Today Focuses on proofs, in the nathernatical sense. Other fields have familiar methods of poof. In science, proof altimately mems experimental verification. In nathematics, it means unassociable logic.

We focus on two knows of proof:

(1) High power proofs: These are proofs that turn big theorems into other results.

det (A-1) = (det (A))-1.

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Suppose that A is an invertible matrix. By proposition 1.54, $det(L) = det(A^{-1} \cdot A) = det(A^{-1}) \cdot det(A)$.

By theorem 1.50, let(I)=1, so we have that

1 = let(A). clet(A).

Hence, we conclude that

let (A") = (let (A))-1.

Proposition 1.54 holds. Proposition 1.54, models in addition, relies on many other results; many, Theorem 1.50. Thus, the above Ba "high power proof."

Example question Using Proposition 1.54, provethet let (AB)=det(BA).

State what assumptions on the stress of A & B in mounts are necessary for this question to make sense.

Class Example

Problem 1.9.12 Prove the determinantal product formula, (1.82); clet (AB) = det(A) det(B).

(a) Let E be an NXNT matrix and B be any NXN matrix. Prove that

det (EB) = det (E) det (B).

PF By Theorem 1.50, we can analyze the cases that Estype 1, 2, or 3.

Suppose E is a type I charactery matrix. The because E is lower triangular with 1's on the main diagonal,

det (E) = 1.

Moreover, multiplying on the left populler MASS by a type 1 matrix corresponds to the appropriate ERO1. This by Theorem 1.50.

Suppose [= is type 2. Then det (E) =-1 by (ii) (noting that

E's obtained from I by one row swap). Moreover, multiply z

B on the left by [= to corresponds to on ERO2, so by another

Copplication of (ii)

det(B) = - det(EB) = det(E) clet(B).

Finally, suppose E is type3; specifically, E=('::::).

Then det(E) = C 1 and be cause multiplying by E is the same as applying on ERO3, so by Theorem 1.50(iii),

det(EB) = c:det(B) = det(E) det(B).

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(b) Prove that if A = E: ... En when E: is an elementary
     motive for 16: EN, then for any B,
      (3) det(AB) = de HA) det(B).
    PF Let A = E, ... EN, where E; one clarentary matrices. Then
    det (AB) = det(E; Ez. Ez. - En. B) = det(E). det(Ez. Ez. .. . En. B)
             = ... = de+ (E,) de+ (B) ... de+ (B).
     Taky B= I, we get
        det (A) = det (G).... det (GN).
        det (AB) = det(E,) · ... det (EN) · det(B) = det(A) det (B).
   Remark The above proof is not quite rigorous. To prove
    a property of this form completely, one must use proof by induction. A complete proof by induction appears below
   Pf (w/ induction) Induction proceeds by proving (3) when No 1,
   and the proofs that if (3) holds for N, then it holds for N+1.
   Base case; N=1 When N=1, Wille equation (3) states
            Let (E, B) = det (E) det 1B) for any B.
    This is proven in part (a).
    N=DN+1; industrested Suppose that (3) holds for N. We
     prove this implies that (3) holds for N+1. By hypothesis,
     let (E, ... En · (Enn · B)) = let (E). ... let (EN) · det(Enn · B).
      By part (a), det (Em. B)=det (Em, ). det (B). This, formy B,
      det (A·B) = det (E). : det (En.,). det (B)
     Taking B = I
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Thus, de+ (A·B) = de+ Lt.)...de+ (Enri) · de+ (B) = de+ (A) de+ (B). I

det (A) = det (G). det (GNH).

(b) Continued Explain why this proves

det (AB) = det (A) det (B) &

for any nonsignal matrix & A and any matrix B.

HE Let A be a nonsequent matrix. Then by definition,

A can be now reduced to an upper triangular matrix with

nonzero entries on the diagonal. It follows that A can

be now reduced to I (using ERO 1, 2, \$3). Thus,

by the correspondence of now operations and multiplication bay

clementary matrices,

A= E:... EN. I = E:... EN.