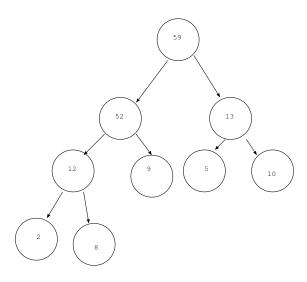
## HW3

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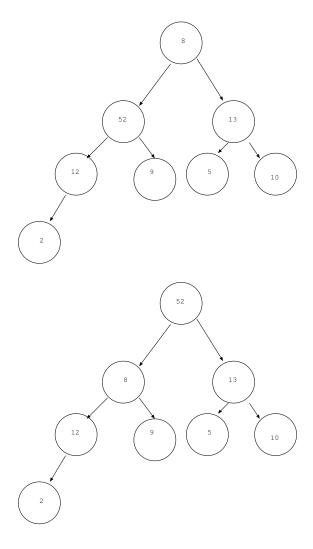
## 1. (a)



First we will replace the removed element with the last element in the tree

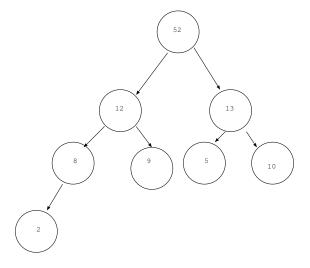
(b):

Then we must heapify the data structure



The max heap is now been heapified.

- 2. Solve the following recurrences. You can not use master theorem to solve them. You must show the steps in your derivation.
  - (a): If we draw out a 3 iterations then we would get



 $[cn]_1 + [\frac{cn}{3} + \frac{2cn}{3}]_2]_2 + [\frac{cn}{9} + \frac{2cn}{9} + \frac{2cn}{9} + \frac{4cn}{9}]_3$  where the brackets represent the layers this can be reduced to  $[cn]_1 + [cn]_2 + \frac{cn}{9} + \frac{2cn}{9} + \frac{2cn}{9} + \frac{2cn}{9} + \frac{4cn}{9}$  $[cn]_3 + ...$  for all iterations. So now we just need to find the number of iterations. Every iteration divides n by 3 or by  $\frac{2}{3}$ . So there should be logarithmic number of iterations. So T(n) = cnlog(n).

(b): First we will draw out a few iterations to find what the infinite series would be.

The sum of the the function would be  $cn + \frac{cn}{5} + \frac{cn}{25} + \dots$ 

Notice that the denominator  $5^i$  where i is the iteration number. So we can write this as an infinite series.

 $\sum_{i=1}^{\infty} \left(\frac{cn}{5}\right)^{i}.$ 

From this we can determine that the total sum is  $\frac{5nc}{4}$ . This means that this function is O(n).

(c): We will draw out the first few iterations again.  $[n^{log_5(7)}]_0 + [\frac{n^{log_5(7)}}{2} + \frac{n^{log_5(7)}}{2}]_1 + [\frac{n^{log_5(7)}}{4} + \frac{n^{log_5(7)}}{2}]_4 + \frac{n^{log_5(7)}}{2}]_4 + \frac{n^{log_5(7)}}{2}]_4]_2 + \dots$  As we can tell this can ben reduced to  $n^{log_5(7)} + n^{log_5(7)} + n^{log_5(7)} + \dots$  so

that means we just have to find the height of the tree.

Notice the denominator's value is  $2^i$  where i is value of the iteration. So we know the last iteration is when  $\frac{n}{2^i} = 1$  so there will be  $log_2(n)$  iterations. So the sum is  $log_2(n) * n^{log_5(7)}$