DCatch: Automatically Detecting Distributed Concurrency Bugs in Cloud Systems

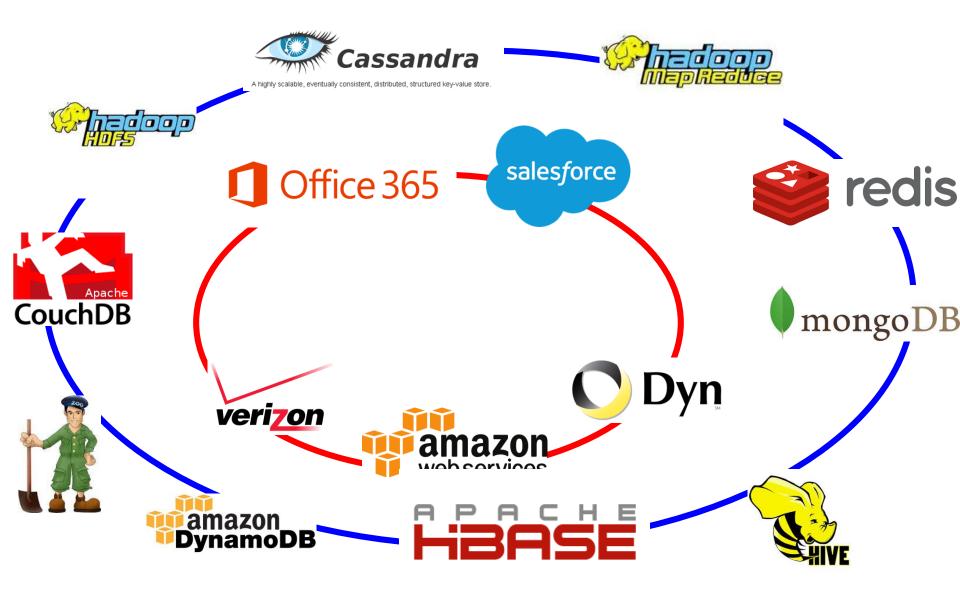
Haopeng Liu, Guangpu Li, Jeffrey Lukman, Jiaxin Li, Shan Lu, Haryadi Gunawi, and Chen Tian*







Cloud systems





Cloud systems







The 10 Biggest Cloud Outages of 2016 redis



www.crn.com/slide-shows/cloud/.../the-10-bigges

Dec 29, 2016 - The **biggest** cloud **outages** of **2016** incl



large, are increasingly vulnerable from downtime.







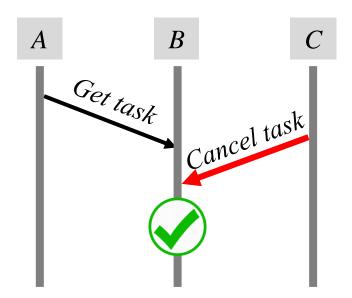




Unexpected timing among distributed operations

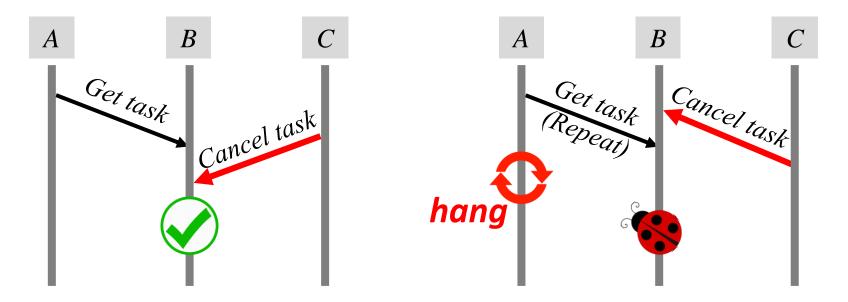


- Unexpected timing among distributed operations
- Example





- Unexpected timing among distributed operations
- Example



MapReduce-3274



- Common in distributed systems [1, 2, 3]
 - 26% failures caused by non-deterministic [1]
 - 6% software bugs in clouds system [2]

^[1] Yuan. Simple Testing Can Prevent Most Critical Failures. In OSDI'14

^[2] Gunawi. What Bugs Live in the Cloud?. In SoCC'14

^[3] Leesatapornwongsa. TaxDC. In ASPLOS'16



- Common in distributed systems [1, 2, 3]
- Difficult to avoid, expose and diagnose

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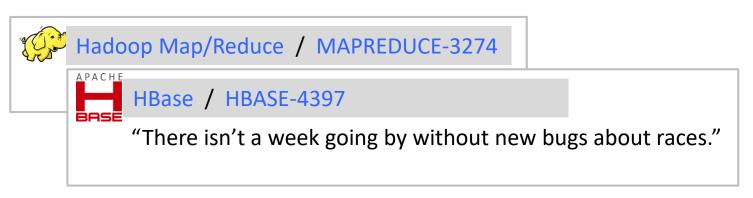


"That is one monster of a race!"

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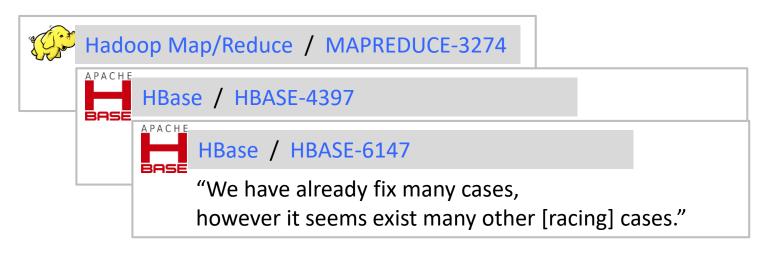
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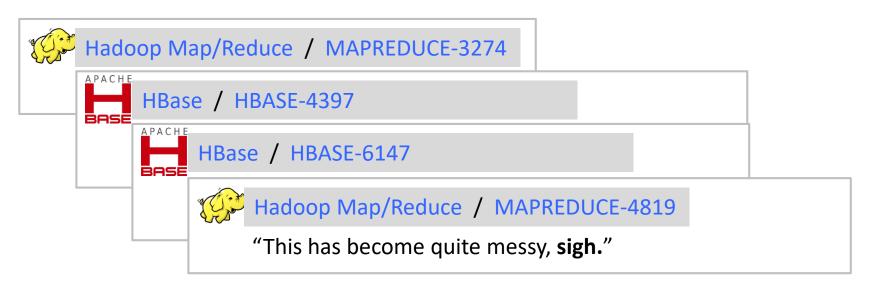
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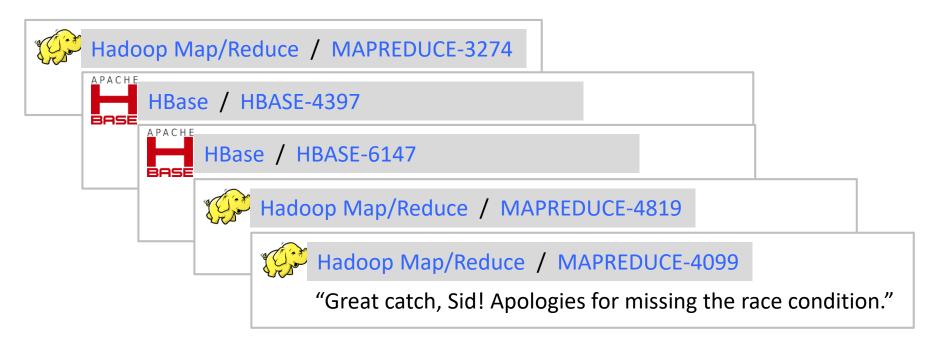
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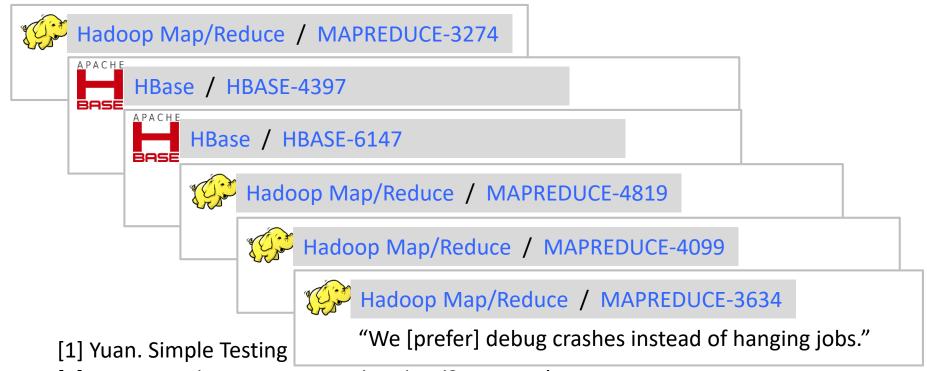
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- Common in distributed systems [1, 2, 3]
- Difficult to avoid, expose and diagnose



Can we detect DCbugs before they manifest?

Hadoop Map/Reduce / MAPREDUCE-3634

[1] Yuan. Simple Testing

"We [prefer] debug crashes instead of hanging jobs."

- [2] Gunawi. What Bugs Live in the Cloud?. In SoCC'14
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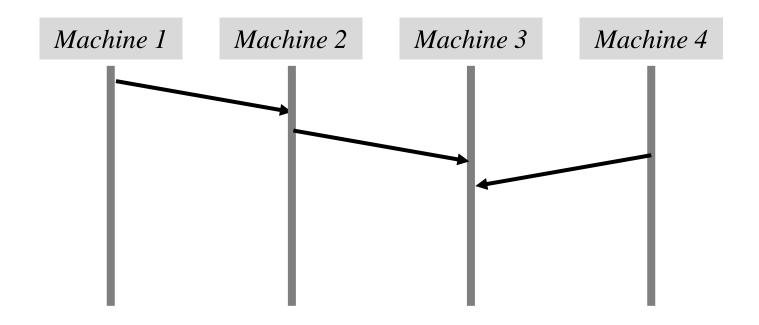


Previous work

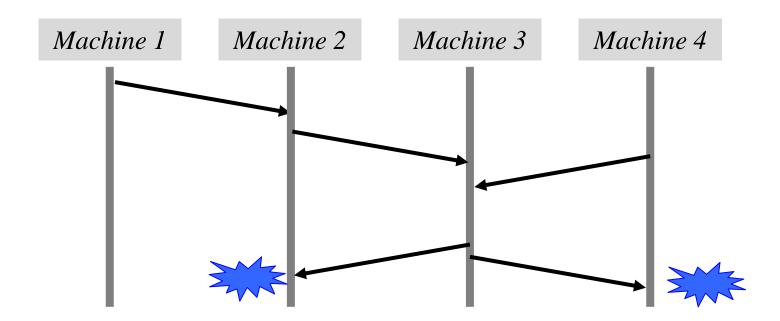
- Model checking
 - Work on abstracted models
 - Face state-space explosion issue



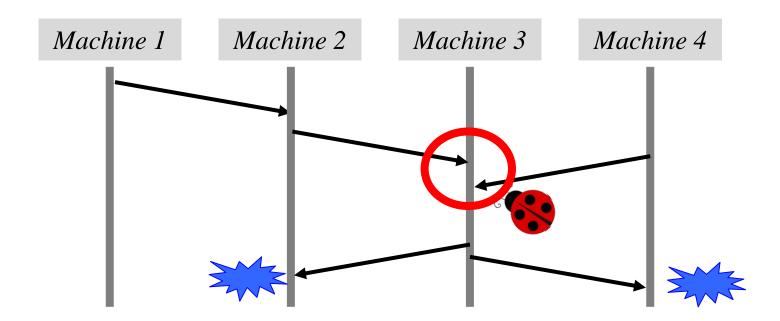






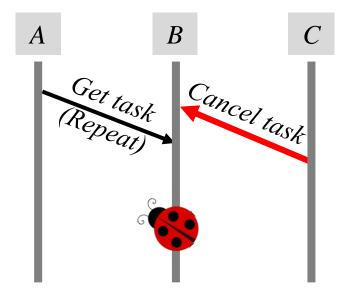






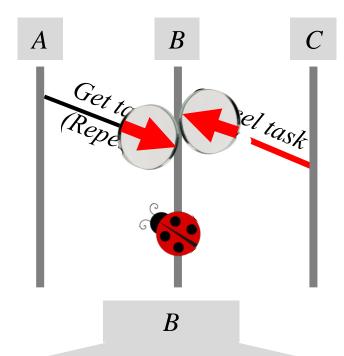


Example





Example







ASPLOS Architectural Support for Programming

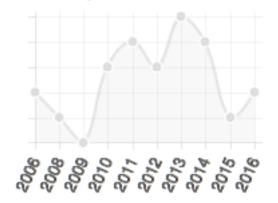
ASPLOS is a multi-disciplinary conference for rese compilers, languages, operating systems, network and engineers to present their latest research fine computer systems innovations of the past two de multiprocessors, clusters and networks-of-workst

This conference occurs at a time when computer processor performance scaling and to new demar increasingly important as boundaries between ha capabilities of computing devices beco expand

Search within ASPLOS: "concurrency bug" "race"

35 results found

Refine by Publication Year



Published Since 2006





ASPLOS Architectural Support for Programming

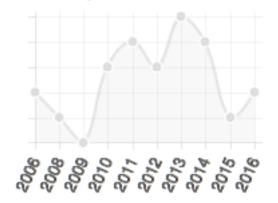
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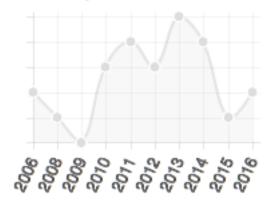
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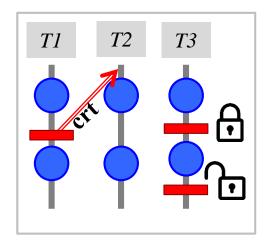
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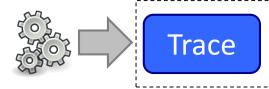


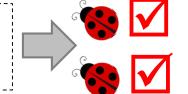
Published Since 2006

Is the problem solved?

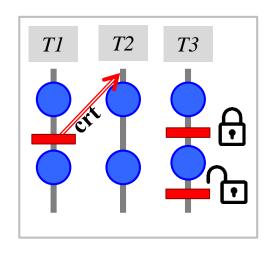


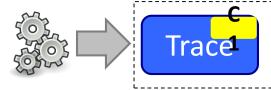


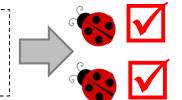








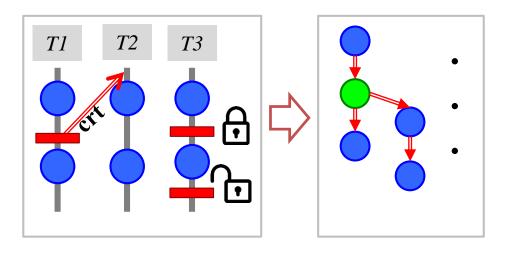




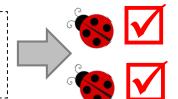
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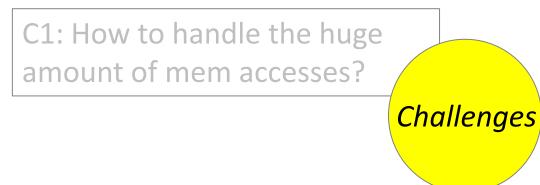
Challenges



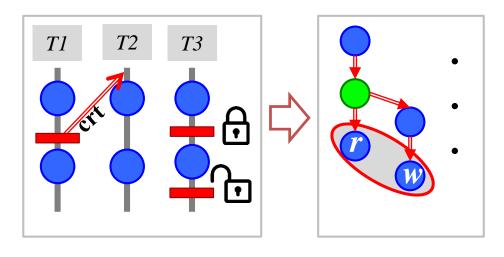




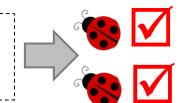








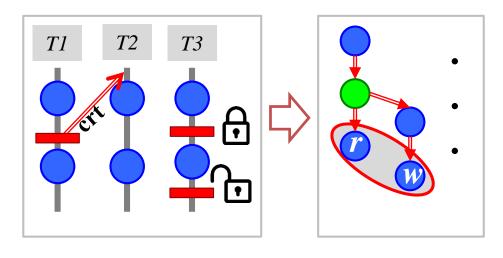




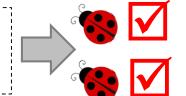
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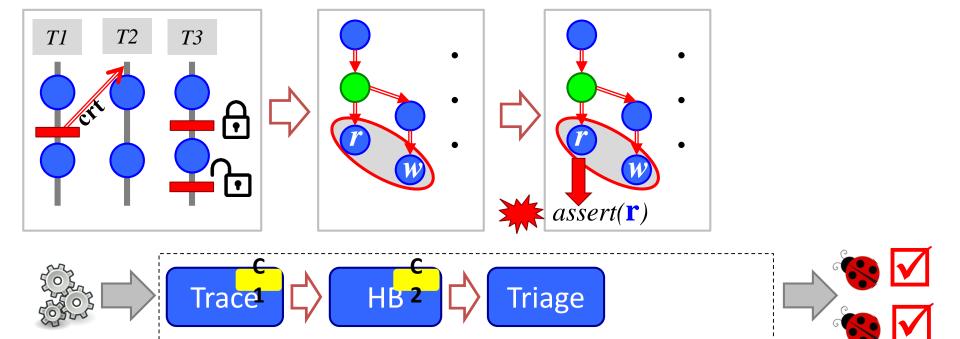


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C2: What's the happens-before model?

Challenges



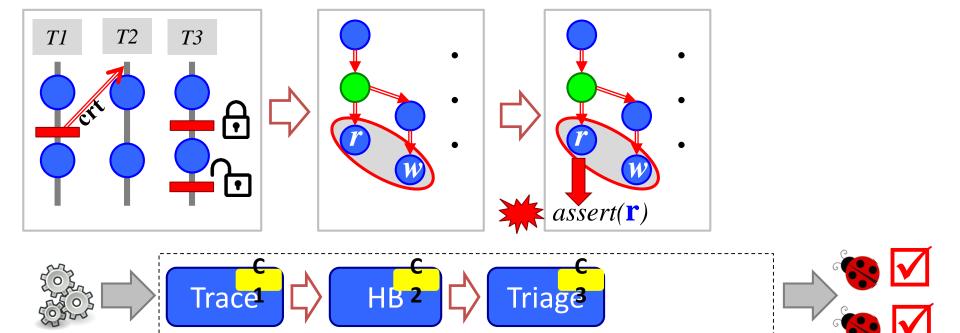


C1: How to handle the huge amount of mem accesses?

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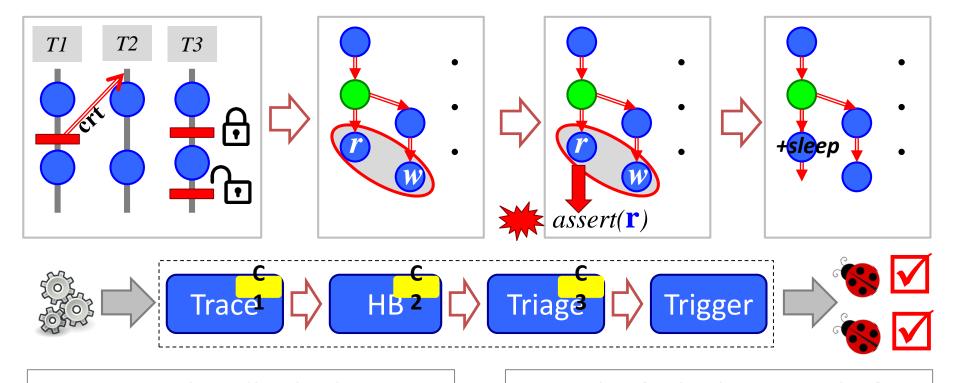
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Challenges

C3: How to estimate the distributed impact of a race?





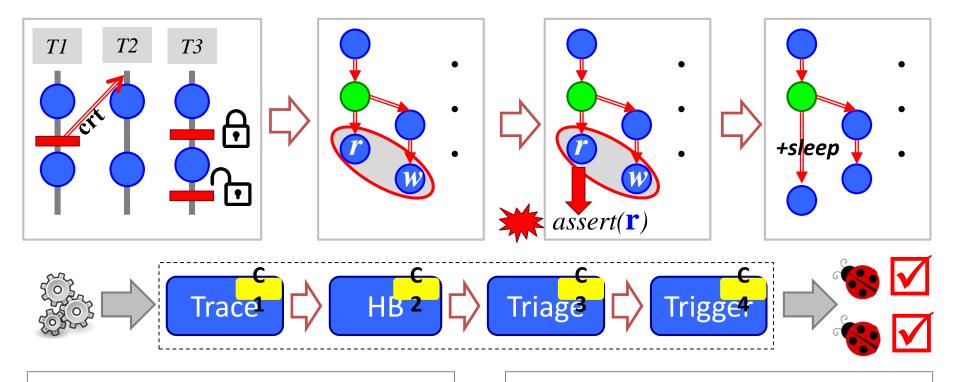
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Challenges

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C4: How to trigger with distributed time manipulation?



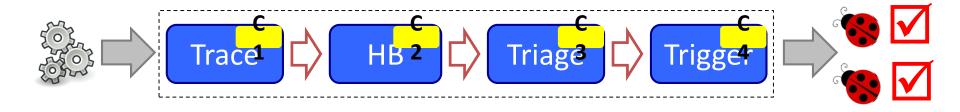
Contribution

A comprehensive HB Model for distributed systems



Contribution

- A comprehensive HB Model for distributed systems
- DCatch tool detects DCbugs from correct runs



C1: How to handle the huge amount of mem accesses?

ChallengesSolved by

C2: What's the happens-before model?

C3: How to estimate the distributed impact of a race?

DCatch C4: How to trigger with

distributed time manipulation?



Contribution

- A comprehensive HB Model for distributed systems
- DCatch tool detects DCbugs from correct runs
- Evaluate on 4 systems









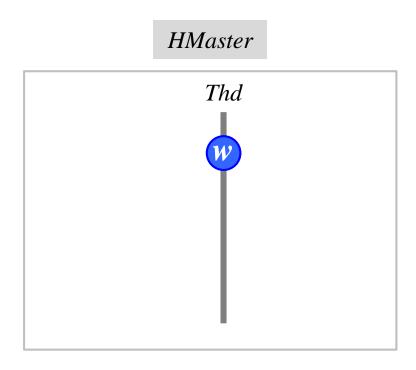
Report 32 DCbugs, with 20 of them being truly harmful



Outline

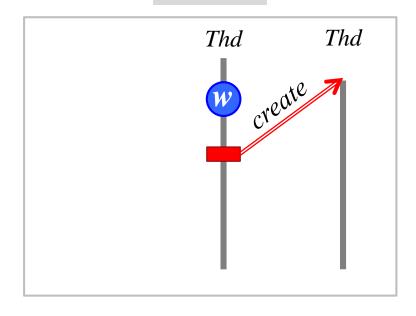
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- Conclusion



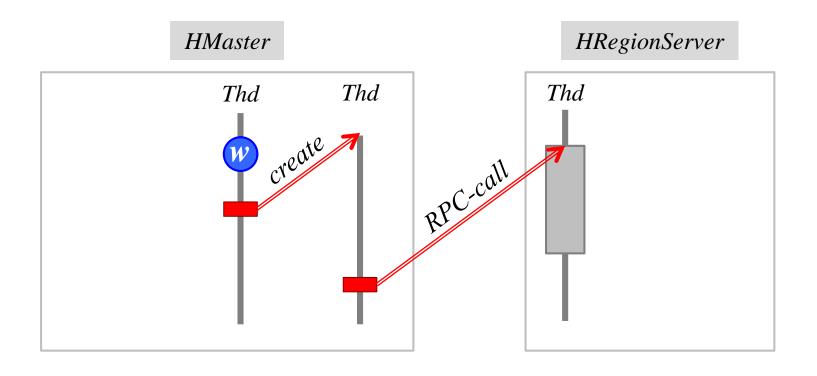




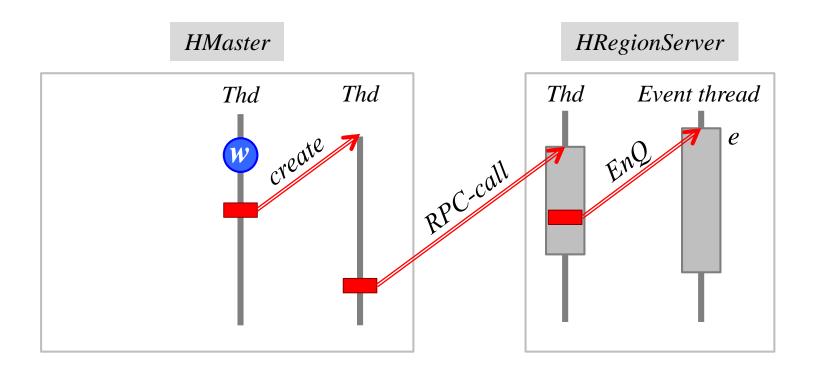
HMaster



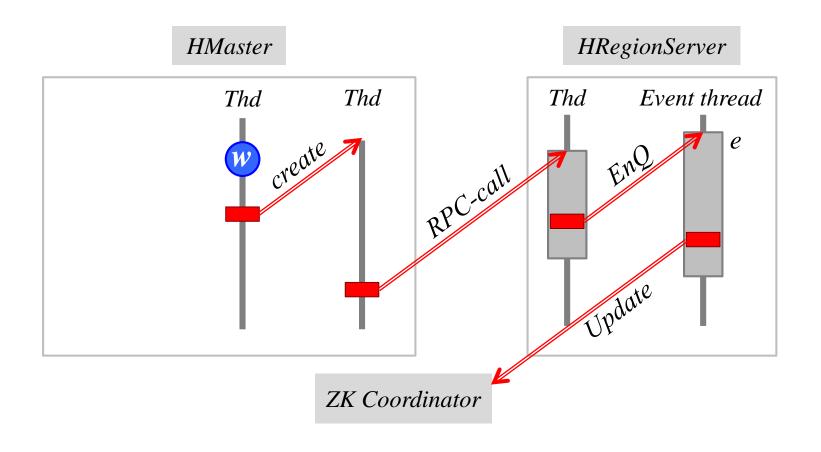




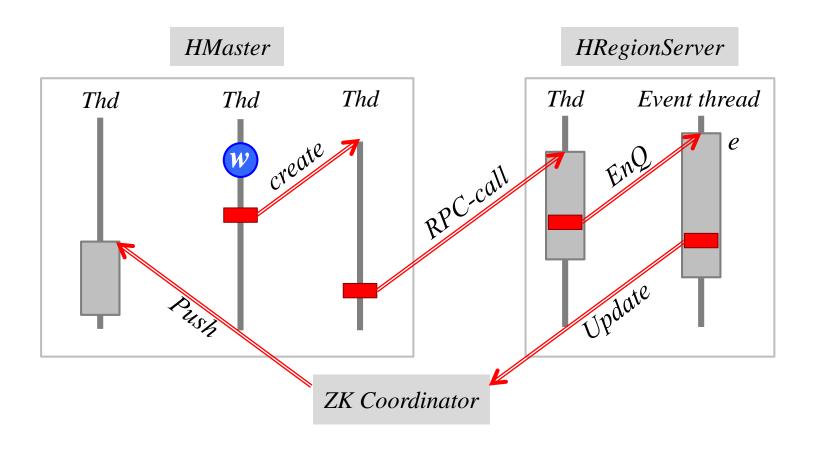




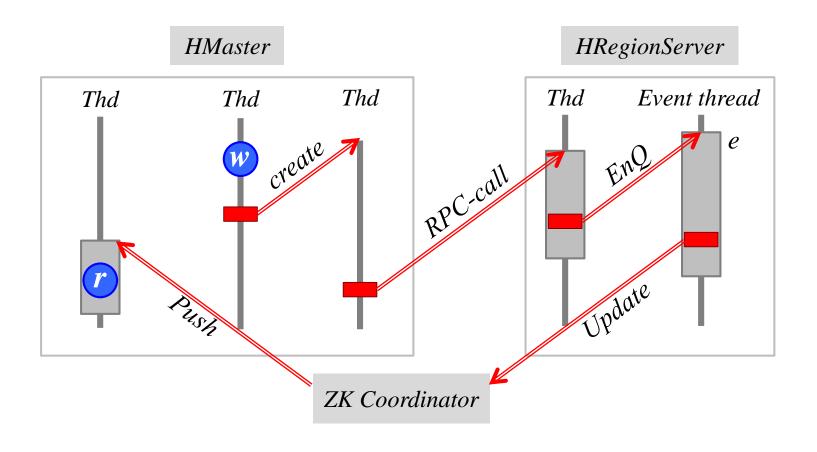




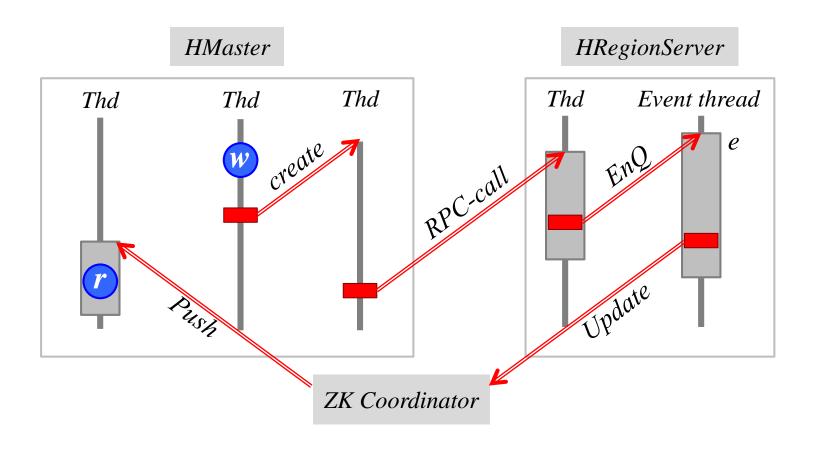






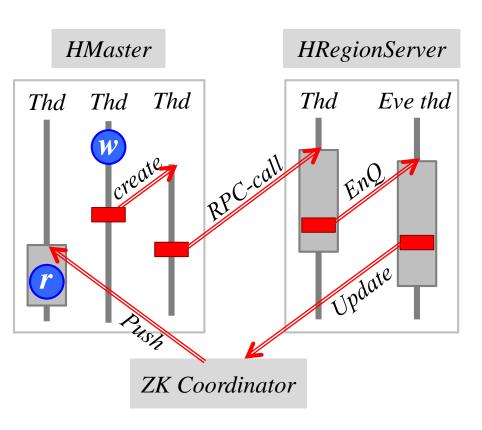




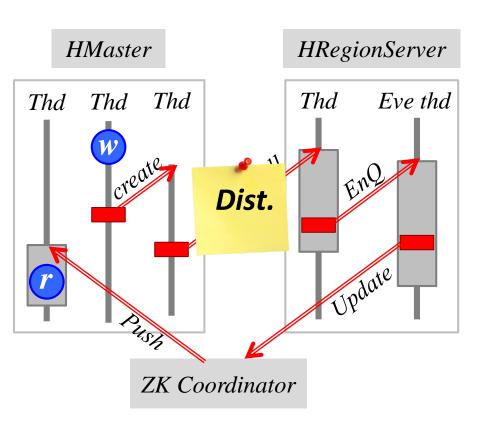


Where is HB model for distributed systems?



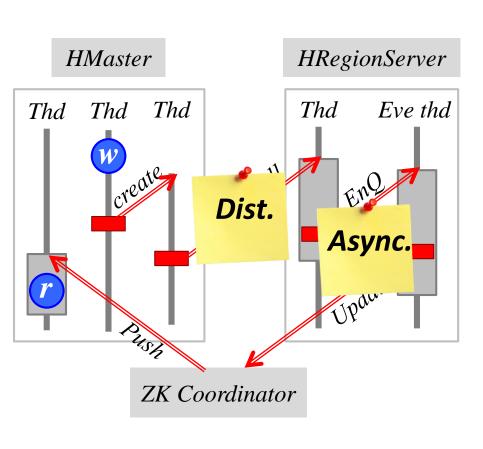


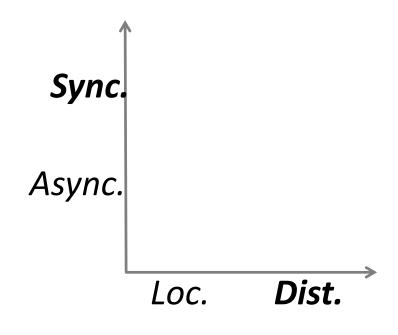




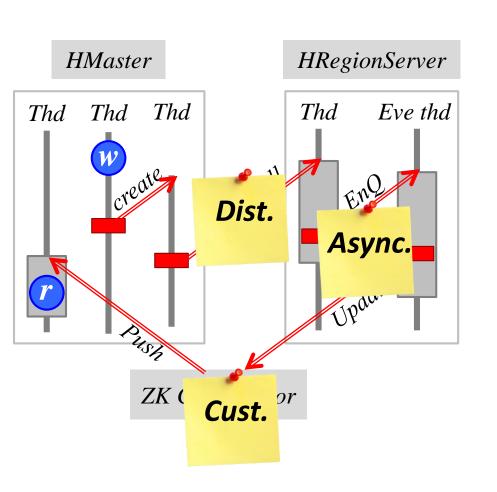
Loc. **Dist.**

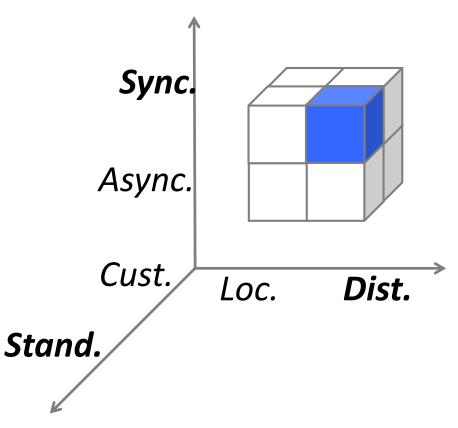




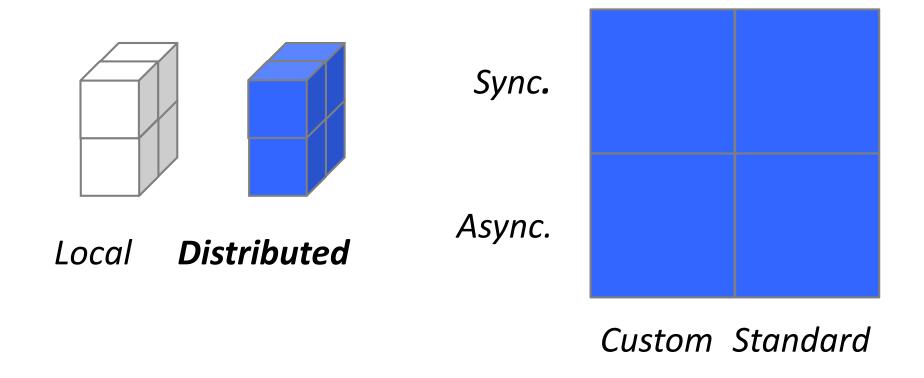




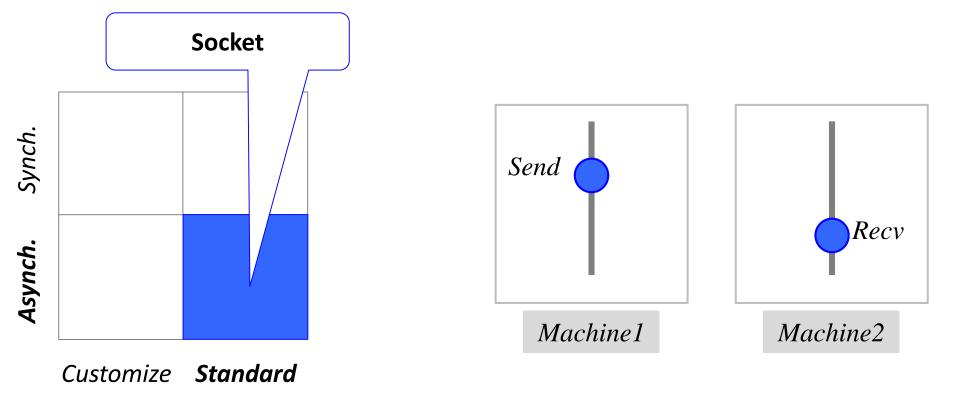






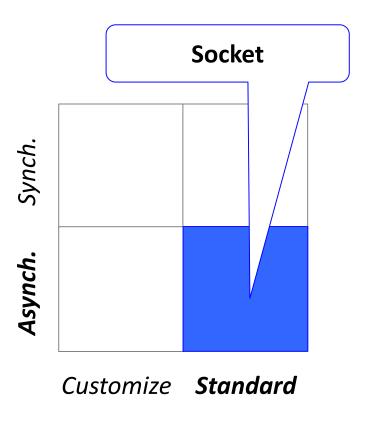


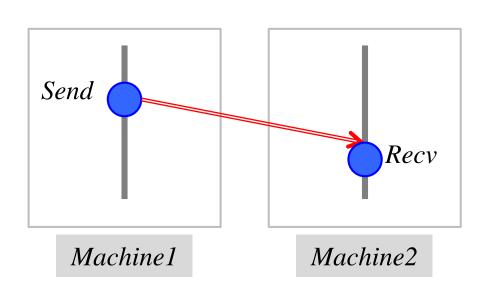




Logical time clock (Leslie Lamport, 1978)

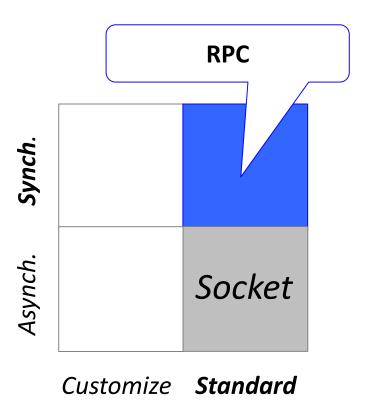


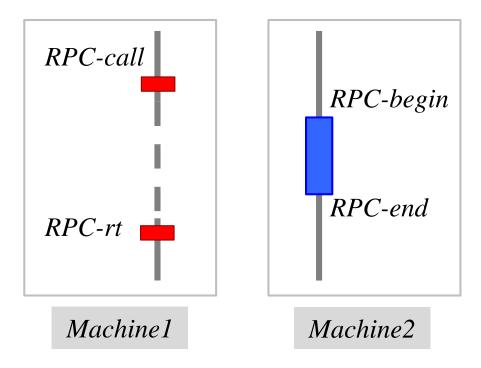




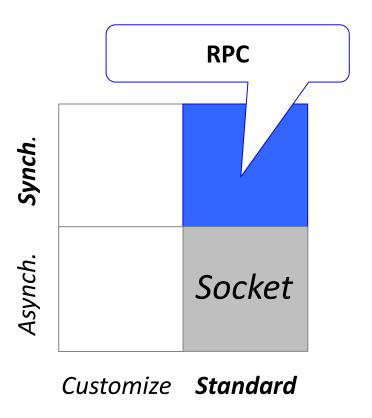
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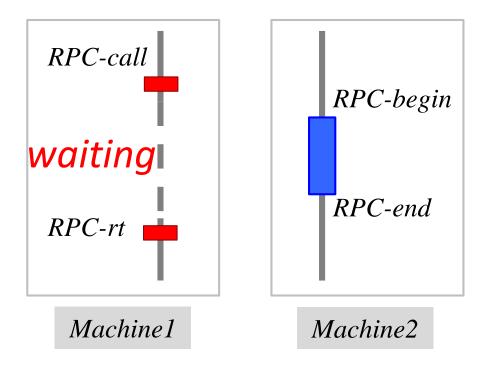




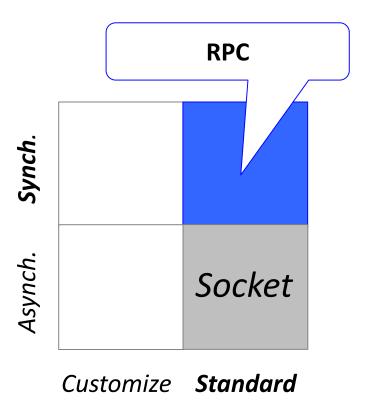


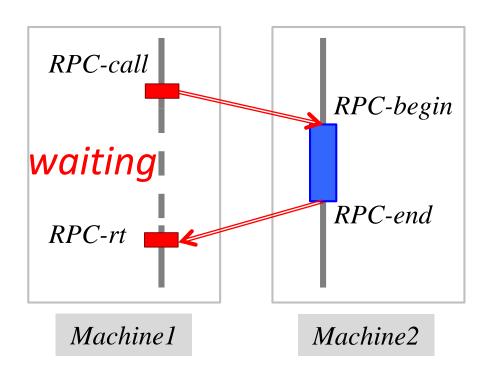




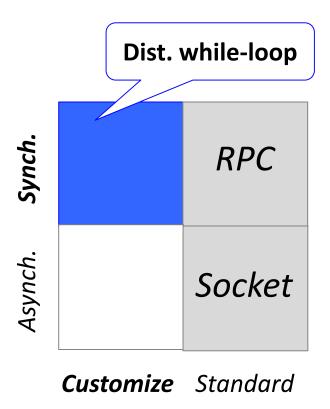






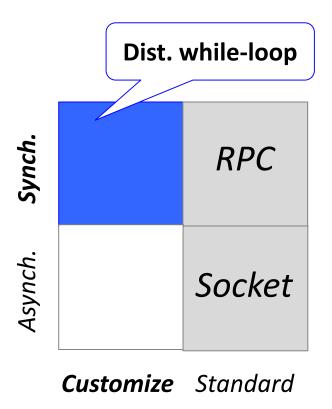






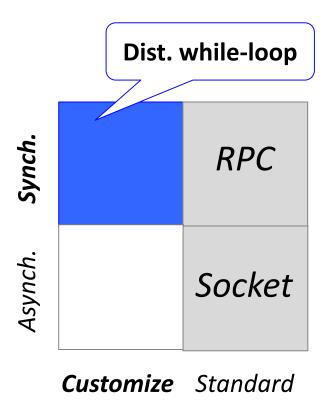
In multi-threaded systems:





In multi-threaded systems:



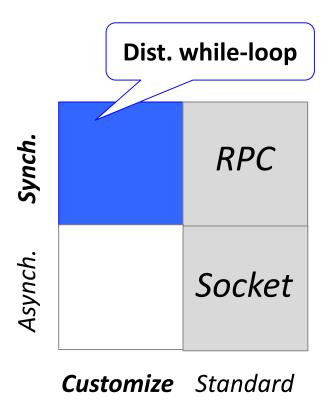


In multi-threaded systems:

In distributed systems:







In distributed systems:

Machine A

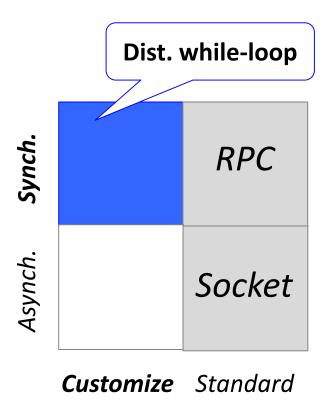
Machine B

```
//Thread1
flag = True;

//Thread2
bool getFlag() {
  return flag;
}
```

```
//Thread
while(!getFlag()){
    ...
}
```





In distributed systems:

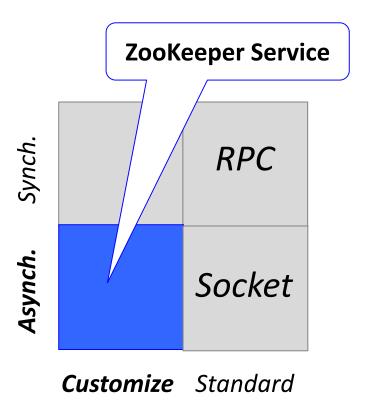
```
Machine A

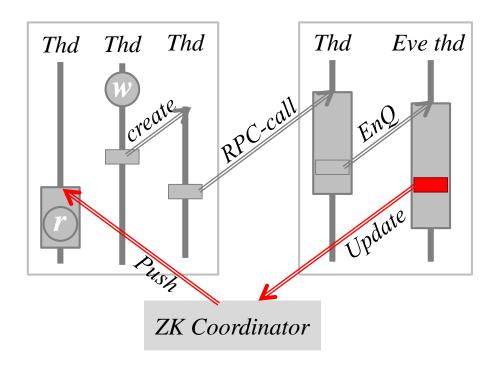
Machine B

//Thread1
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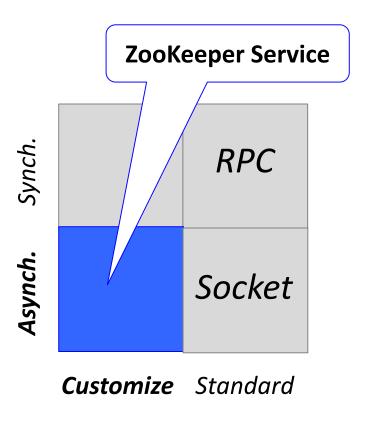
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```

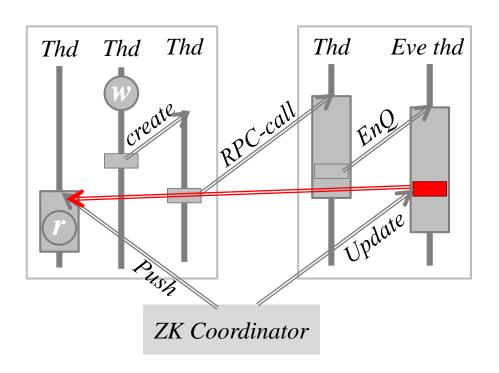








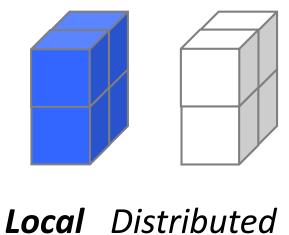




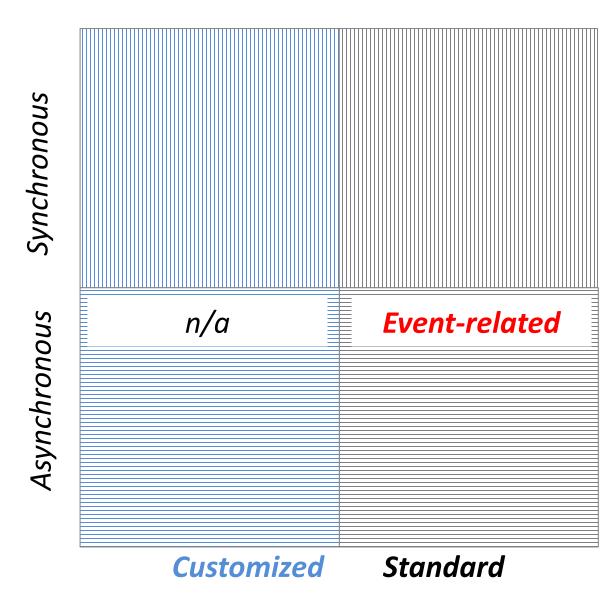


Dist. While-loop **RPC** Synchronous Asynchronous Zookeeper Service Socket Standard **Customized**

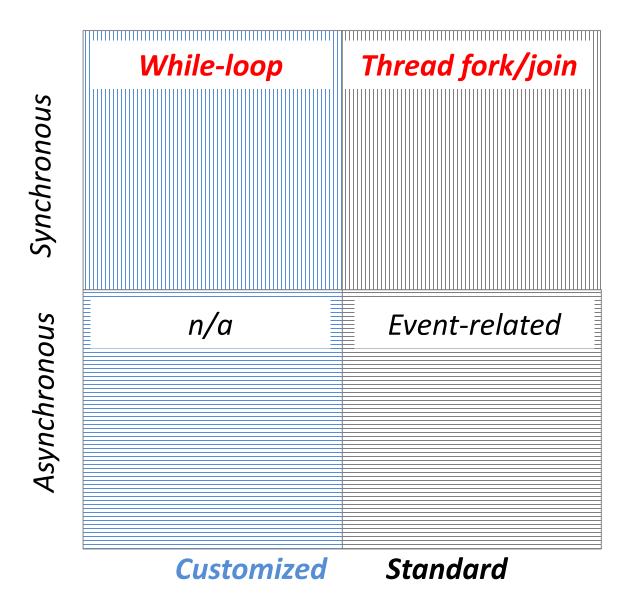








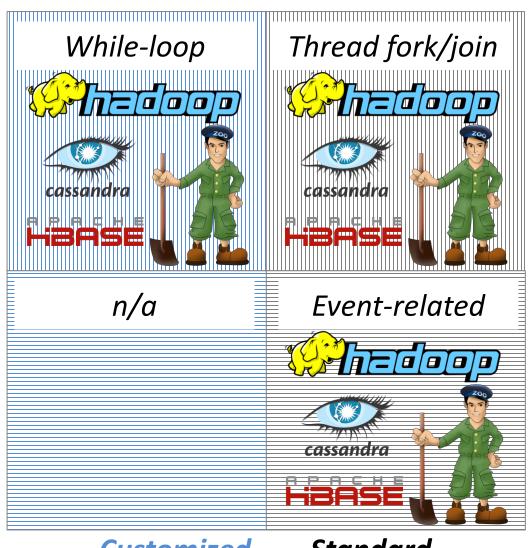






Synchronous

Asynchronous



Customized

Standard

Outline

- Motivation
- DCatch Happens-before Model
- DCatch tool
- Evaluation
- Conclusion





Trace

HB

Triage Trigger

Challenges

C1: How to handle the huge amount of mem accesses?

C2: What's the happens-before model?

C3: How to estimate the distributed impact of a race?

C4: How to trigger with distributed time manipulation?

Selective tracing: only mem accesses in Event/message handlers and their callee.



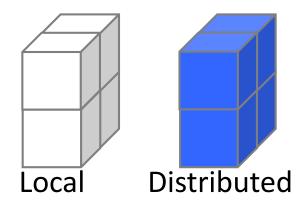
Trace HB Triage Trigger

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HB

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HB

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Machine A

```
while(!getTask(jID)) {
}
```



HB

Triage Trigger

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```
while (!getTask (jID)) {
    hang
```



HB

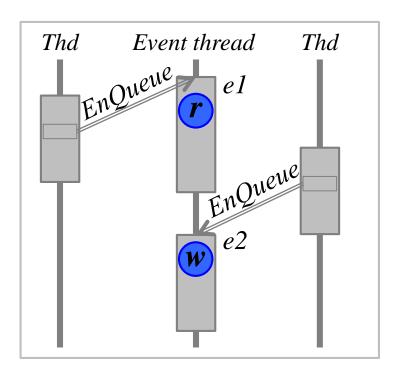
Triage Trigger

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HB Triage

Trigger

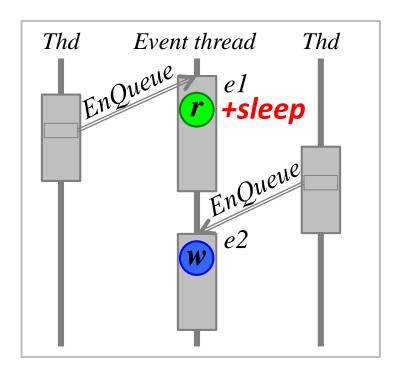
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C2: What's the happens-before model?

C4: How to trigger with distributed time manipulation?





HB

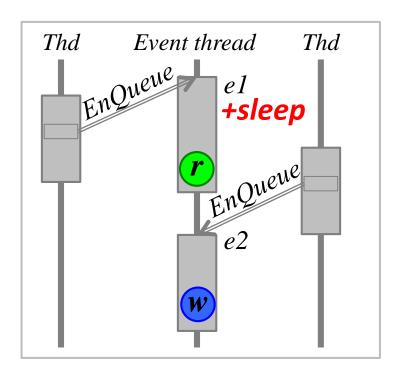
Triage Trigger

C1: How to handle the huge amount of mem accesses?

C3: How to estimate the distributed impact of a race?

C2: What's the happens-before model?

C4: How to trigger with distributed time manipulation?







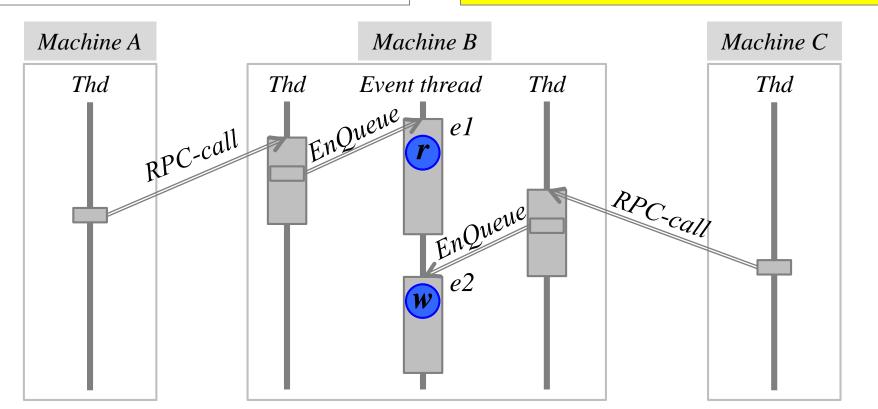
Trace HB Triage Trigger

C1: How to handle the huge amount of mem accesses?

C3: How to estimate the distributed impact of a race?

C2: What's the happens-before model?

C4: How to trigger with distributed time manipulation?





HB Triage

Trigger

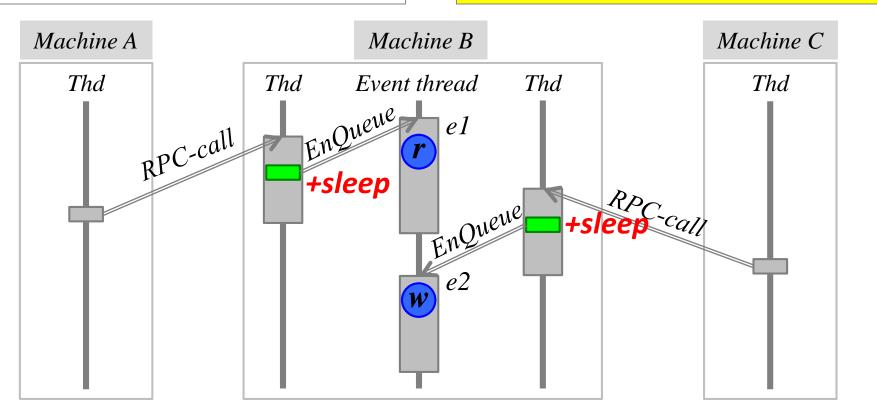
Challenges

C1: How to handle the huge amount of mem accesses?

C3: How to estimate the distributed impact of a race?

C2: What's the happens-before model?

C4: How to trigger with distributed time manipulation?





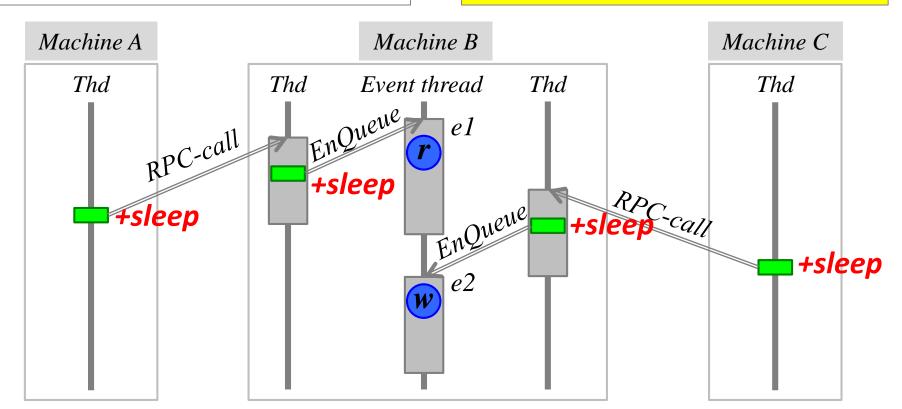
Trace HB Triage Trigger

C1: How to handle the huge amount of mem accesses?

C3: How to estimate the distributed impact of a race?

C2: What's the happens-before model?

C4: How to trigger with distributed time manipulation?



Outline

- Motivation
- DCatch Happens-before Model
- DCatch tool
- Evaluation
- Conclusion



Methodology

- Benchmarks:
 - 7 real-world DCbugs from TaxDC[1]
 - 4 distributed systems











BugID	Detected?	#. Bugs	#. Benign	#. false-pos
CA-1011	✓	3	0	0
HB-4539		3	0	1
HB-4729		4	1	0
MR- 3274		2	0	4
MR- 4637		1	2	4
ZK-1144		5	1	1
ZK-1270	✓	6	2	0



BugID	Detected?	#. Bugs	#. Benign	#. false-pos
CA-1011		3	0	0
HB-4539		3	0	1
HB-4729		4	1	0
MR- 3274		2	0	4
MR- 4637		1	2	4
ZK-1144		5	1	1
ZK-1270		6	2	0



BugID	Detected?	#. Bugs	#. Benign	#. false-pos
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MR- 3274		2	0	4
MR- 4637		1	2	4
ZK-1144		5	1	1
ZK-1270		6 = 1	2 + 8 2	0



BugID	Detected?	#. Bugs	#. Benign	#. false-pos
CA-1011		3	0	0
HB-4539		3	0	1
HB-4729		4	1	0
MR- 3274		2	0	4
MR- 4637		1	2	4
ZK-1144		5	1	1
ZK-1270		6	2	0



Other results in our paper

- Performance overhead
- Trace compositions
- HB model impact
 - False-positive
 - False-negatives

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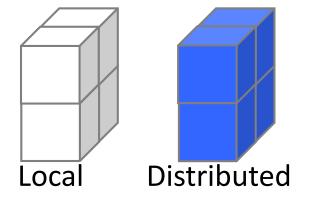
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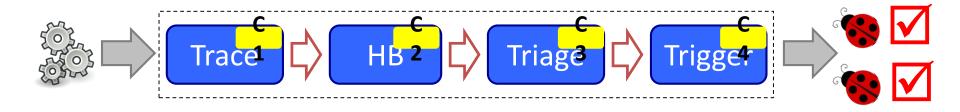


Conclusion

A HB Model for distributed systems



 DCatch detects DCbugs from correct runs with low false positive rates.



Thank you! Q&A