

# Algorithms In A White Box

First, solve the problem. Then, write the code.

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# BITS Manipulation

## Least Significant Set Bit

First thing we need to notice is that when we add 1 to a number  $N$ , what we are doing is just converting the first (right to left) 0-bit into a 1-bit and the 1-bits before get converted to 0-bits because  $1 + 1 = 0$  with carry of 1 in binary, therefore we will be having a carry of 1-bit until we find a 0-bit.

**Example:**

$$00100111 + 1 = 00101000$$

Second thing we need to notice is very simple, lets start by denoting  $\overline{N}$  as  $N$  with all it's bits inverted (1-bits change to 0-bit and viceversa), if we perform an *AND* operation between  $N$  and  $\overline{N}$  we will get all bits in 0 as result.

**Example:**

$$N = 00100111$$

$$\overline{N} = 11011000$$

So, to achieve our main objective which is to extract the least significant bit (rightmost bit) we can just invert  $N$  and add 1 to it that will convert the first 0-bit to 1-bit so if we make an *AND* operation with  $N$  and  $\overline{N}$  we get everything before the lsb as 0-bit and after the lsb we also get everything as 0-bit.

And we can write this as the 2's complement since what we did was just to invert bits and add one, which is just the exact definition of 2's complement.

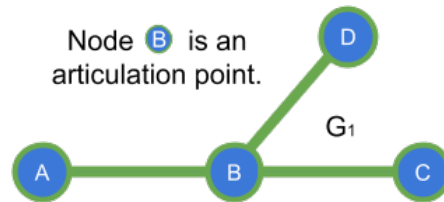
**C++ Code:**

```
1 int lsb(int n) {  
2     return n & -n;  
3 }
```

# Graph Theory

## Articulation Points And Bridges

Let's define what an *articulation point* is. We say that a vertex  $V$  in a graph  $G$  with  $C$  connected components is an *articulation point* if its removal increases the number of connected components of  $G$ . In other words, let  $C'$  be the number of connected components after removing vertex  $V$ , if  $C' > C$  then  $V$  is an *articulation point*.



## How to find articulation points?

•Naive approach  $O(V * (V + E))$

```
1 for every vertex V in the graph G do
2   Remove V from G
3   if the number of connected components increases then V is an articulation point
4   Add V back to G
```