## Ch 8 Textbook Problems

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8.1 2.	a) $P_n = n P_{n-1}$			n	Pn
	Note Po = 1 blc one permutation of	a set w	/	0	= [
	no objects (the empty sequence)				
	$b) P_n = n P_{n-1}$			052	Con 1
	$= n(n-1) P_{n-2}$				3 · 2 · 1
	=				4.3.2.1
	$= n(n-1) \cdots 2 \cdot 1 \cdot P_0$			5	5.4.3.2.1
	= n!				

3. one-dollar coins, \$1 bills, \$5 bills

a) Let an be the # oP ways to deposit n dollars in the vending machine. We must express an in terms of earlier terms in the sequence. If we want to deposit n dollars, we may start w/ a dollar coin and then deposit n-1 dollars. This gives us an -1 ways to deposit n dollars. We can also start w/ a dollar bill and then deposit n-1 dollars. This gives us an -1 more ways to deposit n dollars. Finally, we can deposit a \$5 bill and Pollow that w/ n-5 dollars; there are an-5 ways to do this.

The recurrence relation is an = 2an-1 + an-5. Note this is valid for  $n \geq 5$ , since otherwise an-5 makes no sense.

b) we need initial conditions for all subscripts from 0 to 4. It is clear that  $a_0=1$  (deposit nothing) and  $a_1=2$  (deposit either the dollar coin or the dollar bill). It is also not hard to see that  $a_2=2^2=4$ ,  $a_3=2^3=8$ , and  $a_4=2^4=16$ , since each sequence of a C's and B's corresponds to a way to deposit a dollars -a C meaning to deposit a coin and a B meaning to deposit a bill.

c) 
$$\alpha_{10} = 2\alpha_{10-1} + \alpha_{10-5}$$
  
=  $2\alpha_{q} + \alpha_{5}$   
=  $2(2\alpha_{q-1} + \alpha_{q-5}) + 2\alpha_{5-1} + \alpha_{5-5}$   
=  $2(2\alpha_{8} + \alpha_{4}) + 2\alpha_{4} + \alpha_{6}$   
=  $2(2(2\alpha_{8-1} + \alpha_{8-5}) + 16) + 2(16) + 1$   
=  $2(2(2(2\alpha_{q-1} + \alpha_{q-5}) + \alpha_{3}) + 16) + 33$   
=  $2(2(2(2\alpha_{q-1} + \alpha_{q-5}) + \alpha_{3}) + 16) + 33$   
=  $2(2(2\alpha_{q-1} + \alpha_{q-5}) + \alpha_{q-5}) + \alpha_{q-5}$ 

$$= 2(8\alpha_{6} + 16 + 16 + 16 + 16)$$

$$= 16\alpha_{6} + 96 + 33$$

$$= 16\alpha_{6} + 129$$

$$= 16(2\alpha_{6-1} + \alpha_{6-5}) + 129$$

$$= 32\alpha_{5} + 32 + 129$$

$$= 32(2\alpha_{5-1} + \alpha_{5-5}) + 161$$

$$= 32(32 + 1) + 161$$

$$= 1056 + 161$$

$$= 1056 + 161$$

7. a) Let an be the # of bit strings of length n containing a pair of consecutive 0's.

Case 1) starts with 1

Case 2 starts with O

 $\begin{array}{c}
0 \\
1 \\
-2 \\
contains \\
00
\end{array}$ So  $a_{n-2}$ 

n-2

any bit string

so 2<sup>n-2</sup>

) † 33

Thus, the recurrence relation is

$$a_n = a_{n-1} + a_{n-2} + 2^{n-2}$$
 for  $n \ge 2$ 

b) a = 0, a = 0

ble there are no bit strings of length O or I that contain OD

c) 
$$a_2 = a_1 + a_0 + 2^0 = 0 + 0 + 1 = 1$$
  
 $a_3 = a_2 + a_1 + 2^1 = 1 + 0 + 2 = 3$   
 $a_4 = a_3 + a_2 + 2^2 = 3 + 1 + 4 = 8$   
 $a_5 = a_4 + a_3 + 2^3 = 8 + 3 + 8 = 19$   
 $a_6 = a_5 + a_4 + 2^4 = 19 + 8 + 16 = 43$   
 $a_7 = a_6 + a_5 + 2^5 = 43 + 19 + 32 = 94$ 

11. a) Let an be the # of ways to climb n stairs. In order to climb n states, a person most either start w/ a step of one stair and then climb n-1 stairs (and this can be done in un-1 ways) or else start w/ a step of two stales and then climb n-2 stairs (and this can be done in any ways ). From this analysis we can immediately write down the recurrence relation, valid Per all n = 2: an = an - + an - 2. b) ao = 1 one way to climb no stairs (do nothing)

a, = 1 one way to climb one stair

Note that the recurrence relation is the same as that for the Fibonacci sequence, and the initial conditions are that ao = f, and a, = f2, so it must be that an = fatt for all n.

c)  $a_2 = a_1 + a_0 = 1 + 1 = 2$  $\alpha_3 = \alpha_2 + \alpha_1 = 2 + 1 = 3$ Q4 = Q3 + Q2 = 3 + 2 = 5 as = ay + a3 = 5 + 3 = 8 a6 = as + a4 = 8+5=13 a7 = 06 + as = 13 +8 = 21 ax = az + a6 = 21 + 13 = 34 12. a) Let an be the # of ways to climb n string if the person climbing
the string can take one, two, or three string at a time. In order to climb
n string, a person must either strit w/ a step of one string then climb
n-1 string (this can be done in an-1 ways), strit w/a step of two
string then climb n-2 string (this can be done in an-2 ways), or else
strik w/ a step of three string than climb n-3 string (this can be
done in an-3 ways).

 $\alpha_{n} = \alpha_{n-1} + \alpha_{n-2} + \alpha_{n-3} \quad n \ge 3$ 

b) ao = 1 do nothing

a, = 1 climb one stair

az = 2 two ways to climb two starrs

c)  $a_3 = a_2 + a_1 + a_0 = 2 + 1 + 1 = 4$   $a_4 = a_3 + a_2 + a_1 = 4 + 2 + 1 = 7$   $a_5 = a_4 + a_3 + a_2 = 7 + 4 + 2 = 13$  $a_6 = a_5 + a_4 + a_3 = 13 + 7 + 4 = 24$ 

 $\alpha_7 = \alpha_6 + \alpha_5 + \alpha_4 = 24 + 13 + 7 = 44$   $\alpha_8 = \alpha_7 + \alpha_6 + \alpha_5 = 44 + 24 + 13 = 81$ 

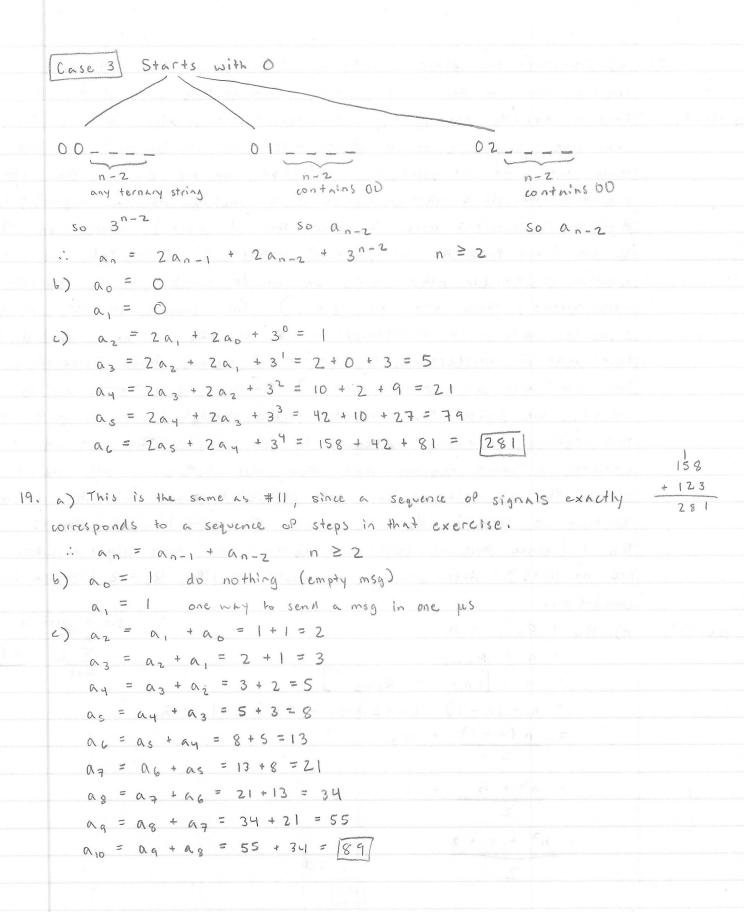
14. a) Let an be the # of ternary strings that contain two consecutive 0's.

[Case 1] starts with 1

n-1
contains 00

so an-1

50 An-1



o clarify

21. a) Consider the plane already divided by n-1 lines into Rn-1 regions. The nth line is now added, intersecting each of the other n-1 lines in exactly one point, n-1 intersections in all. Think of drawing that line beginning at one of its ends (out at "infinity"). As we move toward the 1st point of intersection, we are dividing the unbounded region of the plane through which it is passing into two regions; the division is complete once we reach the 1st point of intersection. Then as we draw from the 1st point of intersection to the 2nd, we cut off another region (in other words we divide another of the regions that were already there into two regions). This process continues as we encounter each point of intersection. By the time we have reached the just point of intersection, the number of regions have increased by n-1 (one for each point of intersection). Finally, as we more off to infinity, we divide the unbounded region through which we pass into two regions, increasing the count by yet I more. Thus there was exactly n more regions than there were before the nth line was added. The analysis we have just completed shows that the rewrence relation we seek is Rn = Rn-1 + n. The initial condition is Ro = 1 ( since there is just one region - the whole plane - when there are no lines). Alternatively, we could specify R, = 2 as the initial condition.

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b) 
$$R_n = R_{n-1} + n$$

=  $n + R_{n-1}$ 

=  $n + [(n-1) + R_{n-2}]$ 

=  $n + (n-1) + (n-2) + ... + 3 + 2 \cdot | + R_n$ 

=  $n + (n+1) + R_n$ 

=  $n + (n+1) + R_n$ 

=  $n + (n+1) + R_n$ 

 $= n^2 + n + 2$ 

7

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(R) (E) (A)
                                     - tiles of the same wolor considered indistinguishable
       red, green, gray
27. We assume that the walkway is one tile in width and n tiles
    long from start to finish. Thus we are talking about ternary sequences
    of length n that do not contain two consecutive R's. Let an represent
                                                       the desired quantity.
     Case 1) Starts with E
               documit contain RA
             50 an-1
     Case 2 starts with A
              doesn't worthin RR
            50 an-1
     Case 3 Starts with R
                                  doesn't contain RR
     discard
                                50 an-2
                                                           So an-7
    a_n = 2a_{n-1} + 2a_{n-2} n \ge 2
    b) as = 1 (empty sequence)
     a_1 = 3
    c) \alpha_2 = 2\alpha_1 + 2\alpha_0 = 6 + 2 = 8
      a3 = 2a2 + 2a1 = 16 + 6 = 22
       ay = 2 a3 + 2 a2 = 44 + 16 = 60
       as = 2ny + 2n3 = 120 +49 = 164
       a6 = 2a5 + 2a4 = 328 + 120 = 448
       a7 = 2a6 + 2a5 = 896 + 328 = 1224
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- no two red files adjacent

EX 3 Let an be the # of bit strings of length n that do not contain OO Case 1) Starts with 1 doesn't contain 00 so anstarts with 0 don't want doesn't contain 00 D says this is  $a_n = a_{n-1} + a_{n-2}$  $n \geq 2$ n 23 (doesn't count empty How many such bit strings of length 5? Chiloss) a = 1 empty string a, = 2 a2 = 3  $a_3 = a_2 + a_1 = 3 + 2 = 5$ ay = a3 + a2 = 5 + 3 = 8

as = ay + a3 = 8 + 5 = 13

9. a) Let an be the # of bit strings of length in that do not contain 000 [ Case 1 | starts with 1 doesn't contain 000 so an-1 starts with O 00 doesn't contain 000 so an-z doesn't contain 000 50 an-3  $a_n = a_{n-1} + a_{n-2} + a_{n-3}$ n ≥ 3 [""} empty string {00,01,10,11} |Total| - |{000}| = 8 - 1 = 7 |Total | = | { 1000, 0000, 0001} = 16 - 3 = 13 ay = 13 c) as = ay + a3 + a2 = 13 + 7 + 4 = 24 a6 = as + a4 + a3 = 24 + 13 + 7 = 44 a7 = a6 + as + a4 = 44 + 24 + 13 = 81

10. a) Let an be the # of bit strings of length n that contain the string 01 [Case 1] starts with 1. contains ol 50 an-1 (Case 2) starts with 0 any bit string 001 any bit string  $a_n = a_{n-1} + 2^{n-1} - 1$   $n \ge 2$  $a_2 = a_1 + 2^{2-1} - 1$ = 0 + 2 - 1 = 1 c)  $\alpha_3 = \alpha_2 + 2^2 - 1 = 1 + 4 - 1 = 4$  $\alpha_{4} = \alpha_{3} + 2^{3} - 1 = 4 + 8 - 1 = 11$ a= = a4 + 2 - 1 = 11 + 16 - 1 = 26 a6 = a5 + 25 - 1 = 26 + 32 - 1 = 57

a7 = a6 + 2 - 1 = 57 + 64 - 1 = [120]

 $2^{2} + 2^{1} + 2^{0} = 4 + 2 + 1 = 7$ 23-1=8-1=7

8.2 Ex 3. 
$$a_{1} = a_{1-1} + 2a_{1-2}$$
  $a_{0} = 2, a_{1} = 7$ 

$$c_{1} = 1 \quad c_{2} = 2$$

$$c^{2} - (1)r - 2 = 0$$

$$(r - 2)(r + 1) = 0$$

$$c^{2} + r - 2r - 2$$

$$c_{1} = 2, c_{2} = 1$$

$$a_{1} = \alpha_{1} \cdot r^{1} + \alpha_{2} \cdot r^{2}$$

$$a_{2} = \alpha_{1} \cdot r^{2} + \alpha_{2} \cdot r^{2}$$

$$a_{3} = \alpha_{1} \cdot r^{2} + \alpha_{2} \cdot r^{2}$$

$$a_{4} = 2 - \alpha_{1} \cdot r^{2}$$

$$a_{5} = 2 - \alpha_{1} \cdot r^{2}$$

$$a_{7} = 2 \cdot r^{2} - 2 - \alpha_{1}$$

$$a_{7} = 3 \cdot r^{2} - 2 - \alpha_{1}$$

$$a_{8} = 3 \cdot r^{2} - 2 - \alpha_{1}$$

$$a_{1} = 3$$

$$a_{1} = 3 \cdot r^{2} - (-1)^{n}$$

$$a_{1} = 3 \cdot r^{2} - (-1)^{n}$$

1. a) yes, degree 3
b) no, 1st weff. not constant (2n)
c) yes, degree 4
d) no, not homogeneous (2)
e) no, not linear (an)
f) yes, degree 2
g) no, no homogeneous

to the 1st power constant coefficients

homogeneous - no terms are not multiples of the a; 's

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3. a) a_n = 2a_{n-1} for n \ge 1, a_0 = 3
    C_1 = 2 C_2 = 0
      r^2 - (z)r - 0 = 0
      r2 - 2r = 0
    1-2=0
   r, = 2
     a_n = \sqrt{2^n}
                          Thm 3 for some constant of
   a = 3 = × 28
  < = 3
   a_n = 3 \cdot 2^n
  b) an = an-1 for n = 1, a0 = 2
  C, = 1 C2 = 0
   r2 - (1)r - 0 = 0
   r2 - r = 0
   1-1=0
   r, = 1
      a_n = \alpha \cdot (1)^n
                   Thm 3 for some constant of
  a = 2 = x · (x)°
     d = 2
   a_n = 2 \cdot (1)^n = 2 for all n
   c) a_n = 5a_{n-1} - 6a_{n-2} for n \ge 2, a_0 = 1, a_1 = 0
   c_1 = 5 c_2 = -6
   r2 - (5)r - (-6) = 0
      r2 - 5r + 6 = 0
   (r-3)(r-2)=0
   r, = 3 r<sub>2</sub> = 2
   an = x, r, + x 2 r2
                            Thm I for some constants x, , xz
   an = x, 3h + x2 . 2h
   a = 1 = x . 36 + x 2 . 26
                             \alpha_1 = 0 = \alpha_1 \cdot 3' + \alpha_2 \cdot 2'
    X = 1 - X 2
                                  0 = 3 \propto 1 + 2 \propto 2
                                   0 = 3(1 - \alpha_2) + 2\alpha_2
      X1 = 1-3 = -2
   a_n = -2 \cdot 3^n + 3 \cdot 2^n
                                   0 = 3 - 3 \times_{2} + 2 \times_{2}
                                   -3 = -\alpha_{2}
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d) an = 4an-1 - 4an-2 for n = 2, a0 = 6, a, = 8
  C, = 4 6, = -4
  r2 - (4)r - (-4) = 0
 r2 - 4r + 4 = 0
 (r-2)(r-2)=0
                          Thm 2
an = x, ron + x2 nron for some constants x, , xz
 \alpha_n = \alpha_1 \cdot 2^n + \alpha_2 \cdot 2^n
 a0=6=x.20+0
                                 \alpha_{1} = 8 = \alpha_{1} \cdot 2' + \alpha_{2} \cdot 1 \cdot 2'
 d, = 6
                                    8 = 2 x, + 2 x z
                                     8 = 2(6) + 2 \times 3
                                    -4 = 2 ×2
                                     X2 = 2
a_n = 6 \cdot 2^n - 2n \cdot 2^n
   = 2^{n} (6 - 2n)
e) an = -4an-1 - 4an-2 for n=2, a0=0, a,=1
 C, = -4 C<sub>2</sub> = -4
 12-(-4)1-(-4)=0
 (2 + 40 + 4 = 0
 (r+2)^2=0
  ro = -2
  an = x, ron + x2 nron Thm 2 For some constants x, , x2
  a_n = \alpha_1 (-2)^n + \alpha_2 n (-2)^n
  \alpha_0 = 0 = \alpha_1(-2)^0 + 0 \alpha_1 = 1 = \alpha_1(-2)^1 + \alpha_2(1)(-2)^1
   \alpha_1 = 0
                                    | = -2 \times, + -2 \times_2
                                     1 = -2(0) + -2×7
                                     1 = -2 ×2
                                     \alpha_2 = -\frac{1}{2}
   \alpha_n = \left(-\frac{1}{2}\right) n \left(-2\right)^n
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f) an = 4an-z for n = 2, a0 = 0, a, = 4
     c, = 0 c<sub>2</sub> = 4
    r2 - (0)r - 4 = 0
    12-4=0
(r+2)(r-2)=0
 r_1 = -2 , r_2 = 2
 an = x, r, n + x2 r2 Thm 1 Por some constants x, , x2
\alpha_n = \alpha_1 (-2)^n + \alpha_2 (2)^n
a = 0 = x (-250 + d = 1250
                                             a, = 4 = x, (-2) + x; 2'
                                                  4 = -2 \, \alpha, \, + 2 \, \alpha
   \alpha' = -\alpha_{7}
     x_1 = -(1) = -1
                                                 4 = -2(-x2) + 2x2
                                                 4 = 4 ~ 2
                                                 d = 1
a_n = -1(-2)^n + 2^n
g) a_n = a_{n-2} for n \ge 2, a_0 = 1, a_1 = 0
  C, = 0 C2 = 4
    r^2 - (0)r - \frac{1}{9} = 0
  r2 - = 0
 \left(\Gamma + \frac{1}{2}\right)\left(\Gamma - \frac{1}{2}\right) = 0
 \Gamma_1 = \frac{-1}{2}, \Gamma_2 = \frac{1}{2}
an = x, r," + x2 r2" Thm I for some constants d, , dz
a_n = \alpha_1 \left(-\frac{1}{2}\right)^n + \alpha_2 \left(\frac{1}{2}\right)^n
                                         \alpha_1 = 0 = \chi_1 \left( \frac{-1}{2} \right)^1 + \chi_2 \left( \frac{1}{2} \right)^1
a = 1 = x, (-1)0 + x = (1)0
  d, = 1 - d2
                                            0 = - 2 0, + 2 0,
  0 = -12 (1-d2) + 1 x x 2
                                              0 = - 1 + 1 x x + 1 x x
\alpha_0 = \frac{1}{2} \left( -\frac{1}{2} \right)^n + \frac{1}{2} \left( \frac{1}{2} \right)^n
                                               = ×2
```