

Integer Types and Literals

Signed?	Literals
Yes	Cast from int: (byte) 3
Yes	None. Cast from int: (short) 4096
No	'a' // (char) 97
	'\n' // newline ((char) 10)
	'\t' // tab ((char) 8)
	'\\' // backslash
	'A', '\101', '\u0041' // == (char) 65
Yes	123
	0100 // Octal for 64
	0x3f, 0xffffffff // Hexadecimal 63, -1 (!)
Yes	123L, 01000L, 0x3fL 1234567891011L

nerals are just negated (positive) literals.
ns that there are 2^N integers in the domain of the type:
range of values is $-2^{N-1} .. 2^{N-1} - 1$.
d, only non-negative numbers, and range is $0..2^N - 1$.

Modular Arithmetic

(mod n) to mean that $a - b = kn$ for some integer k .
inary operation $a \bmod n$ as the value b such that $a \equiv b \pmod n$
 n for $n > 0$. (Can be extended to $n \leq 0$ as well, but
her with that here.) This is *not* the same as Java's %

s: (Here, let a' denote $a \bmod n$).

$$a'' = a'$$
$$a' + b'' = (a' + b)' = a + b'$$
$$(a' - b')' = (a' + (-b)')' = (a - b)'$$
$$(a' \cdot b')' = a' \cdot b' = a \cdot b'$$
$$(a^k)' = ((a')^k)' = (a \cdot (a^{k-1})')', \text{ for } k > 0.$$

Modular Arithmetic and Bits

ound?
tion is the natural one for a machine that uses binary
consider bytes (8 bits):

Decimal	Binary
101	1100101
$\times 99$	1100011
9999	100111 00001111
$- 9984$	100111 00000000
15	00001111

it n , counting from 0 at the right, corresponds to 2^n .
he left of the vertical bars therefore represent multi-
256.
them away is the same as arithmetic modulo 256.

CS61B Lecture #14: Integers

Overflow

w do we handle overflow, such as occurs in 10000*10000*10000?
ges throw an exception (Ada), some give undefined re-
)
; the result of any arithmetic operation or conversion
pes to "wrap around"—*modular arithmetic*.
"next number" after the largest in an integer type is
(like "clock arithmetic").
128 == (byte) (127+1) == (byte) -128

sult of some arithmetic subexpression is supposed to
 T , an n -bit integer type,
ompute the real (mathematical) value, x ,
a number, x' , that is in the range of T , and that is
to x modulo 2^n .
ns that $x - x'$ is a multiple of 2^n .)

Modular Arithmetic: Examples

8) yields 0, since $512 - 0 = 2 \times 2^8$.
2) and (byte) (127+1) yield -128, since $128 - (-128) =$

*99) yields 15, since $9999 - 15 = 39 \times 2^8$.
*13) yields 122, since $-390 - 122 = -2 \times 2^8$.
yields $2^{16} - 1$, since $-1 - (2^{16} - 1) = -1 \times 2^{16}$.

Conversion

Java will silently convert from one type to another if this conversion and no information is lost from value.

Cast explicitly, as in (byte) x.

```
byte b; char aChar; short aShort; int anInt; long aLong;
```

```
aByte; anInt = aByte; anInt = aShort;
aChar; aLong = anInt;
```

(short, might lose information:

```
aLong; aByte = anInt; aChar = anInt; aShort = anInt;
aChar; aChar = aShort; aChar = aByte;
```

special dispensation:

```
13; // 13 is compile-time constant
12+100 // 112 is compile-time constant
```

Bit twiddling

C++ allow for handling integer types as sequences of bits" needed: they already are.

and their uses:

Set	Flip	Flip all
00101100	00101100	
10100111	~ 10100111	~ 10100111
10101111	10001011	01011000

	Arithmetic Right	Logical Right
1 << 3	10101101 >> 3	10101100 >>> 3
0	11110101	00010101

```
1) >>> 29?
<< n?
>> n?
>>> 3) & ((1<<5)-1)?
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= 7.
 $= x \cdot 2^n$.

Negative numbers

representation for -1?

```
1 | 000000012
+ -1 | 111111112
= 0 | 1|000000002
```

in a byte, so bit 8 falls off, leaving 0.

ed bit is in the 2^8 place, so throwing it away gives an r modulo 2^8 . All bits to the left of it are also divisible

types (char), arithmetic is the same, but we choose to ly non-negative numbers modulo 2^{16} :

```
1 | 00000000000000012
+ 216 - 1 | 11111111111111112
= 216 + 0 | 1|00000000000000002
```

Promotion

operations (+, *, ...) promote operands as needed.

just implicit conversion.

operations,

rand is long, promote both to long.

promote both to int.

```
} == (int) aByte + 3 // Type int
} == aLong + (long) 3 // Type long
== (int) 'A' + 2 // Type int
aByte + 1 // ILLEGAL (why?)
```

ely,

```
1; // Defined as aByte = (byte) (aByte+1)
```

mple:

```
e aChar is an upper-case letter
rCaseChar = (char) ('a' + aChar - 'A'); // why cast?
```

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1) >>> 29? = 7.
<< n? = $x \cdot 2^n$.
>> n? = $\lfloor x/2^n \rfloor$ (i.e., rounded down).
>>> 3) & ((1<<5)-1)? 5-bit integer, bits 3-7 of x.

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