

## HW4

Shane Drafahl

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1. Let  $|x|_a$  be the number of occurrences of the symbol  $a$  in the string  $x$ .

Define a context-free grammar for the language  $L = \{ w \in \{ 1, 0 \}^* : |w|_0 = |w|_1 \}$

$$G = \{ (S), (1, 0), S, P \}$$
$$P = \{ S \rightarrow 01 | 10 | 1S0 | 0S1 | SS | \epsilon \}$$

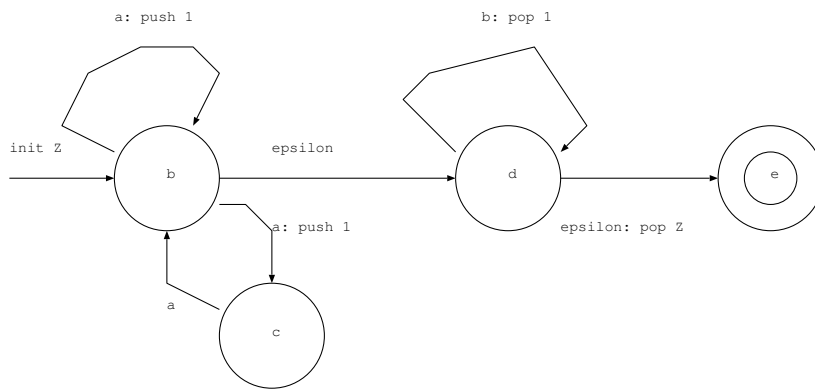
Using an inductive proof I will give a formal proof that the grammar generates  $L$

Basis:  $01 \in \{ 1, 0 \}^*$ ,  $10 \in \{ 1, 0 \}^*$  and  $\epsilon \in \{ 1, 0 \}^*$  and for all cases there are either one 1 and one 0 or zero 0 and zero 1 so they are in the language  $L$  and can be generated by  $G$ .

IH: Suppose that  $A, B$  can be generated from  $G$  and that  $A, B \in L$

Structural:  $A, B$  can be created from transitions from  $S$  so we must prove from  $1S0$ ,  $0S1$ , and  $SS$ . For  $SS$  so for  $AB$  or  $BA$  if we concatenated the strings  $A$  and  $B$  the 1's and 0's would still be equal because they are equal for  $A$  and  $B$ . For the transition  $0S1$  and  $1S0$  or  $0A1$  and  $1A0$ . It is assumed by the IH that  $S$  is in  $L$  and can be generated from  $G$  so  $A$  is in  $L$  and is generated from  $G$ . If we add a 1 and a 0 the number of 1's and 0's still equal because they do for  $A$  and adding a 1 and a 0 is still equal.

2. Define a NPDA for the language  $L = \{ a^n b^m : m, n \in N, m \leq n \leq 2m \}$



3.

