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6.828 2011 Lecture 11: Linux ext3 crash recovery
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Plan
  logging for cash recovery
    xv6: slow and immediately durable
    ext3: fast but not immediately durable
  trade-off: performance vs.safety

topic
  crash recovery
    crash may interrupt a multi-disk-write operation
    leave file system in an unuseable state
  most common solution: logging
  last lecture: xv6 log -- simple but slow
  today: Linux ext3 log -- fast
  theme: speed vs safety
    speed: don't write the disk
    safety: write the disk ASAP

example problem:
  appending to a file
  two writes:
    mark block non-free in bitmap
    add block # to inode addrs[] array
  we want atomicity: both or neither
  so we cannot do them one at a time

why logging?
 goal: atomic system calls w.r.t. crashes
 goal: fast recovery (no hour-long fsck)
 goal: speed of write-back cache for normal operations

xv6
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review of xv6 logging
  [diagram: buffer cache, in-memory log, FS tree on disk, log on disk]
  in-memory log: blocks that must be appended to log, in order
  log "header" block and data blocks
  each system call is a transaction
    begin_op, end_op
  syscall writes in buffer cache and appends to in-memory log
    some opportunity for write absorption
  at end_op, each written block appended to on-disk log
    but NOT yet written to "home" location
    "write-ahead log"
    preserve old copy until sure we can commit
  on commit:
    write "done" and block #s to header block
  then write modified blocks to home locations
  then erase "done" from header blocks
  recovery:
    if log says "done":
      copy blocks from log to real locations on disk

homework
   Q: what does "cat a" produce after panic?
          cannot open a
   Q: how about "ls"
           panic unknown inode type
   Problem:
     dirent is on disk and has an inode#
         that inode hasn't been written to disk to the right place
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        in fact, on-disk it is marked as free
  Q: what does "cat a" produce after recovery?
    empty file.
        recovery wrote inode to the right place
          it is now allocated
          dirent is valid
    create and write are separate transactions
        create made it but write didn't
  modification to avoid reading log in install_trans()
        why is buffer 3 still in buffer cache?

what's wrong with xv6's logging? it is slow!
  all file system operation results in commit
        if no concurrent file operations
  synchronous write to on-disk log
    each write takes one disk rotation time
    commit takes a another
    a file create/delete involves around 10 writes
    thus 100 ms per create/delete -- very slow!
  tiny update -> whole block write
    creating a file only dirties a few dozen bytes
    but produces many kilobytes of log writes
  synchronous writes to home locations after commit
    i.e. write-through, not write-back
    makes poor use of in-memory disk cache

Ext3
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how can we get both performance and safety?
  we'd like system calls to proceed at in-memory speeds
  using write-back disk cache
  i.e. have typical system call complete w/o actual disk writes

Linux's ext3 design
 case study of the details required to add logging to a file system
 Stephen Tweedie 2000 talk transcript "EXT3, Journaling Filesystem"
 ext3 adds a log to ext2, a previous xv6-like log-less file system
 has many modes, I'll start with "journaled data"
   log contains both metadata and file content blocks

ext3 structures:
 in-memory write-back block cache
 in-memory list of blocks to be logged, per-transaction
 on-disk FS
 on-disk circular log file.

what's in the ext3 log?
 superblock: starting offset and starting seq #
 descriptor blocks: magic, seq, block #s
 data blocks (as described by descriptor)
 commit blocks: magic, seq
 |super: offset+seq #|...|Descriptor 4+magic|...metadata blocks...|Commit 4+magic|
|Descriptor 5+magic|...

how does ext3 get good performance despite logging entire blocks?
 batches many syscalls per commit
 defers copying cache block to log until it commits log to disk
 hopes multiple sycalls modified same block
   thus many syscalls, but only one copy of block in log
   much more "write absorbtion" than xv6

sys call:
  h = start()
  get(h, block #)
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    warn logging system we'll modify cached block
      added to list of blocks to be logged
    prevent writing block to disk until after xaction commits
  modify the blocks in the cache
  stop(h)
  guarantee: all or none
  stop() does *not* cause a commit
  notice that it's pretty easy to add log calls to existing code

ext3 transaction
  [circle set of cache blocks in this xaction]
  while "open", adds new syscall handles, and remembers their block #s
  only one open transaction at a time
  ext3 commits current transaction every few seconds (or fsync())

committing a transaction to disk
  open a new transaction, for subsequent syscalls
  mark transaction as done
  wait for in-progress syscalls to stop()
    (maybe it starts writing blocks, then waits, then writes again if needed)
  write descriptor to log on disk w/ list of block #s
  write each block from cache to log on disk
  wait for all log writes to finish
  append the commit record
  now cached blocks allowed to go to homes on disk (but not forced)

is log correct if concurrent syscalls?
  e.g. create of "a" and "b" in same directory
  inode lock prevents race when updating directory
  other stuff can be truly concurrent (touches different blocks in cache)
  transaction combines updates of both system calls

what if syscall B reads uncommited result of syscall A?
  A: echo hi > x
  B: ls > y
  could B commit before A, so that crash would reveal anomaly?
  case 1: both in same xaction -- ok, both or neither
  case 2: A in T1, B in T2 -- ok, A must commit first
  case 3: B in T1, A in T2
    could B see A's modification?
    ext3 must wait for all ops in prev xaction to finish
      before letting any in next start
      so that ops in old xaction don't read modifications of next xaction

T2 starts while T1 is committing to log on disk
  what if syscall in T2 wants to write block in prev xaction?
  can't be allowed to write buffer that T1 is writing to disk
    then new syscall's write would be part of T1
    crash after T1 commit, before T2, would expose update
  T2 gets a separate copy of the block to modify
    T1 holds onto old copy to write to log
  are there now *two* versions of the block in the buffer cache?
    no, only the new one is in the buffer cache, the old one isn't
  does old copy need to be written to FS on disk?
    no: T2 will write it

performance?
  create 100 small files in a directory
    would take xv6 over 10 seconds (many disk writes per syscall)
  repeated mods to same direntry, inode, bitmap blocks in cache
    write absorbtion...
  then one commit of a few metadata blocks plus 100 file blocks
  how long to do a commit?
    seq write of 100*4096 at 50 MB/sec: 10 ms
    wait for disk to say writes are on disk
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    then write the commit record
    that wastes one revolution, another 10 ms
    modern disk interfaces can avoid wasted revolution

what if a crash?
  crash may interrupt writing last xaction to log on disk
  so disk may have a bunch of full xactions, then maybe one partial
  may also have written some of block cache to disk
    but only for fully committed xactions, not partial last one

how does recovery work
  done by e2fsck, a utility program
  1. find the start and end of the log
     log "superblock" at start of log file
     log superblock has start offset and seq# of first transaction
     scan until bad record or not the expected seq #
     go back to last commit record
     crash during commit -> last transaction ignored during recovery
  2. replay all blocks through last complete xaction, in log order

what if block after last valid log block looks like a log descriptor?
  perhaps left over from previous use of log? (seq...)
  perhaps some file data happens to look like a descriptor? (magic #...)

when can ext3 free a transaction's log space?
  after cached blocks have been written to FS on disk
  free == advance log superblock's start pointer/seq

what if block in T1 has been dirtied in cache by T2?
  can't write that block to FS on disk
  note ext3 only does copy-on-write while T1 is commiting
    after T1 commit, T2 dirties only block copy in cache
  so can't free T1 until T2 commits, so block is in log
    T2's logged block contains T1's changes

what if not enough free space in log for a syscall?
  suppose we start adding syscall's blocks to T2
  half way through, realize T2 won't fit on disk
  we cannot commit T2, since syscall not done
  can we free T1 to free up log space?
  maybe not, due to previous issue, T2 maybe dirtied a block in T1
  deadlock!

solution: reservations
  syscall pre-declares how many block of log space it might need
  block the sycall from starting until enough free space
  may need to commit open transaction, then free older transaction
    OK since reservations mean all started sys calls can complete + commit

ext3 not as immediately durable as xv6
  creat() returns -> maybe data is not on disk! crash will undo it.
  need fsync(fd) to force commit of current transaction, and wait
  would ext3 have good performance if commit after every sys call?
    would log many more blocks, no absorption
    10 ms per syscall, rather than 0 ms
  (Rethink the Sync addresses this problem)

no checksum in ext3 commit record
  disks usually have write caches and re-order writes, for performance
    sometimes hard to turn off (the disk lies)
    people often leave re-ordering enabled for speed, out of ignorance
  bad news if disk writes commit block before preceding stuff
  then recovery replays "descriptors" with random block #s!
  and writes them with random content!
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ordered vs journaled
  journaling file content is slow, every data block written twice
  perhaps not needed to keep FS internally consistent
  can we just lazily write file content blocks?
  no:
    if metadata updated first, crash may leave file pointing
    to blocks with someone else's data
  ext3 ordered mode:
    write content block to disk before commiting inode w/ new block #
    thus won't see stale data if there's a crash
  most people use ext3 ordered mode

correctness challenges w/ ordered mode:
  A. rmdir, re-use block for file, ordered write of file,
       crash before rmdir or write committed
     now scribbled over the directory block
     fix: defer free of block until freeing operation forced to log on disk
  B. rmdir, commit, re-use block in file, ordered file write, commit,
       crash, replay rmdir
     file is left w/ directory content e.g. . and ..
     fix: revoke records, prevent log replay of a given block

final tidbit
  open a file, then unlink it
  unlink commits
  file is open, so unlink removes dir ent but doesn't free blocks
  crash
  nothing interesting in log to replay
  inode and blocks not on free list, also not reachably by any name
    will never be freed! oops
  solution: add inode to linked list starting from FS superblock
    commit that along with remove of dir ent
  recovery looks at that list, completes deletions

does ext3 fix the xv6 log performance problems?
  synchronous write to on-disk log -- yes, but 5-second window
  tiny update -> whole block write -- yes (indirectly)
  synchronous writes to home locations after commit -- yes

ext3 very successful
  but: no checksum -- ext4
  but: not efficient for applications that use fsync()
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