

	1000	3000	5000	7000	9000	11000
qs1	10996.72	39617.142	71029.308	103930.394	138403.996	173707.299
qs2	11725.66	40882.576	72446.146	105562.226	139515.856	174076.651
qs3(k = 10)	11171.118	40239.396	72061.51	105390.382	140267.16	175979.937
qs3(k = 100)	22629.122	74411.726	128878.508	185285.804	242794.526	301731.621
qs4(p = .01)	10996.746	39618.536	71034.47	103943.53	138424.906	173747.657
qs4(p = .10)	11278.502	43130.678	79916.844	122885.618	169516.446	231452.878
qs4(p = .25)	16493.414	99938.624	236079.482	454607.886	714832.814	1048181.75

	3	3.47712125	3.69897	3.84509804	3.95424251	4.04139269
qs1	10996.72	39617.142	71029.308	103930.394	138403.996	173707.299
qs2	11725.66	40882.576	72446.146	105562.226	139515.856	174076.651
qs3(k = 10)	11171.118	40239.396	72061.51	105390.382	140267.16	175979.937
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	1000	3000	5000	7000	9000	11000
qs1	629.55926	1968.04088	3136.77685	4389.08507	5688.87745	7090.13237
qs2	351.529268	976.039192	1676.00105	2526.1095	2982.80089	3498.0072
qs3(k = 10)	626.8603	1964.05029	3134.13707	4394.56348	5695.26214	7088.2555
qs3(k = 100)	1136.38994	2497.4235	3683.09574	5065.26027	6460.99512	7806.10196
qs4(p = .01)	629.582017	1969.95002	3143.04005	4400.9583	5726.50892	7134.40529
qs4(p = .10)	986.261965	6750.44536	14435.7288	30900.7447	49404.4222	85973.5849
qs4(p = .25)	5195.88095	49446.6397	130570.956	287212.706	437521.126	655476.061

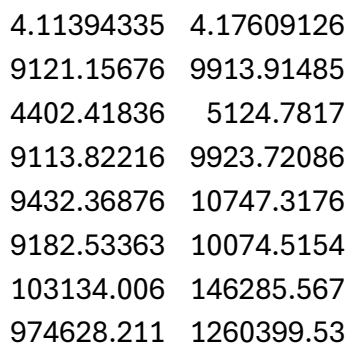
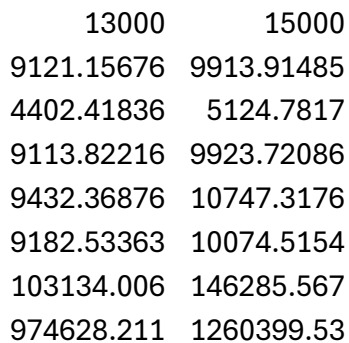
	3	3.47712125	3.69897	3.84509804	3.95424251	4.04139269
qs1	629.55926	1968.04088	3136.77685	4389.08507	5688.87745	7090.13237
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Of the 7 variations of quicksort, in timing of using 1000 elements, QuickSort3 had the fastest quicksort3 was $O(n \log n)$ with a worst case complexity of $O(n^2)$. When running all of the quicksort functions on the chrono time standard the functions ran in speeds respectively as follows: 506306ms, 432490ms, 520808 ms, 492269 ms, and 375993.

As you can see from the following numbers, Quicksort3 with the k=100 had the fastest variation when passing in p=.25. This is due to the decision making of pivot selection. Quicksort4 spec

on the parameter p passed in and quicksort3 for the k respectively. In quicksort3 (our winner), elements helps with avoiding our "worst case scenarios" by comparing the last 3 elements and before sending it to the back to be sorted by that value. In this case, if we had 8,7,9 then 8 would sort the other two in which the median value is chosen, this time adds up as we recursively choose a median when we get down to the last leaves and get down to sizes of 3 for our vector. Asymptotically, quicksort2 due to the lowest standard deviation on average as our n climbed, on the charts with numbers to a log base 10 the quicksort 2 was the fastest by this standard. Overall the basic C++ implementation is faster than the vanilla implementation.

4.11394335	4.17609126
209931.606	246248.62
209660.166	245434.028
212623.296	249351.874
360951.682	420381.17
209973.994	246349.116
283736.726	349051.184
1419401.39	1870341.64



ion followed by QuickSort4
specifically will vary widely based

, the pivot selection of the 3
nd choosing the median value
ould be chosen and already
call the function choosing a
tomatically, my fastest was
ve see when we set the
Quicksort 1 was the slowest due

