B Lecture #3: Values and Containers

mally due at midnight Friday. Last week's is due tonight. ple classes. Scheme-like lists. Destructive vs. non-pperations. Models of memory.

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Recreation

that $|(2+\sqrt{3})^n|$ is odd for all integer $n \ge 0$.

larsky, N. N. Chentzov, I. M. Yaglom, The USSR Olympiad Problem 193), from the W. H. Freeman edition, 1962.]

Structured Containers

tainers contain (0 or more) other containers:

iect Array Object Empty Object

0 1 2
42 17 9

0 42
1 17
2 9

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Values and Containers

umbers, booleans, and pointers. Values never change.

a' true $\frac{\perp}{\bar{z}}$

iners contain values:

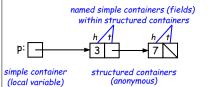
3 L: p: ____

riables, fields, individual array elements, parameters.

Containers in Java

ay be named or anonymous.

simple containers are named, *all* structured containymous, and pointers point only to structured containers. structured containers contain only simple containers).



gnment copies values into simple containers.

Scheme and Python!

has slice assignment, as in x[3:7]=..., which is short-nething else entirely.)

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Pointers

references) are values that reference (point to) con-

ar pointer, called **null**, points to nothing.

uctured containers contain only simple containers, but w us to build arbitrarily big or complex structures any-



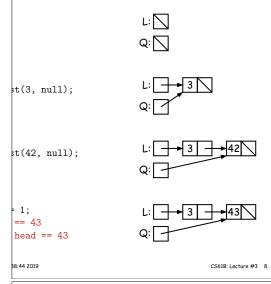
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Primitive Operations



Defining New Types of Object

itions introduce new types of objects.

of integers:

```
s IntList {
uctor function (used to initialize new object)
cell containing (HEAD, TAIL). */
tList(int head, IntList tail) {
ad = head; this.tail = tail;

of simple containers (fields)
G: public instance variables usually bad style!
t head;
tList tail;

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```

nother Way to View Pointers (II)

view:

last: #3

result: #7

5 #3

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native view, you might be less inclined to think that asuld change object #7 itself, rather than just "last".

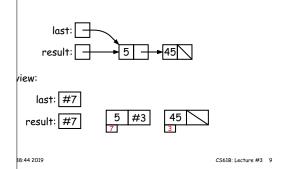
nternally, pointers really are just numbers, but Java as more than that: they have *types*, and you can't just ers into pointers.

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cursion: Another Way to View Pointers

ind the idea of "copying an arrow" somewhat odd.
view: think of a pointer as a *label*, like a street address.
has a permanent label on it, like the address plaque on

ble containing a pointer is like a scrap of paper with a ss written on it.



ondestructive IncrList: Recursive

```
f all items in P incremented by n. */
List incrList(IntList P, int n) {
  null)
  null;
urn new IntList(P.head+n, incrList(P.tail, n));

crList have to return its result, rather than just set-
crList(P, 2), where P contains 3 and 43, which IntList
created first?
```

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Destructive vs. Non-destructive

h a (pointer to a) list of integers, L, and an integer inrn a list created by incrementing all elements of the list

```
f all items in P incremented by n. Does not modify
ng IntLists. */
List incrList(IntList P, int n) {
   /*( P, with each element incremented by n )*/
```

t is non-destructive, because it leaves the input objects nown on the left. A destructive method may modify the o that the original data is no longer available, as shown

```
rList(L, 2): After Q = dincrList(L, 2) (destructive):
L: 5 45
Q: C5618: Lecture #3 11
```

An Iterative Version

rList is tricky, because it is not tail recursive. things first-to-last, unlike recursive version:

```
ncrList(IntList P, int n) {
 last;
st(P.head+n, null);
= null) {
ist(P.head+n, null);
tail;
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```

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                          result:
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```

