Lecture 17 More on Metrix Subspaces

Prop 2.41 The following one exercise for an maxim Matrix A:

(i) Ker(A) = {0}

(ii) rank(A) = n (recall rank(A) = dim(rig(A)))

(iii) The system Ax = b has no free variables
(iv) The system Ax = b has a runique solution frull  $b \in my(A)$ 

DF We will prove that (i) =D(ii) =D(iii) =D(iv) =D(i).

Suppose that Ker(A) = {0}. Then dim(Ker(A))=0, so by the Fondamental Theorem of Linear Algebra, rank (A) = N.

Suppose that rank(A)=n. By definition, there is a proof in every column, so the system has no free variables.

Suppose that Ax=b has no free variables. By chapter 1, the solution to Ax = b is unique when it exists.

Suppose that  $A_x = b$  has a runique solution for all  $b \in rg(A)$ . In particular, this is true for b = 0. Be cause Ao = 0,

Ker (A) = {0}.

Prop 2.42 Suppose that Ais on NXN matrix. The following one ey walent:

(i) A & nonsingular

(ii) A is mostible

(iii) rank (A) = n

(iv) Ker (A) = 20]

(v) my (A) = TR"

(vi) det (A) \$0

(vii) Writing Ain column vector form, A=(v,...vn), the vectors Vi...va forma basis of TR".

## he Superposition Principle

The principle of superposition is the observation that

$$A \times (1) = b(1)$$
 and  $A \times (1) = b(1)$ 

then the system

has the solution

Mon generally, of

it

hes the solution

$$X^* = \sum_{i=1}^{K} c_i x^{(i)}.$$

One important application of this is in solvey the system Ax=b

tor many b. If Ais Mxn, and we need to solve the System Ax = b for the same A but more than m vectors b, it is more efficiento solve

$$A \times = e^{(i)}$$
 for  $i=1...m$   $(e^{(i)} = (\frac{i}{3})e^{ith} position)$ .

Then we write each  $b^{(i)}$  as

and we can struct the solution

as the solution to each equation.

To recap, we use the following techniques for efficient solving:

Ax = 6 once: Use Gaussian Elimination

Ax = b more than once but less than mtims: use a permutal LU factorization Ax = b more than M times: Solve Ax = e(i) for all the M basis we ctors e(i) and reconstruct the rest of our solutions.