Now, for a given degree of accuracy, how many iterations do we actually need? Consider

$$2^{-x} = x$$
 on  $\left[\frac{1}{3}, 1\right]$ 

The bounds are given by

$$|p_n - p| \le k^n \max(p_0 - a, b - p_0)$$

using the initial guess, and by

$$|p_n - p| \le \frac{k^n}{1 - k} |p_1 - p_0|$$

using the initial guess and first iteration. We will look at a few applications of the bounds to the problem above. First notice that  $g'(x) = -\ln(2)2^{-x}$  and  $|g'(x)| \le \ln(2)2^{-1/3} < k = 0.551$ , where k = 0.551 is a bound on the magnitude of g'(x).

The worst possible initial guess would be at one of the endpoints, so we will start there (this maximizes the term  $\max(p_0 - a, b - p_0)$ , which in this case we actually want to do in order to generate a conservative bound). Taking D as the desired accuracy (i.e., an accuracy within  $10^{-D}$ ), this gives,

$$k^{n} \max(p_{0} - a, b - p_{0}) < 10^{-D}$$

$$(0.551)^{n} \left(\frac{2}{3}\right) < 10^{-D}$$

$$(0.551)^{n} < \left(\frac{3}{2}\right) 10^{-D}$$

$$n \log(0.551) < \log\left(\frac{3}{2}\right) - D$$

$$n > \frac{\log\left(\frac{3}{2}\right) - D}{\log(0.551)}$$

In the last line, the inequality has been reversed since we are dividing by a negative. With D=4 this gives n>14.77277 which requires N=15 steps.

For the second bound, we actually need  $p_1$  in addition to  $p_0$ . From  $p_0 = \frac{2}{3}$ , we have  $p_1 = 2^{-1/3}$  (so  $|p_1 - p_0| = |2^{-1/3} - \frac{1}{3}| \approx 0.4604$ ). Similarly, from  $p_0 = 1$ , we have  $p_1 = \frac{1}{2}$  (so  $|p_1 - p_0| = |\frac{1}{2} - 1| = 0.5$ ). We will use the second of these which is larger in value.

$$\frac{k^n}{1-k}|p_1 - p_0| < 10^{-D}$$

$$\frac{(0.551)^n}{1 - 0.551}(0.5) < 10^{-D}$$

$$(0.551)^n < \left(\frac{1 - 0.551}{0.5}\right)10^{-D}$$

$$n\log(0.551) < \log\left(\frac{1 - 0.551}{0.5}\right) - D$$

$$n > \frac{\log\left(\frac{1 - 0.551}{0.5}\right) - D}{\log(0.551)}$$

For consistency, with D=4 this gives n>15.63357 which requires N=16 steps. We have to do at least 16 steps to ensure we are within the bound, though we may satisfy this much more quickly. Notice that this is quite a bit more work than our bound for the Bisection method required.