

- Sorting algorithms: why?
- Insertion Sort.
- Inversions

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- Binary search standard example
- Also supports other kinds of search:
 - Are there two equal items in this set?
 - Are there two items in this set that both have the same value for property X?
 - What are my nearest neighbors?
- Used in numerous unexpected algorithms, such as convex hull (smallest convex polygon enclosing set of points).

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arranges) a sequence of elements to bring them into order, according to some *total order*.

- A total order, \preceq , is:
 - **Total:** $x \preceq y$ or $y \preceq x$ for all x, y .
 - **Reflexive:** $x \preceq x$;
 - **Antisymmetric:** $x \preceq y$ and $y \preceq x$ iff $x = y$.
 - **Transitive:** $x \preceq y$ and $y \preceq z$ implies $x \preceq z$.
- However, our orderings may treat unequal items as equivalent:
 - E.g., there can be two dictionary definitions for the same word. If we sort only by the word being defined (ignoring the definition), then sorting could put either entry first.

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ory.

- *External sorts* process large amounts of data in batches, keeping what won't fit in secondary storage (in the old days, tapes).
- *Comparison-based* sorting assumes only thing we know about keys is their order.
- *Radix sorting* uses more information about key structure.
- *Insertion sorting* works by repeatedly inserting items at their appropriate positions in the sorted sequence being constructed.
- *Selection sorting* works by repeatedly selecting the next larger (smaller) item in order and adding it to one end of the sorted sequence being constructed.

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- The java library provides static methods to sort arrays in the class `java.util.Arrays`.
- For each primitive type `P` other than `boolean`, there are

```
/** Sort all elements of ARR into non-descending
order. */
static void sort(P[] arr) { ... }
```

```
/** Sort elements FIRST .. END-1 of ARR
into non-descending
* order. */
static void sort(P[] arr, int first,
int end) { ... }
```

```
/** Sort all elements of ARR into non-descending
order,
* possibly using multiprocessing for
speed. */
```

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```

/** Sort elements first... End of arr
into non-descending
 * order, possibly using multiprocessing
for speed. */
static void parallelSort(P[] arr, int
first, int end) {...}

```

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- For reference types, *C*, that have a *natural order* (that is, that implement `java.lang.Comparable`), we have four analogous methods (one-argument sort, three-argument sort, and two `parallelSort` methods):

```

/** Sort all elements of ARR stably into
non-descending
 * order. */
static <C extends Comparable<? super
C>> sort(C[] arr) {...}
etc.

```

- And for all reference types, *R*, we have four more:

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```

...
static <R> void sort(R[] arr, Comparator<?
super R> comp) {...}
etc.

```

- **Q:** Why the fancy generic arguments?

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- For reference types, *C*, that have a *natural order* (that is, that implement `java.lang.Comparable`), we have four analogous methods (one-argument sort, three-argument sort, and two `parallelSort` methods):

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/** Sort all elements of ARR stably into
non-descending
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static <C extends Comparable<? super
C>> sort(C[] arr) {...}
etc.

```

- And for all reference types, *R*, we have four more:

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```

...
static <R> void sort(R[] arr, Comparator<?
super R> comp) {...}
etc.

```

- **Q:** Why the fancy generic arguments?
- **A:** We want to allow types that have `compareTo` methods that apply also to more general types.

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two methods similar to the sorting methods for arrays of reference types:

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```

    // Sort all elements of LST stably into
    // non-descending
    // order according to the ordering
    // defined by COMP. */
    static <R> void sort(List<R> lst,
        Comparator<? super R> comp) {...}
    etc.

```

- Also an instance method in the List<R> interface itself:

```

    /** Sort all elements of LST stably into
    non-descending
    * order according to the ordering
    defined by COMP. */
    void sort(Comparator<? super R> comp)
    { ... }

```

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```

import static java.util.Arrays.*;
import static java.util.Collections.*;

```

- Sort X, a String[] or List<String>, into non-descending order:

```

    sort(X);    // or ...

```

- Sort X into reverse order (Java 8):

```

    sort(X, (String x, String y) -> { return
    y.compareTo(x); });
    // or
    sort(X, Collections.reverseOrder());
    // or
    X.sort(Collections.reverseOrder());    //
    for X a List

```

- Sort X[10], ..., X[100] in array or List X (rest unchanged):

```

    sort(X, 10, 101);

```

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- starting with empty sequence of outputs.
- add each item from input, *inserting* into output sequence at right point.

- Very simple, good for small sets of data.
- With vector or linked list, time for find + insert of one item is at worst $\Theta(k)$, where k is # of outputs so far.
- This gives us a $\Theta(N^2)$ algorithm (worst case as usual).
- Can we say more?

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- Consider a typical implementation for arrays:

```

for (int i = 1; i < A.length; i += 1) {
    int j;
    Object x = A[i];
    for (j = i-1; j >= 0; j -= 1) {
        if (A[j].compareTo(x) <= 0) /* (1) */
            break;
        A[j+1] = A[j];            /* (2) */
    }
    A[j+1] = x;
}

```

- #times (1) executes for each $j \approx$ how far x must move.
- If all items within K of proper places, then takes $O(KN)$ operations.
- Thus good for any amount of *nearly sorted* data.

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- Each execution of (2) decreases inversions by 1.

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distant elements:

- First sort subsequences of elements $2^k - 1$ apart:
 - sort items #0, $2^k - 1$, $2(2^k - 1)$, $3(2^k - 1)$, ..., then
 - sort items #1, $1 + 2^k - 1$, $1 + 2(2^k - 1)$, $1 + 3(2^k - 1)$, ..., then
 - sort items #2, $2 + 2^k - 1$, $2 + 2(2^k - 1)$, $2 + 3(2^k - 1)$, ..., then
 - etc.
 - sort items # $2^k - 2$, $2(2^k - 1) - 1$, $3(2^k - 1) - 1$, ...,
 - Each time an item moves, can reduce #inversions by as much as $2^k + 1$.
- Now sort subsequences of elements $2^{k-1} - 1$ apart:

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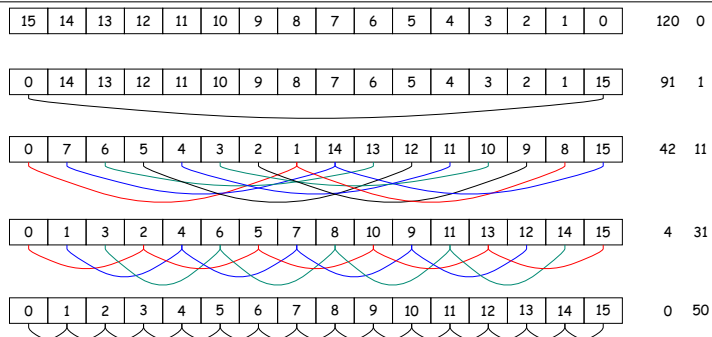
1), $1 + 3(2^{k-1} - 1)$, ...,

- :

- End at plain insertion sort ($2^0 = 1$ apart), but with most inversions gone.
- Sort is $\Theta(N^{3/2})$ (take CS170 for why!).

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I: Inversions left.

C: Cumulative comparisons used to sort subsequences by insertion sort.

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