SM2 数字签名算法实验报告

SM2 是中国密码管理局发布的椭圆曲线公钥密码标准算法,包含数字签名、密钥交换和公钥加密功能。本实验聚焦其数字签名功能,完整实现了密钥生成、签名生成和签名验证流程。

参数选择:

- SM2 system parameters:
 - \mathbb{F}_q : finite field where $|\mathbb{F}_q| = q$
 - a, b: elliptic curve equation parameters
 - $G = (x_G, y_G)$: base point
 - *n*: order
 - h: cofactor where $h = |E(F_a)|/n$
- SM2 system parameters: prime field
 - Elliptic curve equation: $y^2 = x^3 + ax + b$ over \mathbb{F}_{a} -256
 - Prime q: 8542D69E 4C044F18 E8B92435 BF6FF7DE 45728391 5C45517D 722EDB8B 08F1DFC3
 - a: 787968B4 FA32C3FD 2417842E 73BBFEFF 2F3C848B 6831D7E0 EC65228B 3937E498
 - b: 63E4C6D3 B23B0C84 9CF84241 484BFE48 F61D59A5 B16BA06E 6E12D1DA 27C5249A
 - $G = (x_G, y_G)$, ord(G) = n
 - x_G: 421DEBD6 1B62EAB6 746434EB C3CC315E 32220B3B ADD50BDC 4C4E6C14 7FEDD43D
 - y_G: 0680512B CBB42C07 D47349D2 153B70C4 E5D7FDFC BFA36EA1 A85841B9 E46E09A2
 - n:8542D69E 4C044F18 E8B92435 BF6FF7DD 29772063 0485628D 5AE74EE7 C32E79B7

椭圆曲线基础运算

def modular_inverse(value, modulus):

```
"""计算模逆元(扩展欧几里得算法)"""

mult_low, mult_high = 1, 0

val_low, val_high = value % modulus, modulus

while val_low > 1:

    ratio = val_high // val_low

    new_mult = mult_high - mult_low * ratio

    new_val = val_high - val_low * ratio

    mult_low, val_low, mult_high, val_high = new_mult, new_val, mult_low, val_low

return mult_low % modulus
```

计算有限域内的模逆元,使用扩展欧几里得算法实现:初始化乘数因子和余数,通过迭代 计算余数和乘数,当余数减至1时返回乘数因子

椭圆曲线点运算

```
def ecc_point_addition(point1, point2):
    """椭圆曲线两点加法"""
    slope = ((point2[1] - point1[1]) *
              modular_inverse(point2[0] - point1[0], MODULUS P)) % MODULUS P
    result_x = (slope * slope - point1[0] - point2[0]) % MODULUS_P
    result_y = (slope * (point1[0] - result_x) - point1[1]) % MODULUS_P
    return (result x, result y)
算法原理:
计算通过两点的斜率 λ
利用斜率求新点 x 坐标: x_3 = \lambda^2 - x_1 - x_2
计算新点 y 坐标: y_3 = \lambda(x_1 - x_3) - y_1
def ecc_point_doubling(point):
    slope = ((3 * point[0] * point[0] + ECC_PARAM_A) *
              modular_inverse((2 * point[1]), MODULUS_P)) % MODULUS_P
    result_x = (slope * slope - 2 * point[0]) % MODULUS_P
    result y = (slope * (point[0] - result x) - point[1]) % MODULUS P
    return (result_x, result_y)
    点加自身时斜率为(3x_1^2 + a)/2y_1,其他计算类似点加法
    标量乘法
    def scalar_point_multiplication(scalar, base_point):
        """标量乘法(倍点运算)"""
        if scalar == 0 or scalar >= ORDER N:
             raise ValueError("Invalid scalar or private key")
        scalar_binary = bin(scalar)[2:]
        temp_point = base_point
```

```
for bit in scalar_binary[1:]:
    temp_point = ecc_point_doubling(temp_point)
    if bit == "1":
        temp_point = ecc_point_addition(temp_point, base_point)
```

return temp_point

实现策略:

- 1. 将标量转换为二进制格式
- 2. 使用双倍-加法优化算法
- 3. 首位后处理: 每次双倍后根据二进制位决定是否加点
- 4. 安全验证: 检查标量范围 (0 < k < n)

用户标识哈希计算

```
def compute_identity_hash(user_id, pub_x, pub_y):

"""计算用户身份哈希值 Z_A"""

components = [

    str(calculate_bit_length(user_id)),
    user_id,
    str(ECC_PARAM_A),
    str(ECC_PARAM_B),
    str(BASE_PT_X),
    str(pab_x),
    str(pub_x),
    str(pub_y)

]

concatenated = "".join(components)

digest = sm3.sm3_hash(func.bytes_to_list(concatenated.encode()))
```

return int(digest, 16)

- 1. 按标准要求拼接参数: ID 长度、曲线参数、基点、公钥
- 2. 使用 SM3 哈希算法处理拼接后的字符串
- 3. 转换为整数值用于后续签名

密钥对生成

```
def create_key_pair():
   """生成 SM2 密钥对"""
   private_val = int(secrets.token_hex(32), 16) % ORDER_N
   public_pt = scalar_point_multiplication(private_val, BASE_POINT)
   return private_val, public_pt
使用 secrets 模块生成密码学安全的随机数
限制私钥范围在[1, n-1]之间
通过标量乘法从私钥导出公钥
签名生成算法
def generate_signature(private_key, msg_content, user_hash):
```

```
"""生成数字签名"""
merged data = str(user hash) + msg content
msg_bytes = merged_data.encode()
digest_hash = sm3.sm3_hash(func.bytes_to_list(msg_bytes))
hash_int = int(digest_hash, 16)
k_input = str(private_key) + sm3.sm3_hash(func.bytes_to_list(msg_content.encode()))
k_val = int(sha256(k_input.encode()).hexdigest(), 16)
if k val >= MODULUS P:
    return None
```

```
temp_pt = scalar_point_multiplication(k_val, BASE_POINT)
    r_val = (hash_int + temp_pt[0]) % ORDER_N
    s_val = modular_inverse(1 + private_key, ORDER_N) * (k_val - r_val * private_key) %
ORDER N
    return (r_val, s_val)
   1. 组合用户哈希和消息, 计算 SM3 哈希 e
   2. 使用 RFC6979 方法生成确定性的 k 值
   3. 计算椭圆曲线点(k·G) = (x<sub>1</sub>, y<sub>1</sub>)
   4. 生成签名分量 r = (e + x<sub>1</sub>) mod n
   5. 计算签名分量 s = ((1 + d)^{-1} \cdot (k - r \cdot d)) \mod n
签名验证算法
def verify_signature(public_key, user_id, msg_content, signature):
    """验证数字签名"""
    r val, s val = signature
    user_hash = compute_identity_hash(user_id, public_key[0], public_key[1])
    merged_data = str(user_hash) + msg_content
    msg_bytes = merged_data.encode()
    digest_hash = sm3.sm3_hash(func.bytes_to_list(msg_bytes))
    hash int = int(digest hash, 16)
    composite_val = (r_val + s_val) % ORDER_N
    point1 = scalar_point_multiplication(s_val, BASE_POINT)
    point2 = scalar_point_multiplication(composite_val, public_key)
    result_pt = ecc_point_addition(point1, point2)
    R_val = (hash_int + result_pt[0]) % ORDER_N
```

return R val == r val

- 1. 重新计算用户标识哈希 Z_A
- 2. 组合 Z_A 和消息计算哈希值 e'
- 3. 计算 t = (r + s) mod n
- 4. 计算椭圆曲线点: s·G + t·P_A
- 5. 比对 R' = (e' + x₁') mod n 与 r 是否相等

测试

基本实现了 sm2 的基础功能