

## HW4

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1. Let  $|x|_a$  be the number of occurrences of the symbol  $a$  in the string  $x$ .

Define a context-free grammar for the language  $L = \{ w \in \{ 1, 0 \}^* : |w|_0 = |w|_1 \}$

$$G = \{ (S), (1, 0), S, P \}$$
$$P = \{ S \rightarrow 01|10|1S0|0S1|SSS|\epsilon \}$$

Using a inductive proof I will give a formal proof that the grammar generate  $L$

Basis:  $01 \in \{ 1, 0 \}^*$ ,  $10 \in \{ 1, 0 \}^*$  and  $\epsilon \in \{ 1, 0 \}^*$  and for all cases there are either one 1 and one 0 or zero 0 and zero 1.

IH: Suppose that  $w$  can be generated from  $G$  and that  $w \in L$

Structural: Suppose that there are strings  $A, B, C$  and  $A, B, C \in L$  by the IH and they are all generated from  $G$ . Suppose that they are concatenated  $ABC$  or  $CBA$ . All strings have the same number of 1's and 0's so therefore the resulting string  $ABC$  or  $CBA$  would also have the same number of 1's and 0's so  $ABC, CBA \in L$ . One of the transitions in the grammar is  $S \rightarrow SSS$  so  $S \rightarrow ABC$  or  $S \rightarrow CBA$  since we know from the IH that  $A, B, C \in L$  so therefore  $S \rightarrow A, B, C$ . Therefore  $ABC$  and  $CBA$  can be created by the grammar  $G$  and they are in the language  $L$ .

2. Define a NPDA for the language  $L = \{ a^n b^m : m, n \in N, m \leq n \leq 2m \}$

3.

