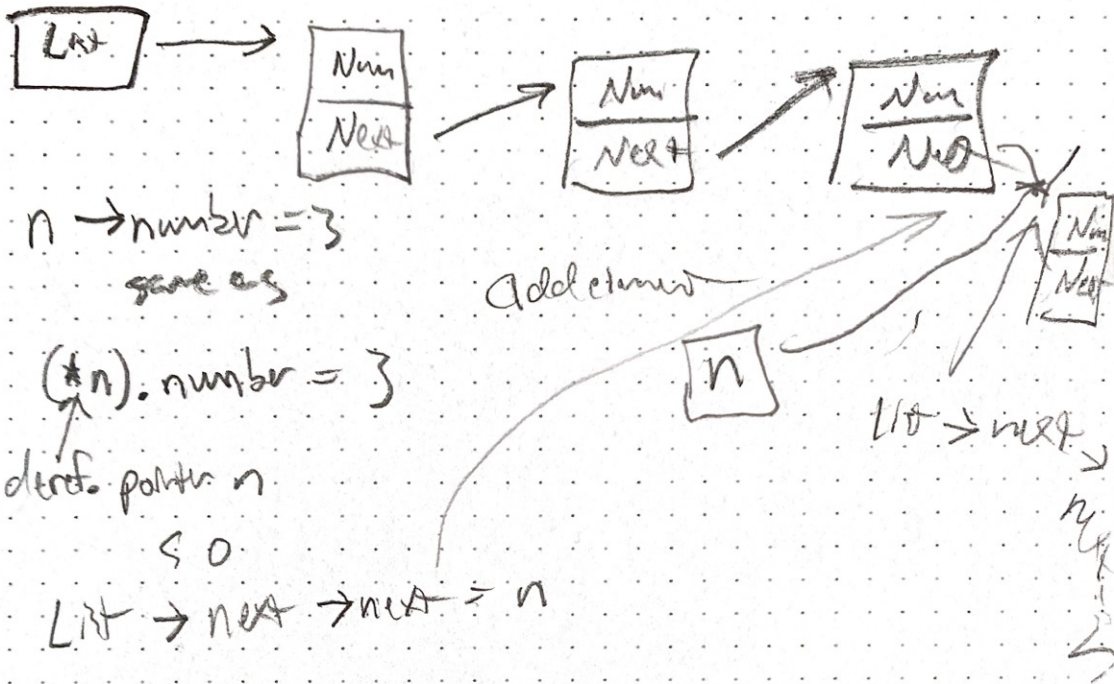


Linked list



$n \rightarrow \text{number} = 3$
same as

$(*n). \text{number} = 3$

def. pointer n

so

$\text{List} \rightarrow \text{next} \rightarrow \text{next} = n$

$(*(*\text{list}). \text{next}). \text{next} = n$

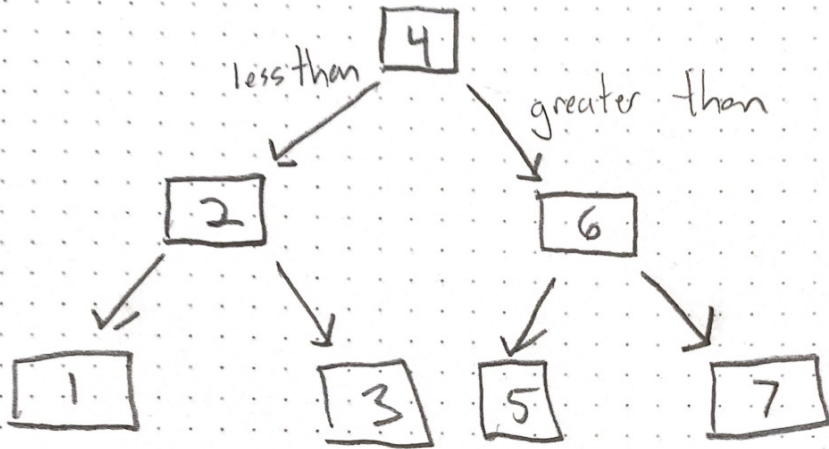
messy notation

Instead, list points to next , which points to next

$\text{list} \rightarrow \text{next} \rightarrow \text{next} = [\text{the address of}]$

Introduce

Binary Trees

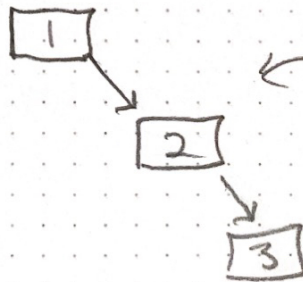


- recursive structure
 - each tree has 2 children
 - each child is a tree

- Search is efficient

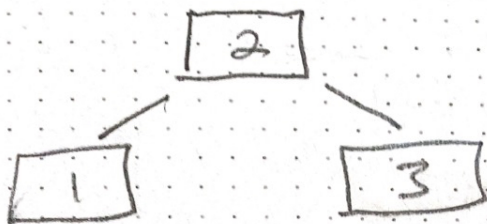
- recursive data structures lend themselves to recursive functions

* a devolution of tree into linked list solution



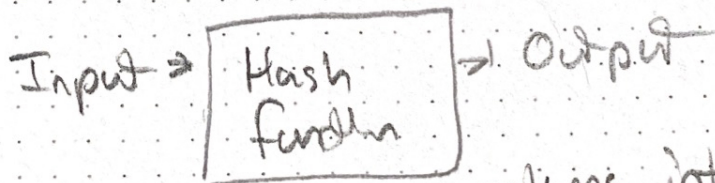
← technically binary tree but also linked list

Make 2 root node



Can we search in constant time
Yes - with Hash tables

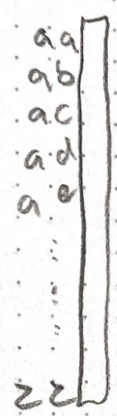
* we use ASCII to an array of
convert index



char *
Zacharias → [] → 25
returns int (index)

To solve inefficiency of Hash search

Solution 1 \rightarrow grow Hash table



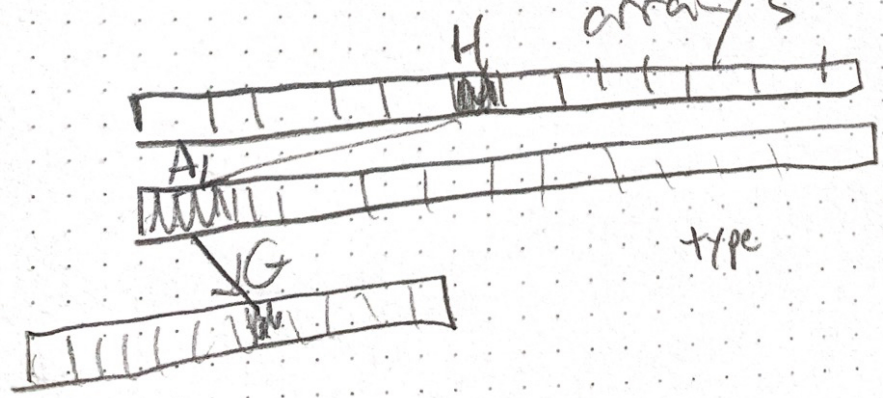
But Harry and Hagrid go to same location

we could use first 3 letters

But we grow the table exponentially
 $26 \rightarrow 26^2 \rightarrow 26^3$ etc

Tries (or prefix tree)

a tree made up of arrays



parent node
 child node

type