CSU34011 Symbolic Programming

Second of Two Assessed Assignments Submit to Blackboard by Nov 24 (Wed)

Problem 1 Exercise 6.6 in Learn Prolog Now describes a street with

(*) three neighbouring houses that all have a different colour, namely red, blue, and green. People of different nationalities live in the different houses and they all have a different pet.

Leaving out all the other constraints mentioned in that exercise, write a DCG that outputs strings

```
[h(Col1,Nat1,Pet1), h(Col2,Nat2,Pet2), h(Col3,Nat3,Pet3)] satisyfing (*), where the nationalities are
```

and the pets are

```
jaguar, snail, zebra.
```

Use the predicate nbd/2 for the 3 houses so that for example,

[20 marks]

Problem 2 For an integer $n \geq 0$, the *n*th *Fibonacci number* F_n is defined as follows

$$F_0 := 0$$
 $F_1 := 1$
 $F_{n+2} := F_n + F_{n+1}$

giving $F_2 = 1$, $F_3 = 2$, $F_4 = 3$, $F_5 = 5$, etc. Define a DCG that generates for every $k \ge 1$, lists $[F_0, F_1, \ldots, F_k]$ so that, for example,

```
?- fib(L,[]).

L = [0,1];

L = [0,1,1];

L = [0,1,1,2];

L = [0,1,1,2,3];

L = [0,1,1,2,3,5];
```

[20 marks]

Problem 3 For each integer n > 0, let

```
L_n := \{s \in \{0,1\}^+ \mid s \text{ ends in a string from } 1(0+1)^{n-1}\}
```

be the set of bit-strings whose n-th to the last bit is 1. That is, L_n is described by the regular expression

$$(0+1)^*1(0+1)^{n-1}$$
.

- (a) Define predicates tran/4 and final/2 so that for every integer n > 0,
 - (i) tran(n,Q1,X,Q2) picks out transitions between states Q1 to Q2 labeled by X, and
 - (ii) final(n, Q) picks out final states Q

for a finite automaton with initial state ${\tt q0}$ accepting L_n according to

For example,

```
?- accept(3,L).
L = [1, 0, 0];
L = [1, 0, 1];
L = [1, 1, 0];
L = [1, 1, 1];
L = [0, 1, 0, 0];
L = [0, 1, 0, 1];
...
```

[20 marks]

(b) Define a DCG for the 3-ary predicate s/3 such that s(n, s, []) is true exactly if s encodes a string in L_n . For example,

```
?- s(3,[A,1,Z],[]).

A = 1, Z = 0;

A = 1, Z = 1;

false.
```

[20 marks]

(c) Define predicates ith/3 and initial/3 such that if ai is the ith string returned by the query s(n,X,[]) then

```
ith(i, n, Z) is true exactly if Z = ai
```

and

```
initial(i, n, Z) is true exactly if Z = [ai, ..., a2, a1].
```

For example, assuming

```
?- s(3,X,[]).

X = [1,0,0];

X = [1,0,1];

X = [1,1,0];

X = [1,1,1];

X = [0,1,0,0];
```

then

```
?- ith(5,3,A). A = [0,1,0,0].
```

and

```
?- initial(5,3,L). 
  L = [[0,1,0,0], [1,1,1], [1,1,0], [1,0,1], [1,0,0]].
```

[20 marks]