

**SEMESTER 2
2022-2023**

**CS603
Rigorous Software Development**

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Time allowed: 3 hours

Answer at least **three** questions

Your mark will be based on your best **three** answers

All questions carry equal marks

Instructions

	Yes	No	N/A
Formulae and Tables book allowed (<i>i.e. available on request</i>)		X	
Formulae and Tables book required (<i>i.e. distributed prior to exam commencing</i>)		X	
Statistics Tables and Formulae allowed (<i>i.e. available on request</i>)		X	
Statistics Tables and Formulae required (<i>i.e. distributed prior to exam commencing</i>)		X	
Dictionary allowed (<i>supplied by the student</i>)		X	
Non-programmable calculator allowed		X	
Students required to write in and return the exam question paper		X	

- 1 (a) Explain how *Design by Contract* guides the verification of: **[25 marks]**
[8 marks]
- i. class methods
 - ii. object constructors, and
 - iii. overridden methods in a class inheritance.

Your explanation should include a coded example of each.

- (b) Use Hoare logic to verify the *partial correctness* of the following program with respect to the specification: [8 marks]

Precondition: $X \geq 0 \wedge Y \geq 0$
Postcondition: $z == X * Y$

Program: `int i = 0;
 int z = 0;
 while (i != Y)
 i = i + 1;
 z = z + X;`

Clearly identify the Hoare logic rules and simplification steps used at each stage of the verification.

- (c) Describe the difference between verifying the *partial* and *total correctness* of a loop. Clearly identify the extra conditions that must be specified and verified. Show how you would verify the total correctness of the program in (b) above. [4 marks]
- (d) Explain what is meant by the term *under-specification* in program verification. Support your solution by providing a correct specification for a program which sorts an array. Explain why your solution is not under-specified. [5 marks]

[25 marks]**2 (a)** With reference to the Dafny code provided below:**[8 marks]**

- i. Provide preconditions and variant functions for `Sum0` and `Sum1` ghost functions.
- ii. Provide a precondition, a variant function and a postcondition for the `Sum2` lemma below.
- iii. Explain the role of ghost code in program verification.
- iv. Explain the role of lemmas in program verification.

```

function F(x: int): int

ghost function Sum0(lo: int, hi: int): int{
  if lo == hi then 0 else F(lo) + Sum0(lo + 1, hi)
}

ghost function Sum1(lo: int, hi: int): int{
  if lo == hi then 0 else
    Sum1(lo, hi - 1) + F(hi - 1)
}

lemma Sum2(lo: int, hi: int){
  if lo != hi {
    PrependSumDown(lo, hi);
  }
}

lemma PrependSumDown(lo: int, hi: int)
  requires lo < hi
  ensures F(lo) + Sum0(lo + 1, hi) == Sum1(lo, hi)
  decreases hi - lo
{
}

```

- 2 (b) The following Dafny method returns the minimum value in an array, or 0 if the array is empty. [8 marks]

- i. Provide a variant and an invariant for the loop verification. Explain how they are used in the verification.
- ii. Describe and explain the expected verification error if the initialisation `min := a[0];` is changed to `min := 0;`
- iii. Describe and explain the expected verification error if the initialisation `min := 0;` is changed to `min := a[0];`

```
method min_elt(a:array<int>) returns(min:int)
ensures
  if (a.Length == 0) then min == 0
  else
    (forall i :: 0<=i<a.Length ==> a[i]>=min) &&
    (exists i :: 0 <=i<a.Length && min == a[i])
{
  if (a.Length == 0) { min := 0;}
  else {
    min := a[0];
    var i:int := 1;
    while (i < a.Length)
    {
      if(a[i] < min) {
        min := a[i];
      }
      i := i + 1;
    }
  }
}
```

- (c) Specify and implement a class `Queue` with the normal LIFO queue behaviour with operations to: [9 marks]

- i. create a queue object.
- ii. add an element to the queue.
- iii. remove an element from the queue.

Use your code to illustrate how verifiers (such as Dafny, OpenJML and Why3) check class invariants, constructor contracts and method contracts.

[25 marks]

- 3 (a) Explain how system modeling and analysis is supported by the *Event-B modelling system*. Your answer should include an explanation of *contexts*, *machines*, *invariants*, *events* and *before-after predicates*. Support your answer with an appropriate example listing at least two proof obligations that must be proved for the model to be correct. [8 marks]
- (b) Define what is meant by the term *Data Refinement*. Support your answer with an example. Explain how both *Dafny* and *Event-B* support data refinement. [6 marks]
- (c) How does the SMT solver Z3 determine if a model exists that satisfies the following set of equations? Explain what property is being checked below and what each line of code achieves. [5 marks]

```
(declare-fun a () Int)
(declare-fun b () Int)
(declare-fun c () Int)
(assert (> a 4))
(assert (< b 0))
(assert (> c 2))
(assert (= (+ (* a a) (* b b)) (* c c)))
(check-sat)
(get-model)
```

- (d) SAT solvers can be classified into those based on *Conflict-Driven Clause Learning (CDCL)* and those based on a *stochastic search*. Explain using an example how the CDCL algorithm operates. What advantage does CDCL have over stochastic search? [6 marks]

[25 marks]

- 4 (a) Explain how the *Spin model checker* verifies temporal properties of a system. Your answer should explain the representation of the system model, the representation of the formula to be checked, expressing invariant properties and using labels to verify properties of the model. [8 marks]

Provide examples to support your answer.

- (b) Specify the following properties in *Linear Temporal Logic (LTL)*. [10 marks]
Explain your answers.

- i. if the coffee level is low in state s_i , $i \leq 0$ then at state s_{i+1} an alarm will ring.
- ii. receiving a good salary and enjoying your job is possible.
- iii. if the door is locked at state s_i , then at some future state s_j , $j \geq i$ the door will unlock.
- iv. if the shop is open and the front of a queue is served at time s_i , then the shop will start serving the queue at some future time s_j , $j \geq i$, and continue serving until the end of the queue is reached at some state s_k , $k > j$.

- (c) Explain the difference between *Linear Temporal Logic (LTL)* and *Computation Tree Logic (CTL)*. Support your explanation by discussing their use when expressing: [7 marks]

- i. *safety properties*.
- ii. *liveness properties*.
- iii. *fairness properties*.