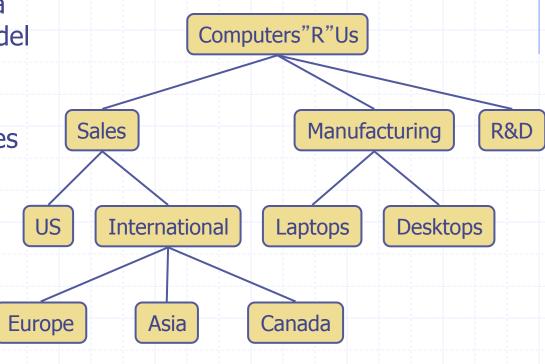


What is a Tree

 In computer science, a tree is an abstract model of a hierarchical structure

A tree consists of nodes with a parent-child relation

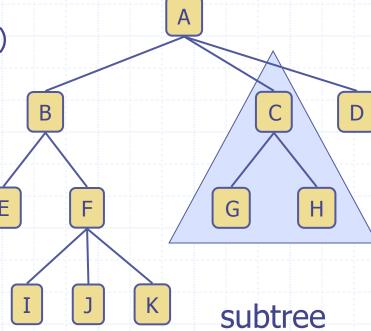
- Applications:
 - Organization charts
 - File systems
 - Programming environments



Tree Terminology

- Root: node without parent (A)
- Internal node: node with at least one child (A, B, C, F)
- <u>External node</u> (a.k.a. leaf): node
 without children (E, I, J, K, G, H, D)
- Ancestors of a node: parent, grandparent, grand-grandparent, etc.
- Depth of a node: number of ancestors
- Height of a tree: maximum depth of any node (3)
- Descendant of a node: child, grandchild, grand-grandchild, etc.

 Subtree: tree consisting of a node and its descendants



Tree ADT

- We use positions to abstract nodes
- Generic methods:
 - Integer len()
 - Boolean is_empty()
 - Iterator positions()
 - Iterator iter()
- Accessor methods:
 - position root()
 - position parent(p)
 - Iterator children(p)
 - Integer num_children(p)

- Query methods:
 - Boolean is_leaf(p)
 - Boolean is_root(p)
- Update method:
 - element replace (p, o)
- Additional update methods may be defined by data structures implementing the Tree ADT

Abstract Nested Position Class in Python

```
#ADT Tree "interface"
class Tree:
    class Position:
        def element( self ):
            pass
        def __eq_ ( self, other ):
            pass
        def __ne__( self, other):
            return not( self == other )
                                            5
                     Trees
```

Abstract Tree Class in Python...

```
#ADT Tree "interface"
class Tree:
    def root( self ):
        pass
    def parent( self, p ):
        pass
    def num children( self, p ):
        pass
    def children( self, p ):
        pass
    def len ( self ):
        pass
    def is root( self, p ):
        return self.root() == p
    def is leaf( self, p ):
        return self.num children() == 0
    def is empty( self ):
        return len( self ) == 0
                   Trees
```

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Abstract Tree Class in Python

```
def depth( self, p ):
      #returns the number of ancestors of p
      if self.is root( p ):
        return 0
      else:
        return 1 + self.depth( self.parent() )
def height1( self, p ):
      #returns the maximum depth of the leaf positions
      return max( self.depth( p ) for p in self.positions() if self.is leaf( p ))
def height2( self, p ):
      #returns the height of the subtree at Position p
      if self.is leaf( p ):
        return 0
      else:
        return 1 + max( self.height2( c ) for c in self.children( p ) )
def height( self, p = None ):
      #returns the height of the subtree rooted at Position p
      #if p is None, then the height of the entire tree
      if p is None:
       p = self.root()
      return self.height2( p )
```

Trees

Preorder Traversal

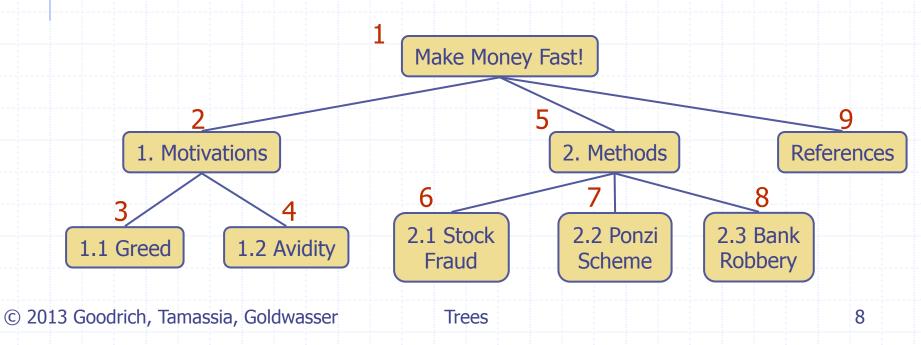
- A traversal visits the nodes of a tree in a systematic manner
- In a preorder traversal, a node is visited before its descendants
- Application: print a structured document

Algorithm preOrder(v)

visit(v)

for each child w of v

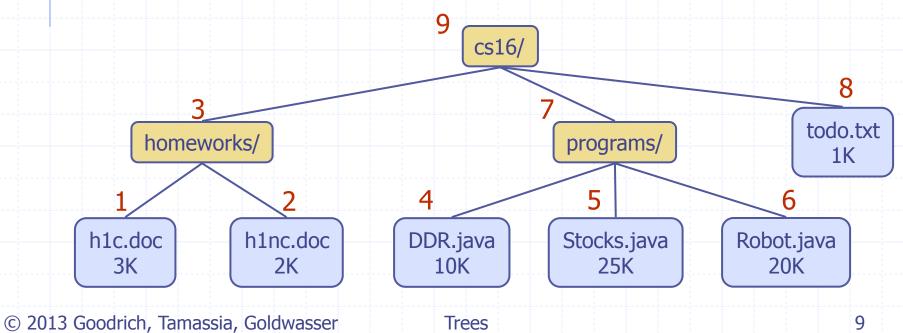
preorder (w)



Postorder Traversal

- In a postorder traversal, a node is visited after its descendants
- Application: compute space used by files in a directory and its subdirectories

Algorithm postOrder(v)
for each child w of v
postOrder (w)
visit(v)



pre- and postorder traversal in Python

```
def preorder_print( self, p ):
    print( p )
    for c in self.children( p ):
        preorder_print( c )

def postorder_print( self, p ):
    for c in self.children( p ):
        postorder_print( c )
    print( p )
```

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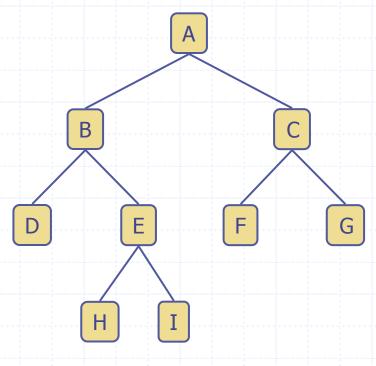
Breadth-first traversal in Python

```
def breadth_first_print( self ):
            Q = ArrayQueue()
            Q.enqueue( self.root() )
                while not Q.is empty():
                      p = Q.dequeue()
                      print( p )
                      for c in self.children( p ):
                           Q.enqueue(c)
                      Make Money Fast!
    1. Motivations
                                 2. Methods
                                                References
                                 2.2 Ponzi
                                           2.3 Bank
                      2.1 Stock
          1.2 Avidity
1.1 Greed
                       Fraud
                                  Scheme
                                            Robbery
                         Trees
                                                     11
```

Binary Trees

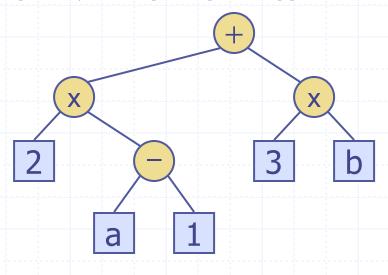
- A binary tree is a tree with the following properties:
 - Each internal node has at most two children (exactly two for proper binary trees)
 - The children of a node are an ordered pair
- We call the children of an internal node left child and right child
- Alternative recursive definition: a binary tree is either
 - a tree consisting of a single node, or
 - a tree whose root has an ordered pair of children, each of which is a binary tree

- Applications:
 - arithmetic expressions
 - decision processes
 - searching



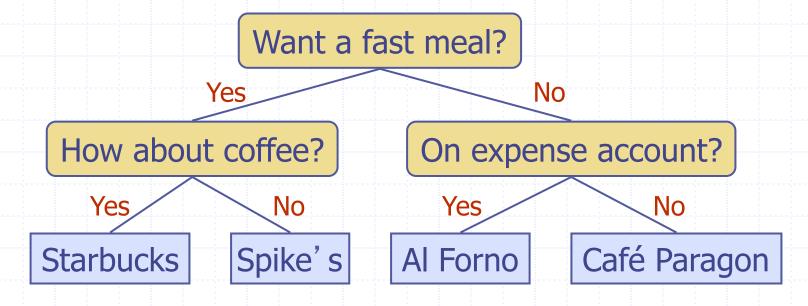
Arithmetic Expression Tree

- Binary tree associated with an arithmetic expression
 - internal nodes: operators
 - external nodes: operands
- Example: arithmetic expression tree for the expression $(2 \times (a 1) + (3 \times b))$



Decision Tree

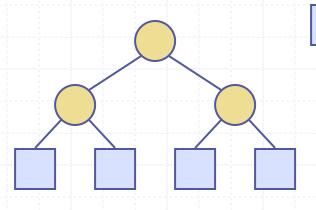
- Binary tree associated with a decision process
 - internal nodes: questions with yes/no answer
 - external nodes: decisions
- Example: dining decision



Properties of Proper Binary Trees

- Notation
 - n number of nodes
 - e number of external nodes
 - i number of internal nodes

h height



Properties:

$$e = i + 1$$

$$n = 2e - 1$$

■
$$h \leq i$$

■
$$h \le (n-1)/2$$

$$e \le 2^h$$

■
$$h \leq \log_2 e$$

BinaryTree ADT

- The BinaryTree ADT extends the Tree
 ADT, i.e., it inherits all the methods of the Tree ADT
- Additional methods:
 - position left(p)
 - position right(p)
 - position sibling(p)

 Update methods may be defined by data structures implementing the BinaryTree ADT

Abstract BinaryTree Class in Python

from Tree import Tree

```
class BinaryTree( Tree ):
    def left( self, p ):
        pass
    def right( self, p ):
        pass
    def sibling( self, p ):
        #return the sibling Position
        parent = self.parent()
        if parent is None:
            return None
        else:
            if p == self.left( parent ):
                return self.right( parent )
            else:
                return self.left( parent )
    def children( self, p ):
         if self.left( p ) is not None:
            yield self.left( p )
        if self.right( p ) is not None:
            yield self.right( p )
```

Inorder Traversal

- In an inorder traversal a node is visited after its left subtree and before its right subtree
- Application: draw a binary tree
 - x(v) = inorder rank of v
 - y(v) = depth of v

Algorithm *inOrder(v)*

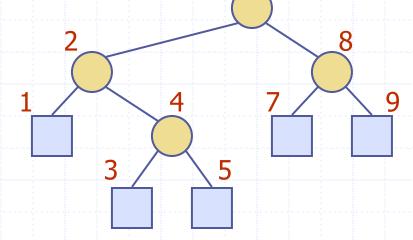
if v has a left child

inOrder(left(v))

visit(v)

if v has a right child

inOrder (right (v))



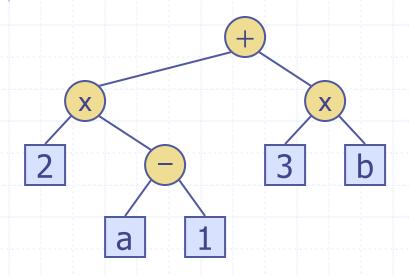
inorder_print in Python

```
def inorder_print( self, p ):
    if self.left( p ) is not None:
        self.inorder_print( self.left( p ) )
    print( p )
    if self.right( p ) is not None:
        self.inorder_print( self.right( p ) )
```

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Print Arithmetic Expressions

- Specialization of an inorder traversal
 - print operand or operator when visiting node
 - print "(" before traversing left subtree
 - print ")" after traversing right subtree



Algorithm printExpression(v)

if v has a left child

print("('')

inOrder (left(v))

print(v.element ())

if v has a right child

inOrder (right(v))

print (")'')

$$((2 \times (a - 1)) + (3 \times b))$$

Evaluate Arithmetic Expressions

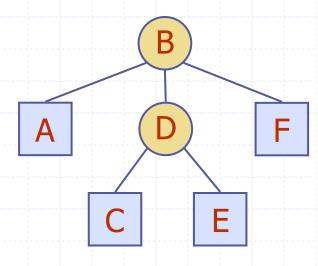
- Specialization of a postorder traversal
 - recursive method returning the value of a subtree
 - when visiting an internal node, combine the values of the subtrees

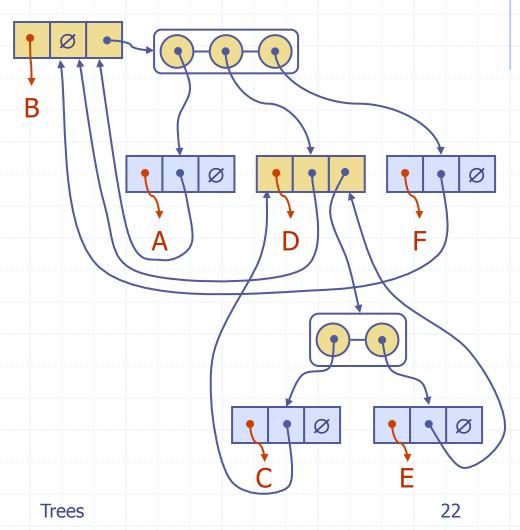
```
x x x 2 2 5 1
```

```
Algorithm evalExpr(v)
if is_leaf (v)
return v.element ()
else
x = evalExpr(left (v))
y = evalExpr(right (v))
op = operator stored at v
return x op y
```

Linked Structure for Trees

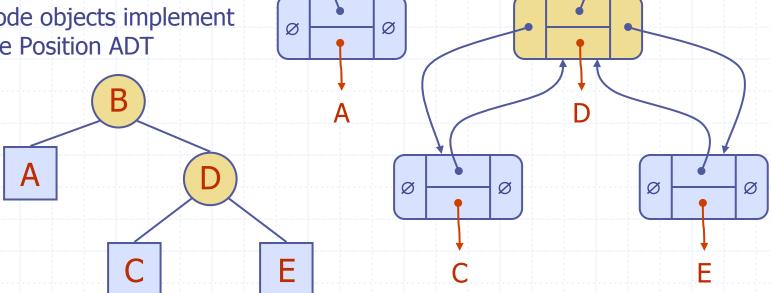
- A node is represented by an object storing
 - Element
 - Parent node
 - Sequence of children nodes
- Node objects implement the Position ADT





Linked Structure for Binary Trees

- A node is represented by an object storing
 - Element
 - Parent node
 - Left child node
 - Right child node
- Node objects implement the Position ADT



B

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Trees

LinkedBinaryTree Class in Python

```
from BinaryTree import BinaryTree
class LinkedBinaryTree( BinaryTree ):
    class Node:
        def __init__( self, element,
                      parent = None,
                      left = None,
                      right = None ):
            self. element = element
            self._parent = parent
            self. left = left
            self. right = right
```

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Nested Abstract Position

```
class Position( BinaryTree.Position ):
 def init ( self, container, node ):
   self._container = container
   self. node = node
 def str ( self ):
   return str( self. node. element )
 def element( self ):
    return self. node. element
 def eq ( self, other ):
   return type( other ) is type( self ) and other._node is self._node
                                Trees
                                                               25
```

Validate and MakePosition

```
def validate( self, p ):
  #return associated node if position is valid
  if not isinstance( p, self.Position ):
    raise TypeError( 'p must be proper Position type' )
  if p. container is not self:
    raise ValueError( 'p does not belong to this container')
  if p. node. parent is p. node:
    raise ValueError( 'p is no longer valid' )
  return p. node
def make position( self, node ):
  #return Position instance for given node (None if no node)
  return self.Position( self, node ) if node is not None else None
```

Trees

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```
def init ( self ):
    #create an initially empty binary tree
    self. root = None
    self. size = 0
def len ( self ):
    return self. size
def root( self ):
    return self. make position( self. root )
def parent( self, p ):
    node = self. validate( p )
    return self. make position( node. parent )
def left( self, p ):
    node = self. validate( p )
    return self. make position( node. left )
def right( self, p ):
    node = self._validate( p )
    return self. make position( node. right )
```

Trees

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```
def num children( self, p ):
    node = self._validate( p )
    count = 0
    if node._left is not None:
        count += 1
    if node._right is not None:
        count += 1
    return count
def _add_root( self, e ):
    if self. root is not None:
      raise ValueError( 'Root exists' )
    self. size = 1
    self._root = self._Node( e )
    return self. make position( self. root )
                    Trees
                                              28
```

```
def add left( self, p, e ):
    node = self. validate( p )
    if node. left is not None:
      raise ValueError( 'Left child exists' )
    self. size += 1
    node._left = self._Node( e, node )
    return self. make position( node. left )
def _add_right( self, p, e ):
    node = self._validate( p )
    if node._right is not None:
      raise ValueError( 'Right child exists' )
    self. size += 1
    node. right = self._Node( e, node )
    return self. make position( node. right )
                    Trees
                                              29
```

```
def replace( self, p, e ):
    node = self. validate( p )
    old = node. element
    node. element = e
    return old
def delete( self, p ):
    #remove node p and replace it with its child if any
    node = self. validate( p )
    if self. num children( p ) == 2:
        raise ValueError( 'p has two children' )
    child = node. left if node. left else node. right
    if child is not None:
        child. parent = node. parent
    if node is self. root:
        self. root = child
    else:
        parent = node. parent
        if node is parent. left:
            parent. left = child
        else:
            parent. right = child
    self. size -= 1
    node. parent = node
    return node. element
                    Trees
                                                         30
```

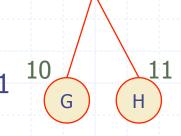
```
def _attach( self, p, t1, t2 ):
#attach trees t1 and t2 as left and right of leaf p
     node = self. validate( p )
     if not self.is leaf( p ):
         raise ValueError( 'position must be leaf' )
     if not type( self ) is type( t1 ) is type( t2 ):
         raise TypeError( 'Tree types must match' )
     self. size += len(t1) + len(t2)
     if not t1.is empty():
         t1._root._parent = node
         node. left = t1. root
         t1._root = None #set t1 to empty
         t1. size = 0
     if not t2.is empty():
         t2._root._parent = node
         node._right = t2._root
         t2._root = None #set t2 to empty
         t2._size = 0
                                                 31
```

Array-Based Representation of Binary Trees

Nodes are stored in an array A



- □ Node v is stored at A[rank(v)]
 - rank(root) = 1
 - if node is the left child of parent(node), rank(node) = 2 x rank(parent(node))
 - if node is the right child of parent(node), rank(node) = 2 x rank(parent(node)) + 1



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