Lecture 6

CIS 341: COMPILERS

Announcements

- HW2: X86lite
 - Due: Weds, February 7th at 11:59:59pm
 - Pair-programming:
 - Register the group on the submission page
 - Submission by any group member counts for the group

see: ir-by-hand.ml, ir<X>.ml

INTERMEDIATE REPRESENTATIONS

Eliminating Nested Expressions

- Fundamental problem:
 - Compiling complex & nested expression forms to simple operations.

```
Source ((1 + X4) + (3 + (X1 * 5)))

Add(Add(Const 1, Var X4),
Add(Const 3, Mul(Var X1,
Const 5)))

IR
```

- Idea: name intermediate values, make order of evaluation explicit.
 - No nested operations.

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Translation to SLL

Given this:

```
Add(Add(Const 1, Var X4),
Add(Const 3, Mul(Var X1,
Const 5)))
```

Translate to this desired SLL form:

```
let tmp0 = add 1L varX4 in
let tmp1 = mul varX1 5L in
let tmp2 = add 3L tmp1 in
let tmp3 = add tmp0 tmp2 in
tmp3
```

- Translation makes the order of evaluation explicit.
- Names intermediate values
- Note: introduced temporaries are never modified

Intermediate Representations

- IR1: Expressions
 - simple arithmetic expressions, immutable global variables
- IR2: Commands
 - global *mutable* variables
 - commands for update and sequencing
- IR3: Local control flow
 - conditional commands & while loops
 - basic blocks
- IR4: Procedures (top-level functions)
 - local state
 - call stack

Basic Blocks

- A sequence of instructions that is always executed starting at the first instruction and always exits at the last instruction.
 - Starts with a label that names the entry point of the basic block.
 - Ends with a control-flow instruction (e.g. branch or return) the "link"
 - Contains no other control-flow instructions
 - Contains no interior label used as a jump target
- Basic blocks can be arranged into a control-flow graph
 - Nodes are basic blocks
 - There is a directed edge from node A to node B if the control flow instruction at the end of basic block A might jump to the label of basic block B.

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