

Dynamic Programming

Garcia):

h a list with an even number of non-negative integers.
er in turn takes either the leftmost number or the
.

get the largest possible sum.

irting with (6, 12, 0, 8), you (as first player) should take
ever the second player takes, you also get the 12, for a

ur opponent plays perfectly (i.e., to get as much as pos-
an you maximize your sum?

s with exhaustive game-tree search.

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Still Another Idea from CS61A

is that we are recomputing intermediate results many

emoize the intermediate results. Here, we pass in an
($N = V.length$) of memoized results, initialized to -1.

```
bestSum(int[] V, int left, int right, int total, int[][] memo) {
    if (left > right)
        return 0;
    if (memo[left][right] != -1)
        return memo[left][right];
    int totalL = total + V[left];
    int totalR = total + V[right];
    memo[left][right] = Math.max(totalL, totalR);
    return memo[left][right];
}
```

emo[left][right];

number of recursive calls to bestSum must be $O(N^2)$, for
length of V , an enormous improvement from $\Theta(2^N)$!

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Longest Common Subsequence

d length of the longest string that is a subsequence of
other strings.

ngest common subsequence of
lls_sea_shells_by_the_seashore" and
ld_salt_sellers_at_the_salt_mines"

ells_the_sae" (length 23)

string, for example.

rsive algorithm:

```
longestCommonSubsequence(S0, k0-1, S1, k1-1) {
    if (k0 == 0 || k1 == 0) return 0;
    if (S0[k0-1] == S1[k1-1]) return 1 + longestCommonSubsequence(S0, k0-1, S1, k1-1);
    return Math.max(longestCommonSubsequence(S0, k0-1, S1, k1), longestCommonSubsequence(S0, k0, S1, k1-1));
}
```

but obviously memoizable.

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Lecture #35

gramming and memoization.

Git.

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Obvious Program

akes it easy, again:

```
bestSum(int[] V) {
    int n = V.length;
    int total = 0;
    for (int i = 0; i < n; i++) total += V[i];
    return total;
}
```

rgest sum obtainable by the first player in the choosing
n the list $V[LEFT \dots RIGHT]$, assuming that $TOTAL$ is the
all the elements in $V[LEFT \dots RIGHT]$. */

```
bestSum(int[] V, int left, int right, int total) {
    if (left > right)
        return 0;
    int totalL = total + V[left];
    int totalR = total + V[right];
    return Math.max(totalL, totalR);
}
```

```
total = total + V[left];
total = total + V[right];
return Math.max(totalL, totalR);
}
```

$C(0) = 1$, $C(N) = 2C(N-1)$; so $C(N) \in \Theta(2^N)$

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Iterative Version

recursive version, but the usual presentation of this
as *dynamic programming*—is iterative:

```
longestCommonSubsequence(S0, S1) {
    int n0 = S0.length, n1 = S1.length;
    int[][] memo = new int[n0][n1];
    for (int i = 0; i < n0; i++)
        for (int j = 0; j < n1; j++)
            memo[i][j] = -1;
    return longestCommonSubsequence(S0, S1, memo, 0, 0);
}
```

emo[0][V.length-1];

figure out ahead of time the order in which the memo-
will fill in memo, and write an explicit loop.

needed to check whether result exists.

ny bother unless it's necessary to save space?

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memoized Longest Common Subsequence

```
longest common subsequence of S0[0..k0-1]
and S1[k1-1] (pseudo Java) */
String S0, int k0, String S1, int k1) {
    new int[k0+1][k1+1];
    : memo) Arrays.fill(row, -1);
    k0, S1, k1, memo);

    int lls(String S0, int k0, String S1, int k1, int[][] memo) {
        k1 == 0) return 0;
        k1 == -1) {
            == S1[k1-1])
        l1 = 1 + lls(S0, k0-1, S1, k1-1, memo);

        l1 = Math.max(lls(S0, k0-1, S1, k1, memo),
                      lls(S0, k0, S1, k1-1, memo));

        l1[k1];
    }
}
```

Will the memoized version be? $\Theta(k_0 \cdot k_1)$

A Little History

by Linus Torvalds and others in the Linux community when
one of their previous, proprietary VCS (Bitkeeper) with-
out a version.
Development effort seems to have taken about 2-3 months,
the 2.6.12 Linux kernel release in June, 2005.
The name, according to Wikipedia,
Linus has quipped about the name Git, which is British
slang meaning "unpleasant person". Torvalds said: "I'm
an arse, and I name all my projects after myself.
'git', now 'git'." The man page describes Git as "the
distributed version tracker."
It is a collection of basic primitives (now called "plumbing")
designed to provide desired functionality.
High-level commands ("porcelain") built on top of these to
provide a convenient user interface.

Conceptual Structure

All components consist of four types of *object*:
Objects typically hold contents of files.
Directory structures of files.
Contain references to trees and additional information
(e.g., date, log message).
References to commits or other objects, with additional
information, intended to identify releases, other important ver-
sion information. (Won't mention further to-

memoized Longest Common Subsequence

```
longest common subsequence of S0[0..k0-1]
and S1[k1-1] (pseudo Java) */
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    new int[k0+1][k1+1];
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    int lls(String S0, int k0, String S1, int k1, int[][] memo) {
        k1 == 0) return 0;
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            == S1[k1-1])
        l1 = 1 + lls(S0, k0-1, S1, k1-1, memo);

        l1 = Math.max(lls(S0, k0-1, S1, k1, memo),
                      lls(S0, k0, S1, k1-1, memo));

        l1[k1];
    }
}
```

Will the memoized version be?

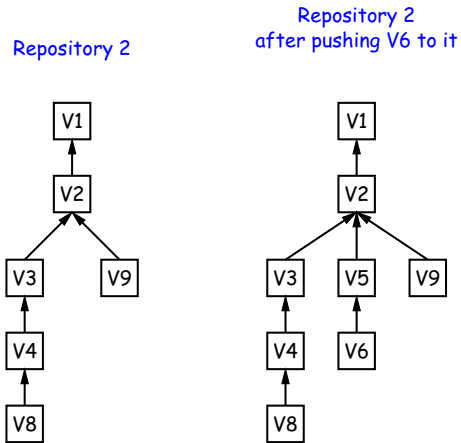
Use Study in System and Data-Structure Design

A distributed version-control system, apparently the most pop-
ular currently.
Git, it stores snapshots (*versions*) of the files and direc-
tory of a project, keeping track of their relationships,
commits, and log messages.
It is *distributed*, in that there can be many copies of a given repos-
itory supporting independent development, with machinery to
reconcile versions between repositories.
Git is extremely fast (as these things go).

Major User-Level Features (I)

Git is a graph of versions or snapshots (called *commits*)
for a project.
The structure reflects ancestry: which versions came from
which.
Each commit contains
a tree of files (like a Unix directory).
Information about who committed and when.
The commit message.
A list of parent commit (or commits, if there was a merge) from which
it was derived.

Version Histories in Two Repositories



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Internals

Repository is contained in a directory.
 It may either be *bare* (just a collection of objects and pointers) or may be included as part of a working directory.
 The repository is stored in various *objects* corresponding to other "leaf" content), trees, and commits.
 For example, data in files is *compressed*.
 We *periodically collect* the objects from time to time to save additional space.

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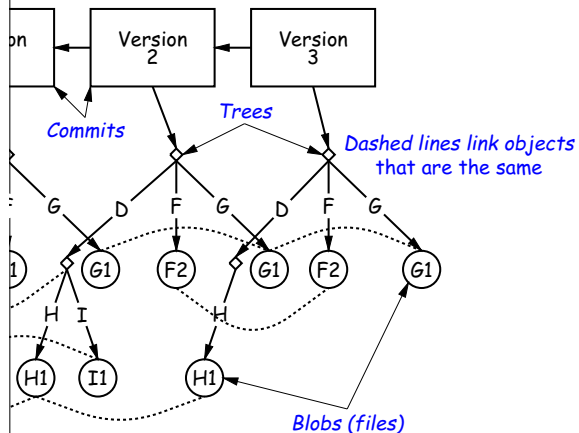
Content-Addressable File System

One way of naming objects that is universal.
 Names, then, as pointers.
 Which objects don't you have?" problem in an obvious way.
 What is invariant about an object, regardless of repository?
 The contents as the name for obvious reasons.
 The *hash of the contents* as the address.
 That doesn't work!
 Use it anyway!!

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Commits, Trees, Files



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Major User-Level Features (II)

Each object has a name that uniquely identifies it to all versions.
 We can transmit collections of versions to each other.
 Making a commit from repository *A* to repository *B* requires transmission of those objects (files or directory trees) that *B* has not yet have (allowing speedy updating of repositories).
 We maintain named *branches*, which are simply identifiers for commits that are updated to keep track of the most recent commits in various lines of development.
 Branches are essentially named pointers to particular commits. The branches in that they are not usually changed.

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The Pointer Problem

How do we represent pointers between objects? Are they files. How should we represent pointers between objects?
 How do we transmit the pointers?
 How do we transfer those objects that are missing in the target repository?
 How do we know which those are?
 How do we maintain a counter in each repository to give each object there a unique identifier?
 But how can that work consistently for two independent repositories?

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SHA1

1 (Secure Hash Function 1).

and with this using the `hashlib` module in Python3.

ames in Git are therefore 160-bit hash codes of con-

.
commit in the shared CS61B repository could be fetched
with

```
ckout e59849201956766218a3ad6ee1c3aab37dfec3fe
```

How A Broken Idea Can Work

o use a hash function that is so unlikely to have a collision, we can ignore that possibility.

ic Hash Functions have relevant property.

ion, f , is designed to withstand cryptanalytic attacks.
, should have

resistance: given $h = f(m)$, should be computationally
to find such a message m .

re-image resistance: given message m_1 , should be infeasible
to find $m_2 \neq m_1$ such that $f(m_1) = f(m_2)$.

esistance: should be difficult to find *any* two messages
such that $f(m_1) = f(m_2)$.

properties, scheme of using hash of contents as name is
likely to fail, even when system is used maliciously.