

# Applied Game Theory

Lecture 7

Sebastian Fest

# Repeated games I

- People and organizations often interact in ongoing relationships:
  - employment relationships
  - international trade
  - competitions and cooperations
- These relationships can be analyzed as repeated games.
- In a repeated game, *history of play* can affect future behavior.

## Repeated games II

- A repeated game is played over a discrete number of periods (period 1, period 2, etc.).
- We let  $t$  denote a given period, and  $T$  denote total number of periods.
- In each period  $t = 1, 2, \dots, T$  players play a *stage game*.

## Repeated games III

- The stage game is a collection  $\langle N, A, u \rangle$ , where  $N$  is the number of players,  $A = A_1 * A_2 * \dots * A_n$  is the set of action profiles and  $u_i(a)$  is player  $i$ 's stage game payoff when  $a$  is played.
- The stage game is played in all periods. For each period  $t$ , the players have observed the history of play, i.e. all actions from first period to period  $t - 1$ .
- The payoff for the whole game is the sum of the stage game payoffs from period 1 to  $T$ .
- When necessary, future payoffs are discounted with a discount factor  $\delta$ .

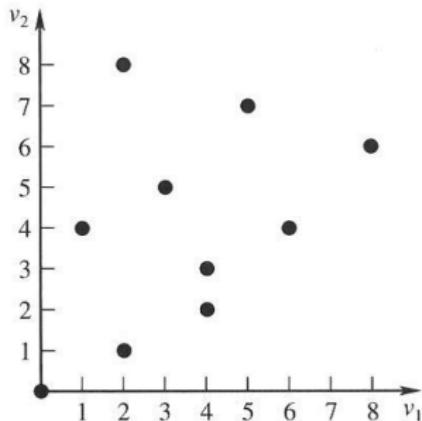
## Repeated games: Example I

An extensive form game tree where Player 1 moves first, choosing between action 1 and 2. Action 1 leads to a terminal node with payoffs (4, 3). Action 2 leads to a decision node for Player 2. At this node, Player 2 chooses between actions X, Y, and Z. Action X leads to payoffs (0, 0), action Y leads to (2, 1), and action Z leads to (1, 4).

	1	2	
	X	Y	Z
A	4, 3	0, 0	1, 4

	B	0, 0	2, 1	0, 0
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- Suppose the stage game is played twice ( $T = 2$ ).
- We have 6 different action profiles in  $t = 1$ , yielding 6 different information sets in  $t = 2$ .
- Since the stage game is also played in  $t = 2$ , we have 36 different strategy profiles and payoff vectors for the whole game.

## Repeated games: Example II

An extensive form game tree for Player 1. Player 1 chooses between 1 and 2. Choosing 1 leads to Player 2 choosing between X, Y, and Z. Choosing 2 leads to Player 2 choosing between A and B. The payoffs are listed as (Player 1 payoff, Player 2 payoff).

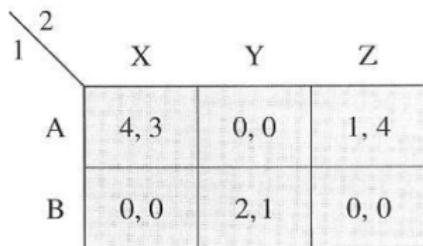
	1	2	
	X	Y	Z
A	4, 3	0, 0	1, 4
B	0, 0	2, 1	0, 0

An extensive form game tree for Player 2. Player 2 chooses between X, Y, and Z. Choosing X leads to Player 1 choosing between A and B. Choosing Y leads to Player 1 choosing between 1 and 2. Choosing Z leads to Player 1 choosing between 1 and 2. The payoffs are listed as (Player 1 payoff, Player 2 payoff).

	1	2	
	X	Y	Z
A	5, 7	1, 4	2, 8
B	1, 4	3, 5	1, 4

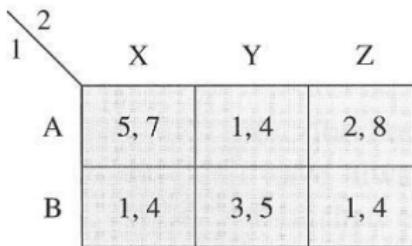
- Subgame perfection requires equilibrium play in every subgame.
- The matrix to the right shows the subgame following (A, Z) in the first stage, yielding the payoff (1, 4) in  $t = 1$ .

## Repeated games: Example III



An extensive form game tree for Player 1. Player 1 chooses between actions A and B. Action A leads to payoffs (4, 3) if Player 2 chooses X, (0, 0) if Player 2 chooses Y, and (1, 4) if Player 2 chooses Z. Action B leads to payoffs (0, 0) if Player 2 chooses X, (2, 1) if Player 2 chooses Y, and (0, 0) if Player 2 chooses Z.

	2
1	X Y Z
A	4, 3 0, 0 1, 4
B	0, 0 2, 1 0, 0



An extensive form game tree for Player 2. Player 2 chooses between actions X, Y, and Z. Action X leads to payoffs (5, 7) if Player 1 chose A, (1, 4) if Player 1 chose B. Action Y leads to payoffs (1, 4) if Player 1 chose A, (3, 5) if Player 1 chose B. Action Z leads to payoffs (2, 8) if Player 1 chose A, (1, 4) if Player 1 chose B.

	2
1	X Y Z
A	5, 7 1, 4 2, 8
B	1, 4 3, 5 1, 4

- A player's preference over actions in the subgame starting at  $t = 2$  are the same as for the stage game. Hence, the Nash equilibria for the subgame and the stage game are identical.
- Sequential rationality requires players to play a Nash equilibrium in the last period.
- We can then have a subgame perfect Nash equilibrium for the two period game if players also play a Nash equilibrium in the first period.

## Repeated games IV

- For any repeated game, any sequence of stage Nash profiles can be supported as the outcome of a subgame perfect Nash equilibrium.
- In addition, other equilibria might exist that don't stipulate the play of stage Nash profiles. Consider the following strategy profile:
  - Select (A,X) in the first period and then, as long as player 2 did not deviate from X in the first period, select (A,Z) in the second period
  - If player 2 deviated by playing Y or Z in the first period, then play (B,Y) in the second period.

## Repeated games V

The diagram shows an extensive form game tree. Player 1 moves first, choosing between X, Y, and Z. Player 2 moves second, choosing between A and B, given Player 1's choice. The payoffs are listed as (Player 1 payoff, Player 2 payoff). The payoffs for Player 1 are in bold.

		X	Y	Z	
		A	4, 3	0, 0	1, 4
		B	0, 0	2, 1	0, 0

- No-one has incentives to deviate from this strategy profile in any of the subgames, hence we have a subgame perfect Nash equilibrium.
- If player 2 deviates in  $t = 1$ , he is a "cheater" and lost reputation. This yields a lower payoff in  $t = 2$  (and a lower payoff overall).
- Hence, in the two period game, subgame perfection requires *stage Nash profiles* in period 2, but there may exist "reputational equilibria" where one does not play stage Nash profile in  $t = 1$ .

# Ininitely repeated games I

	2
1	C      D
C	2, 2      0, 3
D	3, 0      1, 1

- An infinitely repeated game has an infinite horizon, i.e.  $T = \infty$ , which is often useful to model long term relationships.
- A subgame perfect equilibrium can often be achieved through the play of simple strategies, e.g. *trigger strategies*:
  - Play (C,C) (cooperative profile) in each period. If someone deviates, play (D,D) (punishment profile) forever after.

## Ininitely repeated games II

	2
1	C      D
C	2, 2      0, 3
D	3, 0      1, 1

- The players will not deviate from this strategy if the present value from cooperation is greater than the present value from deviating:
- $2 + 2\delta + 2\delta^2 + \dots = \frac{2}{1-\delta} \geq 3 + \delta + \delta^2 + \dots = 3 + \frac{\delta}{1-\delta}$ , i.e.  
 $\delta \geq \frac{1}{2}$

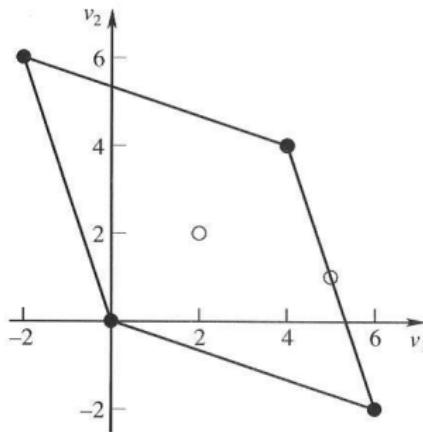
## Ininitely repeated games III

The diagram shows an extensive form game tree. Player 1 moves first, choosing between C and D. Choosing C leads to a terminal node with payoffs (2, 2). Choosing D leads to a decision node for Player 2. At this node, Player 2 can choose C or D. Choosing C leads to payoffs (4, 4). Choosing D leads to payoffs (-2, 6).

	2	
1	C	D
C	4, 4	-2, 6
D	6, -2	0, 0

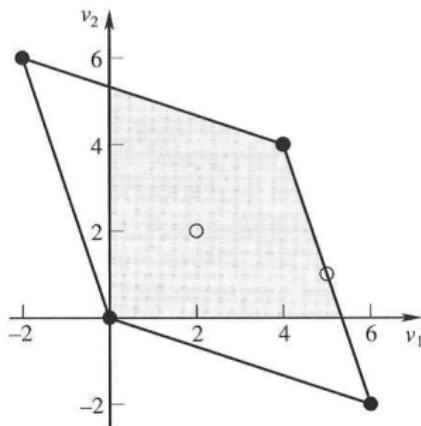
- In this game, the trigger strategy constitutes a subgame perfect equilibrium if  $\frac{4}{1-\delta} \geq 6$ , i.e.  $\delta \geq \frac{1}{3}$ .
- Note that this holds true for all periods, i.e. whether  $t = 1$  or  $t = 1000$  etc.
- The typical case in repeated games: Short term gain from cooperation vs. long term loss from deviation.
- The higher  $\delta$ , the larger is the long term loss, and the easier it is to sustain cooperation in equilibrium.

## Ininitely repeated games IV



- All payoff vectors within the diamond can be obtained as an average stage payoff.
- All payoff vectors within the diamond can be supported in a subgame perfect equilibrium as long as each player gets more than zero (payoff from (D,D)), and the discount factor is sufficiently high.

## Ininitely repeated games V



- The folk theorem: Consider any infinitely repeated game. Suppose there is a stage Nash profile that yields payoff vector  $w$  ( $w_i$  for player  $i$ ,  $i = 1, 2, \dots, n$ ). Let  $v$  be any feasible average per period payoff such that  $v_i > w_i$ . The vector  $v$  can be supported arbitrarily closely by a subgame perfect Nash equilibrium if  $\delta$  is sufficiently close to 1.

## Ininitely repeated games: Cournot

- Consider a Cournot game with  $p = 1 - q_1 - q_2$ . Production costs is 0, so that each player obtain payoff  $(1 - q_i - q_j)q_i$
- This yields Nash equilibrium  $q_1 = q_2 = \frac{1}{3}$  and profits  $\frac{1}{9}$  to each firm.
- Collusion on monopoly quantity where each produces  $\frac{1}{4}$  yields a payoff of  $\frac{1}{8}$  to each. But this cannot be sustained as an equilibrium in the stage game.

## Ininitely repeated games: Cournot

- In an infinitely repeated game, trigger strategies can constitute an equilibrium. Play  $\frac{1}{4}$  in each period as long as both played this in the past. If someone deviates, play stage Nash profil  $\frac{1}{3}$  for ever.
- By deviating firm  $i$  maximizes  $(1 - \frac{1}{4} - q_i)q_i$  which yields  $q_i = \frac{3}{8}$  and payoff  $\frac{9}{64}$  in the first deviation period. But thereafter it obtains  $\frac{1}{9}$  forever.
- Collusion can be sustained as a subgame perfect Nash equilibrium if  $\frac{1}{8(1-\delta)} \geq \frac{9}{64} + \frac{\delta}{9(1-\delta)}$ , i.e.  $\delta \geq \frac{9}{17}$ .

## Self-enforcing contracts

- Collusion is an example of a self-enforcing contract.
- Since collusion is illegal, and thus external enforcement impossible, the parties must rely on self-enforcement.
- Sometimes external enforcement is desirable, but not possible, e.g.:
  - employment contracts (hard to verify relevant variables).
  - international trade (no supranational enforcement).