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Problem: Neighbour (62)

Description: For a given large integer **n**, find the **nearest perfect power** that is equal or smaller than **n**.

N is given once a string(X) is found where sha256(X)[-6:] = *Randomized hex code*

Overview:

The crux of the problem revolved around finding the nearest perfect square(e.g 4, 9, 16...etc) to a large given number. In its simplest form, a perfect power is simply m^k = n, where m and k have to be positive integers and therefore n will be a natural number. However instead of focusing on just one exponent, our search would have to be over a series of different exponents and bases.

Beginning Steps:

Upon connecting to the server, you are presented with the following output.

From this we can extract our formula, which is n-x**y=r, where r>0, x,y>1, and our goal is to find the smallest possible value for r. However, in order to get a value for N and begin the optimization problem, we must first tackle finding a string, such that sha256(x)[-6:]="series of psuedo random hex bytes" (always 6 long). We decided on going for a simple brute force approach. Essentially, we designed a for loop to test a large number of values. Each iteration

```
import hashlib
for i in range(100000000):
    m=hashlib.sha256()
    m.update(str(i))
    code=m.hexdigest()
    if(code[-6:]=="d37879"):
        print("HERE")
        print(i)
        x=rawinput()
```

of the loop would take the iteration number, and then take the sha256 hash of that i'th number (as a string). Then there is a simple comparison between the last 6 digits of the i'th hash digest and the given digest on the server. Since there are many solutions due the short length, we simply break the for loop after the first value is found. This gave us the output of 268684. After the server spits back a large value for a n. Allowing us the begin the hard math.

Part 2:

At first we planned on just doing a simple brute force, with some logical parameters. Since the given n was upwards of 100 digits, (we even got one that 500+ digits), we erroneously assumed the value for y (the exponent) would also have to be relatively large. Therefore, we wrote a quick for loop to have a set x (base) and then an ever increasing exponent value. We would then look at N-x**y and quickly saw that this method had some fatal errors. For starters, the base would also have to change. This can be seen with the following example:

N=12 $x=2 \rightarrow y = \{2,3\}$. This gives the smallest value for R being 4. N=12 $x=3 \rightarrow y=\{2\}$. This gives the smallest value for R being 3, and thus smaller than base of 2.

After reviewing this information, we assumed the new correct path to take would be to use a theorem solver, specifically z3. However, this ended up giving inconclusive results since finding a minimum value using z3 is quite difficult. Therefore we abandoned this path, in favor of a more efficient brute force.

We initially planned to write an algorithm that checks N-x^y, where x is determined in an initial for loop, and y is nested within that – essentially trying every base with every exponent (where x and y are integers). Theoretically this would work, however with these extremely large N numbers the amount of iterations would have to be extraordinary. Therefore, we implemented a binary search to find the correct exponent and base, so as not to exceed the n value/ break the equation. After doing this, finding the nearest perfect power is found very quickly. Since the challenge would give you multiple n values to find different r's, we simply changed the n value in the script each round. This eventually led to the flag which is: ASIS{36812f76cce2753e482ac6f68f9d3012}





