

CS 565 Spring 2022 Homework 3 (Big-Step Semantics)

Your name: _____

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Using the big-step operational semantics of IMP (shown in Figure 1), either (i) fill in the final states produced by each of the following IMP programs and give a derivation justifying your answer **or** (ii) state that the program does not terminate. **Note:** You don't need to supply the derivations for any arithmetic or boolean expressions, just for the IMP statements.

Problem 1 (1 point).

$\emptyset, X := 0; Y := 1 \Downarrow$

Problem 2 (1 point).

$[X \mapsto 2], \text{ if } (X \leq 1) \text{ then } Y := X + 3 \text{ else } Z := 4 \text{ end } \Downarrow$

Problem 3 (1 point).

$[X \mapsto 0]$, **while** $(X \leq 1)$ **do** $Y := Y + 1$ **end** \Downarrow

Problem 4 (1 point).

$[Y \mapsto 0]$, **while** $(Y \leq 1)$ **do** $Y := Y + 1$ **end** \Downarrow

Problem 5 (2 points). Several cutting edge languages, including Perl, Visual Basic, and Pascal include a **repeat c until b** loop construct. These loops differ from the standard **while** loops in two ways:

1. the loop guard is checked /after/ the execution of the body, so the loop always executes at least once.
2. the loop continues executing as long as the condition is false.

Write down the big-step reduction rules for **repeat** loops in IMP.

$$\begin{array}{c}
\frac{}{\sigma, x \Downarrow_A \sigma(x)} \text{(E-VAR)} \quad \frac{}{\sigma, n \Downarrow_A n} \text{(E-CONST)} \quad \frac{\sigma, a_1 \Downarrow_A m \quad a_2 \Downarrow_A n}{\sigma, a_1 + a_2 \Downarrow_A m + n} \text{(E-ADD)} \\
\hline
\frac{}{\sigma, \text{true} \Downarrow_B \text{true}} \text{(E-TRUE)} \quad \frac{}{\sigma, \text{false} \Downarrow_B \text{false}} \text{(E-FALSE)} \quad \frac{\sigma, b_1 \Downarrow_B t \quad b_2 \Downarrow_B v}{\sigma, b_1 \wedge b_2 \Downarrow_B t \wedge v} \text{(E-AND)} \\
\\
\frac{\sigma, b \Downarrow_B \text{true}}{\sigma, \neg b \Downarrow_B \text{false}} \text{(E-NOTT)} \quad \frac{\sigma, b \Downarrow_B \text{false}}{\sigma, \neg b \Downarrow_B \text{true}} \text{(E-NOTF)} \\
\\
\frac{\sigma, a_1 \Downarrow_A n \quad \sigma, a_2 \Downarrow_A n}{\sigma, a_1 = a_2 \Downarrow_B \text{true}} \text{(E-CMPT)} \quad \frac{\sigma, a_1 \Downarrow_A m \quad \sigma, a_2 \Downarrow_A n \quad m \neq n}{\sigma, a_1 = a_2 \Downarrow_B \text{true}} \text{(E-CMPF)} \\
\hline
\\
\frac{}{\sigma, \text{skip} \Downarrow \sigma} \text{(E-SKIP)} \quad \frac{\sigma, a \Downarrow_A n}{\sigma, x := a \Downarrow \sigma[x \mapsto n]} \text{(E-ASS)} \\
\\
\frac{\sigma, c_1 \Downarrow \sigma_1 \quad \sigma_1, c_2 \Downarrow \sigma_2}{\sigma, c_1; c_2 \Downarrow \sigma_2} \text{(E-SEQ)} \quad \frac{\sigma, b \Downarrow_B \text{true} \quad \sigma, c_1 \Downarrow \sigma_1}{\sigma, \text{if } b \text{ then } c_1 \text{ else } c_2 \Downarrow \sigma_1} \text{(E-IFTRUE)} \\
\\
\frac{\sigma, b \Downarrow_B \text{false} \quad \sigma, c_2 \Downarrow \sigma_2}{\sigma, \text{if } b \text{ then } c_1 \text{ else } c_2 \Downarrow \sigma_2} \text{(E-IFFALSE)} \quad \frac{\sigma, b \Downarrow_B \text{false}}{\sigma, \text{while } b \text{ do } c \Downarrow \sigma} \text{(E-WHILEFALSE)} \\
\\
\frac{\sigma, b \Downarrow_B \text{true} \quad \sigma, c \Downarrow \sigma_1 \quad \sigma_1, \text{while } b \text{ do } c \Downarrow \sigma_2}{\sigma, \text{while } b \text{ do } c \Downarrow \sigma_2} \text{(E-WHILETRUE)}
\end{array}$$

Figure 1: Big-step semantics of arithmetic expressions, boolean expressions, and regular IMP.