KAIST,HYU - Freshman Dognodap

Team Note of Freshman Dognodap

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Compiled on July 14, 2022

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ALL BELOW HERE ARE USELESS IF YOU READ THE STATEMENT WRONG

1 Data Structure

1.1 Erasable Priority Queue

```
template<typename T, T inf>
struct pq_set{
 priority_queue<T, vector<T>, greater<T>> in, out; // min heap, inf = 1e18
 // priority_queue<T> in, out; // max heap, inf = -1e18
 pg_set(){ in.push(inf); }
  void insert(T v){ in.push(v); }
 void erase(T v){ out.push(v); }
  while(out.size() && in.top() == out.top()) in.pop(), out.pop();
   return in.top();
  bool empty(){
  while(out.size() && in.top() == out.top()) in.pop(), out.pop();
   return in.top() == inf;
 }
};
1.2 Persistent Segment Tree
struct PSTNode{
 PSTNode *1. *r: int v:
 PSTNode() \{ 1 = r = nullptr; v = 0; \}
PSTNode *root[101010]:
PST(){ memset(root, 0, sizeof root); } // constructor
void init(PSTNode *node, int s, int e){
 if(s == e) return:
 int m = s + e \gg 1;
 node->1 = new PSTNode: node->r = new PSTNode:
 init(node->1, s, m); init(node->r, m+1, e);
void update(PSTNode *prv, PSTNode *now, int s, int e, int x){
 if(s == e){ now->v = prv ? prv->v + 1 : 1; return; }
 int m = s + e >> 1:
 if(x \le m)
   now->1 = new PSTNode; now->r = prv->r;
    update(prv->1, now->1, s, m, x);
 }
  else{
   now->r = new PSTNode: now->l = prv->l:
    update(prv->r, now->r, m+1, e, x);
  int t1 = now->1 ? now->1->v : 0;
  int t2 = now -> r ? now -> r -> v : 0;
 now->v = t1 + t2:
int kth(PSTNode *prv, PSTNode *now, int s, int e, int k){
 if(s == e) return s:
 int m = s + e >> 1, diff = now->l->v - prv->l->v;
 if(k <= diff) return kth(prv->1, now->1, s, m, k);
  else return kth(prv->r, now->r, m+1, e, k-diff);
```

2 Flows, Matching

```
2.1 Hopcroft-Karp Bipartite Matching
```

```
const int MAXN = 50005. MAXM = 50005:
vector<int> gph[MAXN];
int dis[MAXN], 1[MAXN], r[MAXM], vis[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int 1, int r){ gph[1].push_back(r); }
bool bfs(int n){
 queue<int> que;
 bool ok = 0;
 memset(dis, 0, sizeof(dis)):
 for(int i=0; i<n; i++){
   if(l[i] == -1 && !dis[i]){
     que.push(i);
     dis[i] = 1;
 while(!que.empty()){
   int x = que.front();
   que.pop();
   for(auto &i : gph[x]){
     if(r[i] == -1) ok = 1:
     else if(!dis[r[i]]){
       dis[r[i]] = dis[x] + 1:
        que.push(r[i]);
   }
 }
 return ok:
bool dfs(int x){
 if(vis[x]) return 0:
 vis[x] = 1;
 for(auto &i : gph[x]){
   if(r[i] == -1 \mid | (!vis[r[i]] && dis[r[i]] == dis[x] + 1 && dfs(r[i])))
     1[x] = i; r[i] = x;
     return 1:
   }
 }
 return 0;
int match(int n){
 memset(1, -1, sizeof(1));
 memset(r, -1, sizeof(r));
 int ret = 0:
 while(bfs(n)){
   memset(vis, 0, sizeof(vis));
   for(int i=0; i<n; i++) if(l[i] == -1 && dfs(i)) ret++;
 return ret;
bool chk[MAXN + MAXM];
void rdfs(int x, int n){
 if(chk[x]) return;
 chk[x] = 1;
 for(auto &i : gph[x]){
```

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```
chk[i + n] = 1:
   rdfs(r[i], n);
 }
vector<int> getcover(int n, int m){ // solve min. vertex cover
 match(n):
 memset(chk, 0, sizeof(chk));
 for(int i=0; i<n; i++) if(l[i] == -1) rdfs(i, n);
 vector<int> v:
 for(int i=0; i<n; i++) if(!chk[i]) v.push_back(i);</pre>
 for(int i=n; i<n+m; i++) if(chk[i]) v.push_back(i);</pre>
 return v:
2.2 Dinic's Algorithm
const int MAXN = 505;
struct edg{ int pos, cap, rev; };
vector<edg> gph[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int s, int e, int x){
 gph[s].push_back({e, x, (int)gph[e].size()});
 gph[e].push_back({s, 0, (int)gph[s].size()-1});
int dis[MAXN], pnt[MAXN];
bool bfs(int src, int sink){
 memset(dis, 0, sizeof(dis));
 memset(pnt, 0, sizeof(pnt));
 queue<int> que:
 que.push(src);
 dis[src] = 1:
  while(!que.empty()){
   int x = que.front();
   que.pop();
   for(auto &e : gph[x]){
     if(e.cap > 0 && !dis[e.pos]){
       dis[e.pos] = dis[x] + 1;
       que.push(e.pos);
   }
 }
 return dis[sink] > 0:
int dfs(int x, int sink, int f){
 if(x == sink) return f:
 for(; pnt[x] < gph[x].size(); pnt[x]++){</pre>
   edg e = gph[x][pnt[x]];
   if(e,cap > 0 \&\& dis[e,pos] == dis[x] + 1){
     int w = dfs(e.pos, sink, min(f, e.cap));
     if(w){
       gph[x][pnt[x]].cap -= w;
       gph[e.pos][e.rev].cap += w;
       return w:
   }
 return 0;
```

```
lint match(int src, int sink){
 lint ret = 0;
 while(bfs(src, sink)){
   int r:
   while((r = dfs(src, sink, 2e9))) ret += r;
 return ret;
2.3 Min Cost Max Flow
const int MAXN = 100;
struct edg{ int pos, cap, rev, cost; };
vector<edg> gph[MAXN];
void clear(){
 for(int i=0; i<MAXN; i++) gph[i].clear();</pre>
void add edge(int s. int e. int x. int c){
 gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int dist[MAXN], pa[MAXN], pe[MAXN];
bool inque[MAXN];
bool spfa(int src, int sink){
 memset(dist, 0x3f, sizeof(dist));
 memset(inque, 0, sizeof(inque));
 queue<int> que;
 dist[src] = 0;
 inque[src] = 1;
 que.push(src);
 bool ok = 0:
 while(!que.empty()){
   int x = que.front();
   que.pop();
   if(x == sink) ok = 1;
   inque[x] = 0;
   for(int i=0; i<gph[x].size(); i++){</pre>
     edg e = gph[x][i];
     if(e.cap > 0 \&\& dist[e.pos] > dist[x] + e.cost){
       dist[e.pos] = dist[x] + e.cost:
        pa[e.pos] = x;
        pe[e.pos] = i;
        if(!inque[e.pos]){
         inque[e.pos] = 1;
         que.push(e.pos);
     }
   }
 }
 return ok;
int match(int src, int sink){
 int ret = 0:
 while(spfa(src, sink)){
   int cap = 1e9:
   for(int pos = sink; pos != src; pos = pa[pos]){
      cap = min(cap, gph[pa[pos]][pe[pos]].cap);
```

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```
ret += dist[sink] * cap:
   for(int pos = sink; pos != src; pos = pa[pos]){
     int rev = gph[pa[pos]][pe[pos]].rev;
     gph[pa[pos]][pe[pos]].cap -= cap;
     gph[pos][rev].cap += cap;
 }
 return ret;
2.4 Hell-Joseon style MCMF
const int MAXN = 100:
struct edg{ int pos, cap, rev, cost; };
vector<edg> gph[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int s, int e, int x, int c){
 gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int phi[MAXN], inque[MAXN], dist[MAXN];
void prep(int src, int sink){
 memset(phi, 0x3f, sizeof(phi));
 memset(dist, 0x3f, sizeof(dist)):
 queue<int> que;
 que.push(src):
 inque[src] = 1;
  while(!que.empty()){
   int x = que.front();
   que.pop();
   inque[x] = 0;
   for(auto &i : gph[x]){
     if(i.cap > 0 && phi[i.pos] > phi[x] + i.cost){
       phi[i.pos] = phi[x] + i.cost;
       if(!inque[i.pos]){
         inque[i.pos] = 1;
          que.push(i.pos);
     }
   }
 for(int i=0: i<MAXN: i++){</pre>
   for(auto &j : gph[i]){
      if(j.cap > 0) j.cost += phi[i] - phi[j.pos];
   }
 }
 priority_queue<pi, vector<pi>, greater<pi> > pq;
 pq.push(pi(0, src));
 dist[src] = 0;
 while(!pq.empty()){
   auto 1 = pq.top();
   pq.pop();
   if(dist[l.second] != l.first) continue:
   for(auto &i : gph[l.second]){
     if(i.cap > 0 && dist[i.pos] > 1.first + i.cost){
       dist[i.pos] = 1.first + i.cost;
       pq.push(pi(dist[i.pos], i.pos));
```

```
}
}
bool vis[MAXN];
int ptr[MAXN];
int dfs(int pos, int sink, int flow){
 vis[pos] = 1;
 if(pos == sink) return flow;
 for(; ptr[pos] < gph[pos].size(); ptr[pos]++){</pre>
   auto &i = gph[pos][ptr[pos]];
   if(!vis[i.pos] && dist[i.pos] == i.cost + dist[pos] && i.cap > 0){
      int ret = dfs(i.pos, sink, min(i.cap, flow));
      if(ret != 0){
        i.cap -= ret:
        gph[i.pos][i.rev].cap += ret;
        return ret;
   }
 }
 return 0;
int match(int src, int sink, int sz){
 prep(src, sink);
 for(int i=0; i<sz; i++) dist[i] += phi[sink] - phi[src];</pre>
 int ret = 0:
 while(true){
   memset(ptr, 0, sizeof(ptr));
   memset(vis, 0, sizeof(vis));
   int tmp = 0;
    while((tmp = dfs(src, sink, 1e9))){
     ret += dist[sink] * tmp;
     memset(vis, 0, sizeof(vis));
    tmp = 1e9:
   for(int i=0: i<sz: i++){
     if(!vis[i]) continue:
     for(auto &j : gph[i]){
        if(j.cap > 0 && !vis[j.pos]){
          tmp = min(tmp, (dist[i] + j.cost) - dist[j.pos]);
     }
   }
   if(tmp > 1e9 - 200) break;
   for(int i=0; i<sz; i++){</pre>
      if(!vis[i]) dist[i] += tmp;
   }
 }
 return ret;
2.5 Circulation Problem
maxflow mf:
lint lsum;
void clear(){
 1sum = 0:
 mf.clear();
```

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```
void add edge(int s, int e, int l, int r){
 lsum += 1:
 mf.add\_edge(s + 2, e + 2, r - 1);
 mf.add_edge(0, e + 2, 1);
 mf.add_edge(s + 2, 1, 1);
bool solve(int s, int e){
 mf.add\_edge(e+2, s+2, 1e9); // to reduce as maxflow with lower bounds, in circulation
problem skip this line
 return lsum == mf.match(0, 1);
 // to get maximum LR flow, run maxflow from s+2 to e+2 again
2.6 Min Cost Circulation
// Cycle canceling (Dual of successive shortest path)
// Time complexity is ridiculously high (F * maxC * nm^2). But runs reasonably in practice
(V = 70 \text{ in } 1s)
struct edg{ int pos, cap, rev, cost; };
vector<edg> gph[MAXN];
void clear(){ for(int i=0; i<MAXN; i++) gph[i].clear(); }</pre>
void add_edge(int s, int e, int x, int c){
 gph[s].push_back({e, x, (int)gph[e].size(), c});
 gph[e].push_back({s, 0, (int)gph[s].size()-1, -c});
int dist[MAXN], par[MAXN], pae[MAXN];
int negative_cycle(int n){
 bool mark[MAXN] = {};
 memset(dist, 0, sizeof(dist));
 int upd = -1;
 for(int i=0: i<=n: i++){</pre>
   for(int j=0; j<n; j++){
     int idx = 0;
     for(auto &k : gph[j]){
       if(k.cap > 0 \&\& dist[k.pos] > dist[j] + k.cost){
          dist[k.pos] = dist[j] + k.cost;
         par[k.pos] = j;
         pae[k.pos] = idx;
          if(i == n){
            upd = i:
            while(!mark[upd]){
             mark[upd] = 1;
             upd = par[upd];
           }
           return upd;
         }
       }
       idx++:
     }
 }
 return -1;
int match(int n){
 int rt = -1:
 int ans = 0;
 while(~(rt = negative_cycle(n))){
   bool mark[MAXN] = {}:
```

```
vector<pi> cvc:
    while(!mark[rt]){
     cyc.push_back(pi(par[rt], pae[rt]));
     mark[rt] = 1:
     rt = par[rt];
   reverse(cyc.begin(), cyc.end());
   int capv = 1e9;
   for(auto &i : cyc){
     auto e = &gph[i.first][i.second];
      capv = min(capv, e->cap);
   for(auto &i : cyc){
      auto e = &gph[i.first][i.second];
     e->cap -= capv;
     gph[e->pos][e->rev].cap += capv;
     ans += e->cost * capv;
   }
 }
 return ans;
3
   Graph
3.1 \quad 2-\overline{SAT}
strongly_connected scc;
int n: // = number of clauses
void init(int _n){ scc.clear(); n = _n; }
int NOT(int x) { return x \ge n ? (x - n) : (x + n) : }
void add_edge(int x, int y){ // input ~x to denote NOT
 if((x >> 31) & 1) x = (^x) + n;
 if((y >> 31) & 1) y = (~y) + n;
 scc.add_edge(x, y), scc.add_edge(NOT(y), NOT(x));
bool satisfy(vector<bool> &res){
 res.resize(n);
 scc.get_scc(2*n);
 for(int i=0: i<n: i++){
   if(scc.comp[i] == scc.comp[NOT(i)]) return 0;
   if(scc.comp[i] < scc.comp[NOT(i)]) res[i] = 0;</pre>
   else res[i] = 1;
 }
 return 1:
3.2 BCC
void color(int x, int p){
 if(p){
   bcc[p].push_back(x);
   cmp[x].push_back(p);
 for(auto &i : gph[x]){
   if(cmp[i].size()) continue;
   if(low[i] >= dfn[x]){
     bcc[++c].push_back(x);
      cmp[x].push_back(c);
      color(i, c);
```

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```
else color(i, p);
 }
}
3.3 Splay Tree + Link-Cut Tree
// Checklist 1. Is it link cut, or splay?
// Checklist 2. In link cut, is son always root?
void rotate(node *x){
  if(!x->p) return;
  push(x->p); // if there's lazy stuff
 push(x);
  node *p = x->p;
  bool is_left = (p->1 == x);
  node *b = (is_left ? x->r : x->l);
  x->p = p->p;
  if (x-p \&\& x-p-1 == p) x-p-1 = x;
  if (x-p \&\& x-p-r == p) x-p-r = x;
  if(is_left){
   if(b) b \rightarrow p = p;
   p->1 = b;
   p->p = x;
   x->r = p;
  else{
   if(b) b \rightarrow p = p;
   p->r = b;
   p->p = x;
   x->1 = p;
  pull(p); // if there's something to pull up
  pull(x);
  if(!x->p) root = x; // IF YOU ARE SPLAY TREE
  if(p->pp){ // IF YOU ARE LINK CUT TREE
   x->pp = p->pp;
   p->pp = NULL;
void splay(node *x){
  while(x->p){
    node *p = x->p;
   node *g = p \rightarrow p;
    if(g){
      if((p\rightarrow l == x) ^ (g\rightarrow l == p)) rotate(x);
      else rotate(p);
   }
    rotate(x);
 }
}
void access(node *x){
  splay(x);
  push(x);
  if(x->r){
   x->r->pp = x;
   x->r->p = NULL;
   x->r = NULL;
  pull(x);
```

```
while(x->pp){
   node *nxt = x->pp;
   splay(nxt);
   push(nxt);
   if(nxt->r){
     nxt->r->pp = nxt;
     nxt->r->p = NULL;
     nxt->r = NULL;
   nxt->r = x;
   x->p = nxt;
   x->pp = NULL;
   pull(nxt);
   splay(x);
 }
node *root(node *x){
 access(x);
 while (x->1) {
   push(x);
   x = x->1;
 access(x);
 return x;
node *par(node *x){
 access(x);
 if(!x->1) return NULL;
 push(x);
 x = x->1;
 while(x->r){
   push(x);
   x = x->r;
 }
 access(x);
 return x;
node *lca(node *s, node *t){
 access(s);
 access(t);
 splay(s);
 if(s->pp == NULL) return s;
 return s->pp;
void link(node *par, node *son){
 access(par);
 access(son);
 son->rev ^= 1; // remove if needed
 push(son);
 son->1 = par;
 par->p = son;
 pull(son);
void cut(node *p){
 access(p);
 push(p);
```

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```
if(p->1){
   p->1->p = NULL;
   p->1 = NULL;
 pull(p);
3.4 Offline Dynamic MST
int n, m, q;
int st[MAXN], ed[MAXN], cost[MAXN], chk[MAXN];
pi qr[MAXN];
bool cmp(int &a, int &b) { return pi(cost[a], a) < pi(cost[b], b); }
void contract(int s, int e, vector<int> v, vector<int> &must_mst, vector<int> &maybe_mst){
 sort(v.begin(), v.end(), cmp);
 vector<pi> snapshot:
 for(int i=s; i<=e; i++) disj.uni(st[qr[i].first], ed[qr[i].first], snapshot);</pre>
 for(auto &i : v) if(disj.uni(st[i], ed[i], snapshot)) must_mst.push_back(i);
 disj.revert(snapshot);
 for(auto &i : must_mst) disj.uni(st[i], ed[i], snapshot);
 for(auto &i : v) if(disj.uni(st[i], ed[i], snapshot)) maybe_mst.push_back(i);
 disi.revert(snapshot):
void solve(int s, int e, vector<int> v, lint cv){
 if(s == e){}
   cost[qr[s].first] = qr[s].second;
   if(st[qr[s].first] == ed[qr[s].first]){
     printf("%lld\n", cv):
     return;
   }
   int minv = qr[s].second;
   for(auto &i : v) minv = min(minv, cost[i]);
   printf("%lld\n",minv + cv);
   return:
 int m = (s+e)/2:
 vector<int> lv = v. rv = v:
 vector<int> must_mst, maybe_mst;
 for(int i=m+1; i<=e; i++){</pre>
   chk[qr[i].first]--;
   if(chk[qr[i].first] == 0) lv.push_back(qr[i].first);
 vector<pi> snapshot;
 contract(s, m, lv, must_mst, maybe_mst);
  for(auto &i : must_mst) lcv += cost[i], disj.uni(st[i], ed[i], snapshot);
 solve(s, m, maybe_mst, lcv);
 disj.revert(snapshot);
 must_mst.clear(); maybe_mst.clear();
 for(int i=m+1; i<=e; i++) chk[qr[i].first]++;</pre>
 for(int i=s; i<=m; i++){</pre>
   chk[ar[i].first]--:
   if(chk[qr[i].first] == 0) rv.push_back(qr[i].first);
 }
 lint rcv = cv;
```

```
contract(m+1, e, rv, must mst, maybe mst):
 for(auto &i : must_mst) rcv += cost[i], disj.uni(st[i], ed[i], snapshot);
 solve(m+1, e, maybe_mst, rcv);
 disj.revert(snapshot);
 for(int i=s; i<=m; i++) chk[qr[i].first]++;</pre>
int main(){
 scanf("%d %d",&n,&m);
 vector<int> ve;
 for(int i=0: i<m: i++){</pre>
   scanf("%d %d %d".&st[i].&ed[i].&cost[i]):
 scanf("%d",&q);
 for(int i=0; i<q; i++){</pre>
   scanf("%d %d",&gr[i].first,&gr[i].second);
   ar[i].first--:
   chk[qr[i].first]++;
 disj.init(n);
 for(int i=0; i<m; i++) if(!chk[i]) ve.push_back(i);</pre>
 solve(0, q-1, ve, 0);
   Strings
4.1 Aho-Corasick Algorithm
const int MAXN = 100005, MAXC = 26;
int trie[MAXN][MAXC], fail[MAXN], term[MAXN], piv;
void init(vector<string> &v){
 memset(trie, 0, sizeof(trie));
 memset(fail, 0, sizeof(fail));
 memset(term, 0, sizeof(term));
 piv = 0;
 for(auto &i : v){
   int p = 0;
   for(auto &j : i){
     if(!trie[p][j]) trie[p][j] = ++piv;
     p = trie[p][j];
   term[p] = 1;
 queue<int> que;
 for(int i=0; i<MAXC; i++){</pre>
   if(trie[0][i]) que.push(trie[0][i]);
 while(!que.empty()){
   int x = que.front();
   que.pop();
   for(int i=0; i<MAXC; i++){</pre>
     if(trie[x][i]){
       int p = fail[x]:
       while(p && !trie[p][i]) p = fail[p];
       p = trie[p][i];
       fail[trie[x][i]] = p;
       if(term[p]) term[trie[x][i]] = 1;
        que.push(trie[x][i]);
```

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```
}
   }
 }
bool query(string &s){
  int p = 0;
  for(auto &i : s){
    while(p && !trie[p][i]) p = fail[p];
   p = trie[p][i];
   if(term[p]) return 1;
 }
 return 0:
}
4.2 Suffix Array
const int MAXN = 500005;
int ord[MAXN]. nord[MAXN]. cnt[MAXN]. aux[MAXN]:
void solve(int n, char *str, int *sfx, int *rev, int *lcp){
  int p = 1;
  memset(ord, 0, sizeof(ord));
  for(int i=0; i<n; i++){</pre>
   sfx[i] = i:
   ord[i] = str[i]:
 }
  int pnt = 1;
  while(1){
    memset(cnt, 0, sizeof(cnt));
   for(int i=0; i<n; i++) cnt[ord[min(i+p, n)]]++;</pre>
    for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
   for(int i=n-1: i>=0: i--)
      aux[--cnt[ord[min(i+p, n)]]] = i;
    memset(cnt, 0, sizeof(cnt));
    for(int i=0; i<n; i++) cnt[ord[i]]++;</pre>
    for(int i=1; i<=n || i<=255; i++) cnt[i] += cnt[i-1];
    for(int i=n-1; i>=0; i--)
     sfx[--cnt[ord[aux[i]]]] = aux[i]:
    if(pnt == n) break;
    pnt = 1:
    nord[sfx[0]] = 1:
    for(int i=1; i<n; i++){</pre>
     if(ord[sfx[i-1]] != ord[sfx[i]] || ord[sfx[i-1] + p] != ord[sfx[i] + p]){
        pnt++;
      nord[sfx[i]] = pnt;
   memcpy(ord, nord, sizeof(int) * n);
   p *= 2:
 }
  for(int i=0; i<n; i++) rev[sfx[i]] = i;</pre>
  int h = 0:
  for(int i=0; i<n; i++){</pre>
   if(rev[i]){
      int prv = sfx[rev[i] - 1];
      while(str[prv + h] == str[i + h]) h++;
     lcp[rev[i]] = h;
    }
    h = \max(h-1, 0);
```

```
}
4.3 Manacher's Algorithm
const int MAXN = 1000005;
int aux[2 * MAXN - 1];
void solve(int n, int *str, int *ret){
 // *ret : number of nonobvious palindromic character pair
 for(int i=0: i<n: i++){
   aux[2*i] = str[i];
   if(i != n-1) aux[2*i+1] = -1;
 int p = 0, c = 0;
 for(int i=0; i<2*n-1; i++){
   int cur = 0;
   if(i <= p) cur = min(ret[2 * c - i], p - i);
   while(i - cur - 1 >= 0 && i + cur + 1 < 2*n-1 && aux[i-cur-1] == aux[i+cur+1]){
   ret[i] = cur:
   if(i + ret[i] > p){
     p = i + ret[i]:
     c = i;
 }
4.4 Suffix Array (Linear time)
 Should be revised.
class SuffixArray {
public:
 int A[7 * N / 10], B[7 * N / 10], cnt[N + 2], SAV[N];
 int mem[5 * N]; int* mem_pt = mem;
 void clear(int n){
   int *ptr = mem;
   while(ptr != mem_pt){
     *ptr = 0;
     ptr++;
    mem_pt = mem;
   for(int i=0; i<n+10 && i < 7 * N / 10; i++) A[i] = B[i] = 0;
   for(int i=0; i<n+2; i++) cnt[i] = 0;
   for(int i=0; i<n; i++) SAV[i] = 0;
 inline int* mloc(size_t sz) {
   int* ret = mem_pt; mem_pt = mem_pt + sz;
   return ret:
 }
 void rsort(int* a, int* b, int* dat, int n, int k) {
   for (int i = 0; i <= k; i++) cnt[i] = 0;
   for (int i = 0; i < n; i++) SAV[i] = dat[a[i]], cnt[SAV[i]]++;</pre>
   for (int i = 1; i \le k; i++) cnt[i] += cnt[i - 1];
   for (int i = n - 1; i \ge 0; i--) b[--cnt[SAV[i]]] = a[i];
#define I(x) ((x)\%3==1)?((x)/3):((x)/3+num1)
#define I2(x) (x<num1)?(3*x+1):(3*(x-num1)+2)
```

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```
static int cmp(int x, int v, int str[], int A[], int num1) {
   if (x \% 3 == 1) {
     if (y \% 3 == 2) return A[I(x)] < A[I(y)];
     else return str[x] < str[y] \mid | str[x] == str[y] && A[I(x + 1)] < A[I(y + 1)];
   }
   else {
     return str[x] < str[y] \mid \mid str[x] == str[y] && cmp(x + 1, y + 1, str, A, num1);
   }
 void make(int* str, int* sa, int n, int k) {
   if (n == 0) return:
   int num1 = (n + 2) / 3, num2 = n / 3;
   int num = num1 + num2:
   str[n] = str[n + 1] = str[n + 2] = 0:
   int *nsa = mloc(num), *nstr = mloc(num + 3);
   for (int i = 0, j = 0; i < n; i++) if (i % 3) A[j++] = i;
   if (n \% 3 == 1) A [num - 1] = n:
   rsort(A, B, str + 2, num, k): rsort(B, A, str + 1, num, k): rsort(A, B, str, num, k):
   int cnt = 1:
   nstr[I(B[0])] = 1:
   for (int i = 1; i < num; i++) {
     int c = B[i], p = B[i - 1];
     if (str[p] != str[c] || str[p + 1] != str[c + 1] || str[p + 2] != str[c + 2]) cnt++;
     nstr[I(c)] = cnt:
   if (cnt == num) for (int i = 0: i < num: i++) nsa[nstr[i] - 1] = i:
   else make(nstr, nsa, num, cnt);
   for (int i = 0, j = 0; i < num; i++) if (nsa[i] < num1) A[j++] = 3 * nsa[i];
   rsort(A. B. str. num1, k):
   for (int i = 0; i < num; i++) A[nsa[i]] = i, nsa[i] = I2(nsa[i]);</pre>
   A[num] = -1;
   merge(B, B + num1, nsa + (n \% 3 == 1), nsa + num, sa, [&](int x, int y) {
     return cmp(x, v, str. A, num1):
   });
   return:
 }
}sa;
4.5 eertree
int nxt[MAXN][26]:
int par [MAXN], len [MAXN], slink [MAXN], ptr [MAXN], diff [MAXN], series [MAXN], piv;
void clear(int n = MAXN){
 memset(par, 0, sizeof(int) * n):
 memset(len, 0, sizeof(int) * n);
 memset(slink, 0, sizeof(int) * n);
 memset(nxt, 0, sizeof(int) * 26 * n);
 piv = 0:
void init(int n, char *a){
 par[0] = 0:
 par[1] = 1;
 a[0] = -1;
 len[0] = -1:
```

```
piv = 1:
 int cur = 1;
 for(int i=1; i<=n; i++){
   while(a[i] != a[i - len[cur] - 1]) cur = slink[cur]:
   if(!nxt[cur][a[i]]){
     nxt[cur][a[i]] = ++piv:
     par[piv] = cur;
     len[piv] = len[cur] + 2;
     int lnk = slink[cur]:
     while(a[i] != a[i - len[lnk] - 1]){
       lnk = slink[lnk]:
     if(nxt[lnk][a[i]]) lnk = nxt[lnk][a[i]];
     if(len[piv] == 1 || lnk == 0) lnk = 1:
     slink[piv] = lnk;
     diff[piv] = len[piv] - len[lnk];
     if(diff[piv] == diff[lnk]) series[piv] = series[lnk];
     else series[piv] = piv;
   cur = nxt[cur][a[i]];
   ptr[i] = cur;
int query(int s, int e){
 int pos = ptr[e]:
 while(len[pos] \geq e - s + 1){
       if(len[pos] % diff[pos] == (e - s + 1) % diff[pos] &&
          len[series[pos]] <= e - s + 1) return true:</pre>
       pos = series[pos];
       pos = slink[pos]:
 return false:
vector<pi> minimum_partition(int n){ // (odd min, even min)
 vector<pi> dp(n + 1);
 vector<pi> series ans(n + 10):
 dp[0] = pi(1e9 + 1, 0);
 for(int i=1: i<=n: i++){
   dp[i] = pi(1e9 + 1, 1e9);
   for(int j=ptr[i]; len[j] > 0;){
     int slv = slink[series[i]]:
     series_ans[j] = dp[i - (len[slv] + diff[j])];
     if(diff[i] == diff[slink[i]]){
       series_ans[j].first = min(series_ans[j].first, series_ans[slink[j]].first);
       series_ans[j].second = min(series_ans[j].second, series_ans[slink[j]].second);
     auto val = series_ans[i];
     dp[i].first = min(dp[i].first, val.second + 1);
     dp[i].second = min(dp[i].second, val.first + 1);
     j = slv;
 }
 return dp;
```

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4.6 Circular LCS

```
string s1, s2;
int dp[4005][2005];
int nxt[4005][2005];
int n, m;
void reroot(int px){
 int py = 1;
  while(py <= m && nxt[px][py] != 2) py++;
 nxt[px][py] = 1;
  while(px < 2 * n && py < m){
   if(nxt[px+1][py] == 3){
     px++;
     nxt[px][py] = 1;
    else if(nxt[px+1][py+1] == 2){
     py++;
      nxt[px][py] = 1;
    else py++;
  while(px < 2 * n && nxt[px+1][py] == 3){
   px++;
   nxt[px][py] = 1;
 }
}
int track(int x, int y, int e){ // use this routine to find LCS as string
 int ret = 0:
  while(y != 0 \&\& x != e){
   if(nxt[x][y] == 1) y--;
   else if(nxt[x][y] == 2) ret += (s1[x] == s2[y]), x--, y--;
    else if(nxt[x][y] == 3) x--;
 }
  return ret;
int solve(string a, string b){
 n = a.size(), m = b.size();
 s1 = "#" + a + a;
  s1 = '#' + b;
  for(int i=0; i<=2*n; i++){
   for(int j=0; j<=m; j++){
     if(i == 0){
       nxt[i][j] = 3;
        continue:
     }
      if(i == 0){
        nxt[i][j] = 1;
        continue;
      dp[i][j] = -1;
      if(dp[i][j] < dp[i][j-1]){
       dp[i][j] = dp[i][j-1];
       nxt[i][j] = 1;
```

```
if(dp[i][j] < dp[i-1][j-1] + (s1[i] == s2[j])){
        dp[i][j] = dp[i-1][j-1] + (s1[i] == s2[j]);
        nxt[i][j] = 2;
      if(dp[i][j] < dp[i-1][j]){
        dp[i][j] = dp[i-1][j];
        nxt[i][j] = 3;
   }
 }
  int ret = dp[n][m];
 for(int i=1: i<n: i++){
   reroot(i), ret = max(ret, track(n+i, m, i));
 return ret;
5 Geometry
5.1 Buldozer
Point v[2020]:
struct Line{
 ll i, j, dx, dy;
 Line(int i, int j) : i(i), j(j) {
    dx = v[i].x - v[j].x; dy = v[i].y - v[j].y;
 bool operator < (const Line &t) const {</pre>
   ll a = dy * t.dx, b = t.dy * dx;
   return tie(a, i, j) < tie(b, t.i, t.j);
int ccw(Point a, Point b, Point c){}
int ccw(Line a, Line b){
 11 \text{ res} = a.dx*b.dy - b.dx*a.dy;
 if(!res) return 0; return res > 0 ? 1 : -1;
int idx[2020]; vector<Line> line;
void bulldozer(int n){
 sort(v+1, v+n+1); for(int i=1; i<=n; i++) idx[i] = i;
 for(int i=1; i<=n; i++) for(int j=1; j<i; j++) line.emplace_back(i, j);</pre>
 for(int i=0, j=0; i<line.size(); ){</pre>
   int ed = i;
    while(ed < line.size() && !ccw(line[i], line[ed])) ed++;</pre>
   for(int j=i; j<ed; j++){</pre>
     int a = line[j].i, b = line[j].j;
      swap(idx[a], idx[b]); swap(v[idx[a]], v[idx[b]]);
      update(idx[a]); update(idx[b]);
    ans = merge(ans, query()); i = ed;
 }
5.2 Rotating Calibers
Point v[2020];
struct Line{
 ll i, j, dx, dy;
 Line(int i, int j) : i(i), j(j) {
   dx = v[i].x - v[j].x; dy = v[i].y - v[j].y;
```

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```
bool operator < (const Line &t) const {</pre>
   ll a = dy * t.dx, b = t.dy * dx;
   return tie(a, i, j) < tie(b, t.i, t.j);
 }
}:
int ccw(Point a, Point b, Point c){}
int ccw(Line a, Line b){
 11 \text{ res} = a.dx*b.dy - b.dx*a.dy;
 if(!res) return 0; return res > 0 ? 1 : -1;
int idx[2020]: vector<Line> line:
void bulldozer(int n){
  sort(v+1, v+n+1): for(int i=1; i \le n; i++) idx[i] = i:
  for(int i=1; i<=n; i++) for(int j=1; j<i; j++) line.emplace_back(i, j);
  for(int i=0, j=0; i<line.size(); ){</pre>
   int ed = i:
    while(ed < line.size() && !ccw(line[i], line[ed])) ed++;</pre>
   for(int j=i; j<ed; j++){</pre>
     int a = line[j].i, b = line[j].j;
      swap(idx[a], idx[b]); swap(v[idx[a]], v[idx[b]]);
      update(idx[a]); update(idx[b]);
    ans = merge(ans, query()); i = ed;
5.3 Delaunav
using lll = __int128; // using T = ll; (if coords are < 2e4)
// return true if p strictly within circumcircle(a.b.c)
bool inCircle(P p, P a, P b, P c) {
 a = p, b = p, c = p; // assert(cross(a,b,c)>0);
 lll x = (lll)norm(a)*cross(b,c)+(lll)norm(b)*cross(c,a)
     +(lll)norm(c)*cross(a,b);
  return x*(cross(a,b,c)>0?1:-1) > 0;
P arb(LLONG_MAX,LLONG_MAX); // not equal to any other point
using Q = struct Quad*;
struct Quad {
 bool mark; Q o, rot; P p;
 P F() { return r()->p; }
 Q r() { return rot->rot; }
 Q prev() { return rot->o->rot; }
 Q next() { return r()->prev(); }
Q makeEdge(P orig, P dest) {
  Q q[]{new Quad{0.0.0.orig}, new Quad{0.0.0.arb}.
      new Quad{0,0,0,dest}, new Quad{0,0,0,arb}};
 FOR(i,4) q[i] \rightarrow 0 = q[-i \& 3], q[i] \rightarrow rot = q[(i+1) \& 3];
 return *q;
void splice(Q a, Q b) { swap(a->o->rot->o, b->o->rot->o); swap(a->o, b->o); }
Q connect(Q a, Q b) {
  Q q = makeEdge(a->F(), b->p);
  splice(q, a->next()); splice(q->r(), b);
  return q;
```

```
pair<0.0> rec(const vP& s) {
 if (sz(s) \le 3) {
    Q a = makeEdge(s[0], s[1]), b = makeEdge(s[1], s.bk);
    if (sz(s) == 2) return { a, a->r() };
    splice(a->r(), b);
    auto side = cross(s[0], s[1], s[2]):
    Q c = side ? connect(b, a) : 0;
    return {side < 0 ? c > r() : a, side < 0 ? c : b > r() };
#define H(e) e \rightarrow F(), e \rightarrow p
#define valid(e) (cross(e->F().H(base)) > 0)
  Q A, B, ra, rb;
  int half = sz(s) / 2:
  tie(ra, A) = rec({all(s)-half});
  tie(B, rb) = rec({sz(s)-half+all(s)});
  while ((cross(B->p,H(A)) < 0 && (A = A->next())) | |
       (cross(A->p,H(B)) > 0 && (B = B->r()->o)));
  Q base = connect(B->r(), A):
  if (A->p == ra->p) ra = base->r();
  if (B->p == rb->p) rb = base;
#define DEL(e, init, dir) Q e = init->dir; if (valid(e)) \
    while (inCircle(e->dir->F(), H(base), e->F())) { \
      Q t = e \rightarrow dir: \
      splice(e, e->prev()); \
      splice(e->r(), e->r()->prev()); \
      e = t: \
  while (1) {
    DEL(LC, base->r(), o); DEL(RC, base, prev());
    if (!valid(LC) && !valid(RC)) break;
    if (!valid(LC) || (valid(RC) && inCircle(H(RC), H(LC))))
      base = connect(RC, base->r());
    else base = connect(base->r(), LC->r());
 return {ra, rb};
V<AR<P,3>> triangulate(vP pts) {
  sor(pts); assert(unique(all(pts)) == end(pts)); // no duplicates
  if (sz(pts) < 2) return {};
  Q = rec(pts).f; V<Q>q = {e};
  while (cross(e->o->F(), e->F(), e->p) < 0) e = e->o;
#define ADD { Q c = e; do { c\rightarrow mark = 1; pts.pb(c\rightarrow p); \ \ }
 q.pb(c->r()); c = c->next(); } while (c != e); }
  ADD; pts.clear();
  int qi = 0; while (qi < sz(q)) if (!(e = q[qi++])->mark) ADD;
  V<AR<P,3>> ret(sz(pts)/3);
  FOR(i.sz(pts)) ret[i/3][i%3] = pts[i]:
  return ret;
5.4 Dual Graph
const int MV = 101010, ME = 101010; // MAX_V, MAX_E
p pt[MV]; // vertex's coord
vector g[MV]; // g[s].emplace_back(e, edge_id);
vector<int> dual_pt; // coord compress
```

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```
void uf_init(){ iota(par, par+ME*2, 0); }
int find(int v){return v == par[v] ? v : par[v] = find(par[v]);}
void merge(int u, int v){ u = find(u); v = find(v); if(u != v)par[u] = v; }
p base; // sort by angle
bool cmp_angle(const p &_a, const p &_b){
 p = pt[a.x], b = pt[b.x];
 if((a > base) != (b > base)) return a > b;
 return ccw(a, base, b) > 0;
void addEdge(int s, int e, int id){
 g[s].emplace_back(e, id); g[e].emplace_back(s, id);
int out: //outer face
void getDual(int n, int m){
 uf_init();
 for(int i=1: i<=n: i++){
   base = pt[i]; sort(all(g[i]), cmp_angle);
   // up, left : *2+1 / down, right : *2
   for(int j=0; j<g[i].size(); j++){</pre>
     int k = j ? j - 1 : g[i].size()-1;
     int u = g[i][k].y << 1 | 1, v = g[i][j].y << 1;
     p p1 = pt[g[i][k].x], p2 = pt[g[i][j].x];
     if(p1 > base) u = 1;
     if(p2 > base) v = 1;
     merge(u, v);
   }
  int mn_idx = min_element(pt+1, pt+n+1) - pt;
 out = find(g[mn_idx][0].y << 1 | 1);
 for(int i=1; i<=m; i++){</pre>
   dual_pt.push_back(find(i << 1));</pre>
   dual_pt.push_back(find(i << 1 | 1));</pre>
 }
 compress(dual_pt);
 // @TODO coord compress
5.5 Smallest Enclosing Circle / Sphere
namespace cover 2d{
 double eps = 1e-9;
 using Point = complex<double>;
 struct Circle{ Point p; double r; };
 double dist(Point p, Point q){ return abs(p-q); }
  double area2(Point p, Point q){ return (conj(p)*q).imag(); }
  bool in(const Circle& c, Point p){ return dist(c.p, p) < c.r + eps; }</pre>
 Circle INVAL = Circle{Point(0, 0), -1}:
 Circle mCC(Point a, Point b, Point c){
   b -= a: c -= a:
   double d = 2*(conj(b)*c).imag(); if(abs(d)<eps) return INVAL;</pre>
   Point ans = (c*norm(b) - b*norm(c)) * Point(0, -1) / d;
   return Circle{a + ans. abs(ans)}:
 Circle solve(vector<Point> p) {
   mt19937 gen(0x94949); shuffle(p.begin(), p.end(), gen);
   Circle c = INVAL;
   for(int i=0; i<p.size(); ++i) if(c.r<0 ||!in(c, p[i])){
```

int par[ME * 2]: // Union Find

```
c = Circle{p[i], 0}:
     for(int j=0; j<=i; ++j) if(!in(c, p[j])){
        Circle ans{(p[i]+p[j])*0.5, dist(p[i], p[j])*0.5};
        if(c.r == 0) { c = ans; continue; }
        Circle 1, r; 1 = r = INVAL;
        Point pq = p[j]-p[i];
        for(int k=0; k<=j; ++k) if(!in(ans, p[k])) {</pre>
          double a2 = area2(pq, p[k]-p[i]);
          Circle c = mCC(p[i], p[j], p[k]);
          if(c.r<0) continue;
          else if(a2 > 0 && (1.r<0||area2(pq, c.p-p[i])) > area2(pq, 1.p-p[i]))) 1 = c;
          else if(a2 < 0 && (r.r<0)|area2(pq, c.p-p[i]) < area2(pq, r.p-p[i]))) r = c;
        if(1.r<0\&\&r.r<0) c = ans:
        else if(1.r<0) c = r;
        else if(r.r<0) c = 1;
        else c = 1.r<=r.r?1:r:
   }
   return c;
 }
};
namespace cover 3d{
 double enclosing sphere(vector<double> x. vector<double> y. vector<double> z){
   int n = x.size();
   auto hyp = [](double x, double y, double z){
     return x * x + y * y + z * z;
   };
   double px = 0, py = 0, pz = 0;
   for(int i=0; i<n; i++){</pre>
     px += x[i]:
     py += y[i];
     pz += z[i];
   px *= 1.0 / n:
   pv *= 1.0 / n;
   pz *= 1.0 / n:
   double rat = 0.1, maxv;
   for(int i=0; i<10000; i++){
     maxv = -1:
     int maxp = -1;
     for(int j=0; j<n; j++){
       double tmp = hyp(x[j] - px, y[j] - py, z[j] - pz);
       if(maxv < tmp){</pre>
         maxv = tmp;
          maxp = j;
     px += (x[maxp] - px) * rat;
     py += (y[maxp] - py) * rat;
     pz += (z[maxp] - pz) * rat;
     rat *= 0.998;
    return sqrt(maxv);
```

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```
}:
5.6 3D Convex Hull
struct vec3{
 11 x, y, z;
 vec3(): x(0), y(0), z(0) {}
  vec3(11 a, 11 b, 11 c): x(a), y(b), z(c) {}
 vec3 operator*(const vec3& v) const{ return vec3(y*v.z-z*v.y, z*v.x-x*v.z, x*v.y-v*v.x); }
 vec3 operator-(const vec3& v) const{ return vec3(x-v.x, y-v.y, z-v.z); }
 vec3 operator-() const{ return vec3(-x, -y, -z); }
 11 dot(const vec3 &v) const{ return x*v.x+y*v.y+z*v.z; }
struct twoset {
 int a. b:
  void insert(int x) { (a == -1 ? a : b) = x; }
 bool contains(int x) { return a == x || b == x: }
 void erase(int x) { (a == x ? a : b) = -1; }
 int size() { return (a != -1) + (b != -1); }
} E[MAXN][MAXN]: // i < i</pre>
struct face{
  vec3 norm:
 ll disc;
 int I[3]:
};
face make_face(int i, int j, int k, int ii, vector<vec3> &A){ // p^T * norm < disc</pre>
  E[i][j].insert(k); E[i][k].insert(j); E[j][k].insert(i);
 face f; f.I[0]=i, f.I[1]=j, f.I[2]=k;
 f.norm = (A[j]-A[i])*(A[k]-A[i]);
 f.disc = f.norm.dot(A[i]);
 if(f.norm.dot(A[ii])>f.disc){
   f.norm = -f.norm;
   f.disc = -f.disc;
 return f;
vector<face> get_hull(vector<vec3> &A){
  int N = A.size():
  vector<face> faces; memset(E, -1, sizeof(E));
  faces.push_back(make_face(0,1,2,3,A));
  faces.push_back(make_face(0,1,3,2,A));
  faces.push_back(make_face(0,2,3,1,A));
  faces.push_back(make_face(1,2,3,0,A));
  for(int i=4: i<N: ++i){</pre>
   for(int j=0; j<faces.size(); ++j){</pre>
     face f = faces[j];
      if(f.norm.dot(A[i])>f.disc){
        E[f.I[0]][f.I[1]].erase(f.I[2]);
        E[f.I[0]][f.I[2]].erase(f.I[1]):
        E[f.I[1]][f.I[2]].erase(f.I[0]);
        faces[j--] = faces.back();
        faces.pop_back();
   }
```

```
int nf = faces.size():
    for(int j=0; j<nf; ++j){</pre>
     face f=faces[j];
      for(int a=0; a<3; ++a) for(int b=a+1; b<3; ++b){
        int c=3-a-b:
        if(E[f.I[a]][f.I[b]].size()==2) continue:
        faces.push_back(make_face(f.I[a], f.I[b], i, f.I[c], A));
   }
 }
 return faces:
5.7 Dynamic Convex Hull Trick
using line_t = double;
const line_t is_query = -1e18;
struct Line {
 line t m. b:
 mutable function<const Line*()> succ:
 bool operator<(const Line& rhs) const {</pre>
   if (rhs.b != is_query) return m < rhs.m;</pre>
   const Line* s = succ();
   if (!s) return 0;
   line t x = rhs.m:
   return b - s \rightarrow b < (s \rightarrow m - m) * x:
 }
}:
struct HullDynamic : public multiset<Line> { // will maintain upper hull for maximum
 bool bad(iterator y) {
    auto z = next(y);
   if (y == begin()) {
     if (z == end()) return 0;
      return y->m == z->m && y->b <= z->b;
   auto x = prev(y);
    if (z == end()) return y->m == x->m && y->b <= x->b;
   return (x-b-y-b)*(z-m-y-m) >= (y-b-z-b)*(y-m-x-m);
 void insert line(line t m, line t b) {
   auto y = insert({ m, b });
   y->succ = [=] { return next(y) == end() ? 0 : &*next(y); };
   if (bad(v)) { erase(v): return: }
   while (next(y) != end() && bad(next(y))) erase(next(y));
   while (y != begin() && bad(prev(y))) erase(prev(y));
 line_t query(line_t x) {
   auto 1 = *lower bound((Line) { x, is query }):
   return 1.m * x + 1.b;
}H:
5.8 Half-plane Intersection
const double eps = 1e-8;
typedef pair < long double, long double > pi;
bool z(long double x){ return fabs(x) < eps; }</pre>
```

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```
struct line{
 long double a, b, c;
 bool operator<(const line &1)const{</pre>
   bool flag1 = pi(a, b) > pi(0, 0);
   bool flag2 = pi(1.a, 1.b) > pi(0, 0);
   if(flag1 != flag2) return flag1 > flag2;
   long double t = ccw(pi(0, 0), pi(a, b), pi(l.a, l.b));
   return z(t) ? c * hypot(1.a, 1.b) < 1.c * hypot(a, b) : <math>t > 0;
 pi slope(){ return pi(a, b); }
pi cross(line a, line b){
 long double det = a.a * b.b - b.a * a.b;
 return pi((a.c * b.b - a.b * b.c) / det. (a.a * b.c - a.c * b.a) / det):
bool bad(line a, line b, line c){
 if(ccw(pi(0, 0), a.slope(), b.slope()) <= 0) return false;</pre>
 pi crs = cross(a, b);
 return crs.first * c.a + crs.second * c.b >= c.c:
bool solve(vector<line> v, vector<pi> &solution){ // ax + by <= c;</pre>
 sort(v.begin(), v.end());
 deque<line> dq;
 for(auto &i : v){
   if(!dq.empty() && z(ccw(pi(0, 0), dq.back().slope(), i.slope()))) continue;
   while(dq.size() >= 2 && bad(dq[dq.size()-2], dq.back(), i)) dq.pop_back();
   while(dq.size() >= 2 \&\& bad(i, dq[0], dq[1])) dq.pop_front();
   dq.push_back(i);
 }
 while(dq.size() > 2 \&\& bad(dq[dq.size()-2], dq.back(), dq[0])) dq.pop_back();
  while(dq.size() > 2 && bad(dq.back(), dq[0], dq[1])) dq.pop_front();
  vector<pi> tmp;
 for(int i=0; i<dq.size(); i++){</pre>
   line cur = dq[i], nxt = dq[(i+1)%dq.size()];
   if(ccw(pi(0, 0), cur.slope(), nxt.slope()) <= eps) return false;</pre>
   tmp.push back(cross(cur. nxt)):
 }
 solution = tmp:
 return true;
5.9 Point-in-polygon test / Point-to-polygon tangent
// C : counter_clockwise(C[0] == C[N]), N >= 3
// return highest point in C <- P(clockwise) or -1 if strictly in P
// polygon is strongly convex, C[i] != P
int convex_tangent(vector<pi> &C, pi P, int up = 1){
 auto sign = [\&](lint c){ return c > 0 ? up : c == 0 ? 0 : -up: }:
 auto local = [&](pi P, pi a, pi b, pi c) {
   return sign(ccw(P, a, b)) \le 0 && sign(ccw(P, b, c)) >= 0;
 int N = C.size()-1, s = 0, e = N, m;
 if( local(P, C[1], C[0], C[N-1]) ) return 0;
 while(s+1 < e){
   m = (s+e) / 2:
   if( local(P, C[m-1], C[m], C[m+1]) ) return m;
   if ( sign(ccw(P, C[s], C[s+1])) < 0) { // up
     if (sign(ccw(P, C[m], C[m+1])) > 0) e = m;
```

```
else if (sign(ccw(P, C[m], C[s])) > 0) s = m:
     else e = m;
   }
   else{ // down
     if (sign(ccw(P, C[m], C[m+1])) < 0) s = m;
     else if (sign(ccw(P, C[m], C[s])) < 0) s = m;
     else e = m;
   }
 if( s && local(P, C[s-1], C[s], C[s+1]) ) return s;
 if( e != N && local(P, C[e-1], C[e], C[e+1]) ) return e;
 return -1:
5.10 kd-tree
typedef pair<int, int> pi;
struct node{
 pi pnt;
 int spl, sx, ex, sy, ey;
}tree[270000];
pi a[100005]:
int n. ok[270000]:
lint sqr(int x){ return 111 * x * x: }
bool cmp1(pi a, pi b){ return a < b; }</pre>
bool cmp2(pi a, pi b){ return pi(a.second, a.first) < pi(b.second, b.first); }
// init(0, n-1, 1) : Initialize kd-tree
// set dap = INF, and call solve(1, P). dap = (closest point from P)
void init(int s, int e, int p){ // Initialize kd-tree
 int minx = 1e9, maxx = -1e9, miny = 1e9, maxy = -1e9;
 int m = (s+e)/2:
 for(int i=s; i<=e; i++){</pre>
   minx = min(minx, a[i].first);
   miny = min(miny, a[i].second):
   maxx = max(maxx, a[i].first);
   maxy = max(maxy, a[i].second);
 tree[p].spl = (maxx - minx < maxy - miny);</pre>
 sort(a+s, a+e+1, [&](const pi &a, const pi &b){
   return tree[p].spl ? cmp2(a, b) : cmp1(a, b);
 }):
 tree[p] = {a[m], tree[p].spl, minx, maxx, miny, maxy};
 if (s \le m-1) init(s, m-1, 2*p);
 if(m+1 \le e) init(m+1, e, 2*p+1):
lint dap = 3e18;
void solve(int p, pi x){ // find closest point from point x (L^2)
 if(x != tree[p].pnt) dap = min(dap, sqr(x.first - tree[p].pnt.first) + sqr(x.second -
tree[p].pnt.second));
 if(tree[p].spl){
   if(!cmp2(tree[p].pnt, x)){
     if(ok[2*p]) solve(2*p, x);
```

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```
if(ok[2*p+1] \&\& sqr(tree[2*p+1].sv - x.second) < dap) solve(2*p+1, x):
   }
   else{
     if(ok[2*p+1]) solve(2*p+1, x);
     if(ok[2*p] && sqr(tree[2*p].ey - x.second) < dap) solve(2*p, x);
 }
 else{
   if(!cmp1(tree[p].pnt, x)){
     if(ok[2*p]) solve(2*p, x);
     if (ok[2*p+1] \&\& sqr(tree[2*p+1].sx - x.first) < dap) solve(2*p+1, x);
   }
   else{
     if(ok[2*p+1]) solve(2*p+1, x):
     if(ok[2*p] && sqr(tree[2*p].ex - x.first) < dap) solve(2*p, x);</pre>
 }
6
   Math
6.1 FFT / NTT
typedef complex<double> base;
void fft(vector<base> &a, bool inv){
 int n = a.size(), i = 0:
 vector<base> roots(n/2);
 for(int i=1; i<n; i++){</pre>
   int bit = (n >> 1):
   while(j >= bit){
     i -= bit:
     bit >>= 1;
   }
   j += bit;
   if(i < j) swap(a[i], a[j]);</pre>
 double ang = 2 * acos(-1) / n * (inv ? -1 : 1):
 for(int i=0; i<n/2; i++){</pre>
   roots[i] = base(cos(ang * i), sin(ang * i));
 /* In NTT, let prr = primitive root. Then,
 int ang = ipow(prr, (mod - 1) / n);
 if(inv) ang = ipow(ang, mod - 2);
 for(int i=0: i < n/2: i++){
   roots[i] = (i ? (111 * roots[i-1] * ang % mod) : 1);
 XOR Convolution : set roots[*] = 1.
 OR Convolution : set roots[*] = 1, and do following:
   if (!inv) {
       a[j + k] = u + v;
       a[i + k + i/2] = u;
   } else {
       a[i + k] = v:
       a[j + k + i/2] = u - v;
   }
 for(int i=2; i<=n; i<<=1){
   int step = n / i;
```

```
for(int j=0; j<n; j+=i){</pre>
     for(int k=0; k<i/2; k++){
       base u = a[j+k], v = a[j+k+i/2] * roots[step * k];
       a[i+k] = u+v:
       a[j+k+i/2] = u-v:
   }
 }
 if(inv) for(int i=0; i<n; i++) a[i] /= n; // skip for OR convolution.
vector<lint> multiply(vector<lint> &v, vector<lint> &w){
 vector<base> fv(v.begin(), v.end()), fw(w.begin(), w.end());
 int n = 2: while(n < v.size() + w.size()) n <<= 1:
 fv.resize(n); fw.resize(n);
 fft(fv, 0); fft(fw, 0);
 for(int i=0; i<n; i++) fv[i] *= fw[i];
 fft(fv, 1);
 vector<lint> ret(n):
 for(int i=0; i<n; i++) ret[i] = (lint)round(fv[i].real());</pre>
 return ret;
vector<lint> multiply(vector<lint> &v, vector<lint> &w, lint mod){
 int n = 2; while(n < v.size() + w.size()) n <<= 1;</pre>
 vector < base > v1(n), v2(n), r1(n), r2(n):
 for(int i=0; i<v.size(); i++){</pre>
   v1[i] = base(v[i] >> 15, v[i] & 32767);
 for(int i=0; i<w.size(); i++){</pre>
   v2[i] = base(w[i] >> 15, w[i] & 32767);
 fft(v1, 0);
 fft(v2, 0);
 for(int i=0; i<n; i++){</pre>
   int j = (i ? (n - i) : i);
   base ans1 = (v1[i] + conj(v1[j])) * base(0.5, 0);
   base ans2 = (v1[i] - conj(v1[j])) * base(0, -0.5);
   base ans3 = (v2[i] + conj(v2[j])) * base(0.5, 0);
   base ans4 = (v2[i] - conj(v2[j])) * base(0, -0.5);
   r1[i] = (ans1 * ans3) + (ans1 * ans4) * base(0, 1);
   r2[i] = (ans2 * ans3) + (ans2 * ans4) * base(0, 1):
 }
 fft(r1, 1);
 fft(r2, 1):
 vector<lint> ret(n):
 for(int i=0; i<n; i++){
   lint av = (lint)round(r1[i].real());
   lint bv = (lint)round(r1[i].imag()) + (lint)round(r2[i].real());
   lint cv = (lint)round(r2[i].imag());
   av %= mod, bv %= mod, cv %= mod;
   ret[i] = (av << 30) + (bv << 15) + cv;
   ret[i] %= mod;
   ret[i] += mod;
   ret[i] %= mod:
 }
 return ret:
```

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```
6.2 Hell-Joseon style FFT
#include <smmintrin.h>
#include <immintrin.h>
#pragma GCC target("avx2")
#pragma GCC target("fma")
__m256d mult(__m256d a, __m256d b){
  _{m256d} c = _{mm256_{movedup_{pd}(a)}}
  _{m256d} d = _{mm256\_shuffle\_pd(a, a, 15)};
 _{m256d cb} = _{mm256_{mul_pd(c, b)}}
  _{m256d} db = _{mm256} dd, b);
 _{m256d} = _{mm256\_shuffle\_pd(db, db, 5)};
 _{\rm m256d} r = _{\rm mm256\_addsub\_pd(cb, e)};
 return r:
void fft(int n. m128d a[]. bool invert){
 for(int i=1, j=0; i<n; ++i){
   int bit = n >> 1;
   for(; j>=bit; bit>>=1) j -= bit;
   j += bit:
   if(i<j) swap(a[i], a[j]);</pre>
  for(int len=2; len<=n; len<<=1){</pre>
   double ang = 2*3.14159265358979/len*(invert?-1:1);
    _{\text{m256d wlen}}; wlen[0] = cos(ang), wlen[1] = sin(ang);
   for(int i=0; i<n; i += len){</pre>
      _{m256d} w; w[0] = 1; w[1] = 0;
     for(int j=0; j<len/2; ++j){</pre>
        w = _mm256_permute2f128_pd(w, w, 0);
        wlen = _{mm256_{insertf128_{pd}(wlen, a[i+j+len/2], 1)};
        w = mult(w, wlen);
        _{m128d} vw = _{mm256}extractf128_{pd}(w, 1);
        _{\rm m}128d\ u\ =\ a[i+j];
        a[i+j] = _mm_add_pd(u, vw);
        a[i+j+len/2] = _mm_sub_pd(u, vw);
   }
 }
 if(invert){
    _{m128d inv; inv[0] = inv[1] = 1.0/n;
   for(int i=0; i<n; ++i) a[i] = _mm_mul_pd(a[i], inv);</pre>
 }
vector<int64_t> multiply(vector<int64_t>& v, vector<int64_t>& w){
 int n = 2; while(n < v.size()+w.size()) n<<=1;</pre>
  m128d* fv = new m128d[n]:
 for(int i=0; i<n; ++i) fv[i][0] = fv[i][1] = 0;
 for(int i=0; i<v.size(); ++i) fv[i][0] = v[i];</pre>
 for(int i=0; i<w.size(); ++i) fv[i][1] = w[i];
 fft(n, fv, 0); // (a+bi) is stored in FFT
 for(int i=0: i<n: i += 2){
    __m256d a;
   a = _mm256_insertf128_pd(a, fv[i], 0);
   a = _mm256_insertf128_pd(a, fv[i+1], 1);
   a = mult(a, a);
   fv[i] = _mm256_extractf128_pd(a, 0);
```

```
fv[i+1] = mm256 \text{ extractf128 pd(a. 1)}:
 }
 fft(n, fv, 1);
 vector<int64_t> ret(n);
 for(int i=0; i<n; ++i) ret[i] = (int64_t)round(fv[i][1]/2);
 delete∏ fv:
 return ret;
6.3 NTT Polynomial Division
vector<lint> get_inv(int n, const vector<lint> &p){
 vector<lint> q = \{ipow(p[0], mod - 2)\};
 for(int i=2; i<=n; i<<=1){
   vector<lint> res:
   vector<lint> fq(q.begin(), q.end()); fq.resize(2*i);
   vector<lint> fp(p.begin(), p.begin() + i); fp.resize(2*i);
   fft(fq, 0); fft(fp, 0);
   for(int j=0; j<2*i; j++){
     fp[j] *= fq[j] * fq[j] % mod;
     fp[j] %= mod;
   }
   fft(fp, 1);
   res.resize(i);
   for(int j=0; j<i; j++){
     res[j] = mod - fp[j];
     if(j < i/2) res[j] += 2 * q[j];
     res[j] %= mod;
   q = res;
 return q;
vector<lint> poly_divide(const vector<lint> &a, const vector<lint> &b){
 assert(b.back() != 0); // please trim leading zero
 int n = a.size(), m = b.size();
 int k = 2; while (k < n-m+1) k <<= 1;
 vector<lint> rb(k), ra(k);
 for(int i=0; i<m && i<k; ++i) rb[i] = b[m-i-1];</pre>
 for(int i=0; i<n && i<k; ++i) ra[i] = a[n-i-1];
 vector<lint> rbi = get_inv(k, rb);
 vector<lint> res = multiply(rbi, ra);
 res.resize(n - m + 1);
 reverse(res.begin(), res.end());
 return res:
6.4 Black Box Linear Algebra + Kitamasa
vector<int> berlekamp_massey(vector<int> x){
 vector<int> ls, cur;
 int lf, ld;
 for(int i=0; i<x.size(); i++){</pre>
   lint t = 0:
   for(int j=0; j<cur.size(); j++){</pre>
     t = (t + 111 * x[i-j-1] * cur[j]) \% mod;
   if((t - x[i]) \% mod == 0) continue;
   if(cur.empty()){
```

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```
cur.resize(i+1):
     lf = i;
     ld = (t - x[i]) \% mod;
      continue:
   lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod;
   vector<int> c(i-lf-1);
    c.push_back(k);
    for(auto &j : ls) c.push_back(-j * k % mod);
    if(c.size() < cur.size()) c.resize(cur.size());</pre>
    for(int j=0; j<cur.size(); j++){</pre>
      c[j] = (c[j] + cur[j]) \% mod;
   if(i-lf+(int)ls.size()>=(int)cur.size()){
      tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) \% mod);
   }
    cur = c;
  for(auto &i : cur) i = (i % mod + mod) % mod:
  return cur:
int get_nth(vector<int> rec, vector<int> dp, lint n){
 int m = rec.size();
  vector<int> s(m), t(m);
 s[0] = 1:
 if(m != 1) t[1] = 1;
  else t[0] = rec[0];
  auto mul = [&rec](vector<int> v, vector<int> w){
   int m = v.size();
   vector<int> t(2 * m);
   for(int j=0; j<m; j++){
     for(int k=0; k<m; k++){</pre>
       t[j+k] += 111 * v[j] * w[k] % mod;
        if(t[j+k] >= mod) t[j+k] -= mod;
   }
    for(int j=2*m-1; j>=m; j--){
     for(int k=1: k<=m: k++){
       t[j-k] += 111 * t[j] * rec[k-1] % mod;
        if(t[j-k] >= mod) t[j-k] -= mod;
   }
   t.resize(m):
   return t:
 }:
  while(n){
   if(n \& 1) s = mul(s, t);
   t = mul(t, t);
   n >>= 1:
 lint ret = 0;
 for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;
  return ret % mod;
int guess_nth_term(vector<int> x, lint n){
 if(n < x.size()) return x[n]:</pre>
```

```
vector<int> v = berlekamp_massey(x);
  if(v.empty()) return 0;
 return get_nth(v, x, n);
struct elem{int x, y, v;}; // A_(x, y) <- v, O-based. no duplicate please..
vector<int> get_min_poly(int n, vector<elem> M){
 // smallest poly P such that A^i = sum_{j < i} {A^j \times P_j}
 vector<int> rnd1, rnd2;
 mt19937 rng(0x14004);
 auto randint = [&rng](int lb, int ub){
   return uniform_int_distribution<int>(lb, ub)(rng);
 }:
 for(int i=0; i<n; i++){</pre>
   rnd1.push_back(randint(1, mod - 1));
    rnd2.push_back(randint(1, mod - 1));
  vector<int> gobs;
 for(int i=0; i<2*n+2; i++){
   int tmp = 0:
   for(int j=0; j<n; j++){
     tmp += 111 * rnd2[j] * rnd1[j] % mod;
      if(tmp >= mod) tmp -= mod;
    gobs.push_back(tmp);
   vector<int> nxt(n):
   for(auto &i : M){
     nxt[i.x] += 111 * i.v * rnd1[i.y] % mod;
      if(nxt[i.x] >= mod) nxt[i.x] -= mod:
   }
   rnd1 = nxt;
 auto sol = berlekamp_massey(gobs);
 reverse(sol.begin(), sol.end());
 return sol;
lint det(int n. vector<elem> M){
 vector<int> rnd;
 mt19937 rng(0x14004):
 auto randint = [&rng](int lb, int ub){
   return uniform_int_distribution<int>(lb, ub)(rng);
 for(int i=0; i<n; i++) rnd.push_back(randint(1, mod - 1));</pre>
 for(auto &i : M){
   i.v = 111 * i.v * rnd[i.y] % mod;
 auto sol = get_min_poly(n, M)[0];
 if (n \% 2 == 0) sol = mod - sol;
 for(auto &i : rnd) sol = 111 * sol * ipow(i, mod - 2) % mod;
6.5 Gaussian Elimination
int n, inv;
vector<int> basis[505];
lint gyesu = 1;
void insert(vector<int> v){
```

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```
for(int i=0: i<n: i++){</pre>
    if(basis[i].size()) inv ^= 1; // inversion num increases
    if(v[i] && basis[i].empty()){
      basis[i] = v;
      return;
   }
    if(v[i]){
     lint minv = ipow(basis[i][i], mod - 2) * v[i] % mod;
      for(auto &j : basis[i]) j = (j * minv) % mod;
      gvesu *= minv;
      gyesu %= mod;
      for(int j=0; j<basis[i].size(); j++){</pre>
        v[j] += mod - basis[i][j];
        while(v[j] >= mod) v[j] -= mod;
   }
  puts("0");
  exit(0):
// Sample: Calculates Determinant in Z_p Field
int main(){
  scanf("%d",&n);
  for(int i=0: i<n: i++){</pre>
   vector<int> v(n);
   for(int j=0; j<n; j++) scanf("%d",&v[j]);
   if(i \% 2 == 1) inv ^= 1:
   insert(v);
  if(inv) gyesu = mod - gyesu;
  gyesu = ipow(gyesu, mod - 2);
  for(int i=0; i<n; i++) gyesu = gyesu * basis[i][i] % mod;</pre>
  cout << gyesu % mod << endl;</pre>
6.6 Simplex Algorithm
using T = long double;
const int N = 410, M = 30010;
const T eps = 1e-7;
int n, m;
int Left[M], Down[N];
// time complexity: exponential. fast $0(MN^2)$ in experiment. dependent on the modeling.
// Ax <= b, max c^T x. 최댓값: v, 답 추적: sol[i]. 1 based
T a[M][N], b[M], c[N], v, sol[N];
bool eq(T a, T b) { return fabs(a - b) < eps; }</pre>
bool ls(T a, T b)  { return a < b && !eg(a, b); }
void init(int p, int q) {
 n = p; m = q; v = 0;
 for(int i = 1; i <= m; i++){
   for(int j = 1; j <= n; j++) a[i][j]=0;
  for(int i = 1; i \le m; i++) b[i]=0;
  for(int i = 1; i <= n; i++) c[i]=sol[i]=0;</pre>
void pivot(int x,int y) {
  swap(Left[x], Down[y]);
```

```
T k = a[x][v]: a[x][v] = 1:
 vector<int> nz;
 for(int i = 1; i <= n; i++){
   a[x][i] /= k:
   if(!eq(a[x][i], 0)) nz.push_back(i);
 b[x] /= k;
 for(int i = 1: i <= m: i++){
   if(i == x \mid\mid eq(a[i][y], 0)) continue;
   k = a[i][y]; a[i][y] = 0;
   b[i] -= k*b[x]:
   for(int j : nz) a[i][j] -= k*a[x][j];
 if(eq(c[y], 0)) return;
 k = c[v]; c[v] = 0;
 v += k*b[x]:
 for(int i : nz) c[i] -= k*a[x][i];
// 0: found solution, 1: no feasible solution, 2: unbounded
int solve() {
 for(int i = 1; i <= n; i++) Down[i] = i;
 for(int i = 1; i <= m; i++) Left[i] = n+i;
 while(1) { // Eliminating negative b[i]
   int x = 0, y = 0:
   for(int i = 1; i \le m; i++) if (ls(b[i], 0) && (x == 0 || b[i] < b[x])) x = i;
   for(int i = 1; i \le n; i + +) if (ls(a[x][i], 0) && (y == 0 || a[x][i] < a[x][y])) y = i;
   if(y == 0) return 1;
   pivot(x, y);
 }
 while(1) {
   int x = 0, y = 0;
   for(int i = 1; i <= n; i++)
     if (ls(0, c[i]) && (!y || c[i] > c[y])) y = i;
   if(v == 0) break:
   for(int i = 1; i <= m; i++)
     if (ls(0, a[i][y]) && (!x || b[i]/a[i][y] < b[x]/a[x][y])) x = i;
   if(x == 0) return 2;
   pivot(x, y);
 for(int i = 1; i <= m; i++) if(Left[i] <= n) sol[Left[i]] = b[i];</pre>
 return 0:
}
6.7 De Bruijn Sequence
// Create cyclic string of length k^n that contains every length n string as substring.
alphabet = [0, k - 1]
int res[10000000]; // >= k^n
int aux[10000000]; // >= k*n
int de_bruijn(int k, int n) { // Returns size (k^n)
 if(k == 1) {
   res[0] = 0;
   return 1:
 for(int i = 0; i < k * n; i++)
   aux[i] = 0:
```

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```
int sz = 0:
  function<void(int, int)> db = [&](int t, int p) {
   if(t > n) {
     if(n \% p == 0)
       for(int i = 1; i <= p; i++)
          res[sz++] = aux[i]:
   }
    else {
      aux[t] = aux[t - p];
      db(t + 1, p);
     for(int i = aux[t - p] + 1; i < k; i++) {
        aux[t] = i:
        db(t + 1, t);
     }
   }
 };
 db(1, 1);
 return sz;
6.8 Discrete Kth root
* Solve x for x^P = A mod Q
* (P, Q-1) = 1 \rightarrow P^-1 \mod (Q-1) exists
* x has solution iff A^{(0-1)} / P = 1 \mod Q
* PP | (Q-1) -> P < sqrt(Q), solve lgQ rounds of discrete log
* else -> find a s.t. s | (Pa - 1) -> ans = A^a */
using LL = long long;
LL mul(LL x, LL y, LL mod) { return (__int128) x * y % mod; }
LL add(LL x, LL y, LL mod) { return (x + y) % mod; }
LL pw(LL x, LL y, LL mod){
 LL ret = 1, piv = x;
 while(y){
   if(y & 1) ret = mul(ret, piv, mod);
   piv = mul(piv, piv, mod);
   y >>= 1;
 }
 return ret % mod;
void gcd(LL a, LL b, LL &x, LL &y, LL &g){
 if (b == 0) {
   x = 1, y = 0, g = a;
   return;
 }
 LL tx, ty;
 gcd(b, a%b, tx, ty, g);
 x = tv: v = tx - tv * (a / b):
LL P, A, Q, g; // x^P = A \mod Q
const int X = 1e5;
LL base, ae[X], aXe[X], iaXe[X];
unordered_map<LL, LL> ht;
#define FOR(i, c) for (int i = 0; i < (c); ++i)
#define REP(i, 1, r) for (int i = (1); i <= (r); ++i)
void build(LL a) { // ord(a) = P < sqrt(Q)</pre>
 base = a;
 ht.clear();
```

```
ae[0] = 1; ae[1] = a; aXe[0] = 1; aXe[1] = pw(a, X, Q);
  iaXe[0] = 1; iaXe[1] = pw(aXe[1], Q-2, Q);
 REP(i, 2, X-1) {
   ae[i] = mul(ae[i-1], ae[1], Q);
   aXe[i] = mul(aXe[i-1], aXe[1], Q);
   iaXe[i] = mul(iaXe[i-1], iaXe[1], Q);
 FOR(i, X) ht[ae[i]] = i;
LL dis_log(LL x) {
 FOR(i, X) {
   LL iaXi = iaXe[i];
   LL rst = mul(x, iaXi, 0):
   if (ht.count(rst)) return i*X + ht[rst];
 }
}
LL main2() {
   cin >> P >> A >> Q:
 LL t = 0, s = Q-1;
 while (s % P == 0) {
   ++t:
   s /= P;
 }
 if (A == 0) return 0:
 if (t == 0) {
   // a^{P^-1 mod phi(Q)}
   LL x, y, _;
   gcd(P, Q-1, x, y, _);
   if (x < 0) {
     x = (x \% (Q-1) + Q-1) \% (Q-1);
   LL ans = pw(A, x, Q);
   if (pw(ans, P, Q) != A) while(1);
   return ans:
 // A is not P-residue
 if (pw(A, (Q-1) / P, Q) != 1) return -1;
 for (g = 2; g < Q; ++g) {
   if (pw(g, (Q-1) / P, Q) != 1)
     break:
 LL alpha = 0;
 {
   LL y, _;
   gcd(P, s, alpha, y, _);
   if (alpha < 0) alpha = (alpha \% (Q-1) + Q-1) \% (Q-1);
 if (t == 1) {
   LL ans = pw(A, alpha, Q);
   return ans;
 }
 LL a = pw(g, (Q-1) / P, Q);
 build(a):
 LL b = pw(A, add(mul(P%(Q-1), alpha, Q-1), Q-2, Q-1), Q);
 LL c = pw(g, s, Q);
```

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3

997

168

12

99999999989

```
LL h = 1:
  LL e = (Q-1) / s / P; // r^{t-1}
  REP(i, 1, t-1) {
   e /= P:
   LL d = pw(b, e, Q);
   LL i = 0:
    if (d != 1) {
      j = -dis_log(d);
      if (j < 0) j = (j % (Q-1) + Q-1) % (Q-1);
    b = mul(b, pw(c, mul(P%(Q-1), j, Q-1), Q), Q);
   h = mul(h, pw(c, j, Q), Q);
    c = pw(c, P, Q);
  return mul(pw(A, alpha, Q), h, Q);
6.9 Miller-Rabin Test + Pollard Rho Factorization
namespace miller rabin {
   lint mul(lint x, lint y, lint mod){ return (__int128) x * y % mod; }
  lint ipow(lint x, lint y, lint p){
   lint ret = 1, piv = x \% p;
    while(y){
      if(y&1) ret = mul(ret, piv, p);
      piv = mul(piv, piv, p);
     y >>= 1;
   }
    return ret;
  bool miller_rabin(lint x, lint a){
   if(x % a == 0) return 0;
   lint d = x - 1;
    while(1){
     lint tmp = ipow(a, d, x);
      if(d&1) return (tmp != 1 && tmp != x-1);
      else if(tmp == x-1) return 0;
      d >>= 1:
   }
  bool isprime(lint x){
   for(auto &i : {2, 3, 5, 7, 11, 13, 17, 19, 23, 29, 31, 37}){
      if(x == i) return 1:
      if (x > 40 \&\& miller_rabin(x, i)) return 0;
    if(x \le 40) return 0;
    return 1;
 }
}
namespace pollard rho{
 lint f(lint x, lint n, lint c){
    return (c + miller_rabin::mul(x, x, n)) % n;
  void rec(lint n. vector<lint> &v){
   if(n == 1) return;
    if(n \% 2 == 0){
     v.push_back(2);
```

```
rec(n/2, v):
     return;
   }
   if(miller_rabin::isprime(n)){
     v.push_back(n);
     return:
   }
   lint a, b, c;
   while(1){
     a = rand() \% (n-2) + 2;
     b = a:
     c = rand() \% 20 + 1;
     do{
       a = f(a, n, c):
       b = f(f(b, n, c), n, c);
     \frac{1}{2} while \left( \gcd(abs(a-b), n) == 1 \right);
     if(a != b) break;
   lint x = gcd(abs(a-b), n);
   rec(x, v);
   rec(n/x, v);
 vector<lint> factorize(lint n){
   vector<lint> ret;
   rec(n, ret);
   sort(ret.begin(), ret.end());
   return ret;
 }
};
6.10 Highly Composite Numbers, Large Prime
 < 10^k
                 number divisors 2 3 5 71113171923293137
                     6
                                  4 1 1
                    60
                                 12 2 1 1
 3
                   840
                                 32 3 1 1 1
                  7560
                                 64 3 3 1 1
                                128 3 3 1 1 1
                  83160
                 720720
                8648640
                                448 6 3 1 1 1 1
 8
               73513440
                                768 5 3 1 1 1 1 1
 9
              735134400
                               1344 6 3 2 1 1 1 1
 10
             6983776800
                               2304 5 3 2 1 1 1 1 1
                               4032 6 3 2 2 1 1 1 1
 11
            97772875200
 12
           963761198400
                                     6 4 2 1 1 1 1 1 1
 13
          9316358251200
                              10752
                                     6321111111
 14
         97821761637600
                              17280
        866421317361600
                              26880
                                     6 4 2 1 1 1 1 1 1 1 1
 16
       8086598962041600
                              41472 8 3 2 2 1 1 1 1 1 1 1
 17
      74801040398884800
                              64512 6 3 2 2 1 1 1 1 1 1 1 1
     897612484786617600
                             103680 8 4 2 2 1 1 1 1 1 1 1 1
 < 10 k prime # of prime
                                      < 10^k
 1
              7
                                      10
                                                  9999999967
 2
             97
                          25
                                      11
                                                 9999999977
```

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9999999999971	13	1229	9973	4
9999999999973	14	9592	99991	5
99999999999989	15	78498	999983	6
999999999999937	16	664579	9999991	7
9999999999999997	17	5761455	99999989	8
99999999999999989	18	50847534	999999937	9

NTT Prime:

```
998244353 = 119 \times 2^{23} + 1. Primitive root: 3.
985661441 = 235 \times 2^{22} + 1. Primitive root: 3.
1012924417 = 483 \times 2^{21} + 1. Primitive root: 5.
```

7 Miscellaneous

7.1 Mathematics

- Tutte Matrix. For a simple undirected graph G, Let M be a matrix with entries $A_{i,j} = 0$ if $(i,j) \notin E$ and $A_{i,j} = -A_{j,i} = X$ if $(i,j) \in E$. X could be any random value. If the determinants are non-zero, then a perfect matching exists, while other direction might not hold for very small probability.
- Cayley's Formula. Given a degree sequence $d_1, d_2 \cdots, d_n$ for each labeled vertices, there exists $\frac{(n-2)!}{(d_1-1)!(d_2-1)!\cdots(d_n-1)!}$ spanning trees. Summing this for every possible degree sequence gives n^{n-2} .
- Kirchhoff's Theorem. For a multigraph G with no loops, define Laplacian matrix as L = D A. D is a diagonal matrix with $D_{i,i} = deg(i)$, and A is an adjacency matrix. If you remove any row and column of L, the determinant gives a number of spanning trees.
- Green's Theorem. Let C is positive, smooth, simple curve. D is region bounded by C. $\oint_C (Ldx + Mdy) = \iint_D (\frac{\partial M}{\partial x} \frac{\partial L}{\partial y})$

To calculate area, $\frac{\partial M}{\partial x} - \frac{\partial L}{\partial y} = 1$, common selection is $M = \frac{1}{2}x$, $L = -\frac{1}{2}y$.

Line integral of circle parametrized by $(x,y) = (x_C + r_C \cos \theta, \ y_C + r_C \sin \theta)$, when $\theta = t\theta_i + (1-t)\theta_f$, is given as follows:: $\frac{1}{2}(r_C(x_C(\sin \theta_f - \sin \theta_i) - y_C(\cos \theta_f - \cos \theta_i)) + (\theta_f - \theta_i)r_C^2)$.

Line integral of line parametrized by $(x, y) = t(x_1, y_1) + (1 - t)(x_2, y_2)$ is given as follows:: $\frac{1}{2}(x_1y_2 - x_2y_1)$.

• Burnside's lemma / Pólya enumeration theorem. let G and H be groups of permutations of finite sets X and Y. Let $c_m(g)$ denote the number of cycles of length m in $g \in G$ when permuting X. The number of colorings of X into |Y| = n colors with exactly r_i occurrences of the i-th color is the coefficient of $w_1^{r_1} \dots w_n^{r_n}$ in the following polynomial:

$$P(w_1, \dots, w_n) = \frac{1}{|H|} \sum_{h \in H} \frac{1}{|G|} \sum_{g \in G} \prod_{m \ge 1} (\sum_{h^m(b) = b} (w_b^m))^{c_m(g)}$$

When $H = \{I\}$ (No color permutation): $P(w_1, ..., w_n) = \frac{1}{|G|} \sum_{g \in G} \prod_{m > 1} (w_1^m + ... + w_n^m)^{c_m(g)}$

Without the occurrence restriction:

$$P(1,...,1) = \frac{1}{|G|} \sum_{g \in G} n^{c(g)}$$

where c(g) could also be interpreted as the number of elements in X that are fixed up to g.

Pick's Theorem. A = i + b/2 - 1, where: P is a simple polygon whose vertices are grid points, A is area of P, i is # of grid points in the interior of P, and b is # of grid points on the boundary of P. If h is # of holes of P (h + 1 simple closed curves in total), A = i + b/2 + h - 1.

```
// number of (x, y) : (0 <= x < n && 0 < y <= k/d x + b/d)

// argument should be positive

11 count_solve(11 n, 11 k, 11 b, 11 d) {
    if (k == 0) {
        return (b / d) * n;
    }
    if (k >= d || b >= d) {
        return ((k / d) * (n - 1) + 2 * (b / d)) * n / 2 + count_solve(n, k % d, b % d, d);
    }
    return count_solve((k * n + b) / d, d, (k * n + b) % d, k);
}

• Xudyh Sieve. F(n) = \sum_{d|n} f(d)
S(n) = \sum_{i \le n} f(i) = \sum_{i \le n} F(i) - \sum_{d=2}^{n} S\left(\left\lfloor \frac{n}{d} \right\rfloor\right)
Preprocess S(1) to S(M) (Set M = n^{\frac{2}{3}} for complexity)
S(n) = \sum f(i) = \sum_{i \le n} \left[F(i) - \sum_{j|i,j \ne i} f(j)\right] = \sum F(i) - \sum_{i/j = d = 2}^{n} \sum_{dj \le n} f(j)
S(n) = \sum if(i) = \sum_{i \le n} i \left[F(i) - \sum_{j|i,j \ne i} f(j)\right] = \sum iF(i) - \sum_{i/j = d = 2}^{n} \sum_{dj \le n} djf(j)
\sum_{d|n} \varphi(d) = n \qquad \sum_{d|n} \mu(d) = if (n > 1) \text{ then 0 else 1} \qquad \sum_{d|n} \mu(n^{\frac{n}{d}}) \sum_{e|d} f(e)) = f(n)
```

7.2 Popular Optimization Technique

- CHT. DnC optimization. Mo's algorithm trick (on tree). IOI 2016 Aliens trick. IOI 2009 Regions trick.
- Knuth's $O(n^2)$ Optimal BST : minimize $D_{i,j} = Min_{i \leq k < j}(D_{i,k} + D_{k+1,j}) + C_{i,j}$. Quadrangle Inequality : $C_{a,c} + C_{b,d} \leq C_{a,d} + C_{b,c}$, $C_{b,c} \leq C_{a,d}$. Now monotonicity holds.
- Sqrt batch processing Save queries in buffer, and update in every sqrt steps (cf : IOI 2011 Elephant. hyea calls it "ainta technique")
- Dynamic insertion in static set (Make $O(\log n)$ copy. Merge like binomial heap.)
- Offline insertion / deletion in insert-only set (Pair insertion-deletion operation, and regard it as range query)
- Atcoder Median Pyramid: Reduce the input to binary, and solve the easier problem.
- LP Duality. max $c^T x$ sjt to $Ax \leq b$. Dual problem is min $b^T x$ sjt to $A^T x \geq c$. By strong duality, min max value coincides.

7.3 Fast LL Division / Modulo

```
inline void fasterLLDivMod(unsigned long long x, unsigned y, unsigned &out_d, unsigned &out_m) {
  unsigned xh = (unsigned)(x >> 32), xl = (unsigned)x, d, m;
#ifdef __GNUC__
  asm(
   "divl %4; \n\t"
  : "=a" (d), "=d" (m)
  : "d" (xh), "a" (xl), "r" (y)
  );
#else
  __asm {
  mov edx, dword ptr[xh];
  mov eax, dword ptr[xl];
  div dword ptr[y];
  mov dword ptr[d], eax;
  mov dword ptr[m], edx;
};
```

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```
#endif
  out_d = d; out_m = m;
}
//x < 2^32 * MOD !
inline unsigned Mod(unsigned long tong x){
  unsigned v = mod:
 unsigned dummy, r;
 fasterLLDivMod(x, y, dummy, r);
 return r;
7.4 Bit Twiddling Hack
int __builtin_clz(int x);// number of leading zero
int __builtin_ctz(int x);// number of trailing zero
int __builtin_clzll(long long x);// number of leading zero
int __builtin_ctzll(long long x);// number of trailing zero
int __builtin_popcount(int x);// number of 1-bits in x
int __builtin_popcountll(long long x);// number of 1-bits in x
lsb(n): (n & -n); // last bit (smallest)
floor(log2(n)): 31 - __builtin_clz(n | 1);
floor(log2(n)): 63 - __builtin_clzll(n | 1);
// compute next perm. ex) 00111, 01011, 01101, 01110, 10011, 10101...
long long next_perm(long long v){
 long long t = v \mid (v-1);
 return (t + 1) \mid (((^t \& -^t) - 1) >> (_builtin_ctz(v) + 1));
7.5 Fast Integer IO
static char buf[1 << 19]: // size : any number geg than 1024
static int idx = 0;
static int bytes = 0:
static inline int _read() {
 if (!bytes || idx == bytes) {
   bytes = (int)fread(buf, sizeof(buf[0]), sizeof(buf), stdin);
   idx = 0;
  return buf[idx++];
static inline int _readInt() {
 int x = 0, s = 1;
 int c = _read();
 while (c \le 32) c = read():
  if (c == '-') s = -1, c = _read();
  while (c > 32) x = 10 * x + (c - '0'), c = _read();
 if (s < 0) x = -x:
  return x;
7.6 OSRank in g++
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace "gnu pbds;
typedef
tree<int, null_type, less<int>, rb_tree_tag, tree_order_statistics_node_update> ordered_set;
```

```
ordered set X:
X.insert(1); X.insert(2); X.insert(4); X.insert(8); X.insert(16);
cout<<*X.find_by_order(1)<<endl; // 2</pre>
cout<<*X.find_by_order(2)<<endl; // 4</pre>
cout<<*X.find_by_order(4)<<endl; // 16</pre>
cout<<(end(X)==X.find_by_order(6))<<endl; // true</pre>
cout<<X.order_of_key(-5)<<endl; // 0</pre>
cout<<X.order_of_key(1)<<endl; // 0</pre>
cout<<X.order_of_key(3)<<endl; // 2</pre>
cout<<X.order_of_key(4)<<endl; // 2</pre>
cout<<X.order_of_key(400)<<endl; // 5</pre>
7.7 Nasty Stack Hacks
// 64bit ver.
int main2(){ return 0; }
int main(){
 size t sz = 1 << 29: // 512MB
  void* newstack = malloc(sz);
  void* sp_dest = newstack + sz - sizeof(void*);
  asm __volatile__("movq %0, %%rax\n\t"
  "movg %%rsp , (%%rax)\n\t"
 "movq %0, %%rsp\n\t": : "r"(sp_dest): );
 main2();
 asm __volatile__("pop %rsp\n\t");
 return 0;
7.8 C++ / Environment Overview
// vimrc : set nu sc ci si ai sw=4 ts=4 bs=2 mouse=a svntax on
// compile : g++ -o PROB PROB.cpp -std=c++11 -Wall -02
// options : -fsanitize=address -Wfatal-errors
struct StupidGCCCantEvenCompileThisSimpleCode{
 pair<int, int> array[1000000];
}; // https://gcc.gnu.org/bugzilla/show_bug.cgi?id=68203
// how to use rand (in 2018)
mt19937 rng(0x14004);
int randint(int lb, int ub){ return uniform_int_distribution<int>(lb, ub)(rng); }
// comparator overload
auto cmp = [](seg a, seg b){ return a.func() < b.func(); };</pre>
set<seg, decltype(cmp)> s(cmp);
map<seg, int, decltype(cmp)> mp(cmp);
priority_queue<seg, vector<seg>, decltype(cmp)> pq(cmp); // max heap
// hash func overload
struct point{
    int x, y;
   bool operator == (const point &p)const{ return x == p.x && y == p.y; }
};
struct hasher {
    size_t operator()(const point &p)const{ return p.x * 2 + p.y * 3; }
unordered_map<point, int, hasher> hsh;
```