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Homework #5

1. Ford-Fulkerson Algorithm

Credit: <http://www.geeksforgeeks.org/ford-fulkerson-algorithm-for-maximum-flow-problem/>

a) Pseudo code:

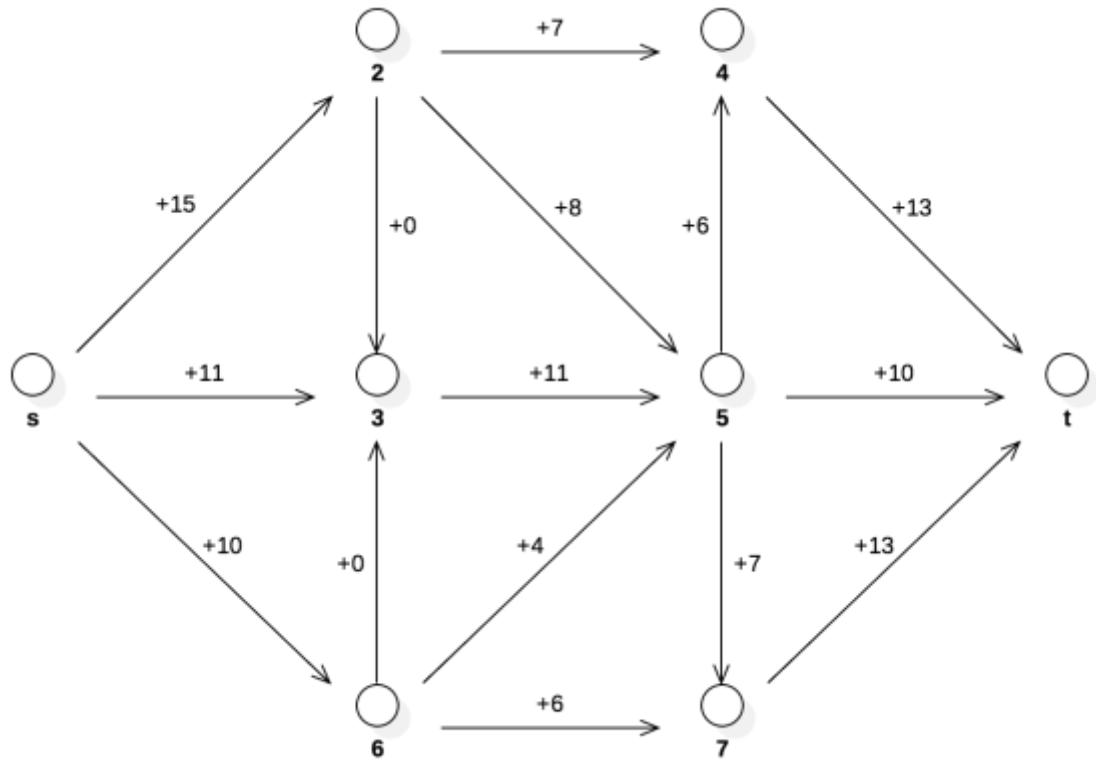
```
Let G be the node – link graph
Let rG be the residual graph and let rG = G initially
Let E be the number of edges
Let f(e) be the flow of edge e
Let C(e) be the capacity of edge e
Let maxFlow be the maximum possible flow of graph
Let f(e) = 0 for all edges and let maxFlow = 0 initially

While there is an s – t path P from s to t using BFS
    An s – t path exists if f(e) < C(e) for every edge e on the path
    If no path found
        return maxFlow
    Else
        Find minimum residual capacity of the edges
        along the path P filled by BFS and assign to pathFlow

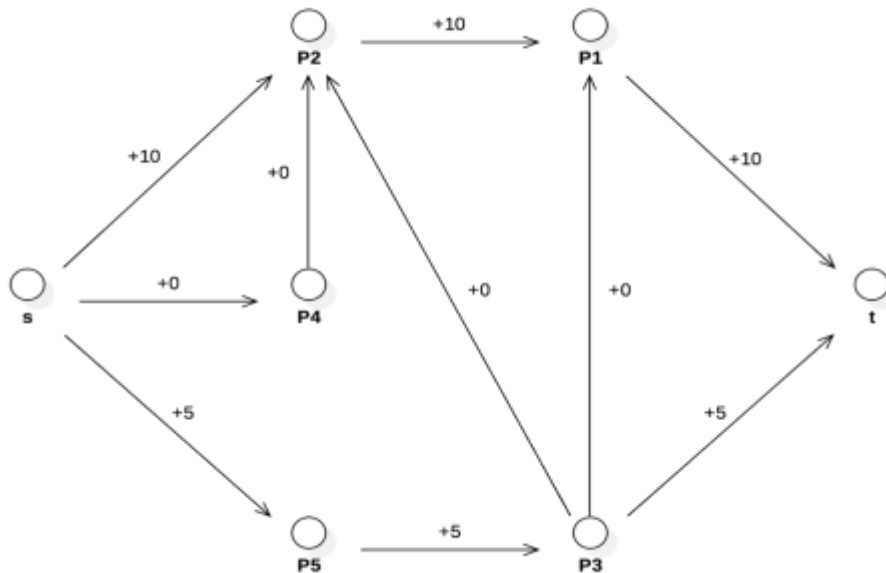
        Update residual capacities of the edges and
        reverse edges along the path
        For each edge in path P
            Let u be the start node of the edge
            Let v be the end node of the edge
            rG[u][v] = rG[u][v] – pathFlow
            rG[v][u] = rG[v][u] + pathFlow
        Endfor
    Endif
    maxFlow = maxFlow + pathFlow
Endwhile
return maxFlow
```

b) Let F_m be the max flow, E is the total number of nodes in graph. When we iterate while there is still augmenting path, the worst case is that we only add 1 unit in every iteration, as a result, the time complexity is bounded by $O(F_m * E)$.

c) The maximum possible flow is 36, the flow graph is as following:



2. We model the problem using network flow diagram. The network flow of profit is as following:



The maximum profit we can make is 15, we will take P1, P2, P3, P5 to generate this profit.

3. a)

Problem:

- n : number of doctors
- D : set of vacation days; $|D| = d$
- $S_i \subseteq D$: set of vacation-days doctor i can work
- k : vacation periods
- $D_j \subseteq D$: days in period j

Give a polynomial algorithm which returns assignment of doctors to vacation days, or reports if no such assignment exists; and satisfies following constraints:

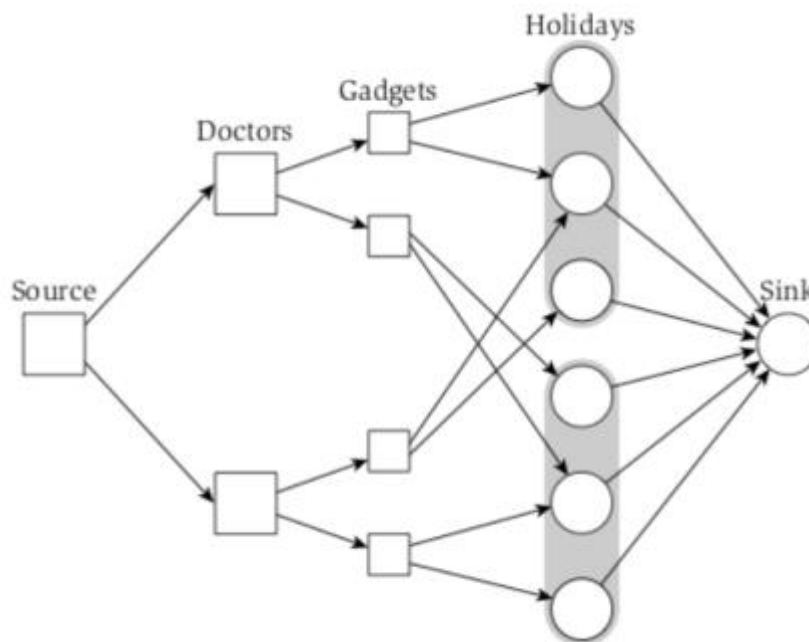
- For a given parameter c , each doctor should be assigned to work at most c vacation days total, and only days when he or she is available.
- For each vacation period j , each doctor should be assigned to work at most one of the days in the set D_j .

Solution:

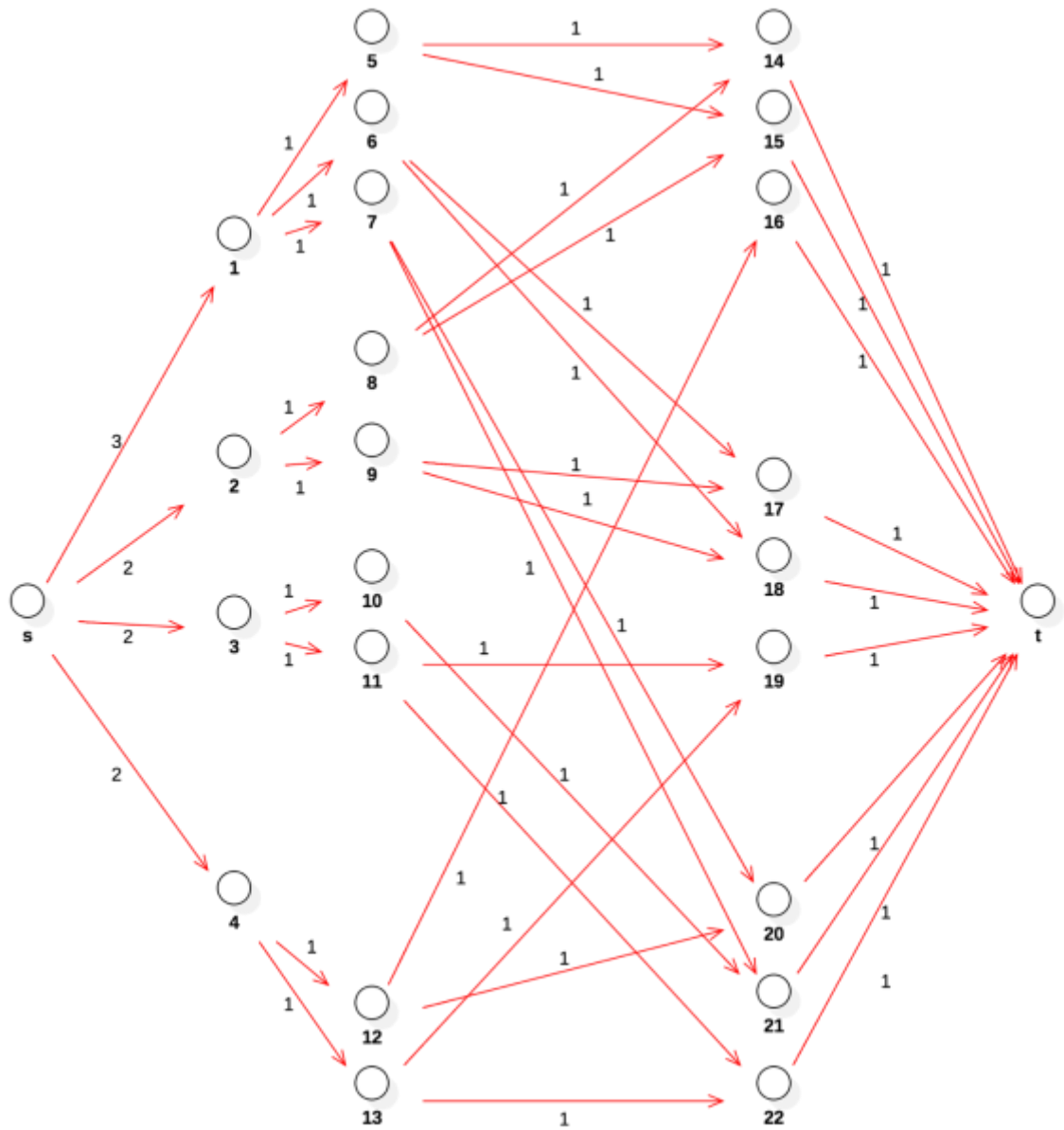
We use network flow model to solve this problem. We have a node u_i representing each doctor attached to a node v_i representing each day when he or she can work; each edge has capacity of 1.

We attach a super-source s to each doctor node u_i by an edge of capacity c , and we attach each day node v_i to a super-sink t by an edge with upper and lower bounds of 1. Suppose there are d vacation days total; we put a demand of $+d$ on the sink and $-d$ on the source.

We include a new node w_{ij} with an incoming edge of capacity 1 from the doctor node u_i , and with outgoing edges of capacity 1 to each day in vacation period j when doctor i is available to work. The network flow model is as following



b) Problem instance with 4 doctors and 3 vacation periods. The network flow diagram is as following



- Vertices from 1 to 4 represent 4 doctors.
- Vertices from 5 to 13 represent “gadgets” for 4 doctors, these gadgets guarantee each doctor only cover 1 day in each vacation period.
- Vertices from 14 to 22 represent vacation days, each 3-day block represents each vacation period.

c) We encode the network flow diagram into a matrix with values of each edge as following

```
new int[][]{ /* s 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 t */
/* s */ {0, 3, 2, 2, 2, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0},
/* 1 */ {0, 0, 0, 0, 0, 1, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0},
/* 2 */ {0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0},
/* 3 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0},
/* 4 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0, 0},
/* 5 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0},
/* 6 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0},
/* 7 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0},
/* 8 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0, 0, 0},
/* 9 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 1, 0, 0, 0, 0, 0, 0},
/* 10 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 0},
/* 11 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 1, 0, 0},
/* 12 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 0, 1, 0, 0, 0},
/* 13 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0, 0, 1, 0, 0, 0},
/* 14 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1, 0},
/* 15 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 16 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 17 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 18 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 19 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 20 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 21 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* 22 */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
/* t */ {0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 0, 1},
};
```

When applying to question 1 implementation, the result is as following:

Maximum flow is 9, which is all vacation days are covered, the details are as following:

- Day 1 of vacation period 1 is covered doctor 1
- Day 2 of vacation period 1 is covered doctor 2
- Day 3 of vacation period 1 is covered doctor 4

- Day 1 of vacation period 2 is covered doctor 1
- Day 2 of vacation period 2 is covered doctor 2
- Day 3 of vacation period 2 is covered doctor 3

- Day 1 of vacation period 3 is covered doctor 1
- Day 2 of vacation period 3 is covered doctor 3
- Day 3 of vacation period 3 is covered doctor 4

4. a)

Problem:

n users of the website

k groups from $\{G_1, \dots, G_k\}$ of the demographic groups, which users belong to

m advertisers

Give an efficient algorithm to decide if there is a way to show a single ad to each user so that the site's contracts with each of the m advertisers is satisfied for this minute? (That is, for each $i = 1, 2, 3, \dots, m$, can at least r_i of the n users, each belonging to at least one demographic group in $X_i \subseteq \{G_1, \dots, G_k\}$, be shown an ad provided by advertiser i?), and if so, to choose an ad to show each user.

Solution:

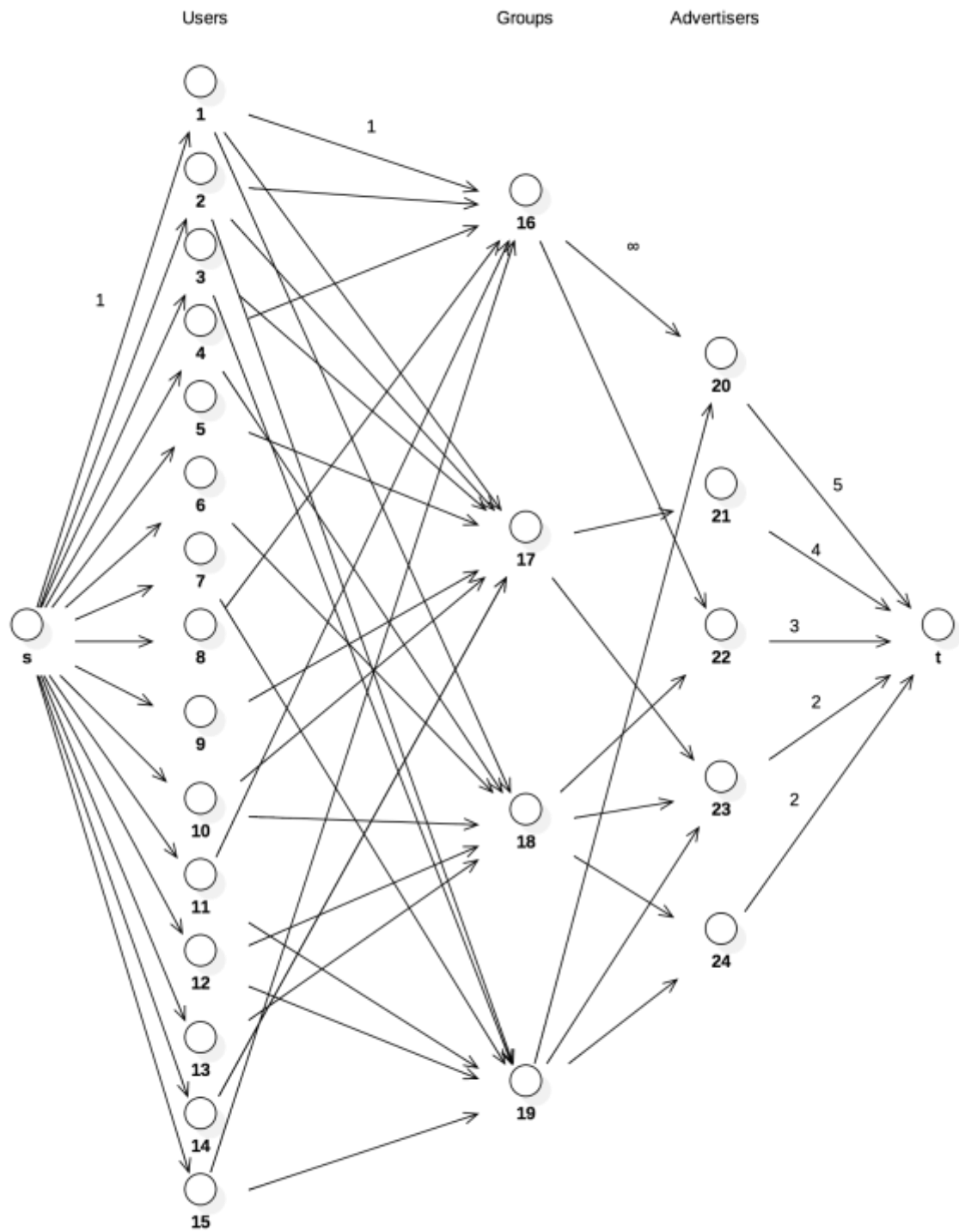
We design a flow network $G = (V, E)$ as following:

- There is a source s
- Vertices u_1 to u_n for n users
- Vertices g_1 to g_k for k demographic groups
- Vertices v_1 to v_m for m advertisers
- There is a sink t
- There is an edge of capacity 1 from u_i to each group g_j for which user i belongs to a demographic group.
- There is an edge of capacity ∞ from each group g_j to at least one advertiser v_l .
- There is an edge with a capacity of 1 from s to each u_i for each i .
- There is an edge with a capacity of the site contracts with each of m advertisers which is r_j of the n users from v_j to t for each j .
- The source s has supply of $-n$
- The sink t has demand of n

If there is a valid circulation in this graph, meaning that advertiser v_l shows ads to a subset of vertices of $\{g_1, \dots, g_k\}$, in which a subset of $\{u_1, \dots, u_n\}$ users belong to, which mean that advertiser shows ads to appropriate people. We will find out the maximum flow in this graph, if this maximum flow is equal to n , it means advertisers show their ads to all site's users with appropriate content.

b) Problem instance with $k=4$, $n=15$, $m=5$ is represented as following network flow diagram

- Each edge from s to vertices 1 to 15 has capacity of 1 (source to each user vertex)
- Each edge from vertices 1 to 15 to vertices 16 to 19 has capacity of 1 (user vertices to groups)
- Each edge from vertices 16 to 19 to vertices 20 to 24 has capacity of 15 (we choose 15 because n is 15 therefore the maximum flow to one edge is 15).
- Each edge from vertices 20 to 24 to sink t has capacity of site's contract to that advertiser (advertisers to sink t).



[illegible]

The maximum flow is 15, which means all users are shown appropriate ads.

- Node 20 to sink t has flow of $5/5$, which means the site's contract between advertiser 20 has been met.
- Node 21 to sink t has flow of $4/4$, which means the site's contract between advertiser 21 has been met.
- Node 22 to sink t has flow of $3/3$, which means the site's contract between advertiser 22 has been met.
- Node 23 to sink t has flow of $2/2$, which means the site's contract between advertiser 23 has been met.
- Node 24 to sink t has flow of $1/2$, which means the site's contract between advertiser 24 has NOT been met.