

Center for Information Services and High Performance Computing (ZIH)

An Introduction to Scout, a Vectorizing Source-to-Source Transformator

ACCU 2012, Oxford, UK

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Introducing myself

software engineer

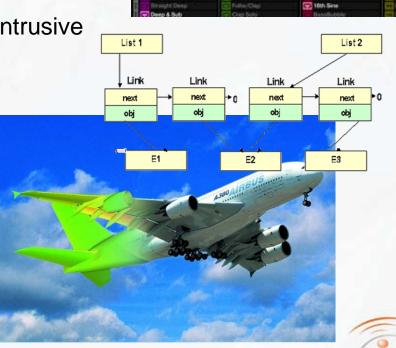
over 10 years in the "software industry"

author of the initial version of boost::intrusive

Kudos to Ion Gaztañaga!

back to the university in 2009

focus of ZIH on HPC research



Rattlesnake Vision



4 124.00

Motivation

Project HiCFD

Highly efficient Implementation of CFD-Codes for Many-Core-Architectures

Improve the runtime and parallel efficiency of computational fluid dynamics codes on HPC-Many-Core architectures by exploiting all levels of parallelism.

→ "all levels" includes the SIMD features of modern CPUs

starting point 2009: **3300 h** for a typical flight maneuver of one minute (40 million points, time resolution 0,01 sec, 10.000 CPUs)

MPI • many-core

OpenMP • multi-core

SIMD • single core





Agenda today

- 1. Motivation
 - where we come from: the HiCFD project
 - state of the art of (auto) vectorization
- 2. Introducing Scout
 - overview
 - background: using the clang framework for source-to-source transformation
 - features, capabilites, configuration
- 3. Applying Scout
 - measurements of two real-world CFD codes
- 4. Advanced vectorization techniques
 - Register blocking
 - Gather and scatter operations
- 5. Discussion

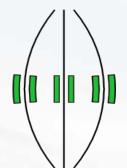




SIMD Vectorization: The Task

scalar code

vectorized code



source code (e.g. C)

source code (e.g. C)

intermediate code (e.g. llvm)

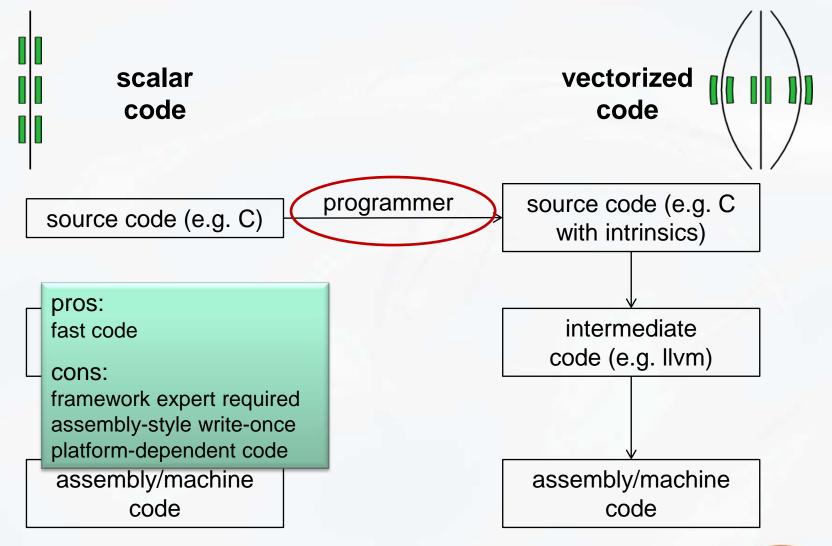
intermediate code (e.g. llvm)

assembly/machine code

assembly/machine code

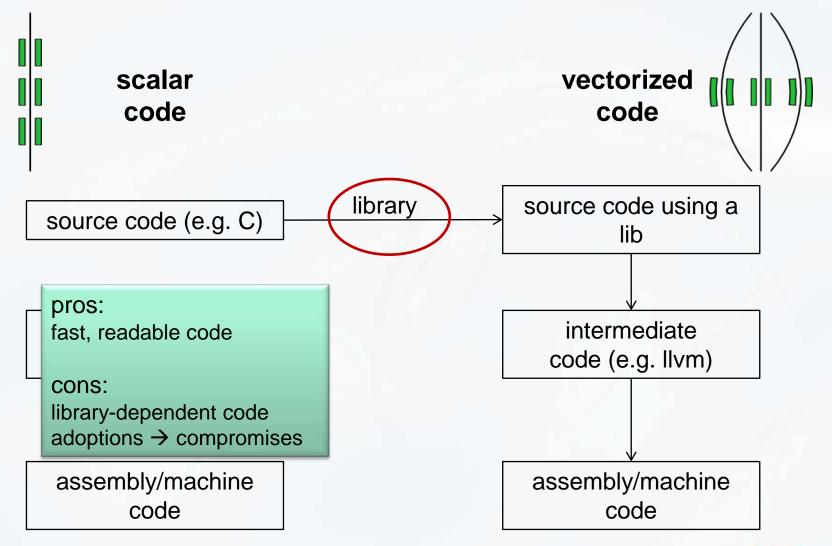






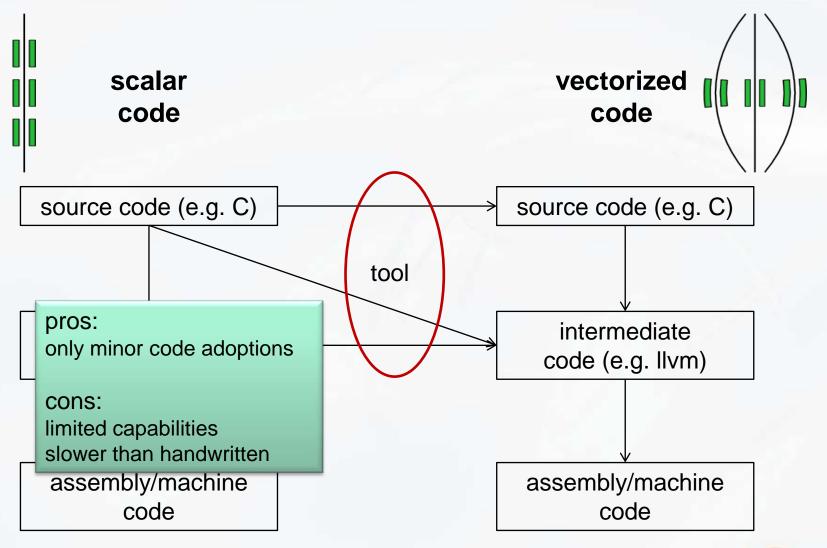






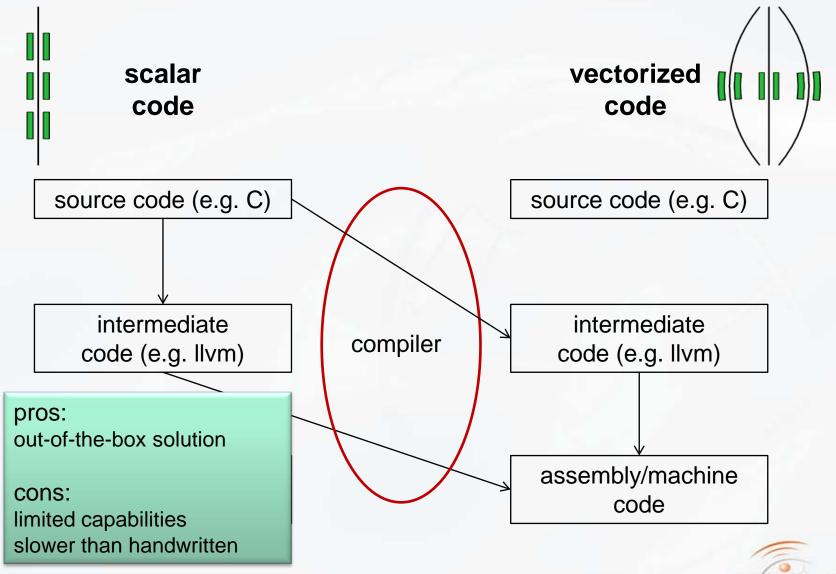














SIMD Vectorization: Compilers

- automatic dependency analysis incapable
 - e.g. no way to reason about indirectly indexed arrays
 - → auto-vectorization without meta-information sometimes impossible
 - Maleki, S. et. al.: An Evaluation of Vectorizing Compilers, PACT 2011
- pragma-based vectorization reveals new pecularities
 - e.g. requirement of an unsigned loop index
 - vectorization of real world application loops ranged from hard to impossible
- considerable improvements by icc V12 (#pragma simd)





SIMD Vectorization: Tools

- some tools are lost in a former millennium
 - techniques still remain
- academic tools:
 - Hohenauer et.al.: A SIMD optimization framework for retargetable compilers
 - Pokam et.al.: SWARP: a retargetable preprocessor for multimedia instructions



- commercial tools:
 - HMPP by CAPS Enterprise (source-to-source compiler)

List certainly incomplete...





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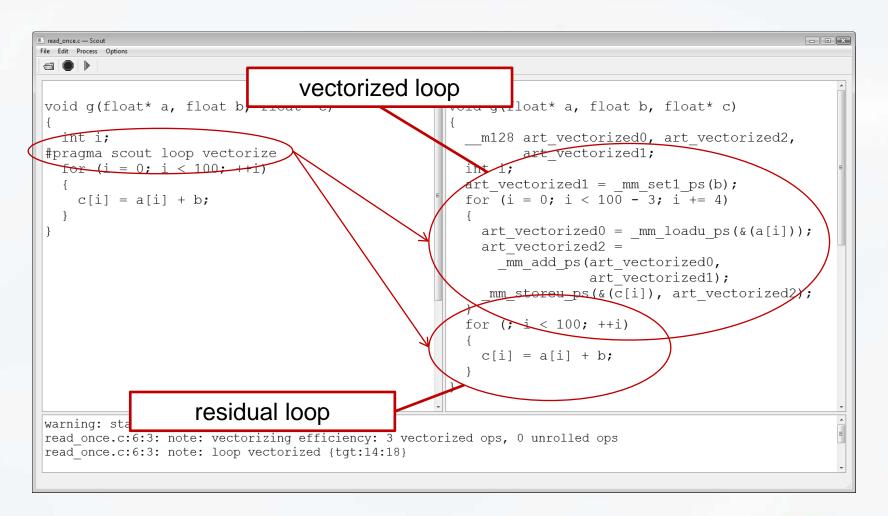
Scout: Introduction

Scout: A Source-to-Source Transformator for SIMD-Optimizations

- C as input and output source code
- focus on pragma-annotated loop vectorization
 - loop body simplification by function inlining and loop unswitching
 - unroll or vectorize inner loops
- User-Interface: GUI and command-line version
- Open source software tool: http://scout.zih.tu-dresden.de/
 - used in the production to accelerate code without much effort
 - a means to investigate new vectorization techniques



Scout: Impression





Scout: Command Line

CLI for an automatic build process via Makefile:

```
> scout -scout:configuration=./config/avx2.cpp -scout:extension=sc
-scout:prolog=./config/prolog.inc -I/usr/include file.c
```

- compiler arguments passed to clang
 - clang uses mostly gcc-like syntax
- Scout-specific arguments start with <u>-scout</u>:

```
-scout:configuration=file [req]:configuration file
```

-scout:preprocess=file [opt]: source file(s) containing definitions of

inlined functions

-scout:extension=text [opt]: target file extension

-scout:prolog=file [opt]: file content inserted in the target file





Scout: Command Line

Effect of -scout:prolog=./prolog.inc -scout:extension=sse:

```
prolog.inc:
  /* prolog start */
#include "ia32intrin.h"
  /* prolog end */
```

```
source.c:
#include "user.h"

void foo(float* a, float b)
{
#pragma scout loop vectorize
  for (i = 0; i < 100; ++i) {
    a[i] += b;
  }
}</pre>
```

```
source.c.sse:
#include "user.h"

/* prolog start */
#include "ia32intrin.h"

/* prolog end */

void foo(float* a, float b)
{
   __m128 av, bv, tv;
  bv = _mm_set1_ps(b);
  for (i = 0; i < 100; i += 4)</pre>
```

 $tv = _mm_loadu_ps(a + i);$

 $av = _mm_add_ps(tv, bv);$

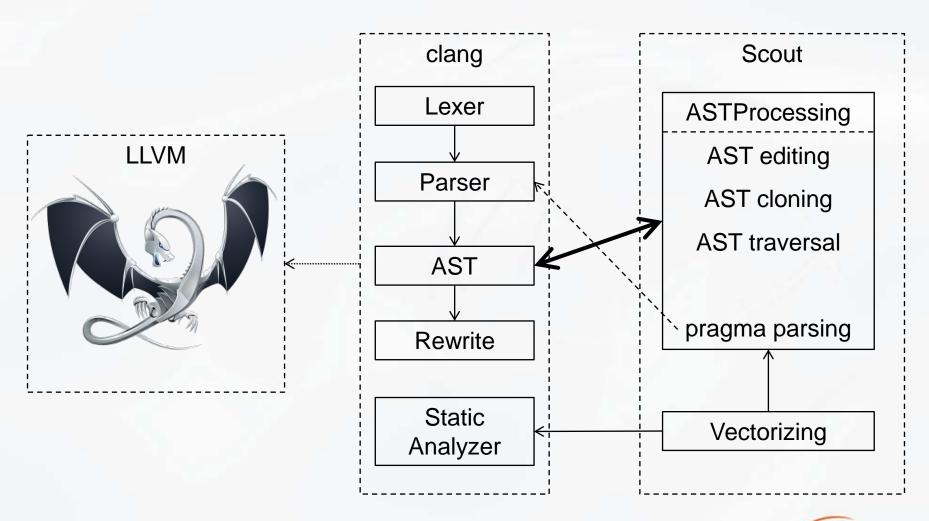
_mm_storeu_ps(a + i, av);

- content of prolog file inserted before the first meaningful non-include line
- → target file directly compilable





Background: Scout and clang







Background: Scout and clang

- manipulation of clang's abstrac syntax tree:
 - actually immutable
 - actually more than an AST
 - → actually a tricky and sometimes hacky approach
- nevertheless scout::ASTProcessing works:



Background: Scout and clang

- parsing expressions in pragma arguments
 - forecast: #pragma scout loop vectorize aligned(a.array, a.ptr)
 - not supported by clang → brute force patch
 - but they are going to need something similar for OpenMP
 - better solution: configurable C++11 attributes
- clang::StaticAnalyzer for alias analysis

```
double* b = //...
for (int i=0; ...) {
  double* ptr = b;
  // b[i] → Node1
  // ptr[i] → Node2
}
```





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- configuration file written in C++:
 - no need to learn another configuration syntax
 - usual preprocessing means (conditional compilation, includes) available
 - somewhat stretched semantics
- replace expressions by their vectorized intrinsic counterparts
 - even complex expressions like a+b*c (fmadd) or a<b?a:b (min)
 - and function calls like sqrt
- currently available configurations:
 - SSE2, SSE4 (blending): double, float
 - AVX, AVX2 (fmadd, gather): double, float
 - ARM NEON: float
- other architectures easy to provide
 - MIC (sorry, under NDA)
 - even upcoming or experimental ones





- specialized template named config in the namespace scout:
 - type parameter: base type of the vector instruction set
 - integral parameter: vector size (number of vector lanes)
- name of target SIMD type introduced by a typedef named type
 - probably that name needs to be introduced
 - → no need to include the header of the appropriate SIMD architecture

```
namespace scout {
  template < class T, unsigned size > struct config;
  template <> struct config < float, 4 > {
    typedef float __m128 __attribute__ ((__vector_size__ (16)));
    typedef __m128 type;
    enum { align = 16 };
  };
} // end of namespace
```



- specialized template named config in the namespace scout:
 - type parameter: base type of the vector instruction set
 - integral parameter: vector size (number of vector lanes)
- name of target SIMD type introduced by a typedef named <u>type</u>
 - probably that name needs to be introduced
 - → no need to include the header of the appropriate SIMD architecture
- alignment requirement by a enum named align

```
namespace scout {
  template < class T, unsigned size > struct config;
  template < > struct config < float, 4 > {
    typedef float __m128 __attribute__ ((__vector_size__ (16)));
    typedef __m128 type;
    enum { align = 16 };
};
} // end of namespace
```



- instruction definition by static member functions
 - last statement of a member function body always a string literal
 - that string is inserted in the target code
 - arguments are expanded using boost::format

```
template<> struct config<float, 4> {
  typedef float base_type;
  typedef float __m128 __attribute__ ((__vector_size__ (16)));
  typedef __m128 type;

static type load_aligned(base_type*)
  {
    "_mm_load_ps(%1%)";
  }
  static void store_aligned(base_type*, type)
  {
    "_mm_store_ps(%1%, %2%)";
  }
};
```



picking the instruction by pre-defined method name:

```
type load [un]aligned(base type* p)
                                                   return *p as vector
void store [un]aligned(base type* p, type v)
                                                   *p := v
void store nt packed [un]aligned
                                                   v = : q^*
                                                   (no cache pollution)
  (base type* p, type v)
type splat(base_type x)
                                                   return { x, x,..., x }
type broadcast(base_type* p)
                                                   return {*p, *p,..., *p }
                                                   return { xn, ... x1 }
type set(base type x1, ..., base type xn)
base_type extract(type v, unsigned int i)
                                                   return v[i]
void insert(type v, base type x, unsigned int i) v[i] := x
```

additional names for gather/scatter support:

```
gs_index_type get_uniform_gs_index(int)

type gather(base_type*, gs_index_type)

void scatter(base_type*, gs_index_type, type)
```





- picking the instruction by expressions or functions:
 - method name ignored
 - method signature uses fundamental types
 - string literal preceded by a list of expressions and/or function declarations



Example: complex expression mapping with SSE4

- computation of all lanes
 - → turn off fp exceptions (if not already done)
- blending -> good acceleration



- simplify loop bodies
 - source-level inlining
 - loop unswitching (remove loop-invariant conditions)
- vectorize assignments
 - don't change the arrays-of-structs data layout
 composite load / store operations required
 - complex loops with conditions and mixed data types
- modern SMD architectures provide rather short vector registers
 - whole-loop transformations better suited for traditional vector machines
- → vectorization by *unroll-and-jam* approach



Unroll-and-Jam:

```
#pragma scout loop vectorize
for (i = si; i < ei; ++i)
{
    s<sub>i</sub>;
    t<sub>i</sub>;
}
```

```
for (i=si; i<ei-vs+1; i+=vs)
{
    s<sub>i</sub>;
    s<sub>i+1</sub>;
    ...
    s<sub>i+vs</sub>;
    t<sub>i</sub>;
    t<sub>i+1</sub>;
    ...
    t<sub>i+1</sub>;
    ...
    t<sub>i+vs</sub>;
}
// residual loop
```

first step: unroll all statements according to the vector size



Unroll-and-Jam:

```
#pragma scout loop vectorize
for (i = si; i < ei; ++i)
{
    s<sub>i</sub>;
    t<sub>i</sub>;
}
```

```
for (i=si; i<ei-vs+1; i+=vs)
{
    s<sub>i:i+vs</sub>;
    t<sub>i</sub>;
    t<sub>i+1</sub>;
    ...
    t<sub>i+vs</sub>;
}
// residual loop
```

- first step: unroll all statements according to the vector size
- second step: jam vectorizeable statements
- unvectorizeable statements remain unrolled
 - but they don't inhibit vectorization





unrolling of non-inlined functions:

```
double a[100], d[100];
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
   a[i] = foo(2.0 * d[i]);
}</pre>
```



```
__m128d cv, av1, av2, av3;
double a[100], d[100];
cv = _mm_set1_pd(2.0);
for (i = 0; i < 100; i += 2) {
  av1 = _mm_loadu_pd(d + i);
  av2 = _mm_mul_pd(cv, av1);
  av3 = _mm_set_pd(
  foo(_mm_extract_pd(av2,1)),
  foo(_mm_extract_pd(av2,0)));
  _mm_storeu_pd(a + i, av3);
}</pre>
```

no support for functions with out-arguments





unrolling of if-statements:

```
double a[100], b[100], d[100];
#pragma scout loop vectorize
for (i = 0; i < 100; ++i)
 y = a[i];
 if (y < 0)
   b[i] = d[i] + y;
 else
   b[i] = d[i] - y;
 a[i] = b[i] * d[i];
```

```
__m128d yv, bv, tv, dv;
double a[100], b[100], d[100];
for (i = 0; i < 100; i += 2)
  yv = mm loadu pd(a + i);
  if ( mm extract pd(yv,0)<0)
    b[i] = d[i] +
      _mm_extract_pd(yv, 0);
  else
   b[i] = d[i] -
      _mm_extract_pd(yv, 0);
  if ( mm extract pd(yv,1)<0)</pre>
    b[(i + 1)] = d[(i + 1)] +
      _mm_extract_pd(yv, 1);
  else
   b[(i + 1)] = d[(i + 1)] -
      _mm_extract_pd(yv, 1);
  bv = mm loadu pd(b + i);
  dv = mm loadu pd(d + i);
  tv = _{mm\_mul\_pd(bv, dv)};
  mm storeu pd(a + i, tv);
```

only necessary for loop-variant conditions





Scout: Scalar Transformations

Loop unswitching:

```
double a[100], b, c[100];
int mode = /*...*/

#pragma scout loop vectorize scalar
for (i = 0; i < 100; ++i)
{
   if (mode == 0)
     c[i] = a[i] + b;
   else
     c[i] = a[i] - b;
}</pre>
```

```
double a[100], b, c[100];
int mode = /*...*/
if (mode == 0) {
  for (i = 0; i < 100; ++i)
    c[i] = a[i] + b;
} else {
  for (i = 0; i < 100; ++i)
   c[i] = a[i] - b;
```

- moving of loop-invariant conditions outside of loops
- code bloat outweighted by vectorization gains





Scout: Scalar Transformations

Loop unswitching:

```
double a[100], b, c[100];
int mode = /*...*/
if (mode == 0) {
  for (i = 0; i < 100; ++i)
    c[i] = a[i] + b;
} else {
  for (i = 0; i < 100; ++i)
    c[i] = a[i] - bi
```

- moving of loop-invariant conditions outside of loops
- code bloat outweighted by vectorization gains
- also used for conditional expressions





Scout: Capabilities

Mixed data types:

```
float a[100];
double b[100], x;
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
 x = a[i];
 x = x / b[i];
 a[i] = x;
```

```
float a[100];
double b[100];
 m128 av;
 m128d xv1, xv2, bv1, bv2;
for (i = 0; i < 100; i += 4)
  av = _mmloadu_ps (a + i);
  xv1 = _mm_cvtps_pd (av);
  xv2 = _mm_cvtps_pd (
    _mm_movehl_ps (av, av));
  bv1 = mm loadu pd (b+i);
  bv2 = \underline{mm} loadu pd (b+i+2);
  xv1 = _mm_div_pd (xv1, dv1);
  xv2 = mm_div_pd(xv2, dv2);
  av = mm movelh ps (
    _mm_cvtpd_ps (xv1),
    _mm_cvtpd_ps (xv2));
  _mm_storeu_ps (a + i, av);
```

- vectorized according to the largest vector size
- by-product of the rather local scope of the unroll-and-jam approach



Scout: Capabilities

Arbitrary constant loop stride:

```
int k = /*...*/i
double a[100], b[100];
#pragma scout loop vectorize
for (i = 0; i < 100; i += k)
 a[i] = b[i] * b[i];
```

```
int k = /*...*/;
__m128d av1, av2;
double a[100], b[100];
for (i = 0;
     i < 100 - 1;
     i += k * 2) {
  av1 = _mm_set_pd(b[(i + k)],
                   b[i]);
  av2 = _mm_mul_pd(av1, av1);
  a[i] =
    _mm_extract_pd(av2, 0);
  a[(i + k)] =
    _mm_extract_pd(av2, 1);
/* residual loop */
```

- array-of-structure data layout need composite load and store operations
 this nets that feature anyway
- coming later (but today): gather and scatter



Partial vectorization:

```
double a [100], c[100];
int d [100];
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
  int j = d[i];
 double b = a[j];
  // computations
  // introduces an inner-loop
  // dependency if
  // d[i]==d[i+1]:
#pragma scout vectorize unroll
 c[i] += b;
```

```
m128d b_v;
int j_v [2];
double a [100], c[100];
int d [100];
for (i = 0; i < 100; i += 2) {
  j_v[0] = d[i];
  j_v[1] = d[(i + 1)];
  b_v = _mm_set_pd (
    a[j_v[0]], a[j_v[1]]);
  // vectorized computations
  // compute every element
  // separately:
  c[j v[0]] += mm extract pd
                   (b v , 0);
  c[j v[1]] += mm extract pd
                   (b v , 1);
```

- enables the vectorization of loops in the presence of dependencies
- transform the vectorizeable part and leave dependencies intact





Inner loop vectorization:

```
double *a, *b, c, *d;
#pragma scout loop vectorize
for (i = 0; i < 100; ++i)
{
   d[i] = a[i];
   for (j = 0; j < k; ++j)
   {
      d[i] += b[j * k + i];
   }
   d[i] *= c;
}</pre>
```

```
double *a, *b, c, *d;
__m128d dv, cv, bv;
for (i = 0; i < 100; i += 2)
{
    dv = _mm_loadu_pd(a + i);
    for (j = 0; j < k; ++j)
    {
        bv = _mm_loadu_pd(b+j*k+i);
        dv = _mm_add_pd(dv, bv);
    }
    dv = _mm_mul_pd(dv, cv);
    _mm_storeu_pd(d + i, dv);
}</pre>
```

only for constant inner loop range





Inner loop vectorization:

```
double *a, *b, c, *d;
#pragma scout loop vectorize
for (i = 0; i < 100; ++i)
 d[i] = a[i];
 for (j = 0; j < k; ++j)
   d[i] += b[i * k + j];
 d[i] *= c;
```

```
double *a, *b, c, *d;
__m128d dv, cv, bv;
for (i = 0; i < 100; i += 2)
  dv = mm loadu pd(a + i);
  for (j = 0; j < k; ++j)
    bv = mm set pd(
      b[(i * k + j + k)],
      b[i * k + j]);
    dv = mm \text{ add } pd(dv, bv);
  dv = mm \ mul \ pd(dv, cv);
  _mm_storeu_pd(d + i, dv);
```

only for constant inner loop range





Inner loop vectorization:

```
double *a, *b, c, *d;
#pragma scout loop vectorize
for (i = 0; i < 100; ++i)
{
   d[i] = a[i];
   for (j = 0; j < k; ++j)
   {
     d[i] += b[i];
   }
   d[i] *= c;
}</pre>
```

```
double *a, *b, c, *d;
__m128d dv, cv, bv;
for (i = 0; i < 100; i += 2)
{
    dv = _mm_loadu_pd(a + i);
    for (j = 0; j < k; ++j)
    {
        bv = _mm_loadu_pd(b + i);
        dv = _mm_add_pd(dv, bv);
    }
    dv = _mm_mul_pd(dv, cv);
    _mm_storeu_pd(d + i, dv);
}</pre>
```

- only for constant inner loop range
- no displacement of inner-loop-invariant expressions by Scout





Reductions:

```
float *a, x, y;

#pragma scout loop vectorize
for (i = 0; i < 100; ++i)
{
    x += a[i];
    y = MIN(y, a[i]);
}</pre>
```

```
float *a, x, y;
 m128 av, xv, yv;
xv = mm set1 ps(0);
yv = mm set1 ps(y);
for (i = 0; i < 100; i += 2) {
  av = _mm_loadu_ps(a + i);
 xv = _mm_add_ps(xv, av);
 yv = mm min ps(yv, av);
for (ti = 0U; ti < 4U; ++ti) {
 x = x + mm extract ps(xv, ti);
for (ti = 0U; ti < 4U; ++ti) {
  y = y<_mm_extract_ps(yv, ti) ?
         mm extract ps(yv, ti);
```

- note the numerical instability due to a different computation order
- TODO: merging the loops and introduce horizontal operations





Scout: Fine-Tuning the Vectorization

list of additional #pragma scout vectorize arguments:

noremainder	don't generate a residual loop
size(N)	unroll N times (N >= VS && N % VS==0)
scalar	no vectorization (only inlining and unswitching)
aligned(expr)	use aligned loads/stores
nontemporal(expr)	use nontemporal loads/stores (avoid cache pollution)
align	automatic alignment of the most often accessed memory location

additional pragmas:

<pre>#pragma scout loop vectorize unroll</pre>	used in loop bodies: next statement remains unrolled
#pragma scout function expand	used in front of a function definition: all calls in that function are recursively inlined
#pragma scout loop unroll	used in front of a loop with constant iteration range: loop gets completely unrolled





Using aligned and nontemporal:

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
#pragma ... aligned(a, b)
for (i = 0; i < 100; ++i)
{
   a.array[i] = b[i];
   a.ptr[i] = b[i+1];
}</pre>
```

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
__m128 tv
for (i = 0; i < 100; i += 4)
{
   tv = _mm_load_ps(b + i);
   _mm_store_ps(a.array + i, tv);
   tv = _mm_load_ps(b + i+1);
   _mm_storeu_ps(a.ptr + i, tv);
}</pre>
```

accesses to regions and their direct subregions are aligned



Using aligned and nontemporal:

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
#pragma ... aligned(a, b)
for (i = 0; i < 100; ++i)
{
   a.array[i] = b[i];
   a.ptr[i] = b[i+1];
}</pre>
```

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
__m128 tv
for (i = 0; i < 100; i += 4)
{
   tv = _mm_load_ps(b + i);
   _mm_store_ps(a.array + i, tv);
   tv = _mm_load_ps(b + i+1);
   _mm_storeu_ps(a.ptr + i, tv);
}</pre>
```

- accesses to regions and their direct subregions are aligned
- no reasoning about the index



Using aligned and nontemporal:

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
#pragma ... aligned(a, b[i])
for (i = 0; i < 100; ++i)
{
   a.array[i] = b[i];
   a.ptr[i] = b[i+1];
}</pre>
```

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
__m128 tv
for (i = 0; i < 100; i += 4)
{
   tv = _mm_load_ps(b + i);
   _mm_store_ps(a.array + i, tv);
   tv = _mm_loadu_ps(b + i+1);
   _mm_storeu_ps(a.ptr + i, tv);
}</pre>
```

- accesses to regions and their direct subregions are aligned
- no reasoning about the index -> denote the access directly
 - all names must be declared before the pragma line





Using aligned and nontemporal:

```
typedef struct {
   float* ptr;
   float array[100];
} A;

A a;
float b[100];
__m128 tv
for (i = 0; i < 100; i += 4)
{
   tv = _mm_loadu_ps(b + i+1);
   _mm_stream_ps(a.ptr + i, tv);
}</pre>
```

- accesses to regions and their direct subregions are aligned
- no reasoning about the index -> denote the access directly
 - all names must be declared before the pragma line
- SSE nontemporal streaming requires aligned data





Agenda today

1. Motivation

- where we come from: the HiCFD project
- state of the art of (auto) vectorization
- 2. Introducing Scout
 - overview
 - background: using the clang framework for source-to-source transformation
 - features, capabilites, configuration

3. Applying Scout

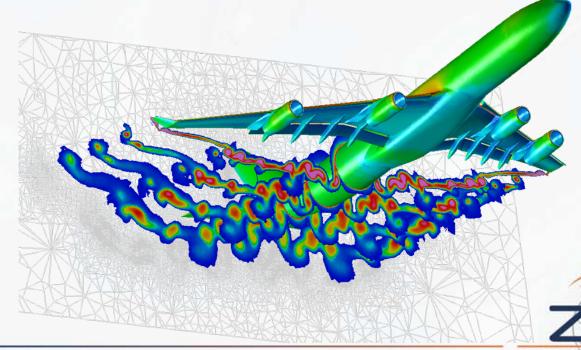
- measurements of two real-world CFD codes
- 4. Advanced vectorization techniques
 - Register blocking
 - Gather and scatter operations
- 5. Discussion





first CFD computation kernel computes flow around air planes

- unstructured grid → indirect indexing
- arrays-of-structure data layout
- partial vectorization
 - enforcing compiler auto-vectorization by pragmas lead to incorrect results
- double precision, SSE platform → two vector lanes





Compiler: Intel 11.1, Windows 7, Intel Core 2 Duo, 2.4 MHz

```
int d [100];
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
  int j = d[i];
  // first computations with
  // a[j], b[j] aso.
}
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
  int j = d[i];
  // second computations with
  // a[j], b[j] aso.
}</pre>
```

speedup relation	original to vectorized	ideally a value near 2.0
Grid 1	1.070	ideally a value near 2.0
Grid 2	1.075	→ unsatisfying result





lot of consecutive loops traverses over the same data structures:

```
int d [100];
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
  int j = d[i];
  // first computations with
  // a[j], b[j] aso.
}
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
  int j = d[i];
  // second computations with
  // a[j], b[j] aso.
}</pre>
```



```
int d [100];
#pragma scout loop vectorize
for (i = 0; i < 100; ++i) {
  int j = d[i];
  // first computations with
  // a[j], b[j] aso.

// second computations with
  // a[j], b[j] aso.
}</pre>
```

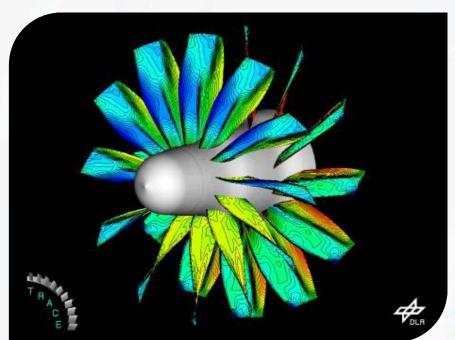
speedup relation	original to vectorized	merged to merged+vect.	original to merged+vect.
Grid 1	1.070	1.391	1.489
Grid 2	1.075	1.381	1.484





second CFD computation kernel computes interior flows of jet turbines

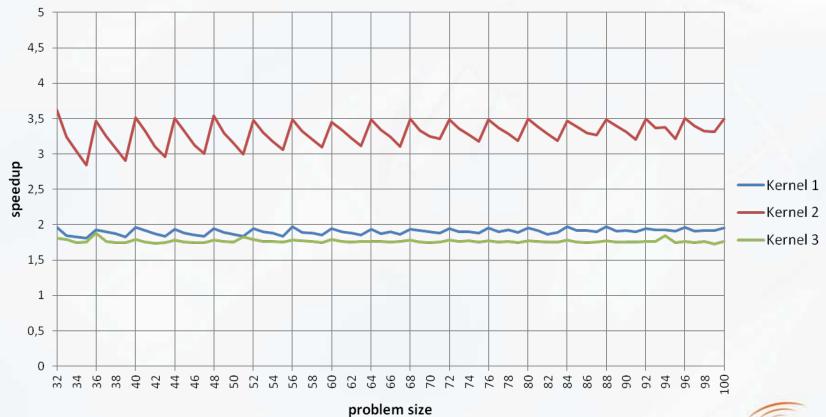
- structured grid → direct indexing
- arrays-of-structure data layout
- divided in three sub-kernels
- vectorization of complete loops





a CFD computation kernel computing interior flows of jet turbines

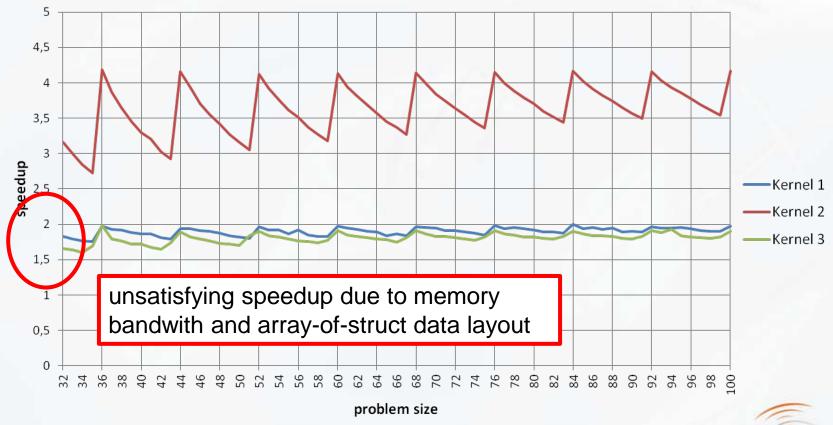
- three-dimensional grid \rightarrow (problem size)³ cells to compute
- single precision, target architecture: SSE → 4 vector lanes





a CFD computation kernel computing interior flows of jet turbines

single precision, target architecture: AVX → 8 vector lanes





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- 5. Discussion





derived from loop blocking (aka loop tiling):

- treats register file like a zero-level cache
- example: SIMD architecture provides 4 vector lanes

```
#pragma ... size(8)
for (i = si; i < ei; ++i)
{
    s<sub>i</sub>;
    t<sub>i</sub>;
}

for (i = si; i < ei-7; i += 8)
{
    s<sub>i:i+3</sub>;
    s<sub>i+4:i+7</sub>;
    t<sub>i:i+3</sub>;
    t<sub>i:i+3</sub>;
    t<sub>i:i+3</sub>;
}
```

vectorized statements are logically blocked, technically unrolled





derived from loop blocking (aka loop tiling):

- treats register file like a zero-level cache
- example: SIMD architecture provides 4 vector lanes

```
float b[100], d[100], x;

#pragma ... size(8)
for (i = 0; i < 100; ++i)
{
    x = d[i];
    for (j = 0; j < k; ++j)
    {
        x += b[j];
    }
    d[i] = x;
}</pre>
```

```
float b[100], d[100], x;

for (i=0; i < 100 - 7; i += 8)
{
    dv0 = _mm_loadu_ps(d+i);
    dv1 = _mm_loadu_ps(d+i+4);
    for (j = 0; j < k; ++j)
    {
        tv0 = _mm_set1_ps(b[j]);
        dv0 = _mm_add_ps(dv0,tv0);
        dv1 = _mm_add_ps(dv1,tv0);
    }
    _mm_storeu_ps(d+i, dv0);
    _mm_storeu_ps(d+i+4, dv1);
}</pre>
```

number of loads for tv0 halved



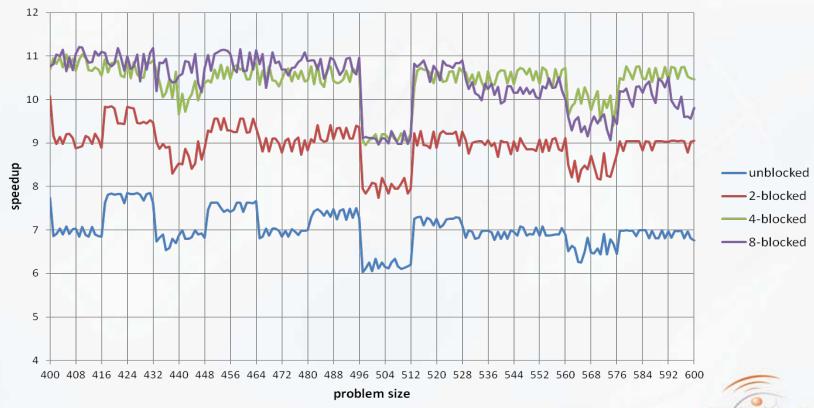


- useful for loop-invariant variables
 - e.g. variables in inner loops not depending on the outer-loop index
- test case derived from production code:

```
#pragma scout loop vectorize size(BLOCK_SIZE) align(a,b,c)
for (i = 0; i < S; ++i)
{
   for (j = 0; j < G; j++)
   {
     float x = a[j];
     for (d = 0; d < D; d++)
     {
        x += b[j*D+d] * c[d*S+i];
     }
     output[i * G + j] = x;
}</pre>
```

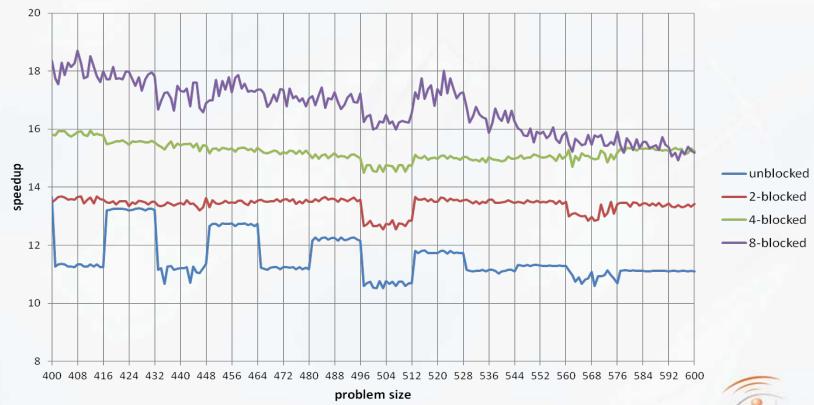


- ICC12, SSE, 8 registers, no scalar residual loop:
 - no spillings in up to 4-blocked loops
 - 8-blocked loop: 2 spillings in innermost and 8 in 2nd innermost loop



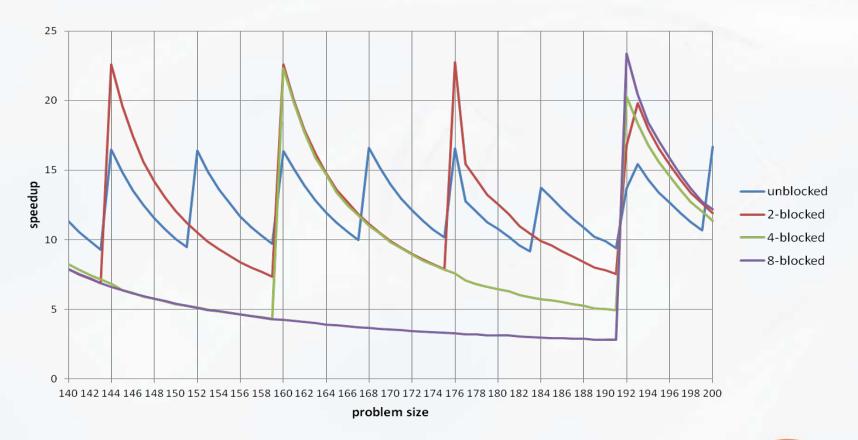


- ICC12, AVX, 16 registers, no scalar residual loop:
 - no spillings in up to 4-blocked loops
 - 8-blocked loop: 5 spillings in the 2nd innermost loop (no spilling in the innermost loop)





- beware of simple scalar residual loops!
 - same nice speedup only at the peaks







- array-of-struct data layout forces the use of composite loads and stores
 - composite intrinsics (e.g. _mm_set_pd(x,y)) result in a series of scalar assembly instructions
 - broader vector registers makes the problem even worse

Thus either

→ rearrange your data layout

or

→ use gather and scatter operations



AVX example with comile-time constant gather distance:

```
struct V { float x, y; };
struct A { V v[5]; };
A a[100];

#pragma scout loop vectorize
for (i=0; i<S; ++i) {
   s = a[i].v[1].x;
}</pre>
```

```
struct V { float x, y; };
struct A { V v[5]; };
A a[100];
for (i=0; i < S-7; i += 8) {
   s_v = _mm256_set_ps(
   a[i+7].v[1].x, a[i+6].v[1].x,
   a[i+5].v[1].x, a[i+4].v[1].x,
   a[i+3].v[1].x, a[i+2].v[1].x,
   a[i+1].v[1].x, a[i].v[1].x);
}</pre>
```



AVX2 example with comile-time constant gather distance:

```
struct V { float x, y; };
struct A { V v[5]; };
A a[100];
#pragma scout loop vectorize
for (i=0; i<S; ++i) {
  s = a[i].v[1].x;
```

```
struct V { float x, y; };
struct A { V v[5]; };
A a[100];
m256i d v =
 mm256 set epi32(
  sizeof(A)*7,sizeof(A)*6,
  sizeof(A)*5, sizeof(A)*4,
  sizeof(A)*3, sizeof(A)*2,
  sizeof(A) ,0);
for (i=0; i < S-7; i += 8) {
 s_v = _mm256_i32gather_ps(
  &a[i].v[1].x, dist_v, 1);
```

scalar inititialization outside of the loop



SDE of Intel gives some estimations of the effect

	total # of instructions		parallel portion	
	AVX	AVX 2	AVX	AVX 2
Kernel 1	1244	793	43%	90%
Kernel 2	2451	2232	37%	45%
Kernel 3	2885	2666	51%	57%

- Kernel 1: only scatter is missing
- Kernel 2: unrolled loop-variant condition
- Kernel 3: indirect indexing → needs a recursive gather → TODO



Remember the promise I made: gather with arbitrary constant loop stride

```
int k = /*...*/;
double a[100], b[100];
#pragma scout loop vectorize
for (i = 0; i < 100; i += k) {
   a[i] = b[i] * b[i];
}</pre>
```

```
int k = /*...*/i
double a[100], b[100];
int kd = k * sizeof(double);
m128d av1, av2;
 m128d ki = mm set pi(
       kd * 3, kd * 2, kd, 0);
for (i = 0; i < 100 - 1;
     i += k * 4) {
  av1 = mm256 i32gather pd(
   b + i, art vect1, 1);
  av2 = _mm_mul_pd(av1, av1);
  // ...
/* scalar epilog */
```





Coding Advices

- allowed statements in loop bodies:
 - (compound) assign expressions
 - including copy assignments of records
 - function calls
 - if- and for-statements
 - TODO: while, switch/case
- better use ?: expressions for loop-variant conditions
 - vectorized blend operations are much faster than unrolled if-statements
- inlined functions must use SESE style
 - otherwise the inliner generates gotos
- functions configured for direct vectorization are not inlined
 - allows for integration of user-implemented vector code



SIMD is a rather cheap way to exploit data-parallelism

- even a simple application of Scout nets remarkable speedups
 - just augment your code with pragmas and put Scout in your makefile
- overall speedup of hybird-parallelized CFD production codes due to Scout:
 - TAU: 8 10 %
 - TRACE: 20 80 % (1-Proc/Node AVX/Sandy Bridge EP)
- out-of-the-box performance results:
 - exploit hardware development without programming effort
- great flexibility by using a source-to-source approach



Future

- C++ support:
 - improved function inlining
 - will require -fno-access-control or the like
- folllow-up project waits in the wings:
 - investigate automatic data layout transformations



Questions?

Code Snippets?

http://scout.zih.tu-dresden.de/



