Purposes of Sorting

orts searching

h standard example

s other kinds of search:

two equal items in this set?

two items in this set that both have the same value for

my nearest neighbors?

rous unexpected algorithms, such as convex hull (smallolygon enclosing set of points).

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rithms: why?

rt.

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Classifications

ts keep all data in primary memory.

ts process large amounts of data in batches, keeping it in secondary storage (in the old days, tapes).

based sorting assumes only thing we know about keys is

g uses more information about key structure.

rting works by repeatedly inserting items at their apsitions in the sorted sequence being constructed.

rting works by repeatedly selecting the next larger m in order and adding it to one end of the sorted seconstructed.

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Some Definitions

porithm (or sort) permutes (re-arranges) a sequence of brings them into order, according to some total order.

```
Ի, <u>Ճ</u>, is:

otin y \le x
otin y = x
otin x, y.
otin y.
 : x \prec x:
etric: x \leq y and y \leq x iff x = y.
x \prec y and y \prec z implies x \prec z.
```

orderings may treat unequal items as equivalent:

e can be two dictionary definitions for the same word. t only by the word being defined (ignoring the definih sorting could put either entry first.

at does not change the relative order of equivalent enpared to the input) is called stable.

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lys of Reference Types in the Java Library

```
te types, C, that have a natural order (that is, that im-
a.lang.Comparable), we have four analogous methods
ht sort, three-argument sort, and two parallelSort
all elements of ARR stably into non-descending
 extends Comparable<? super C>> sort(C[] arr) {...}
eference types, R, we have four more:
all elements of ARR stably into non-descending order
ding to the ordering defined by COMP. */
> void sort(R[] arr, Comparator<? super R> comp) {...}
fancy generic arguments?
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```

ays of Primitive Types in the Java Library

ary provides static methods to sort arrays in the class rrays.

```
mitive type P other than boolean, there are
```

```
all elements of ARR into non-descending order. */
id sort(P[] arr) { ... }
elements FIRST .. END-1 of ARR into non-descending
id sort(P[] arr, int first, int end) { ... }
all elements of ARR into non-descending order,
bly using multiprocessing for speed. */
id parallelSort(P[] arr) { ... }
elements FIRST .. END-1 of ARR into non-descending
, possibly using multiprocessing for speed. */
id parallelSort(P[] arr, int first, int end) {...}
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```

forting Lists in the Java Library

va.util.Collections contains two methods similar to

```
hethods for arrays of reference types:
all elements of LST stably into non-descending
 . */
 extends Comparable<? super C>> sort(List<C> lst) {...}
all elements of LST stably into non-descending
 according to the ordering defined by COMP. */
> void sort(List<R> , Comparator<? super R> comp) {...}
nce method in the List<R> interface itself:
all elements of LST stably into non-descending
 according to the ordering defined by COMP. */
(Comparator<? super R> comp) {...}
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```

lys of Reference Types in the Java Library

```
te types, C, that have a natural order (that is, that im-
a.lang.Comparable), we have four analogous methods
ht sort, three-argument sort, and two parallelSort
all elements of ARR stably into non-descending
 extends Comparable<? super C>> sort(C[] arr) {...}
eference types, R, we have four more:
all elements of ARR stably into non-descending order
ding to the ordering defined by COMP. */
> void sort(R[] arr, Comparator<? super R> comp) {...}
fancy generic arguments?
to allow types that have compare To methods that apply
general types.
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```

Sorting by Insertion

```
lith empty sequence of outputs.
tem from input, inserting into output sequence at right
good for small sets of data.
or linked list, time for find + insert of one item is at
where k is # of outputs so far.
a \Theta(N^2) algorithm (worst case as usual).
ore?
```

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Examples

```
atic java.util.Arrays.*;
atic java.util.Collections.*;
ing[] or List<String>, into non-descending order:
   // or ...
everse order (Java 8):
String x, String y) -> { return y.compareTo(x); });
lollections.reverseOrder()); // or
llections.reverseOrder()); // for X a List
 ..., X[100] in array or List X (rest unchanged):
0, 101);
 ..., L[100] in list L (rest unchanged):
iblist(10, 101));
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```

Shell's sort

```
insertion sort by first sorting distant elements:
bsequences of elements 2^k - 1 apart:
\sharp #0, 2^k - 1, 2(2^k - 1), 3(2^k - 1), ..., then
\sharp #1, 1+2^k-1, 1+2(2^k-1), 1+3(2^k-1), ..., then
\sharp #2, 2+2^k-1, 2+2(2^k-1), 2+3(2^k-1), ..., then
\mathbf{5} \# 2^k - 2, 2(2^k - 1) - 1, 3(2^k - 1) - 1, ....
an item moves, can reduce #inversions by as much as
beguences of elements 2^{k-1} - 1 apart:
5 #0. 2^{k-1} - 1, 2(2^{k-1} - 1), 3(2^{k-1} - 1), ..., then
\sharp #1, 1 + 2<sup>k-1</sup> - 1, 1 + 2(2<sup>k-1</sup> - 1), 1 + 3(2<sup>k-1</sup> - 1), ...
insertion sort (2^0 = 1 apart), but with most inversions
^{/2}) (take CS170 for why!).
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```

Inversions

```
N) comparisons if already sorted.
pical implementation for arrays:
: i < A.length: i += 1) {
; j >= 0; j -= 1) {
ompareTo(x) <= 0) /* (1) */
                  /* (2) */
[j];
kecutes for each j \approx how far x must move.
within K of proper places, then takes O(KN) operations.
r any amount of nearly sorted data.
of unsortedness: # of inversions: pairs that are out
when sorted, N(N-1)/2 when reversed).
ion of (2) decreases inversions by 1.
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```

