



UNIVERSITÄT ZU LÜBECK

Fully Dynamic Bin Packing Revisited

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Fully Dynamic Bin Packing = Online + Removal + Repacking

INSERT: $a/0.2$



Fully Dynamic Bin Packing = Online + Removal + Repacking

INSERT: $a/0.2$, **INSERT:** $b/0.7$



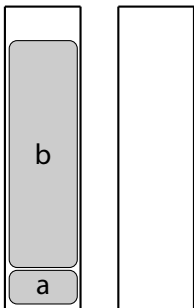
Fully Dynamic Bin Packing = Online + Removal + Repacking

INSERT: $a/0.2$, **INSERT:** $b/0.7$, **INSERT:** $c/0.4$



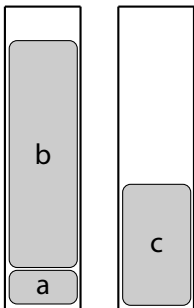
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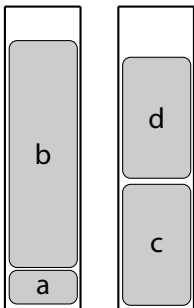
Fully Dynamic Bin Packing = Online + Removal + Repacking

INSERT: $a/0.2$, **INSERT:** $b/0.7$, **INSERT:** $c/0.4$, **INSERT:** $d/0.4$



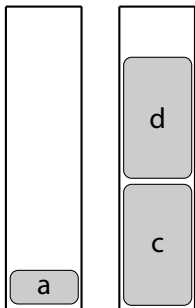
Fully Dynamic Bin Packing = Online + Removal + Repacking

INSERT: $a/0.2$, **INSERT:** $b/0.7$, **INSERT:** $c/0.4$, **INSERT:** $d/0.4$, **REMOVE:** $b/0.7$



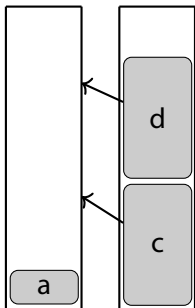
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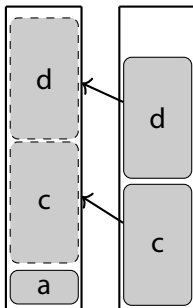
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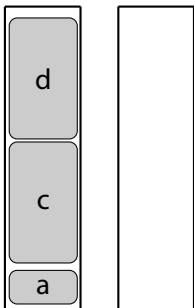
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$$\text{Migration Factor} = \frac{\text{c} + \text{d}}{\text{b}} = \frac{\text{SIZE(moved)}}{\text{SIZE(new/removed)}}$$

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$$\text{Migration Factor} = \frac{\text{c} + \text{d}}{\text{b}} = \frac{\text{SIZE(moved)}}{\text{SIZE(new/removed)}}$$

$$\text{Shifting Moves} = 2 = \# \text{ moved}$$

Lower Bound

MF of $\Omega(1/\varepsilon)$ is necessary for rate $1 + \varepsilon$

$$a = 1/2 - b/3, \quad b = 1/MF$$

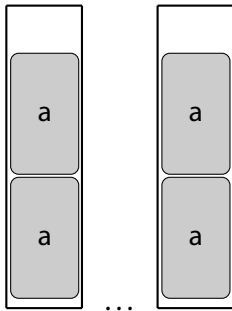
INSERT: a^+

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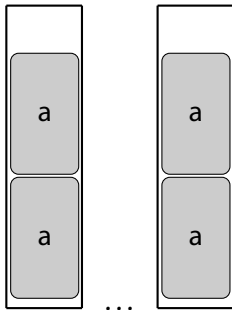
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INSERT: a^+ ,

INSERT: b^+



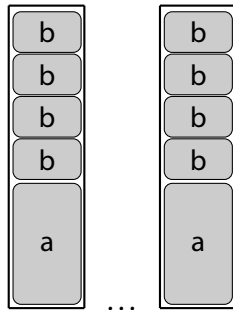
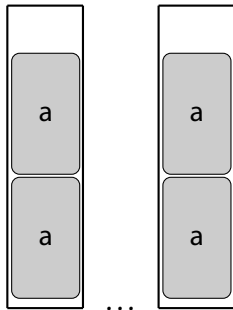
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Fully Dynamic Bin Packing

Known Results

Rate	Shifting Moves	Migration Factor	Authors
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Fully Dynamic Bin Packing

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Rate	Shifting Moves	Migration Factor	Authors
$5/4$	7 (amortized)	∞	Ivković, Lloyd (1998)

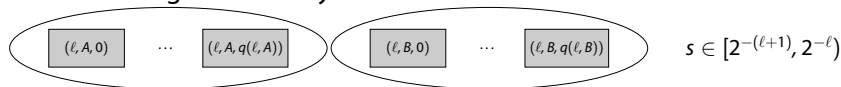
Fully Dynamic Bin Packing

Known Results

Rate	Shifting Moves	Migration Factor	Authors
$5/4$	7 (amortized)	∞	Ivković, Lloyd (1998)
$1 + \varepsilon$	$\mathcal{O}(1/\varepsilon^4 \cdot \log(1/\varepsilon))$	$\mathcal{O}(1/\varepsilon^4 \cdot \log(1/\varepsilon))$	<i>B., Jansen, Klein (2015)</i>

Large Items

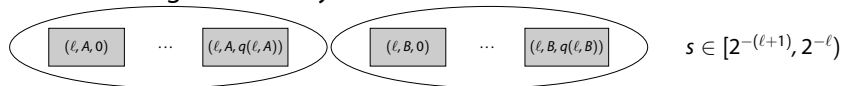
Round Items geometrically:



■ $|(\ell, A, \cdot)| = 2^\ell \cdot \text{SIZE}, \quad |(\ell, B, \cdot)| = 2^\ell \cdot (\text{SIZE} - 1)$

Large Items

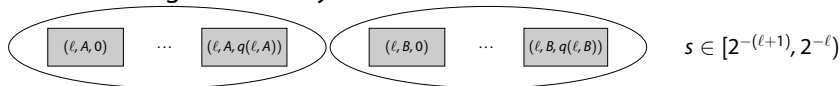
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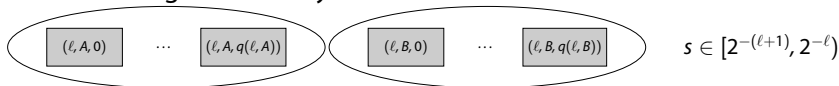
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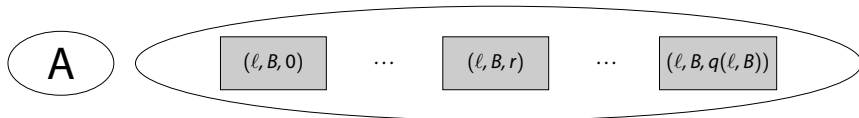
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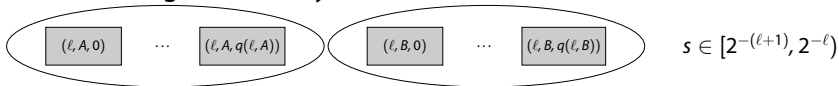


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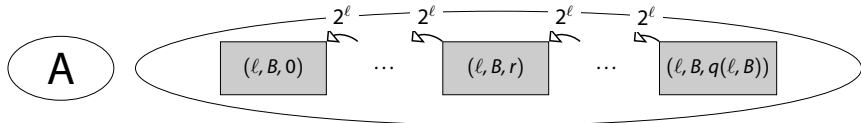


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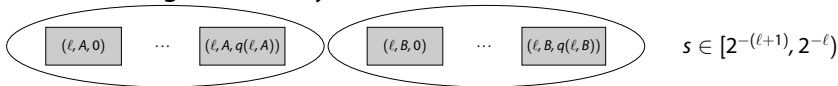


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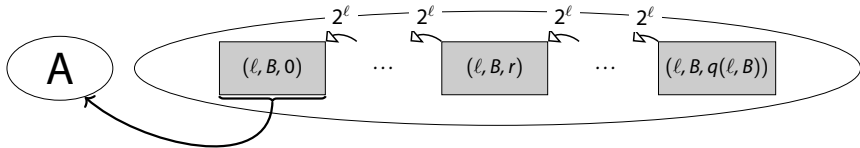


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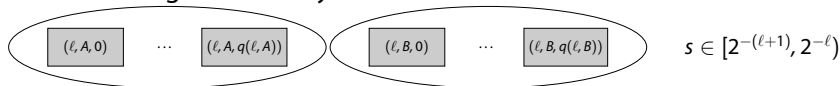


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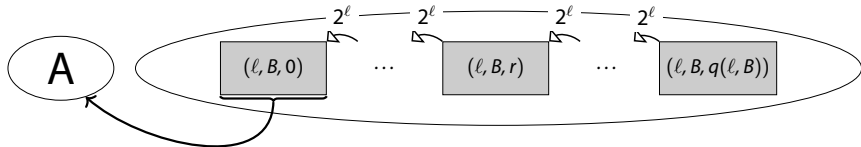


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Idea : Shifting Items to the left complies with the packing!

Small Items

Greedy fails: $\varepsilon \gg b \gg a$

Small Items

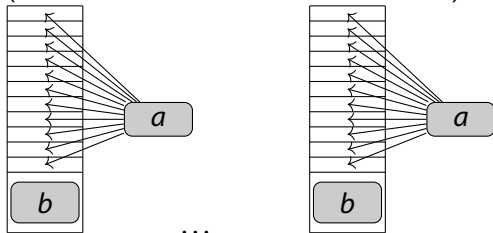
Greedy fails: $\varepsilon \gg b \gg a$

(INSERT: b , INSERT: a , . . . , INSERT: a)⁺

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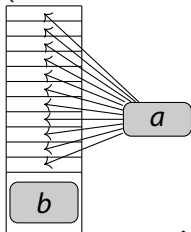
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Small Items

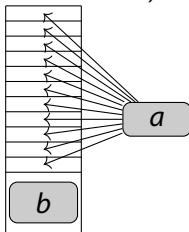
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...

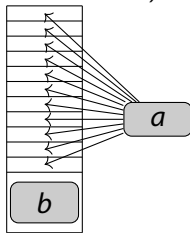
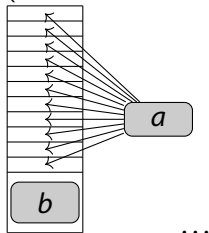
(REMOVE: a , REMOVE: a , ...)⁺



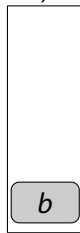
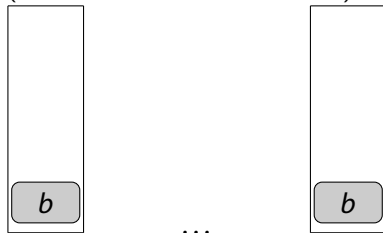
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Idea: "Sort" small items from left to right



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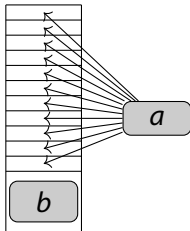
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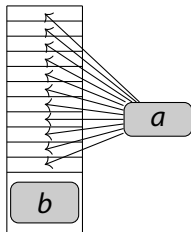
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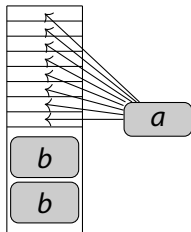
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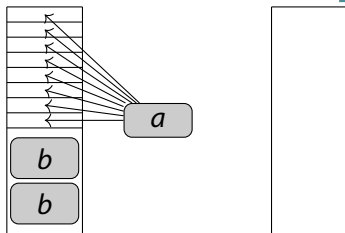
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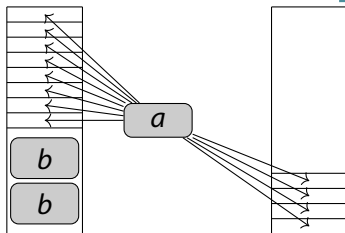
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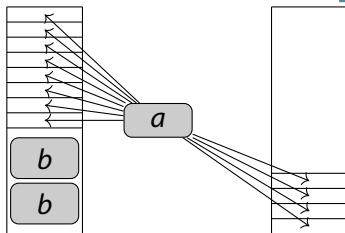
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Stop at every $1/\varepsilon$ -th bin (buffer bin) to bound MF

Everything

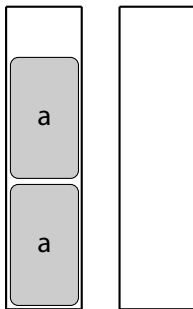


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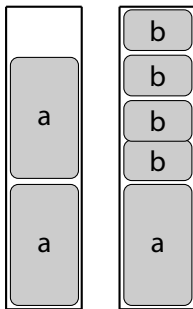
Large

Everything



Large

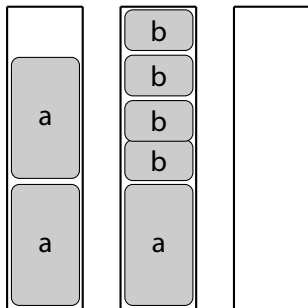
Everything



Large

Mixed

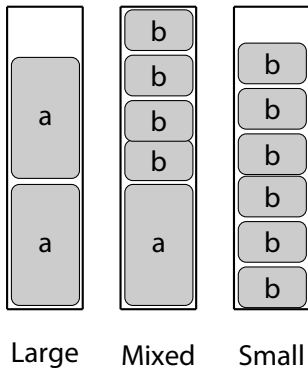
Everything



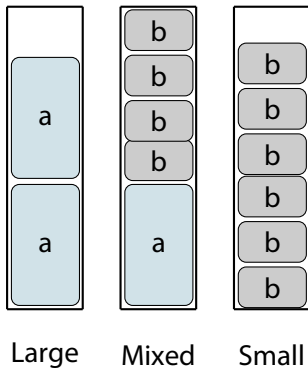
Large

Mixed

Everything

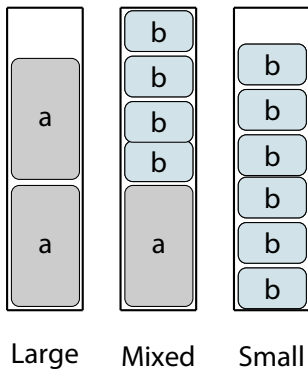


Everything



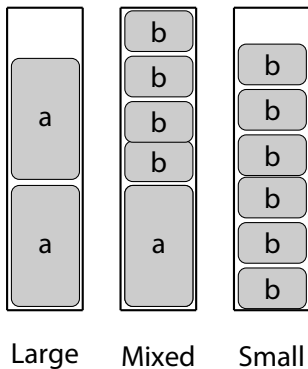
■ *Pack via LP*

Everything



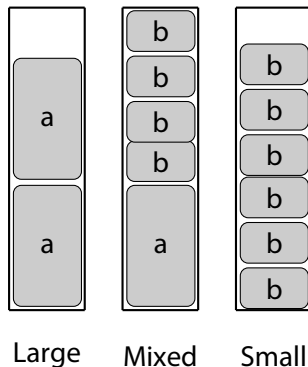
- Pack via LP
- *Pack via "Sorting"*

Everything



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- *Small bins \Rightarrow little free space in other bins*

Everything



- Pack via LP
- Pack via "Sorting"
- Small bins \Rightarrow little free space in other bins
- *Relate nearly full / nearly empty bins via potential function*

Lookout

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- *Simplify handling of small items*

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- Simplify handling of small items
- *Adapt techniques to other problems*