

HW <sup>7</sup> Sean Poston

1. a) If you replace all variables with True/False, does the formula resolve to True? If yes, it's satisfiable.
- b)  $O(n)$ , if only one variable per clause,  $O(n^2)$  if variables = clauses.  
We only need to prove that not all clauses contain a literal and its negation.
- ```
if (!clause[0].contains(n, !n)) {  
    break; // the formula is DNF-satisfiable at this point.  
}
```

c) take  $(a \vee b) \wedge (a \vee c)$  in CNF form  
We can use the distributive law to say:  
 $a \vee (b \wedge c)$  which is DNF form

d) Because Dr. Love didn't solve for ALL possible formulas + K-SAT.  
Or maybe he's just in it for the love of the game.

2. decisions: Given the vertices in  $G$ , are any of them connected by an edge?

NP: Check every vertex. If the edges extending from the vertices touch another vertex in  $G$ , we reject it. Otherwise, we accept it.