CS61B Lecture #8: Object-Oriented Mechanis

Today:

- New in this lecture: the bare mechanics of "object-orient gramming."
- The general topic is: Writing software that operates on main of data.

Overloading

Problem: How to get System.out.print(x) to print x, regard type of x?

- In Scheme or Python, one function can take an argument of a and then test the type (if needed).
- In Java, methods specify a single type of argument.
- Partial solution: overloading—multiple method definitions v same name and different numbers or types of arguments.
- E.g., System.out has type java.io.PrintStream, which defi

```
void println() Prints new line.
void println(String s) Prints 5.
void println(boolean b) Prints "true" or "false"
void println(char c) Prints single character
void println(int i) Prints I in decimal
etc.
```

 Each of these is a different function. Compiler decides whice on the basis of arguments' types.

Generic Data Structures

Problem: How to get a "list of anything" or "array of anything

- Again, no problem in Scheme or Python.
- But in Java, lists (such as IntList) and arrays have a single element.
- First, the short answer: any reference value can be convetype java.lang.Object and back, so can use Object as the '(reference) type":

And Primitive Values?

- Primitive values (ints, longs, bytes, shorts, floats, doubles and booleans) are not really convertible to Object.
- Presents a problem for "list of anything."
- So Java introduced a set of wrapper types, one for each p type:

Prim.	Ref.	Prim.	Ref.	Prim.	Re
byte	Byte		Short	int	Inte
J	Long Float		Character Double	boolean	Bool

• One can create new wrapper objects for any value (boxing):

```
Integer Three = new Integer(3);
Object ThreeObj = Three;
and vice-versa (unboxing):
  int three = Three.intValue();
```

Autoboxing

Boxing and unboxing are automatic (in many cases):

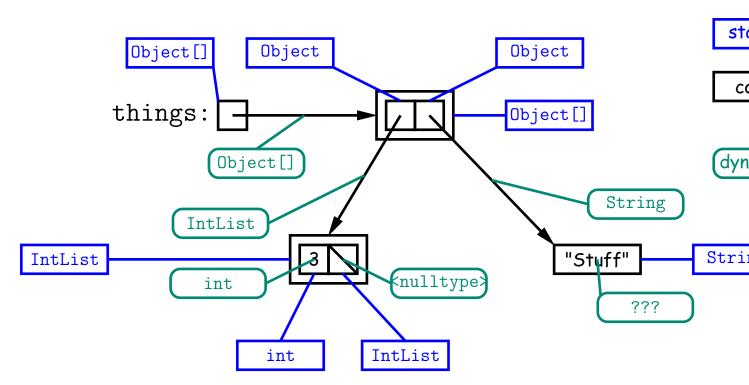
```
Integer Three = 3;
int three = Three;
int six = Three + 3;

Integer[] someInts = { 1, 2, 3 };
for (int x : someInts) {
    System.out.println(x);
}
System.out.println(someInts[0]);
// Prints Integer 1, but NOT unboxed.
```

Dynamic vs. Static Types

- Every value has a type—its dynamic type.
- ullet Every container (variable, component, parameter), literal, call, and operator expression (e.g. x+y) has a type—its stati
- Therefore, every expression has a static type.

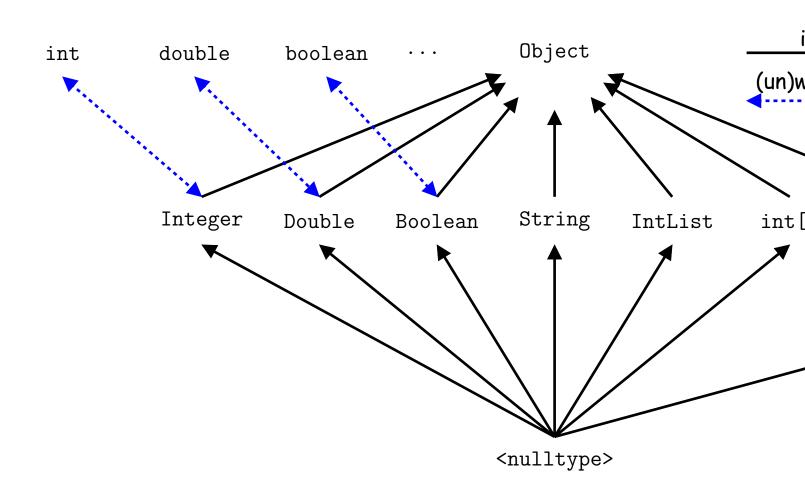
```
Object[] things = new Object[2];
things[0] = new IntList(3, null);
things[1] = "Stuff";
```



Type Hierarchies

- A container with (static) type T may contain a certain value that value "is a" T—that is, if the (dynamic) type of the value subtype of T. Likewise, a function with return type T may only values that are subtypes of T.
- All types are subtypes of themselves (& that's all for primitive
- Reference types form a type hierarchy; some are subtypes ers. null's type is a subtype of all reference types.
- All reference types are subtypes of Object.

Java Library Type Hierarchy (Partial)



The Basic Static Type Rule

- Java is designed so that any expression of (static) type 7
 yields a value that "is a" T.
- Static types are "known to the compiler," because you declar as in

or they are pre-declared by the language (like 3).

ullet Compiler insists that in an assignment, L = E, or function combined where

```
void f(SomeType L) { ... },
```

E's static type must be subtype of L's static type.

Similar rules apply to E[i] (static type of E must be an ariother built-in operations.

Coercions

- The values of type short, for example, are a subset of tint (shorts are representable as 16-bit integers, ints as integer)
- But we don't say that short is a subtype of int, because the quite behave the same.
- Instead, we say that values of type short can be coerce verted) to a value of type int.
- Leads to a slight fudge: compiler will silently coerce "smalle ger types to larger ones, float to double, and (as just se tween primitive types and their wrapper types.
- So.

```
short x = 3002;
int y = x;
```

works without complaint.

Consequences of Compiler's "Sanity Checks"

This is a conservative rule. The last line of the following, when might think is perfectly sensible, is illegal:

```
int[] A = new int[2];
Object x = A; // All references are Objects
A[i] = 0; // Static type of A is array...
x[i+1] = 1; // But not of x: ERROR
```

Compiler figures that not every Object is an array.

- \bullet Q: Don't we *know* that x contains array value!?
- A: Yes, but still must tell the compiler, like this:

```
((int[]) x)[i+1] = 1;
```

- Defn: Static type of cast (T) E is T.
- \bullet Q: What if x *isn't* an array value, or is null?
- A: For that we have runtime errors—exceptions.

Overriding and Extension

- Notation so far is clumsy.
- Q: If I know Object variable x contains a String, why can't x.startsWith("this")?
- A: startsWith is only defined on Strings, not on all Objects compiler isn't sure it makes sense, unless you cast.
- But, if an operation were defined on all Objects, then you need clumsy casting.
- Example: .toString() is defined on all Objects. You can alv x.toString() if x has a reference type.
- The default .toString() function is not very useful; on an I would produce string like "IntList@2f6684"
- But for any subtype of Object, you may override the defaunction.

Overriding to String

- For example, if s is a String, s.toString() is the identity: (fortunately).
- For any type you define, you may supply your own definitiexample, in IntList, could add

```
public String toString() {
    StringBuffer b = new StringBuffer();
    b.append("[");
    for (IntList L = this; L != null; L = L.tail)
        b.append(" " + L.head);
    b.append("]");
    return b.toString();
}
```

- If x = new IntList(3, new IntList(4, null)), then x.te is "[3 4]".
- Conveniently, the "+" operator on Strings calls .toString who
 to append an Object, and so does the "%s" formatter for pr
- With this trick, you can supply an output function for any t define.

Extending a Class

To say that class B is a direct subtype of class A (or A is superclass of B), write

```
class B extends A { ... }
```

- By default, class ... extends java.lang.Object.
- The subtype inherits all fields and methods of its direct support (and passes them along to any of its subtypes).
- In class B, you may override an instance method (not a static by providing a new definition with same signature (name, type, argument types).
- I'll say that a method and all its overridings form a dynamic set.
- The Point: If f(...) is an instance method, then the call x calls whatever overriding of f applies to the dynamic type of gardless of the static type of x.

Illustration

```
class Worker {
  void work() {
    collectPay();
  }
}
```

Lesson: For instance methods (only), select method based on type. Simple to state, but we'll see it has profound consequence

What About Fields and Static Methods?

```
class Parent {
  int x = 0;
  static int y = 1;
  static void f() {
    System.out.printf("Ahem!%n");
  }
  static int f(int x) {
    return x+1;
  }
}
class Child extends Parent {
    String x = "no";
    static String y = "way";
    static void f() {
        System.out.printf("I wa
        }
    }
    return x+1;
}
```

Lesson: Fields *hide* inherited fields of same name; static r *hide* methods of the same signature.

Real Lesson: Hiding causes confusion; so understand it, but do

What's the Point?

- The mechanism described here allows us to define a kind of method.
- A superclass can define a set of operations (methods) that a mon to many different classes.
- Subclasses can then provide different implementations o common methods, each specialized in some way.
- All subclasses will have at least the methods listed by the class.
- So when we write methods that operate on the superclass, tautomatically work for all subclasses with no extra work.