

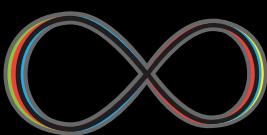
D F I N I T Y



# The Decentralized Cloud

## A Brief Tour of DFINITY

Doc vers. 3rd January 2017



D F I N I T Y

# Mission

# Extend Ethereum Ecosystem With Cloud Network



ethereum



---

Extreme availability  
Casper

Scaling : Speed : Public-Private IOP  
crypto:3



Ungoverned  
The Code is Law

Governed  
Blockchain Nervous System

# **A Boon To Ethereum**

# Virtual Machine Battle

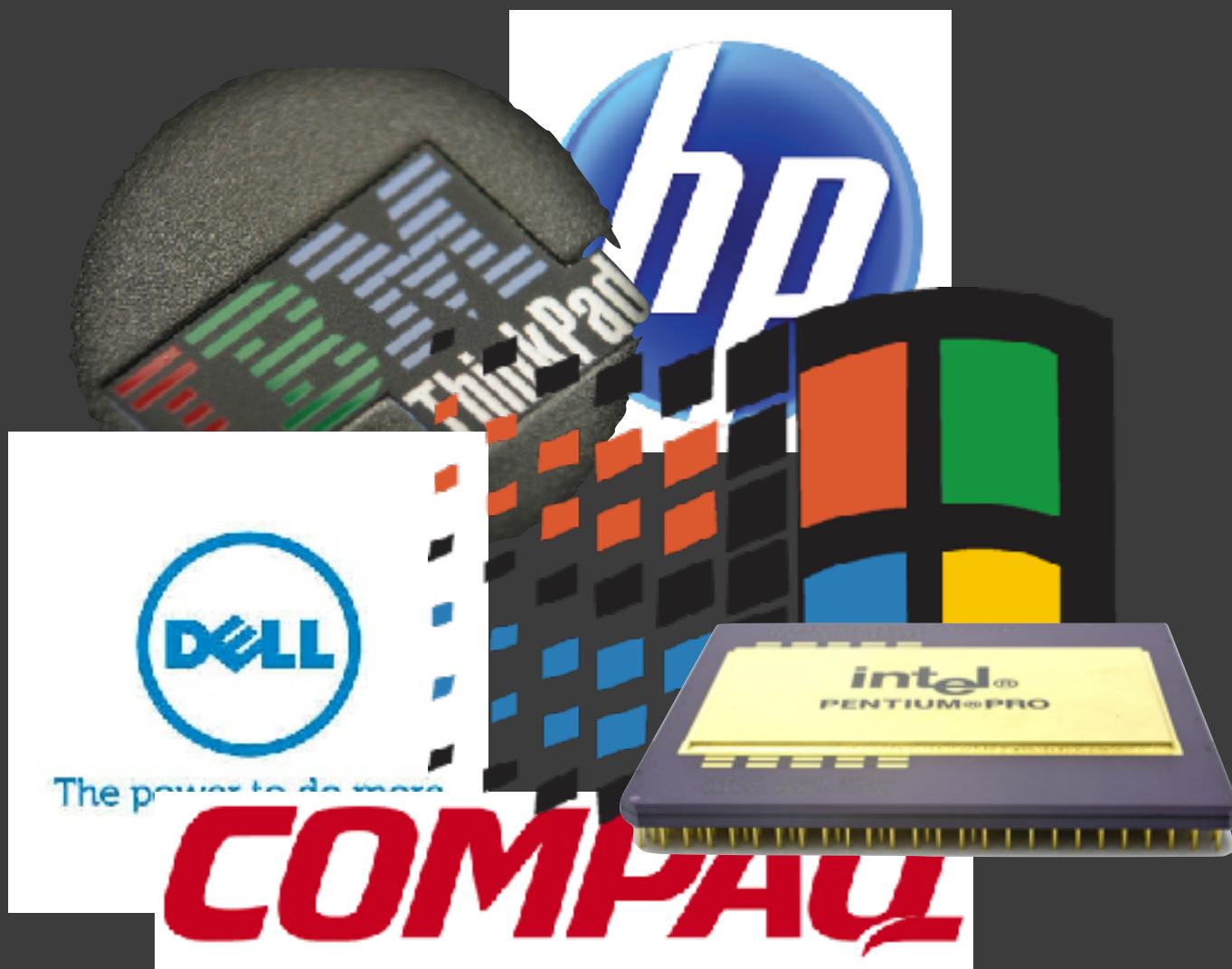
PowerPC



The Macintosh

8086

Microsoft, Dell, IBM, HP... Apple



Sparc

Sun Microsystems



SPARCstation 5

Compatibility More Popular!

# Virtual Machine Battle

Not compatible

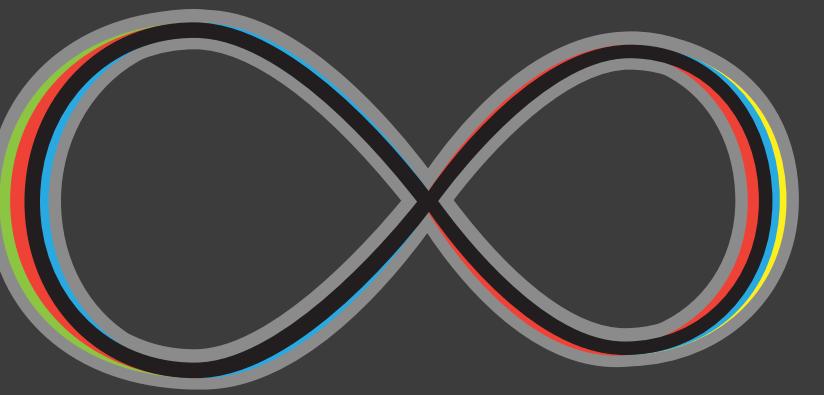


EVM

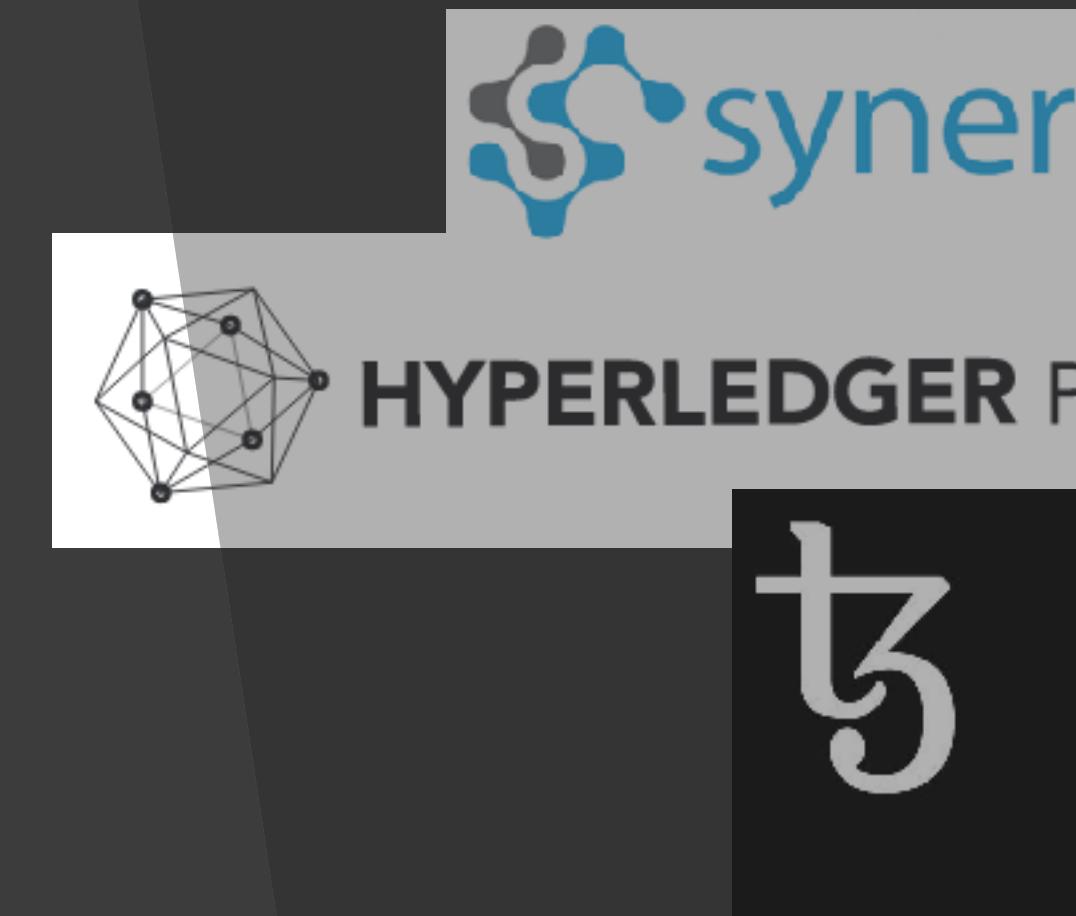
The  
Code is Law



The  
Decentralized Cloud



Not compatible



Compatibility More Popular!

# DFINITY Foundation

- ★ Enable decentralized cloud computing
- ★ Promote compatibility for Dapp developers
- ★ Advance the Ethereum ecosystem
- ★ Fund and advance core Ethereum technologies  
e.g. Ethereum Virtual Machine, Solidity, state channels, etc
- ★ Fund and advance novel DFINITY technologies  
e.g. crypto:3 protocols, Blockchain Nervous System, etc

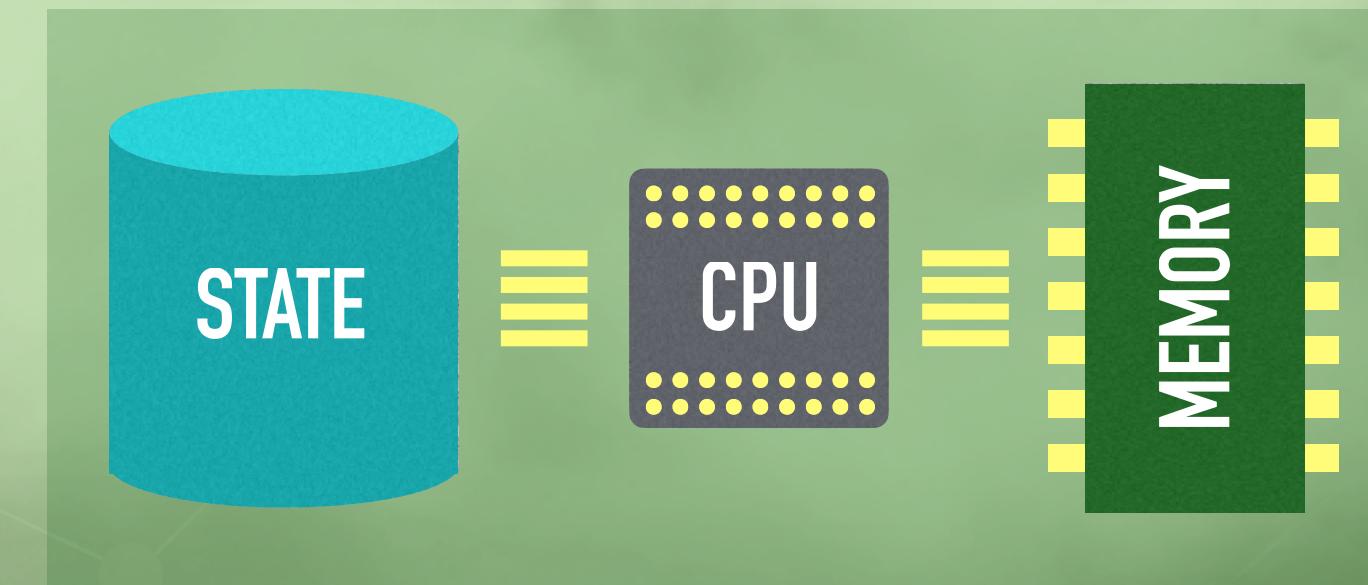
# **DFINITY**

Core Vision

# The Decentralized Cloud

**A Revolutionary Limitless Virtual Computer**

**created by a decentralized network protocol rather than servers**

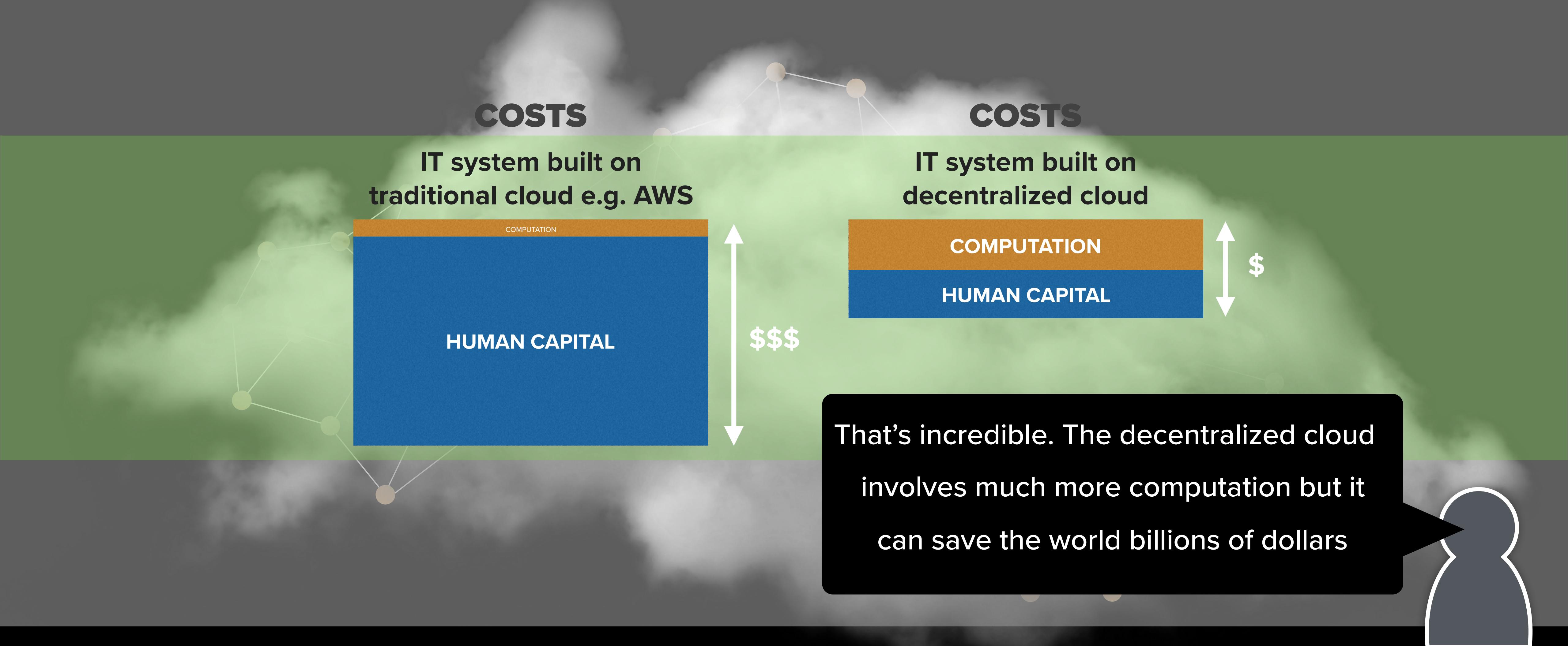


Hmm... doesn't that mean software systems running  
on the virtual machine might be protocols too?



D F I N I T Y

# The Decentralized Cloud Can Crush IT Costs



# IT Systems Should Not Be Rube Goldberg Machines



# Code Business Logic Not Systems

```
1 pragma solidity ^0.4.4;
2
3 import "ConvertLib.sol";
4
5 // This is just a simple example of
6 // It is not standards compatible
7 // coin/token contracts. If you want
8 // token, see: https://github.com/
9
10 contract MetaCoin {
11     mapping (address => uint) balances;
12
13     event Transfer(address _from,
14
15     function MetaCoin() {
16         balances[tx.origin] = 1000;
17     }
18
19     function sendCoin(address receiver, uint amount) returns (bool)
20     {
21         if (balances[msg.sender] > amount)
22             balances[msg.sender] -= amount;
23         balances[receiver] += amount;
24         Transfer(msg.sender, receiver);
25         return true;
26     }
27
28     function getBalanceInEth(address account) returns (uint)
29     {
30         return ConvertLib.convert(balances[account]);
31     }
32
33     function getBalance(address account) returns (uint)
34     {
35         return balances[account];
36     }
37
38     window.onload = function() {
39         web3.eth.getAccounts(function(err, accs) {
40             accounts = accs;
41             account = accounts[0];
42
43             setStatus("Connected to Metacoin");
44         });
45     };
46
47     function setStatus(string message) {
48         var status = document.getElementById("status");
49         status.innerHTML = message;
50     }
51
52     function refreshBalance() {
53         var meta = MetaCoin.deployed();
54
55         meta.getBalance.call(account, {from: account}).then(function(value) {
56             var balance_element = document.getElementById("balance");
57             balance_element.innerHTML = value.valueOf();
58         }).catch(function(e) {
59             console.log(e);
60             setStatus("Error getting balance; see log.");
61         });
62     }
63
64     function sendCoin() {
65         var meta = MetaCoin.deployed();
66
67         var amount = parseInt(document.getElementById("amount").value);
68         var receiver = document.getElementById("receiver").value;
69
70         setStatus("Initiating transaction... (please wait)");
71
72         meta.sendCoin(receiver, amount, {from: account}).then(function() {
73             setStatus("Transaction complete!");
74             refreshBalance();
75         }).catch(function(e) {
76             console.log(e);
77             setStatus("Error sending coin; see log.");
78         });
79     }
80
81     window.onload = function() {
82         web3.eth.getAccounts(function(err, accs) {
83             accounts = accs;
84             account = accounts[0];
85
86             setStatus("Connected to Metacoin");
87         });
88     };
89 }
```

## Abstraction Wins

In the GUI code directly against persistent smart contract objects as though they were libraries. No server. No database. No scaling. Just simple UX and business logic...

# Decentralized Cloud Properties

**Autonomous** | systems w/o dependencies

**Verifiable** | know the code you interact with

**Unstoppable** | no servers to fail

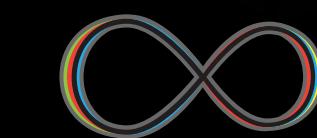
**Cyberspace** | no servers so no geography

**Simple** | code processing not components

**Shareable** | share code and governance

**Tamperproof** | no servers to meddle with

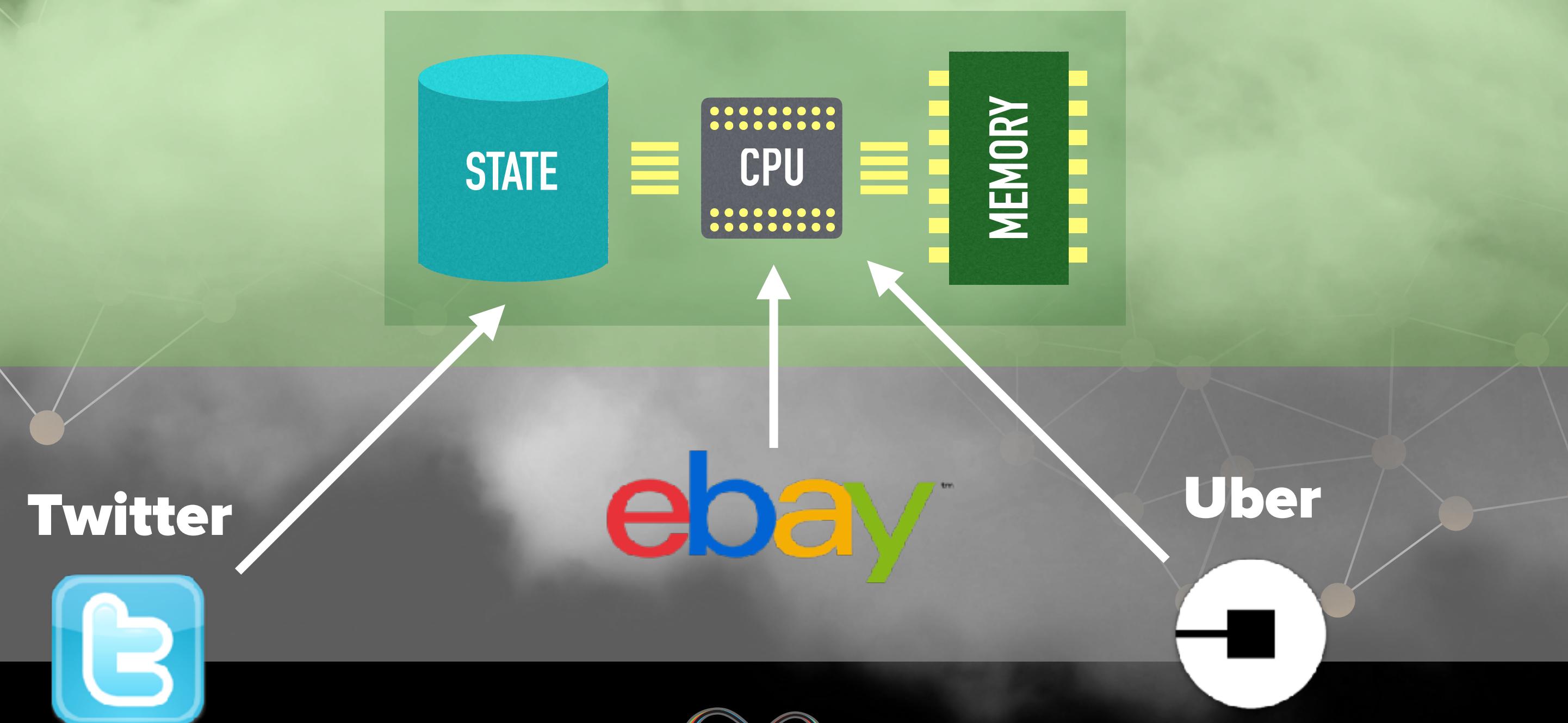
**Interoperability** | no server boundaries



D F I N I T Y

# Decentralized Core Public Services

**“Generic” Business Systems And Their Data Embedded  
as open protocols within fabric of Internet**



D F I N I T Y

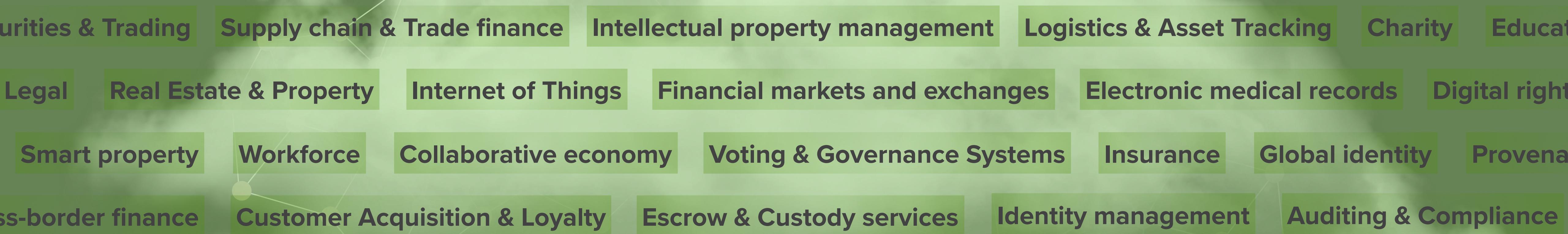
# Decentralized Core Public Services



D F I N I T Y

# Root Level Business Reengineering

**Combine DFINITY with Interoperable Private Cloud Computing Networks**



# Decentralization of Core Functions



## PHI Decentralized Commercial Banking

Similarly to commercial banks, PHI originates new loans by creating new stable cryptocurrency that piggybacks the local currency. Loans granted using a “random validator chain” technique. Can originate loans more judiciously at a fraction of the cost...



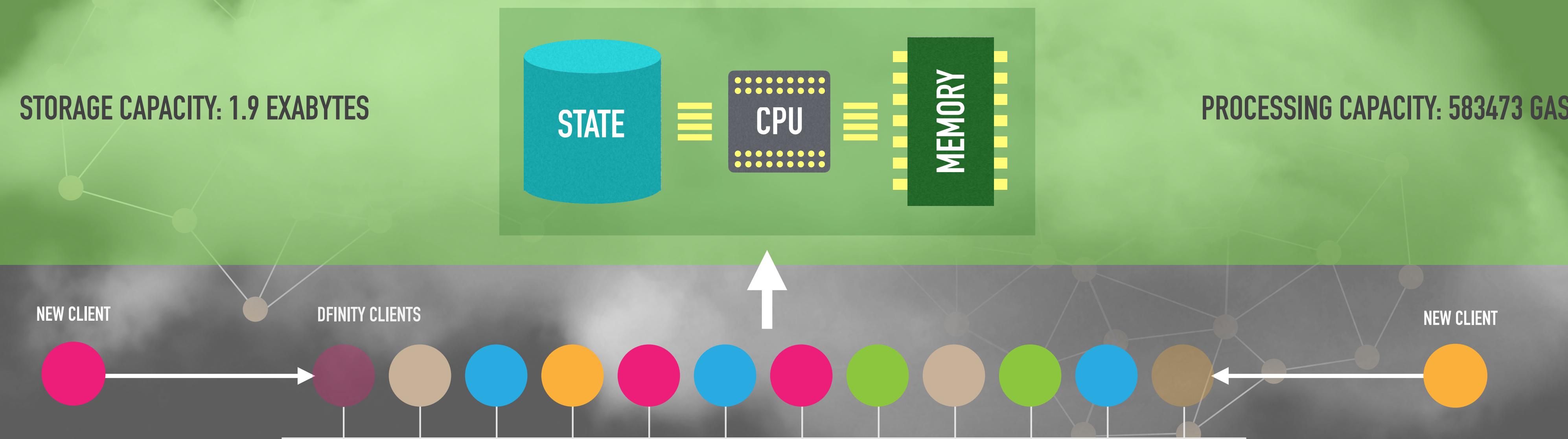
# **Scale-out**

By applying crypto:3 techniques

# crypto:3 Scale-out

A Network of DFINITY Mining Clients Run crypto:3 Protocols

virtual computer formed from protocol interactions

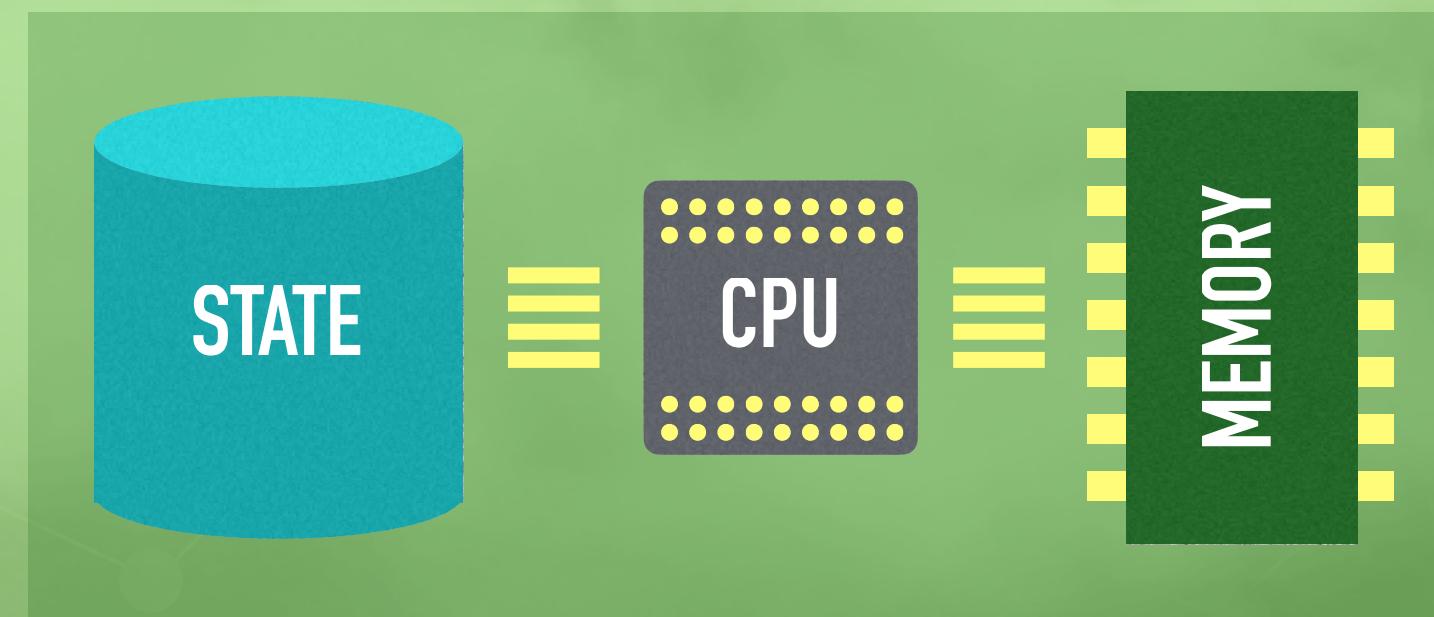


# crypto:3 Scale-out

**As New Mining Clients Join The Network...**

**computational and storage capacity grows indefinitely**

STORAGE CAPACITY: 2.1 EXABYTES



PROCESSING CAPACITY: 731011 GAS

DFINITY CLIENTS



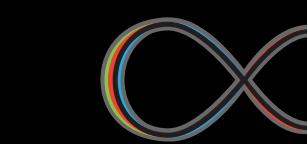
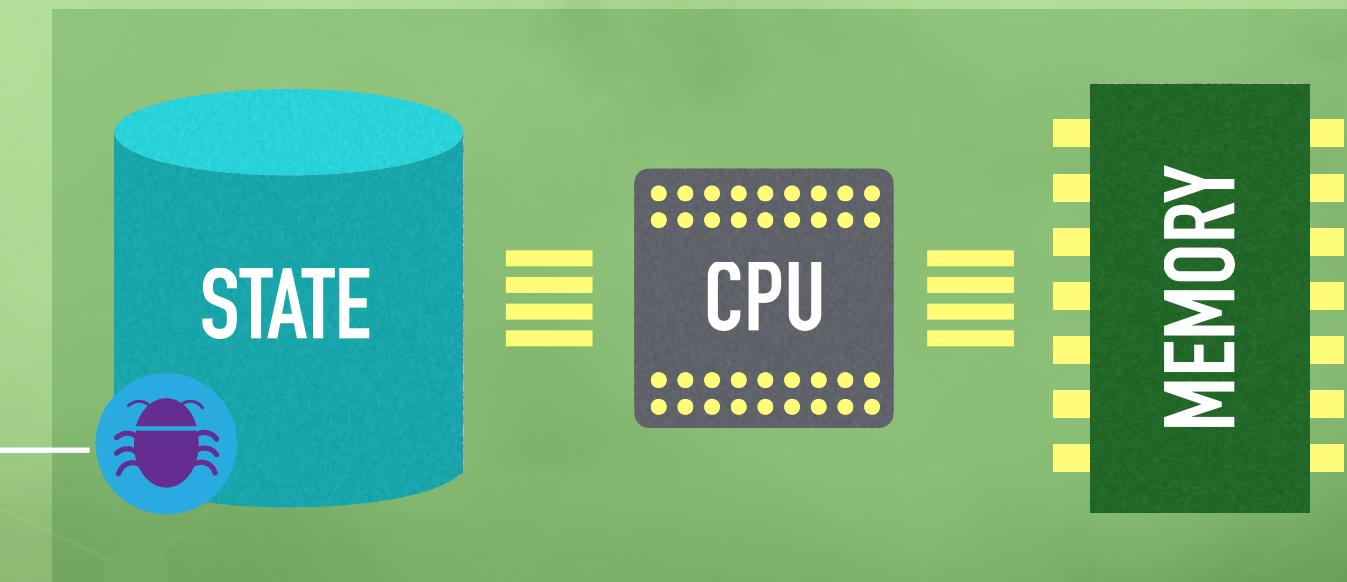
DFINITY

# **Blockchain Nervous System**

Decentralized governance

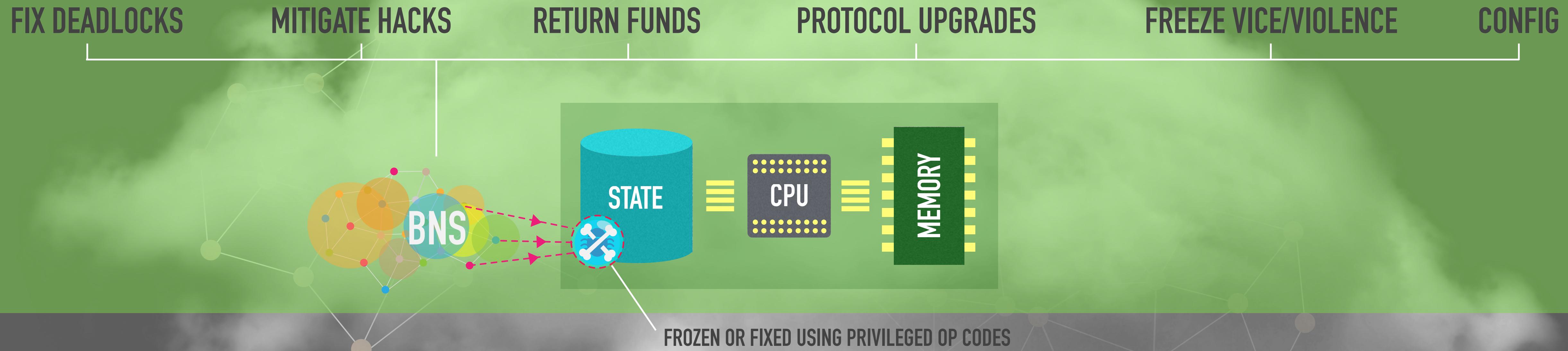
# Blockchain Nervous System AI

**“Code is Law” Now Contingent Upon a Decentralized Intelligence  
this governance system uses privileged machine instructions...**



D F I N I T Y

# Blockchain Nervous System AI



D F I N I T Y

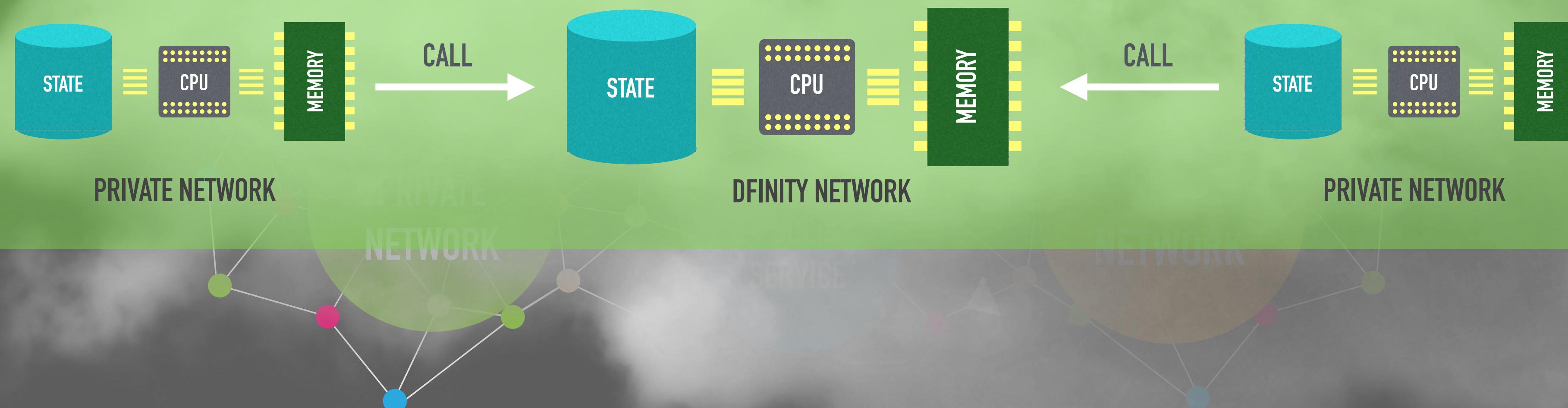
# **Public-Private Cloud System Interoperability**

Made possible by crypto:3 network technology

# Public-Private Interoperability

**Private Network Software System Calls Into Public Software System**

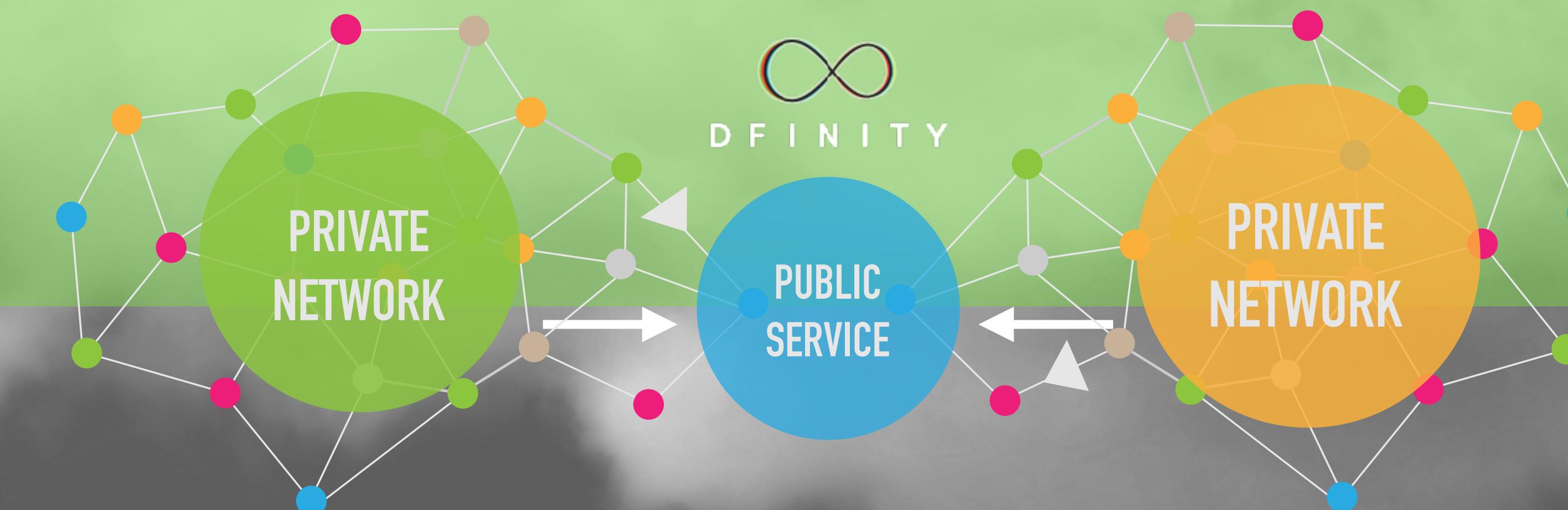
**ZINC RELEASE - one top-level private method can be twinned with a public method**



D F I N I T Y

# Public-Private Interoperability

**Corporates Build Quickly Using Open Business Systems Like Software Libraries**  
e.g. incorporate stable coin, identity, arbitration, interact through an exchange...



## **Review Technicalities...**

- 1 Threshold Relay**
- 2 PSP blockchain (Threshold Relay)**
- 3 Applications of Randomness**
- 4 Blockchain Nervous System**
- 5 Public-Private Interoperability**

1

## **Threshold Relay**

Unmanipulable, unpredictable randomness on  
demand in decentralized networks

This is the crypto:3 technique  
you need to know about first !



## Protocol Design Challenge

**How to organize 1M+ mining clients to  
produce 1 massive virtual computer (DFINITY cloud)?**

**Composed 1M+ servers**



Google Cloud Platform

**Composed 1M+ servers**



**Start with randomness...**

## OBSERVATION

**Randomness is  
the fundamental  
engine used to  
drive stateful  
decentralized  
networks**



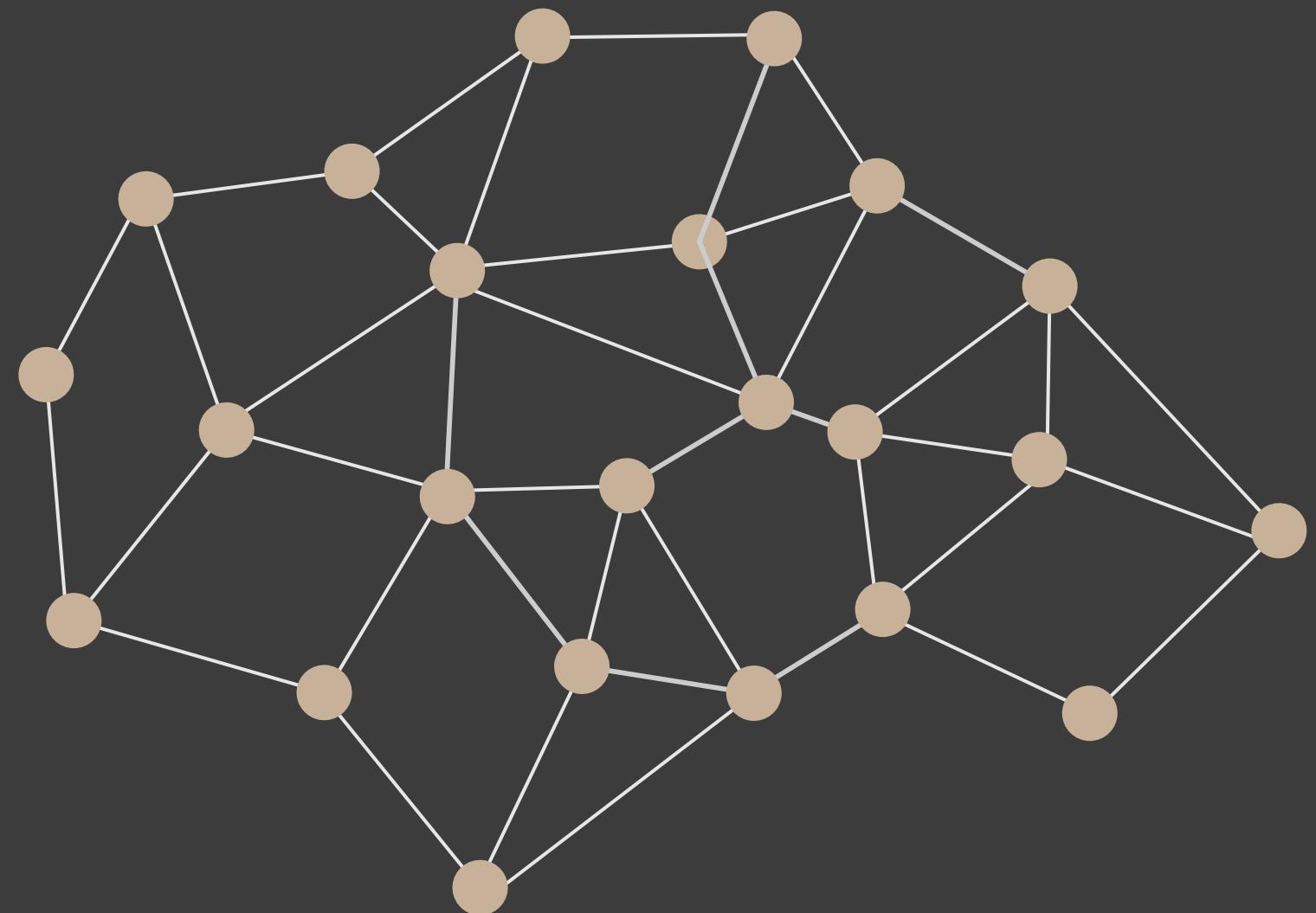
## FOR EXAMPLE...

### Proof of Work



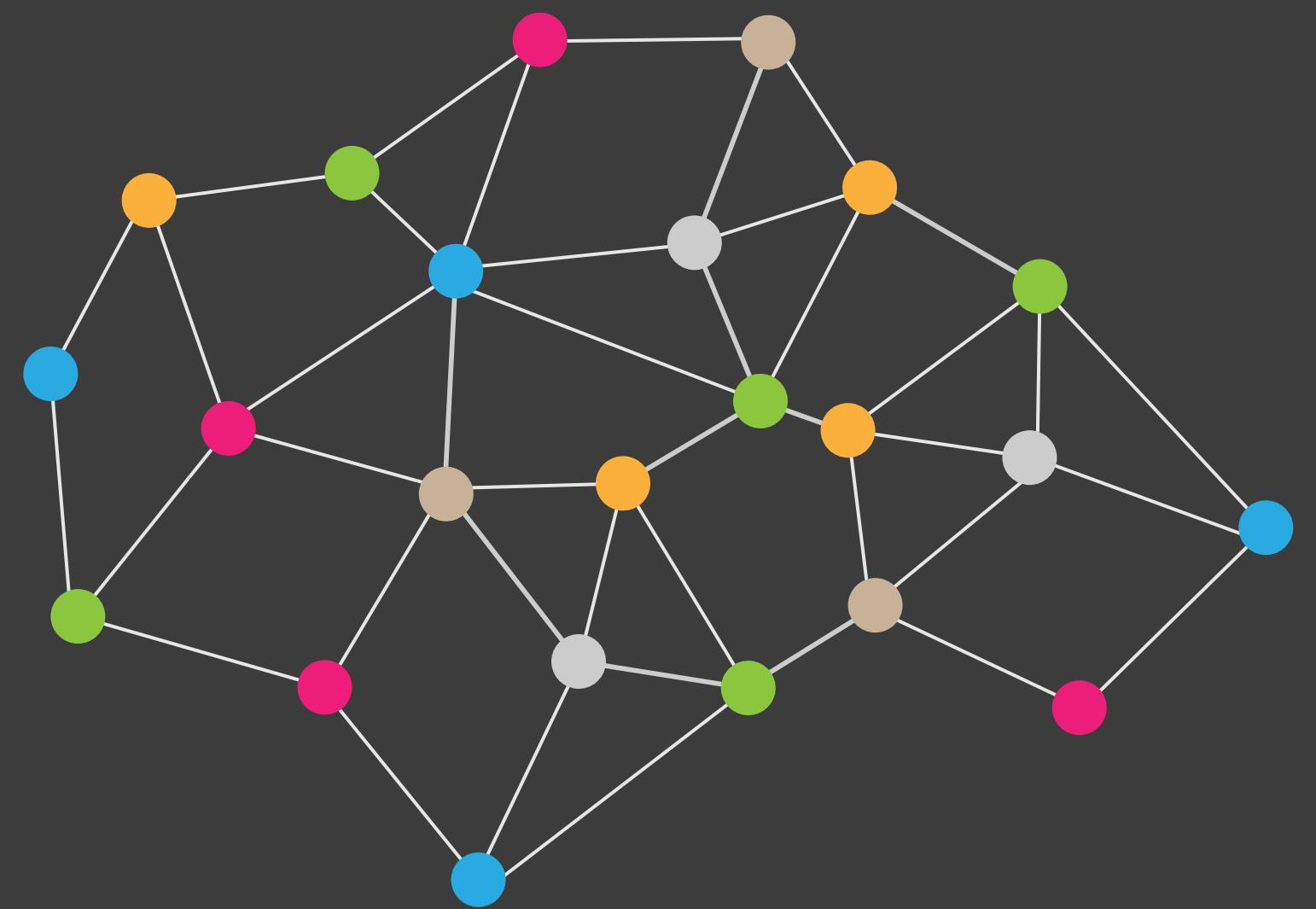
- Miners race to solve a current puzzle
- Solutions found randomly in Poisson distribution
- Winner appends block to blockchain
- In effect was elected a temporary “leader”
- Next leader is unknown
- Cannot DOS, manipulate etc.
- Honest majority
- Chain functions correctly if majority of leaders are honest (~ since selfish mining...)
- Adversary cannot control chain

# A network of processes...



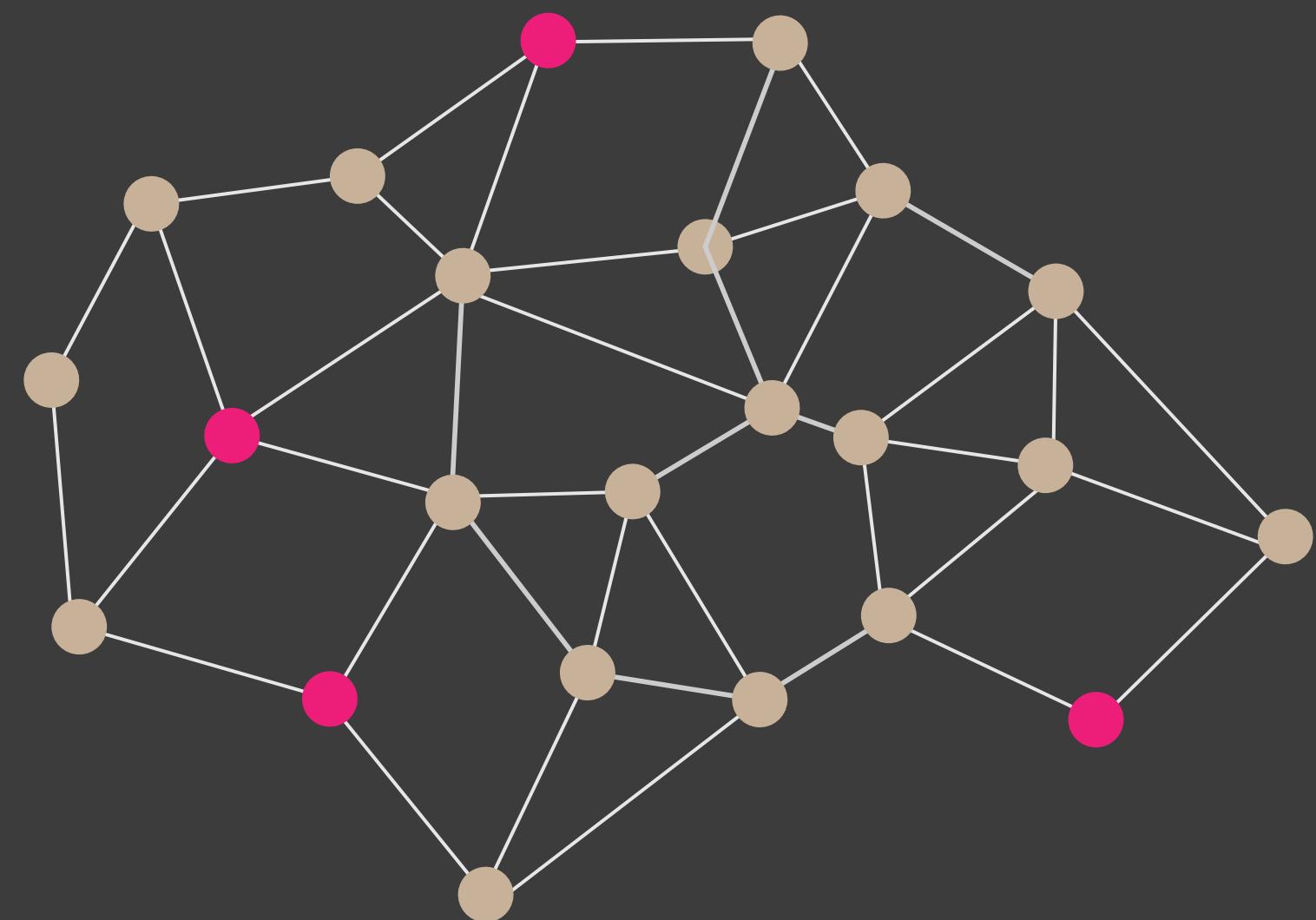
- **Mining “client software” is a process**
  - Fundamental unit of computational resource
- **Connected in P2P broadcast net**
  - E.g. gossip based. Can use Kademlia structure
- **Each process has “mining identity”**
  - Public key with meta data attached
- **IDs mediate participation**
  - Private network: trusted dealer defines list
  - Public network: CC security deposit, USCIDs
- **Massive network size**

# is organized into random groups...



- **Random members**
  - Each process is a member of multiple groups
  - Groups intersect, have e.g. 400 members
- **Groups setup threshold crypto**
  - Run VSS secret sharing protocol
  - 51% threshold e.g. 201+/400 needed create signature
- **BLS signature scheme**
  - **Math magic...** If 51% of group members broadcast “signature shares” on a message, these can be combined to create the group’s threshold signature.  
This will always be the same irrespective of which 51% subset signs (the system is “unique and deterministic”)

# current group signs... signature of previous group



- **Signature *is* random number**
  - Otherwise it would be predictable/insecure
- **Number selects next group**
  - $g = G[r \% |G|]$
- **Next group use prev no. as message**
  - Thus sequence is entirely deterministic
- **Verifiable Random Function**
  - Numbers verifiable using group public key
  - New values produced on threshold agreement
  - Unmanipulable, unpredictable...

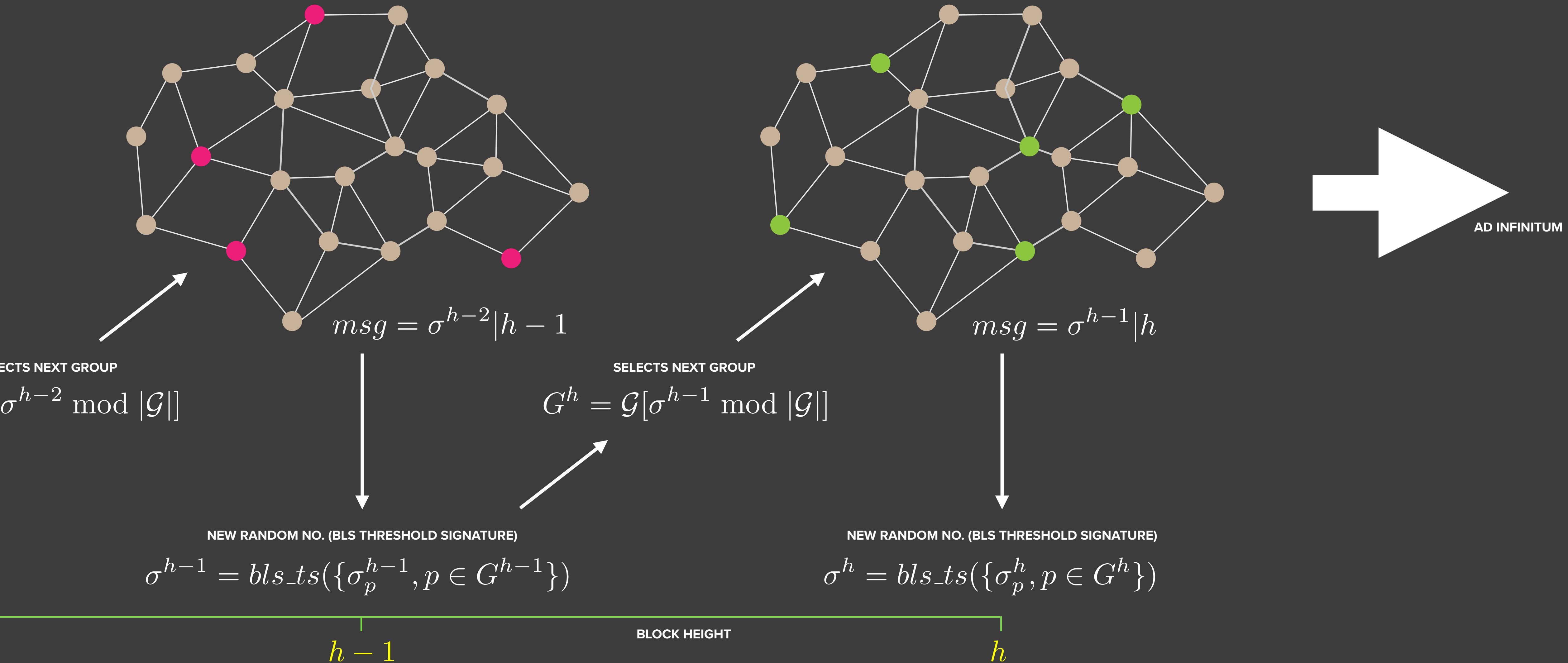
## LEGEND BY EXAMPLE

$\sigma^h$  Threshold signature at height  $h$  ( $h-1$  signature used as message)

$\sigma_p^h$  Signature “share” on  $h-1$  threshold signature by process  $p$

$G^h$  Threshold group that will sign at height  $h$

$|\mathcal{G}|$  The total number of threshold groups in network



# Overhead Example

## MESSAGE FORMAT

Process ID	20 bytes
Signature share	32 bytes
Signature on comms	32 bytes
<b>Total</b>	84 bytes

## GROUP SIZE

Group size	400
Threshold	201

## COMMUNICATION OVERHEAD

<b>Maximum</b>	<i>only 34 KB</i>
----------------	-------------------

**When a group must sign the previous signature, each member process creates a signature share using it as the message. This must be broadcast together with some other information**

**If all group members are active, a total of 34 KB messages will be created each round. In practice, broadcast halts as soon as the 32 byte group signature is broadcast (requires 17 KB of messages)**

# Resilience Example

NETWORK MAKEUP

Processes	10,000
Faulty	3,000
(Correct)	7,000
Group Size	400
Threshold	201

Note: our example assumes almost one third mining processes in the network are faulty. In practice this would be an extreme situation where professional mining is involved.

Probability that 200 or more processes in randomly selected group are faulty, preventing production of signature:

**1e-17**

Calculate odds using any hypergeometric probability calculator  
<http://www.geneprof.org/GeneProf/tools/hypergeometric.jsp>

Nb. groups expire to address “adaptive” adversaries

2

## **Threshold Relay Blockchain**

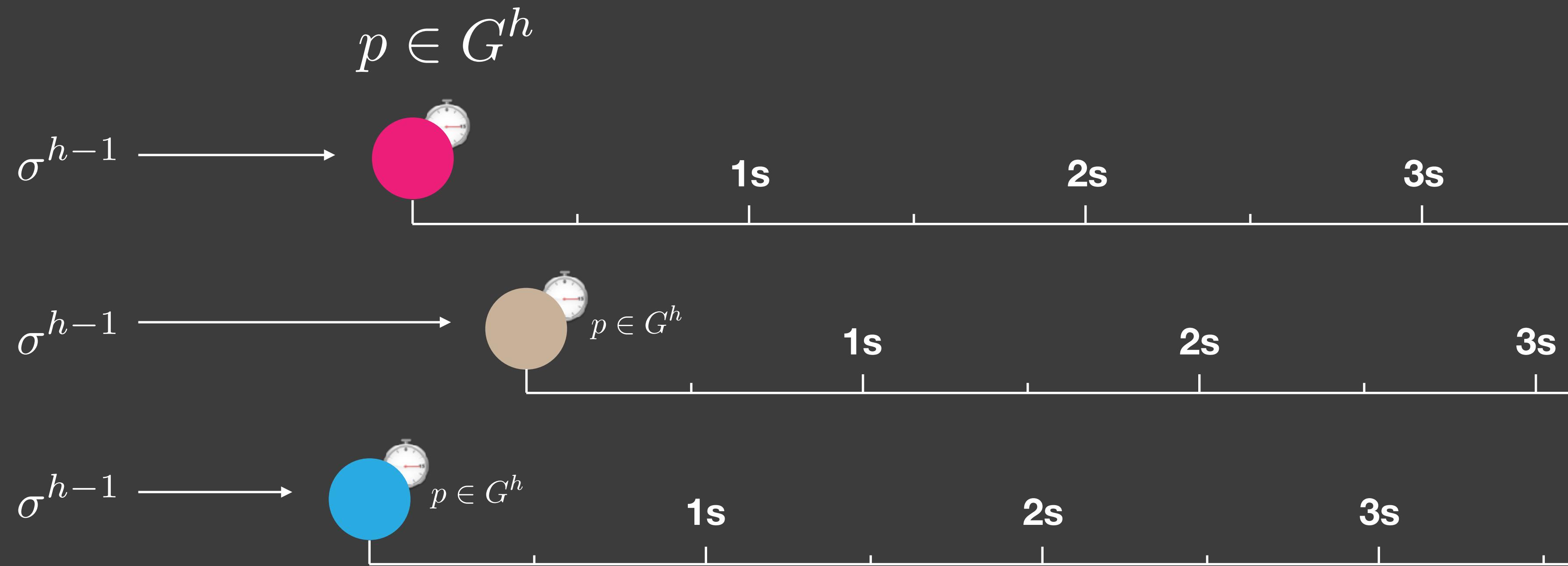
Probabilistic Slot Protocol (PSP)

# Upon Selection Group Members Start Their Timers

and wait for the current Block Time to expire...

## Waiting For The “Block Time”

When a new group is selected by the previous group's threshold signature (the “verifiable random number”) the member's start their stopwatches and wait for the “Block Time” to expire. It's fine that the stopwatch of each member is slightly out of sync.



# Processes Assigned To Priority Forger List

Randomness at  $h-1$  orders processes into new list at  $h$

SLOT	Publish	Points	$h$	→	$h+2$	$h+3$	$h+4$	$h+5$
0	5+ secs	1 pt	P <sub>2313</sub>	→	P <sub>4792</sub>	P <sub>1101</sub>	P <sub>3883</sub>	P <sub>877</sub>
1	6+ secs	1/2 pt	P <sub>3493</sub>	→	P <sub>103</sub>	P <sub>2993</sub>	P <sub>4632</sub>	P <sub>2840</sub>
2	7+ secs	1/4 pt	P <sub>939</sub>	→	P <sub>9291</sub>	P <sub>7520</sub>	P <sub>4486</sub>	P <sub>9811</sub>
3	8+ secs	1/8 pt	P <sub>93843</sub>	→	P <sub>382</sub>	P <sub>4207</sub>	P <sub>7723</sub>	P <sub>9837</sub>

$d = 6s \ t = +1s \ q = 0.5$

Random but shared ordering of all processes in network...

Blocks published by process “2313” recognized after 5 seconds.

Blocks published by process “3493” recognized after 6 seconds.

The nominal “Block Time” here is 5 seconds. Staggering the times at which published forger blocks are recognized is only a network performance optimization and the protocol does not depend upon it for correctness of safety.

# Processes Assigned To Priority Forger List

A new list is created at each block height

SLOT	Publish	Points	h	h+1	h+2	h+3	h+4	h+5	→
0	5+ secs	1 pt	P <sub>2313</sub>	P <sub>4792</sub>	P <sub>1101</sub>	P <sub>3883</sub>	P <sub>877</sub>	P <sub>7615</sub>	
1	6+ secs	$\frac{1}{2}$ pt	P <sub>3493</sub>	P <sub>103</sub>	P <sub>2993</sub>	P <sub>4652</sub>	P <sub>2840</sub>	P <sub>5620</sub>	
2	7+ secs	$\frac{1}{4}$ pt	P <sub>939</sub>	P <sub>9291</sub>	P <sub>9742</sub>	P <sub>8746</sub>	P <sub>7221</sub>	P <sub>66</sub>	
3	8+ secs	$\frac{1}{8}$ pt	P <sub>93843</sub>	P <sub>382</sub>	P <sub>4207</sub>	P <sub>7723</sub>	P <sub>9837</sub>	P <sub>2811</sub>	↓

**d = 6s t = +1s q = 0.5**

# Chain Choice Rules

Correct processes try to build on the highest scoring chain

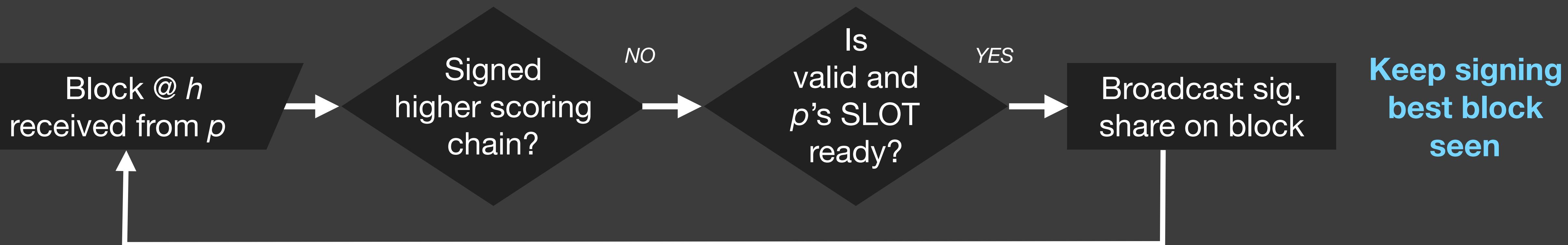
SLOT	Publish	Points	h	h+1	h+2	h+3	h+4	h+5	→
0	5+ secs	1 pt	P <sub>2313</sub>	P <sub>4792</sub>	P <sub>1101</sub>	P <sub>3883</sub>	P <sub>877</sub>	P <sub>7615</sub>	<b>BEST PARENT + 4 points</b>
1	6+ secs	$\frac{1}{2}$ pt	P <sub>3493</sub>	P <sub>103</sub>	P <sub>2993</sub>	P <sub>4652</sub>	P <sub>2840</sub>	P <sub>5620</sub>	+ $2\frac{3}{4}$ points
2	7+ secs	$\frac{1}{4}$ pt	P <sub>939</sub>	P <sub>9291</sub>	P <sub>9742</sub>	P <sub>8746</sub>	P <sub>7221</sub>	P <sub>66</sub>	
3	8+ secs	$\frac{1}{8}$ pt	P <sub>93843</sub>	P <sub>382</sub>	P <sub>4207</sub>	P <sub>7723</sub>	P <sub>9837</sub>	P <sub>2811</sub>	↓

**d = 6s t = +1s q = 0.5**

# Threshold Block Notarization

Group members sign until  $\geq 1$  block gets threshold signature

The signing behavior of a group member at block height  $h$



Thresh. sig. on block  
at  $h$  received

Broadcast sig.  
share on  $\sigma^{h-1}$

STOP

Threshold  
relay  
and halt

# Threshold Block Notarization

Of course, generally only 1 block gets signed....

**Before the Block Time  
expires, received blocks  
placed in priority queue**

Group members place received blocks into a priority queue while waiting for the Block Time to expire. If a block from the process in SLOT 0 is received, it will be placed at the head of the queue.

**Group members only sign  
blocks if not seen higher  
priority block...**

After expiry of the Block Time, a member will usually have a block from SLOT 0 waiting in the priority queue. This is signed first, so no others will be signed\*. Generally, we expect groups to sign only 1 block.

**When the chain score  
drops, the protocol  
lengthens the Block Time**

Imagine if SLOT 0 was produced in Europe and SLOT 1 in America. Both might get signed if Block Time too low. If score best chain drops, the protocol lengthens the Block Time to accommodate network speed.

\* Notwithstanding equivocation. Always sign equal highest priority prevent deadlock

# Notarization Drives Rapid Convergence (Consistency)

Only blocks with group signatures are valid...

SLOT	Publish	Points	h	h+1	h+2	h+3	h+4	h+5	→
0	5+ secs	1 pt	P <sub>2313</sub> σ	P <sub>4792</sub> σ	P <sub>1101</sub> σ	FAULTY P <sub>3883</sub>	P <sub>877</sub> σ	P <sub>7615</sub> σ	
1	6+ secs	1/2 pt	P <sub>3493</sub>	P <sub>103</sub> σ	P <sub>2993</sub>	P <sub>4652</sub> σ	P <sub>2840</sub>	P <sub>5620</sub>	
2	7+ secs	1/4 pt	P <sub>939</sub>	P <sub>9291</sub>	P <sub>9742</sub>	P <sub>8746</sub>	P <sub>7221</sub>	P <sub>66</sub>	
3	8+ secs	1/8 pt	P <sub>93843</sub>	P <sub>382</sub>	P <sub>4207</sub>	P <sub>7723</sub>	P <sub>9837</sub>	P <sub>2811</sub>	↓

**d = 6s t = +1s q = 0.5**

# Timestamping Benefits

Threshold power!

**SUPER FAST - during normal operation expect overwhelming probability of transaction irreversibility (“finality”) in:**

**7.5s**



**Threshold groups notarizing blocks resolves several security challenges...**

- Nothing At Stake
- Equivocation
- Selfish Mining

**Quantifiable risk**  
Presence multiple notarizations, chain score, allow est. statistical risk

**SPV**  
Prove to a light client that only has Merkle root of groups

# Waiting for Casper... PoW On Ethereum

- Currently 50+% of blocks mined are empty...
- Proof-of-work's "Poission distribution" is cause
- Sooner you publish, greater chance of being "confirmed"
- Publishing an empty block without delaying to validate transactions is more profitable...
- Building on an empty block that does not involve validation delay is more profitable...
- To combat problem, per block gas limit set to tiny levels



Bitcoin Could Consume as Much Electricity as Denmark by 2020,

Motherboard 3/29/2016  
Ridiculous!!!

# Relative Performance

Threshold power!



## Block time

Average 10 mins  
*varies wildly*

Average 20 secs  
*varies wildly*

Average 5 secs  
*low variance*

## “TX finality” (speed)

6 confirmations  
avg. 1 hr

25 confirmations  
avg. 6 mins

2 confirmations  
**avg. 7.5 secs**

## Gas available

- - -

Low due to  
Poisson distribution

**25-50X Ethereum**

*( Unlimited scale-out achieved  
by applying randomness in  
following techniques... )*

3

## **Applications of Randomness**

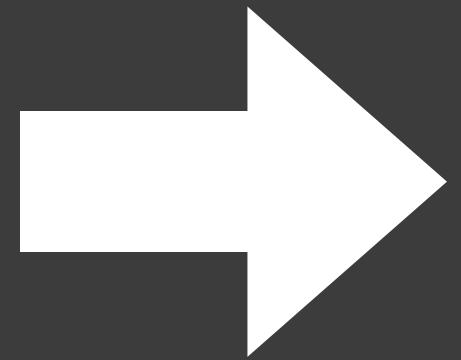
# **Scale-out virtual computer**

Handling any volume of computation and state

# Separate and decouple concerns

## Proof-of-Work Blockchain

*Sybil resistance*  
*Validation*  
*State storage*  
*Consensus*



## DFINITY

Consensus  
———  
Validation  
———  
State storage

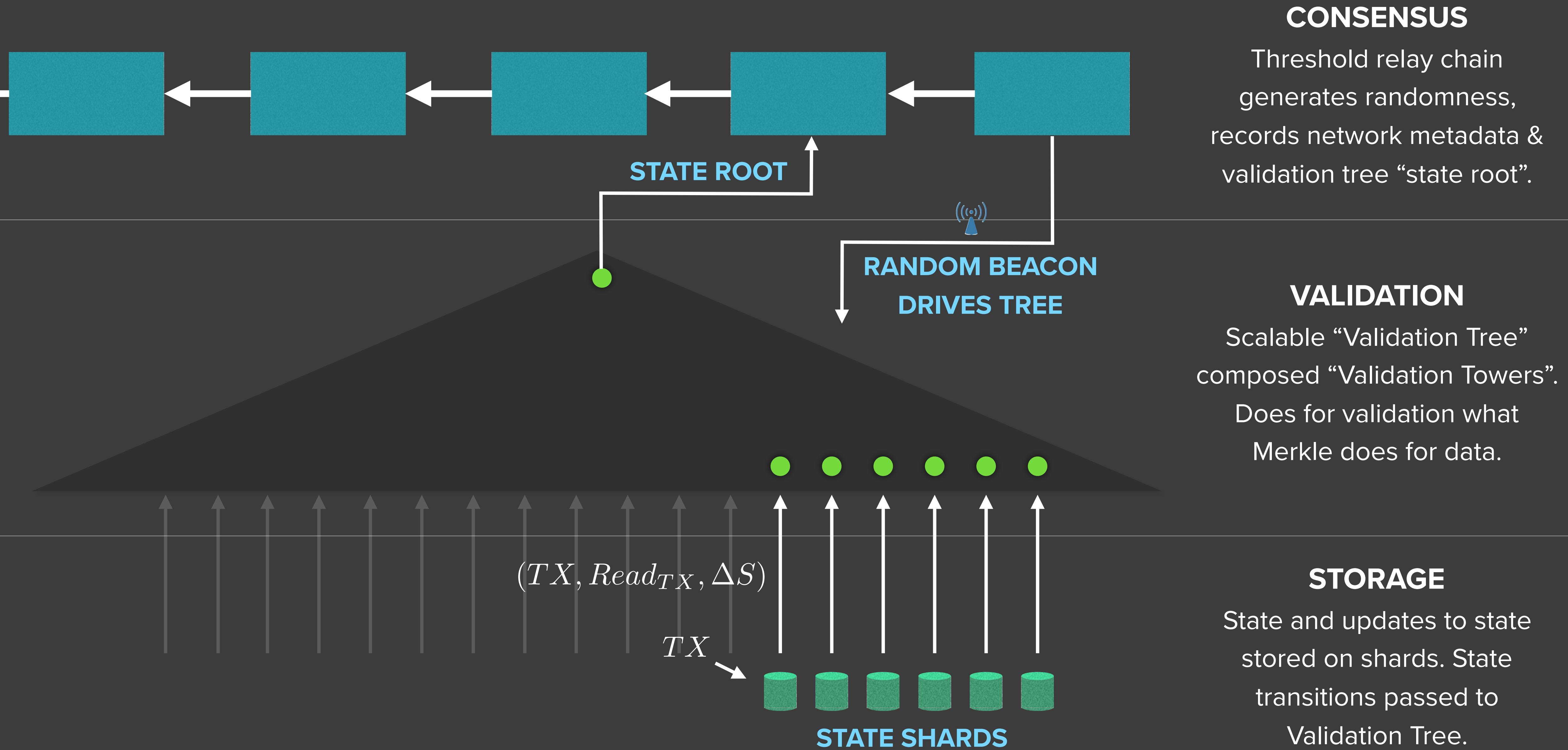
Sybil  
resistance

TCP/IP

Application  
———  
Transport  
———  
Internet  
———  
Network Access

Good old-fashioned Computer Science!

# Three Layer “Scale-out” Architecture



# Example Applications Of Randomness

1

**SCALING**

## VALIDATION TOWER

***Validate network state changes using few processes***

Each additional level of the tower validates new state transitions applied to some storage shard, and is built by processes selected by the random beacon. These processes must also validate levels beneath them to some validation depth  $d$ . Once a level has been added, a process becomes economically inactive until it has been buried by  $d$  further levels. It cannot predict who will build these levels, and thus it becomes computationally infeasible to collude with shards and have bad transitions validated.

2

**USER EXPERIENCE**

## LAZY VALIDATION

***Create instant Web search results***

Some queries of data on the virtual computer are relatively inexpensive and do not warrant the cost of a validation in a Validation Tower. For example, a search of a Web index is a low value operation. Nonetheless, validation is still necessary, since otherwise miners might insert advertising into search results. To address this situation, query results are validated 1. only occasionally upon direction of the random beacon and 2. after the query has been returned to the user.

3

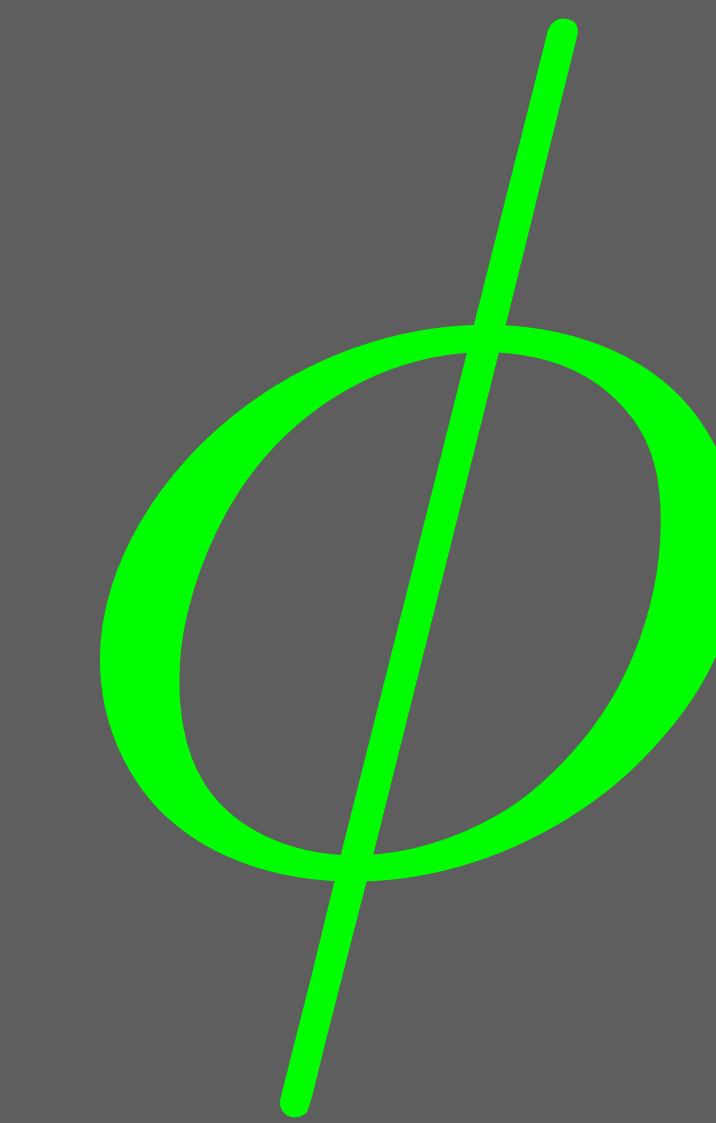
**COST REDUCTION**

## LOTTERY CHARGING

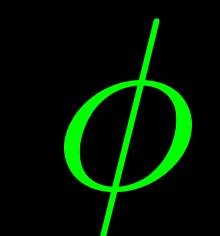
***Create inexpensive Web search results***

While we can use lazy validation to reduce the cost of low value operations such as search queries, any operation that involves currency (in this case “dfinities”) must necessarily be fully validated - after all, lots of small frauds make a big fraud!!! A problem thus exists, because the computer cannot make any operation free or it will be DOSed. The solution is to multiply charges by e.g. 1000, and then only apply them 1/1000 times as directed by the random beacon.

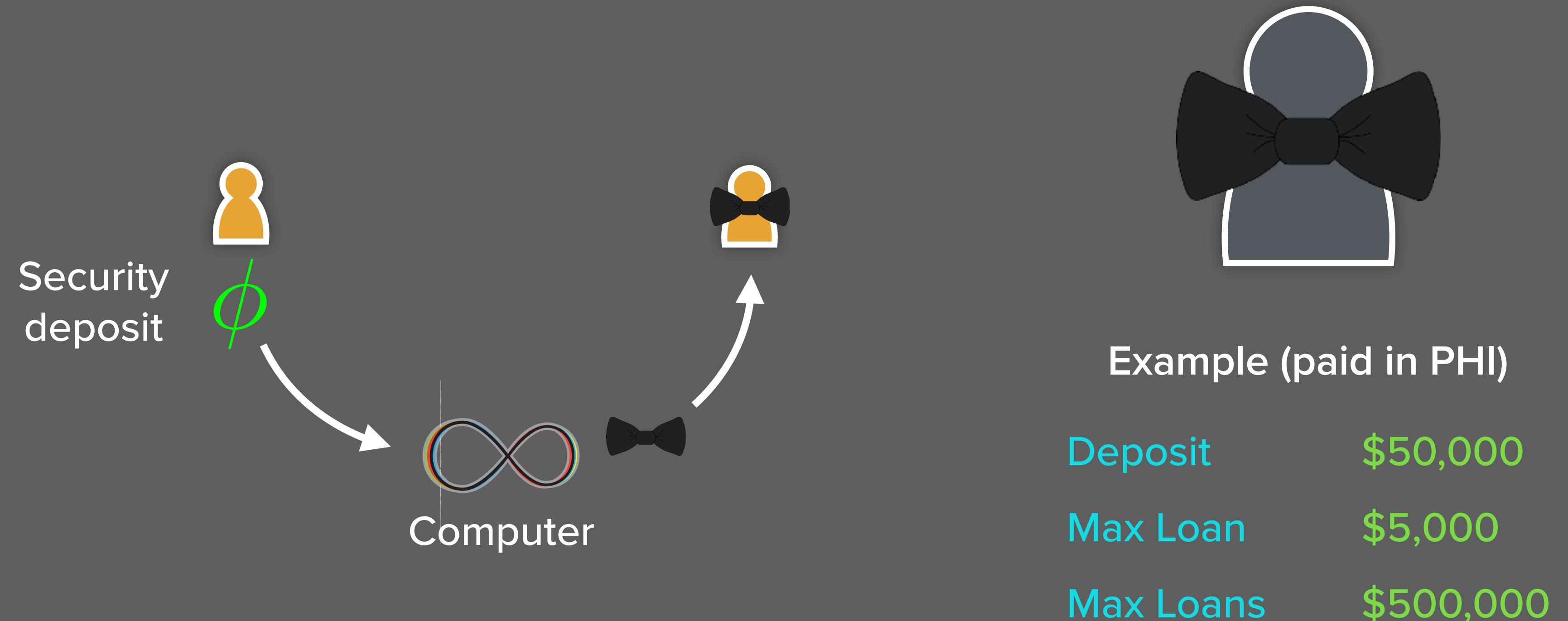
#### 4. Extract from DEVCN2 PHI presentation demonstrating use of randomness to originate loans algorithmically



Decentralized commercial banking and stable currency \_



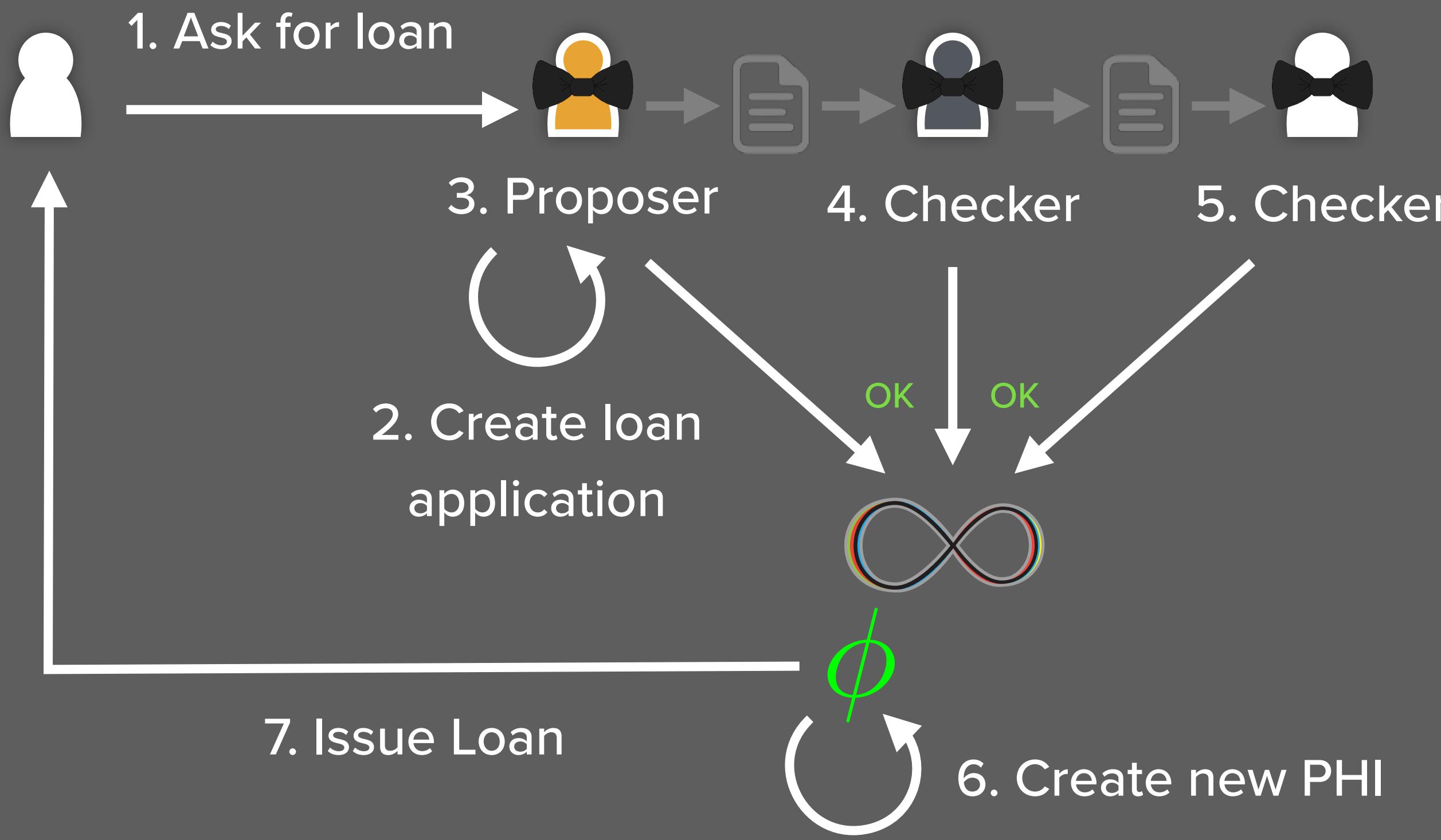
# Anyone Can Be A PHI Validator



Anyone can become a PHI Validator by making a security deposit to the computer. If a loan you approve becomes delinquent the computer takes compensation from your deposit.

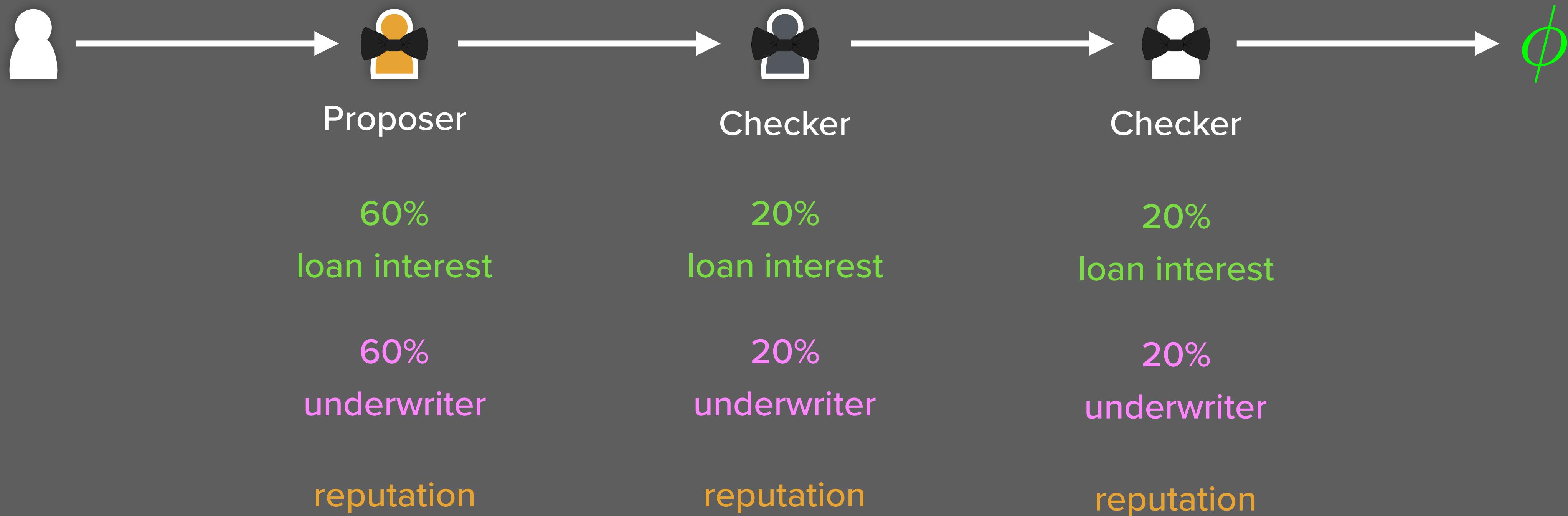
$\phi$

# How Computer Issues Loans



- ▶ Random sequence validators
    - Nobody knows who's next
    - Nobody knows length
  - ▶ Choice validators
    - Size of their deposit
    - Reputation
  - ▶ Loan application
    - Format is open standard

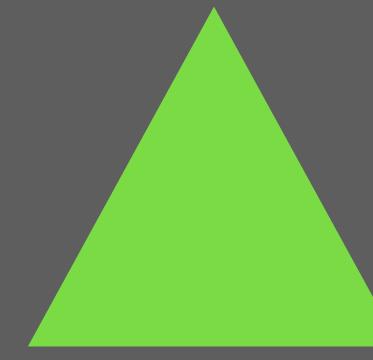
## Validator Incentives



EXAMPLE

$\phi$

# Validator Reputation

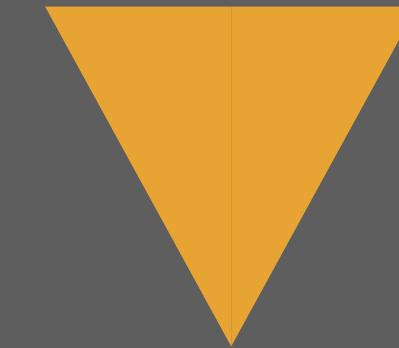


Loans  
performing

Deposit size  
maintained

Incompetence &  
game playing are  
losing strategies

The lower your  
reputation falls the  
longer the validation  
chains and the  
harder it is to make  
returns from security  
deposit, which can  
get “frozen”

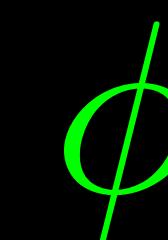


Loans have  
delinquent payments

Deposit size  
decreasing

Other validators  
reject loans

*If validation decision rejected by others, can appeal and computer decides using new random sequence*

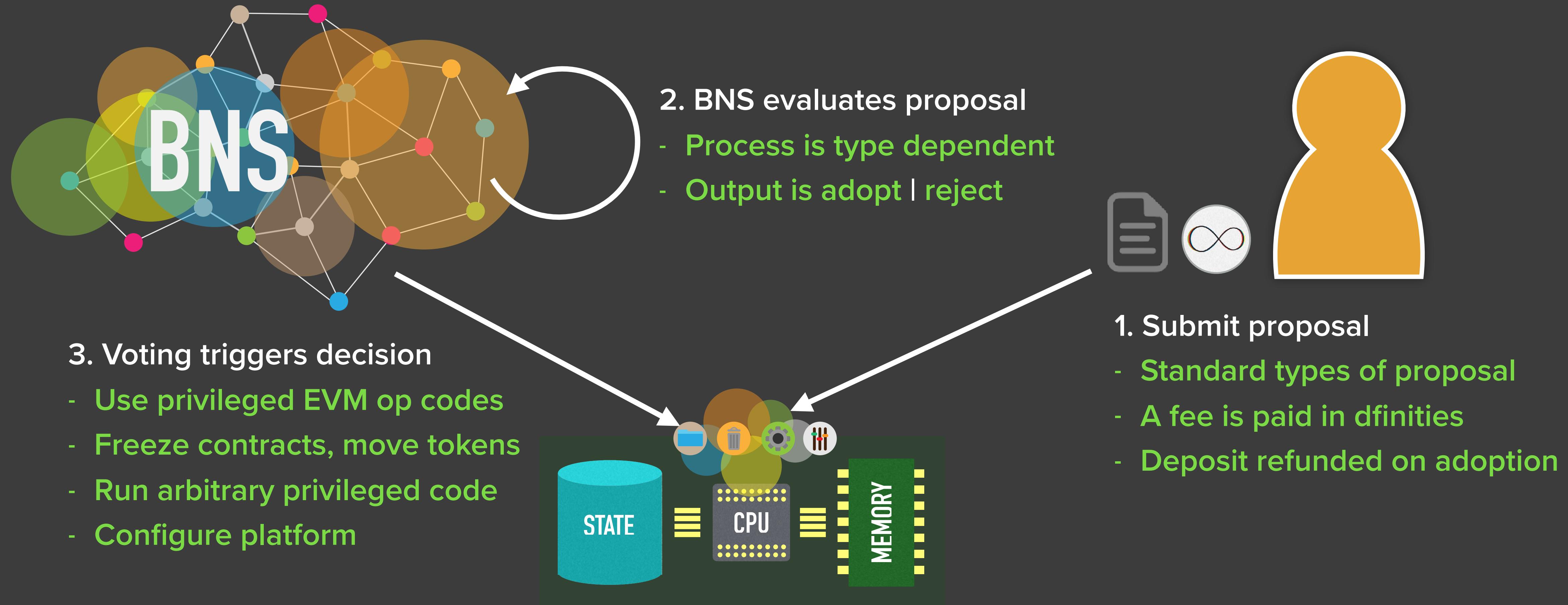


4

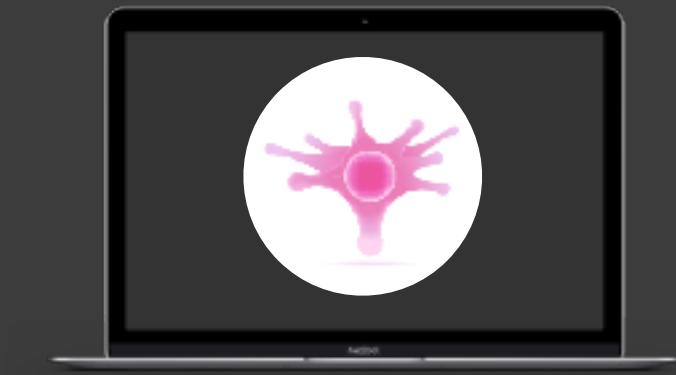
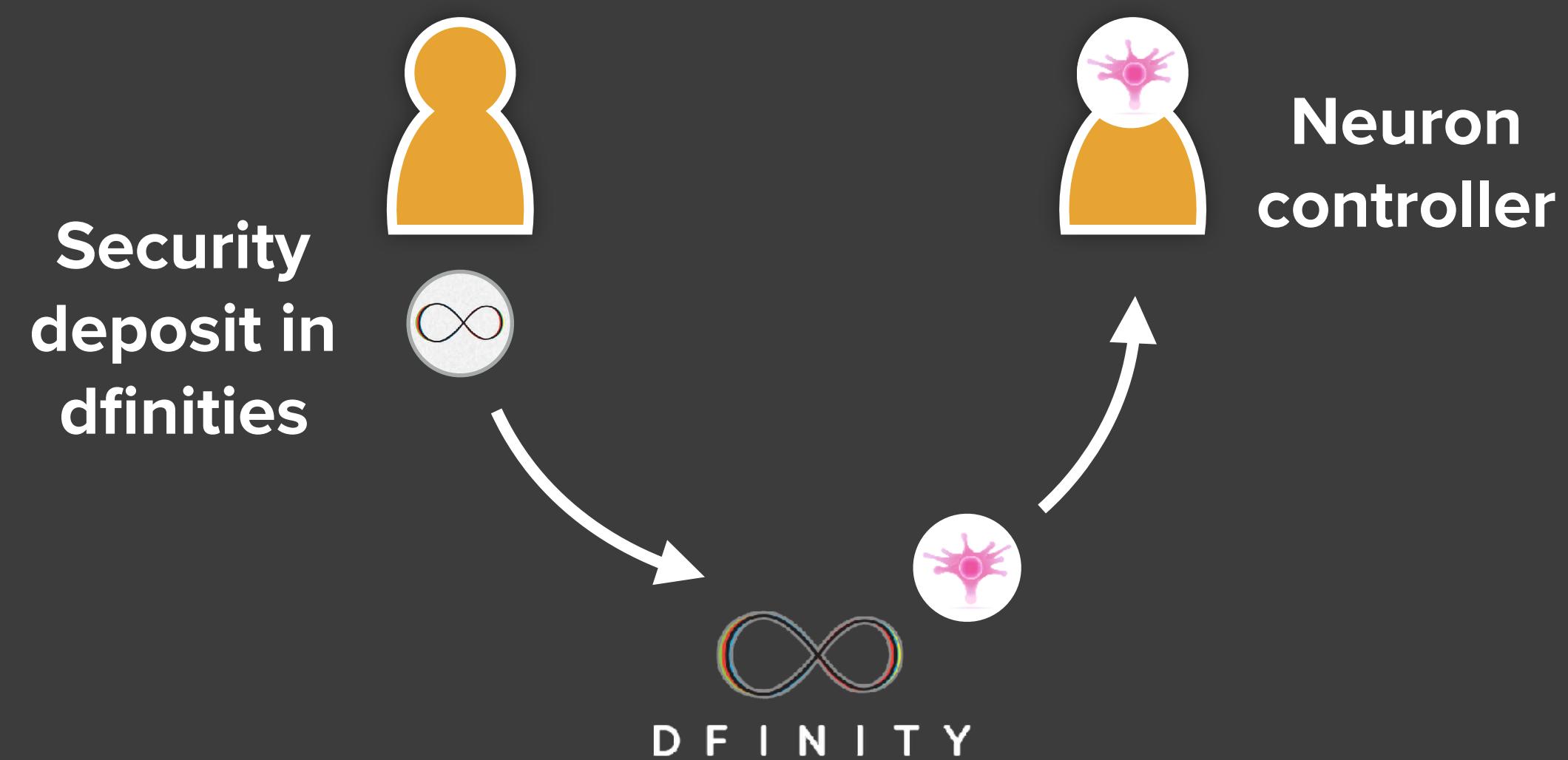
## **Blockchain Nervous System**

Decentralized governance

# Proposal Processing



# People Create Neurons

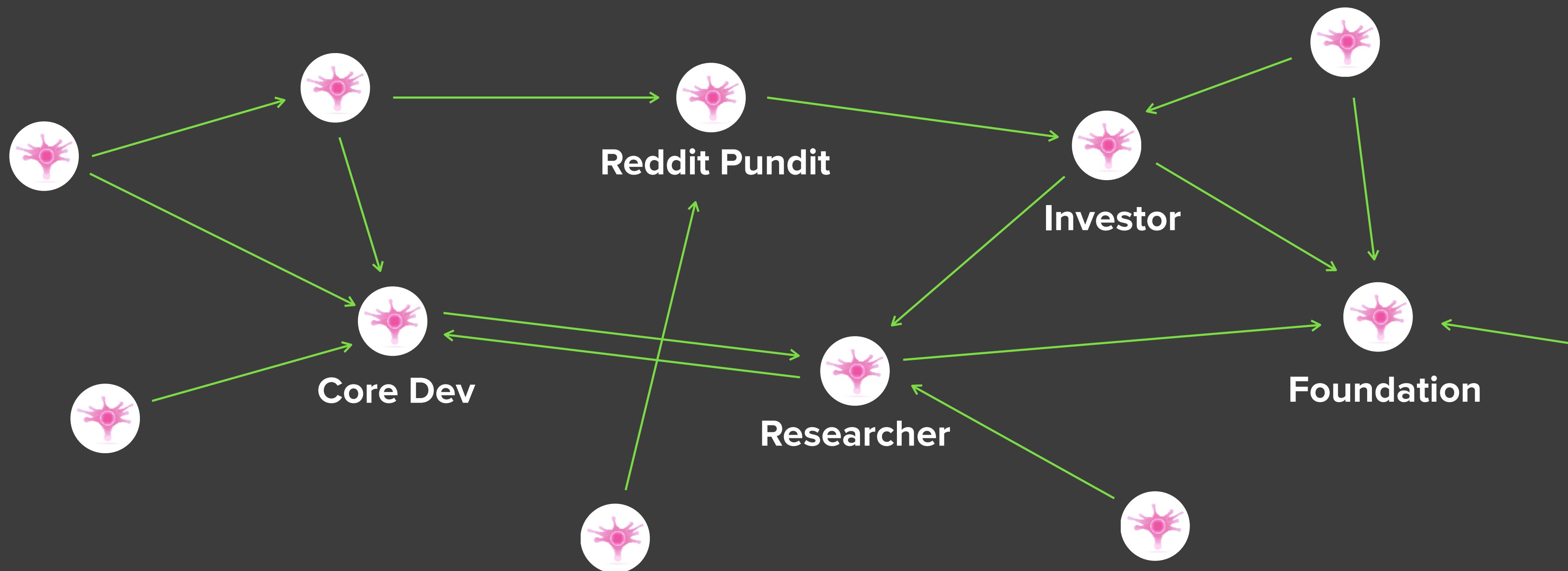


Neuron key pair  
configured into laptop  
or smartphone client

Neuron's voting power  
equals dfinities  
deposited

Neurons are created by making a security deposit to a special DFINITY smart contract.  
It takes 3 months to dissolve a neuron and retrieve the deposit, incentivizing good decision making

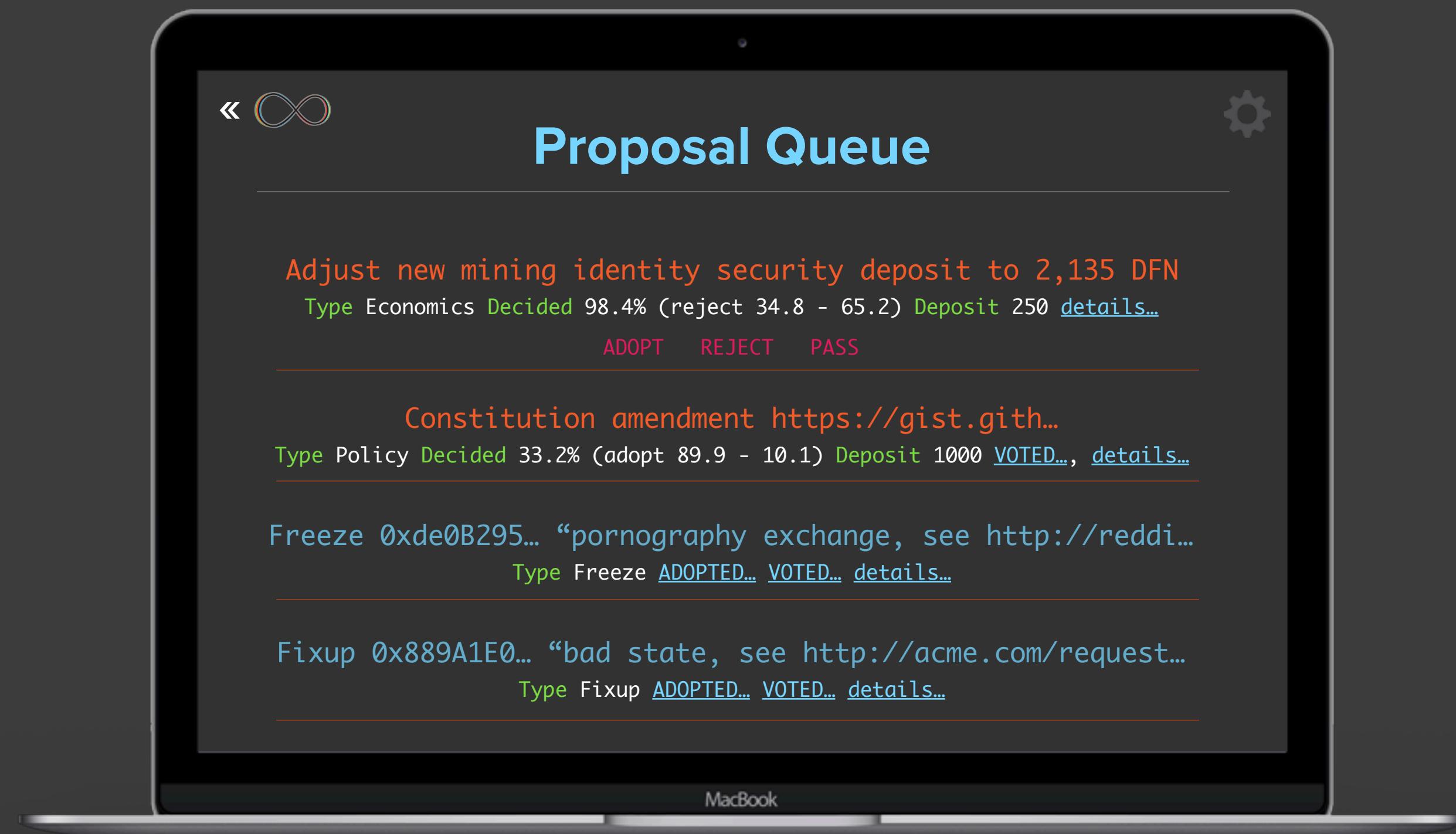
# Neurons Follow Neurons...



Neurons follow other neurons. This enables them to make decisions without benefiting from the input of their direct human controller. People can advertise the address of their neuron



# Neuron Client Software

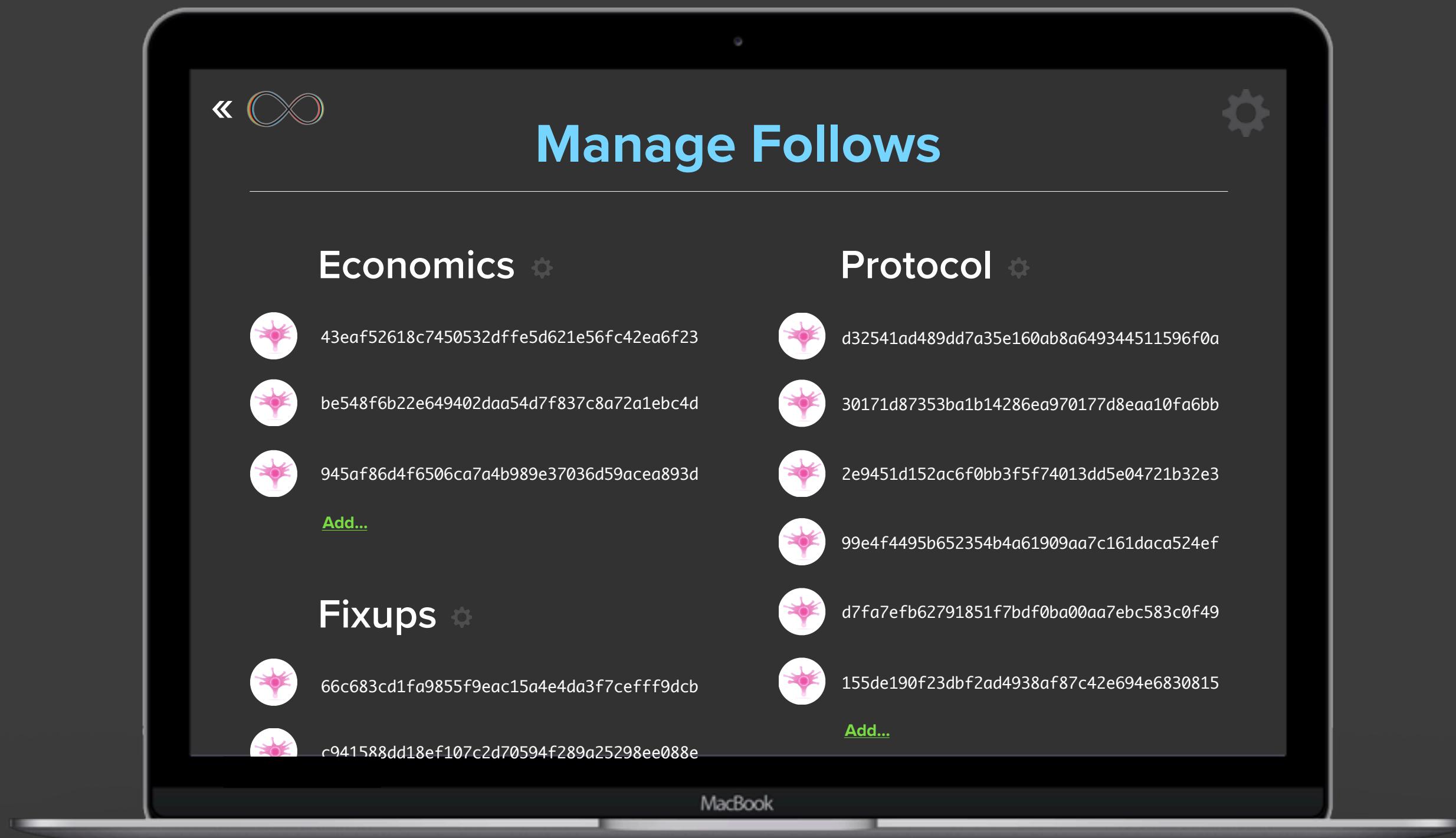


**Proposals are submitted with a non-refundable fee and deposit that is refunded if the proposal is adopted. Before a timeout, the neuron controller (owner) can manually vote on proposals...**

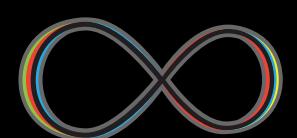


DFINITY

# Neuron Client Software



**When the neuron's controller (owner) doesn't vote on a proposal before the timeout, the neuron client software examines the follow list configured for the proposal's type...**



D F I N I T Y

# BNS Properties

## Decentralized intelligence

### Captures wisdom of crowds

- Community expertise incorporated
- Opaque liquid democracy

### Neurons cascade to make decisions on proposals

- Non-deterministic  
(depends timing)

### Highly resilient

- Trust graph (follow relationships) exist on edges and are unknowable...
- Difficult to kidnap, extort or influence “key holders”
- Government, or “frozen” cos cannot capture, sue etc...
- Good incentives system

### The BNS learns

- Follow relationships are dynamic
- Improves over time

### Thought mining

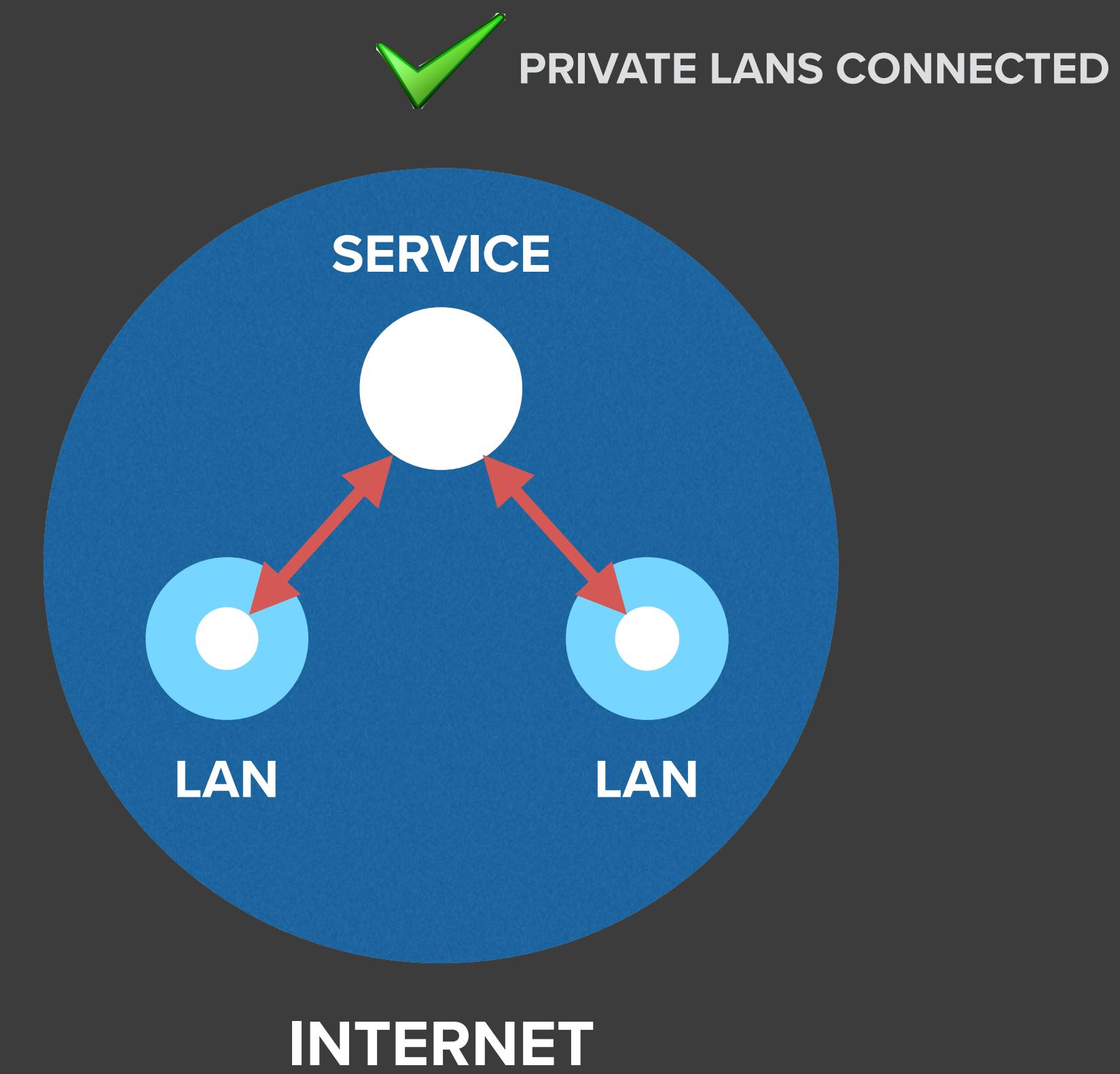
- Neurons earn money
- Factor down by proportion votes missed

# 5

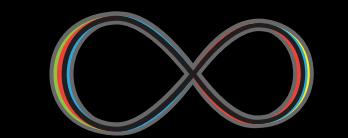
## **Public - Private Interoperability**

Possible using high consistency of crypto:3 networks

# Internet Network

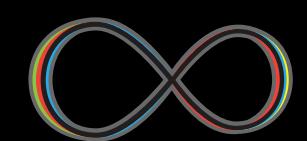
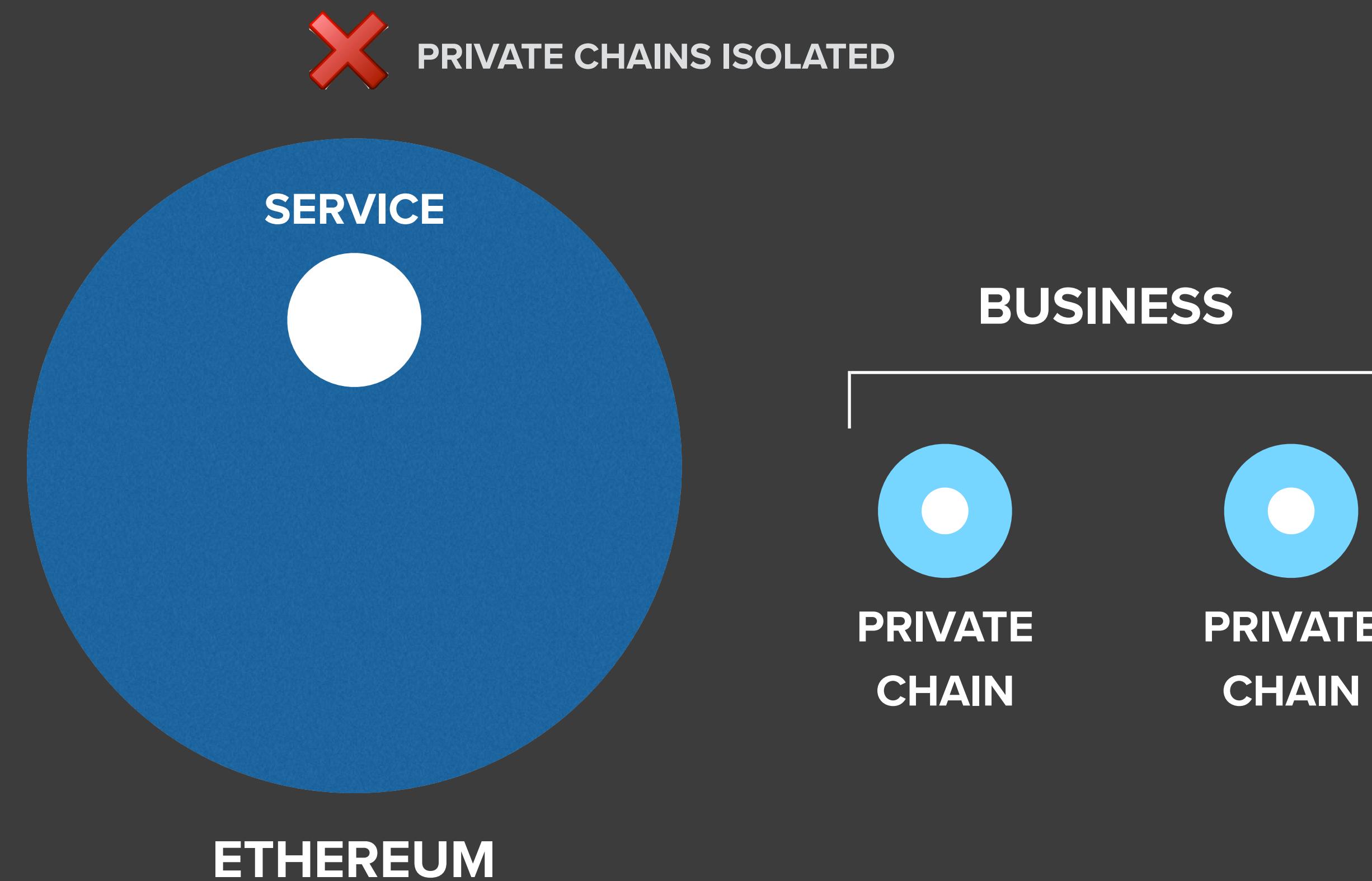


\* LAN = Local Area Network, run in office or home



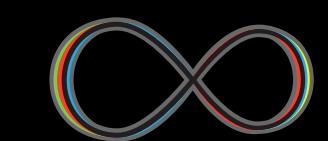
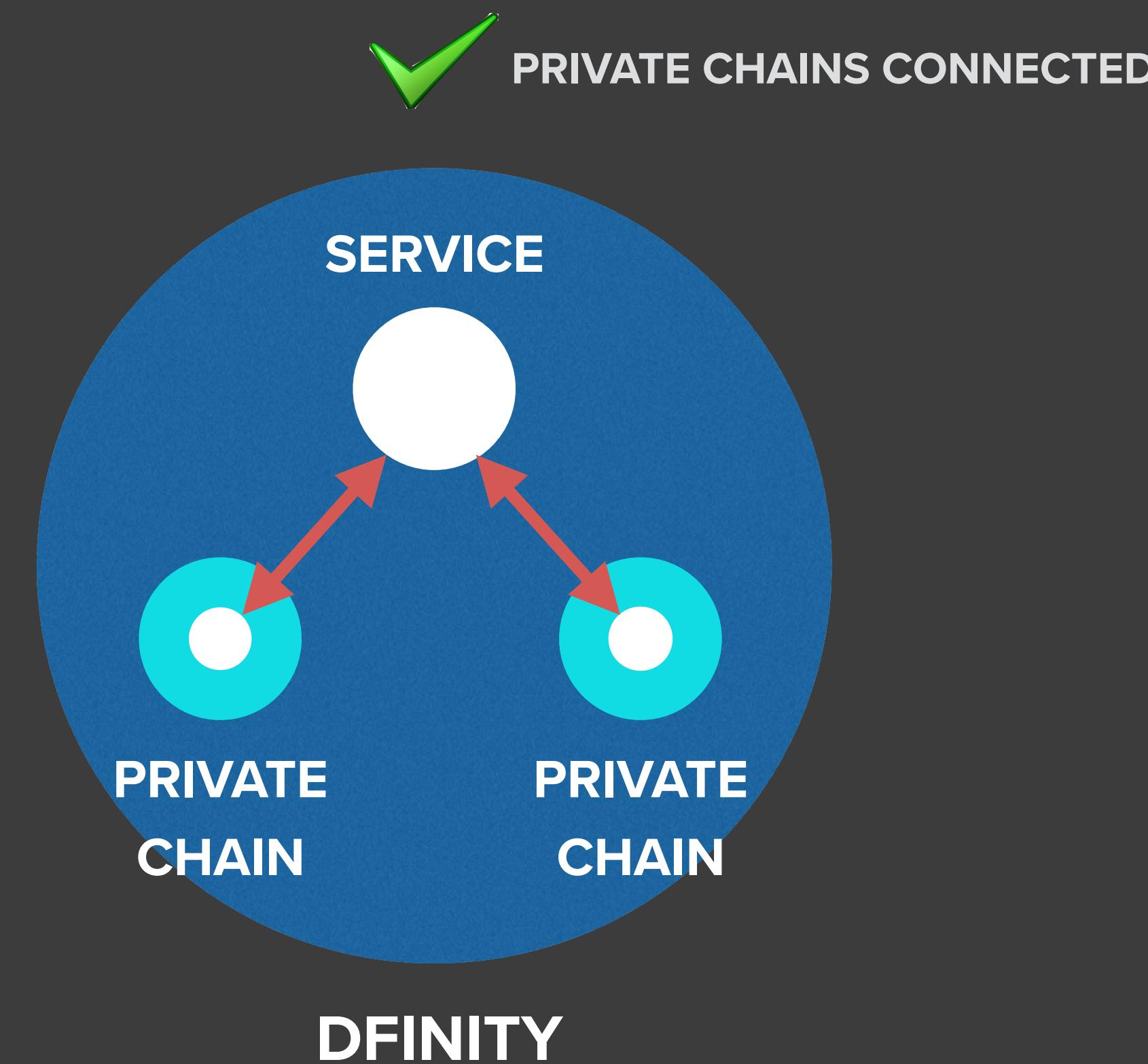
D F I N I T Y

# Ethereum Network



D F I N I T Y

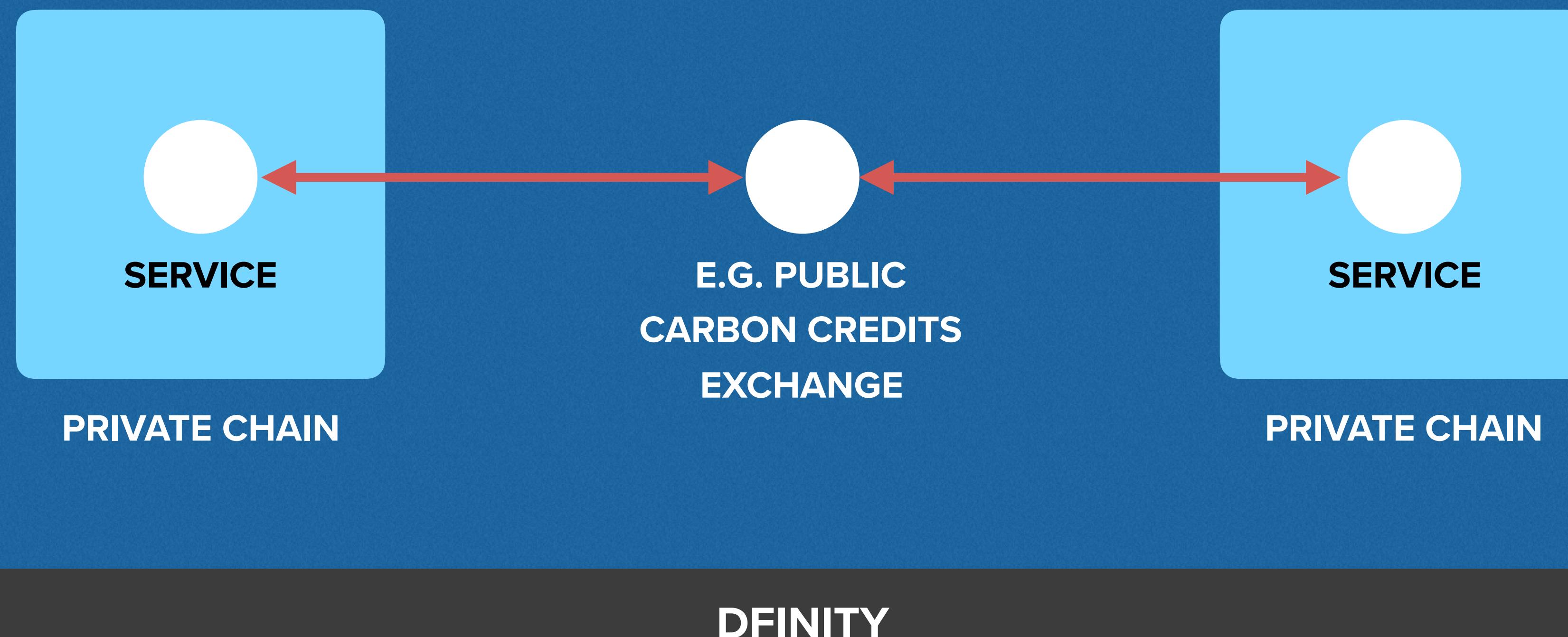
# DFINITY Network



D F I N I T Y

# Private Chains Interact

## EXAMPLE

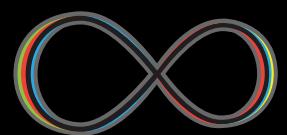


# Public Systems Are Building Blocks

## EXAMPLE

```
Include "dfinity:StableCoin.sol"  
Include "dfinity:Arbitration.sol"  
Include "dfinity:Identity.sol"  
Include "dfinity:Haulage.sol"  
  
-
```

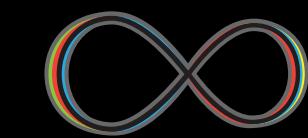
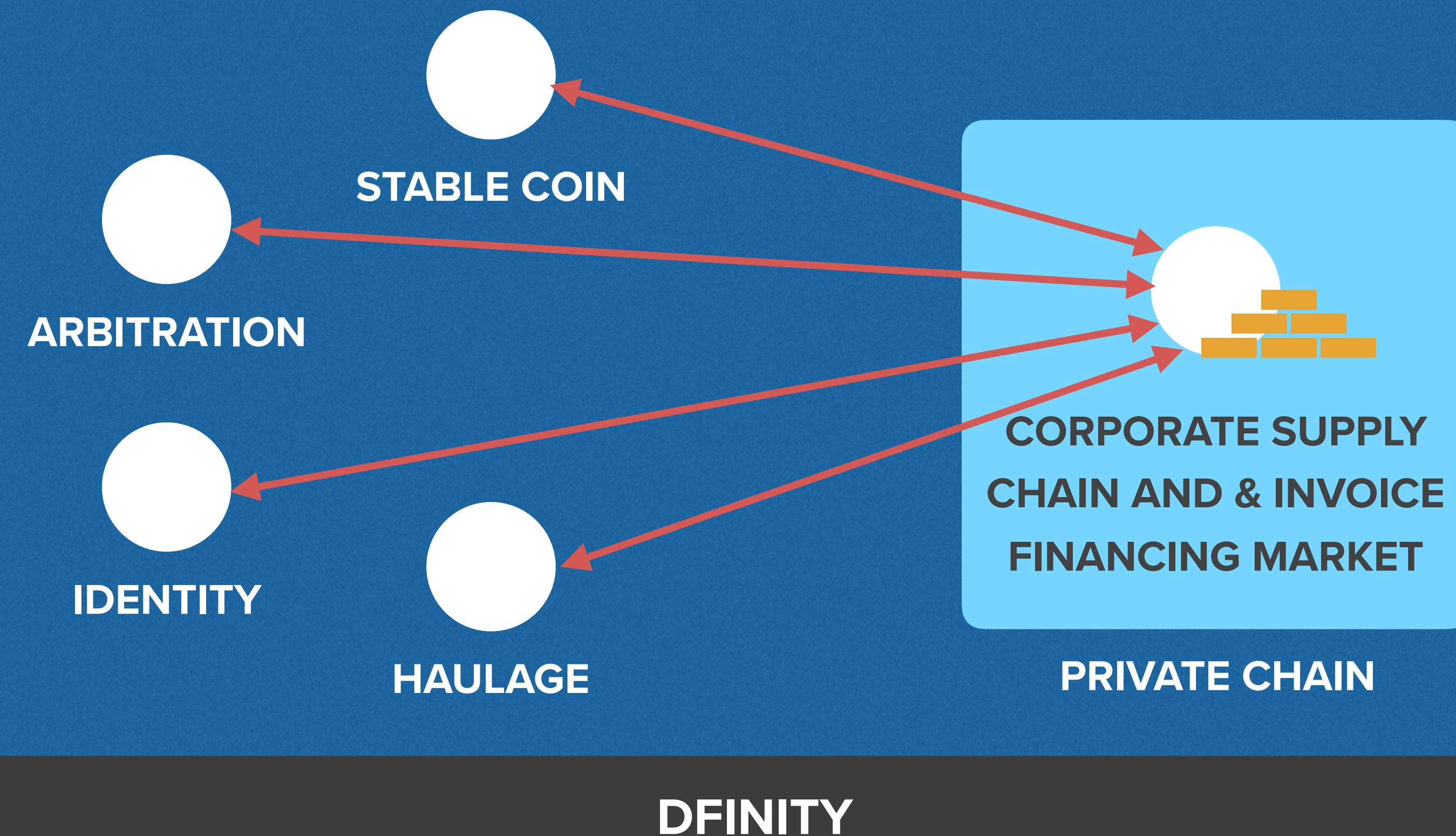
PRIVATE SMART CONTRACT CODE



D F I N I T Y

# Public Systems Are Building Blocks

## EXAMPLE



D F I N I T Y

6

**Release Schedule...**

# Release Schedule

Objective: race to Tungsten < 1 year

## 1| COPPER

- Threshold Relay Chain
- Blockchain Nervous System (BNS)
- Security deposits
- State-root-only-chain

## 2| ZINC

- Private - Public IOP

## 3| TUNGSTEN

- State sharding (basic)
- Validation Towers (basic)
- Asynchronous model for cross-shard programming
- USCIDs (Unique State Copy IDs)
- Advancements in BNS

## 4| LITHIUM

- Full “Validation Tree” architecture
- Micro-sharding
- Advanced economic models
- zkSNARKs, and other privacy enhancements

**More TBA...**