

# CS345 Homework 4

5.

$$p \wedge C_1 \{S\} q$$

$$p \wedge \neg C_1 \wedge C_2 \{S_2\} q$$

$$p \wedge \neg C_1 \wedge \neg C_2 \wedge \dots \wedge C_{n-1} \{S_n\} q$$

6. I think I'm misunderstanding the tautology assertion because if the initial assertion were always true, then  $y$  would be  $-2$ . Otherwise, it would take the second stage of the above:

$$p \wedge \neg C_1 \wedge C_2 \{y := 2\} y = 2$$

$$2. i \leq n \{ \text{power} := \text{power} * x \wedge i := i + 1 \} i > n$$

$$12. r := a \} S_1$$

$$q := 0$$

$$\text{while } r \neq d \} S_2$$

$$r := r - d$$

$$q := q + 1$$

$$(a \wedge d) \wedge \{S_1\} p \wedge (q = 0)$$

$$r \geq d \{ r := r - d \wedge q := q + 1 \} (q = dq + r) \wedge (0 \leq r < d)$$

This is true. If  $a$  is 2 and  $d$  is 1, then it will loop twice, leaving  $r = 0$ ,  $d = 1$ ,  $q = 2$ ,  $a = 2$   
 $a = dq + r \Rightarrow 2 = 1(2) + 0 \Rightarrow 2 = 2 \checkmark$   
 $0 \leq r < d \Rightarrow 0 \leq 0 < 1 \checkmark$

228 1. a.) Yes

b.) No

c.) Yes

d.) Yes

e.) No

f.) No

$$229 \quad 30. a.) O(x) = \Omega(x)$$

$$b.) O(\log x) = \Omega(\log x)$$

$$c.) O(x^2) = \Omega(x^2)$$

$$d.) O(\log x) = \Omega(\log x)$$

$$e.) O(x) = \Omega(x)$$

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11. a.) while  $i < n$

b.) It would be  $O(n^2)$

while  $j < n_2$

if  $k$

true

242 12 For  $i$  to  $n \Rightarrow n$

For  $j$  to  $n \Rightarrow n$

For  $k$  to  $j \Rightarrow n$

$m[i][j] = \min[m[i][j], a[k]] \Rightarrow \text{assignment}$

$n^3 + \text{assign}$

$O(n^3)$

b.)