

Exact and approximate algorithms for resultant polytopes

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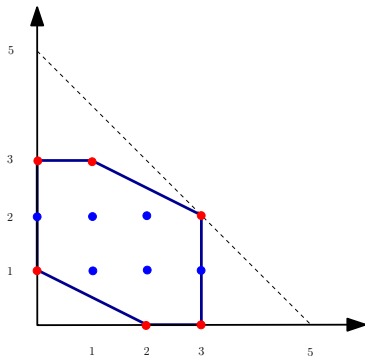


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Polynomials and Newton polytopes

- The **support** of a polynomial f is the set of exponents of its monomials with non-zero coefficient.
- The **Newton polytope** of f is the convex hull of its support.

$$f(x, y) = 8y + xy - 24y^2 - 16x^2 + 220x^2y - 34xy^2 - 84x^3y + 6x^2y^2 - 8xy^3 + 8x^3y^2 + 8x^3 + 18y^3$$



Polynomial systems

We study polynomials that expresses the solvability of polynomial systems.

Given a system of $n + 1$ **linear** polynomials f_0, f_1, \dots, f_n , on n variables the **determinant** is a polynomial on the coefficients which is zero iff the system has a common solution.

Polynomial systems

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$$f_0 = ax + by + c = 0$$

$$f_1 = dx + ey + f = 0$$

$$f_2 = gx + hy + i = 0$$

$$\begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix}$$

Polynomial systems

Given a system of $n + 1$ **linear** polynomials f_0, f_1, \dots, f_n , on n variables the **determinant** is a polynomial on the coefficients which is zero iff the system has a common solution.

$$f_0 = 4x + y + 2 = 0$$

$$f_1 = x + 2y + 1 = 0$$

$$f_2 = x + y + 8 = 0$$

$$\begin{vmatrix} 4 & 1 & 2 \\ 1 & 2 & 1 \\ 1 & 1 & 8 \end{vmatrix} = 51$$

Polynomial systems

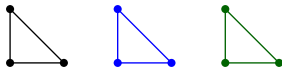
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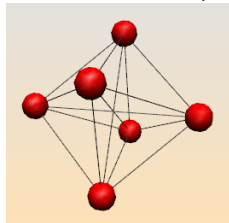
$$f_1 = dx + ey + f = 0$$

$$f_2 = gx + hy + i = 0$$

supports



Newton polytope of **determinant** (**Birkhoff** polytope)



$$\begin{vmatrix} a & b & c \\ d & e & f \\ g & h & i \end{vmatrix}$$

Polynomial systems

Given a system of $n + 1$ ~~linear~~ general polynomials f_0, f_1, \dots, f_n , on n variables the ~~determinant~~ resultant is a polynomial on the coefficients which is zero iff the system has a common solution.

$$f_0 = 4xy^2 + x^4y + 2 = 0$$

$$f_1 = x + 2y = 0$$

$$f_2 = 3x^2 + y + 8 = 0$$

Polynomial systems

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hard to compute
the resultant

Polynomial systems

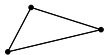
Given a system of $n + 1$ linear general polynomials f_0, f_1, \dots, f_n , on n variables the determinant resultant is a polynomial on the coefficients which is zero iff the system has a common solution.

$$f_0 = 4xy^2 + x^4y + 2 = 0$$

$$f_1 = x + 2y = 0$$

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supports



Newton polytope of determinant resultant
(Birkhoff resultant polytope II)

hard to compute
the resultant



Polynomial systems

Given a system of $n + 1$ **linear** **general** polynomials f_0, f_1, \dots, f_n , on n variables the **determinant** **resultant** is a polynomial on the coefficients which is zero iff the system has a common solution.

supports



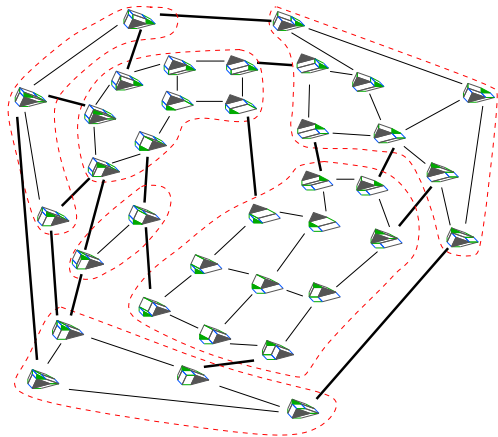
Newton polytope of **determinant** **resultant**
(**Birkhoff** **resultant** polytope Π)



The idea of the algorithm

The input supports define a pointset $\mathcal{A} \in \mathbb{Z}^{2n}$

Naive method: compute the secondary polytope $\Sigma(\mathcal{A})$ to compute Π

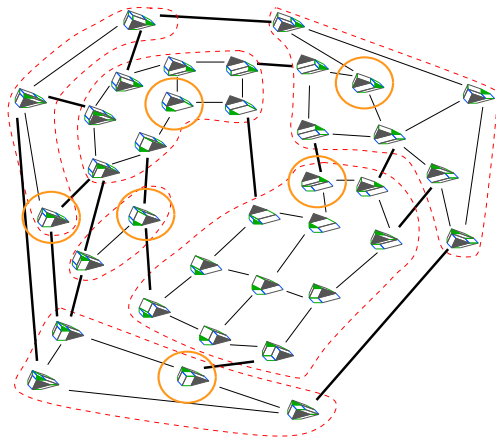


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Idea: incrementally construct Π using an **oracle** that given a direction produces vertices of Π

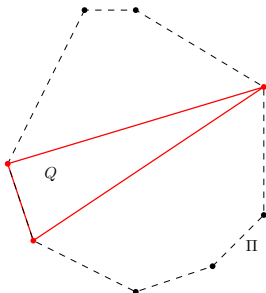


Incremental Algorithm

Input: \mathcal{A}

Output: H-rep. Q_H , V-rep. Q_V of $Q = \Pi$

1. initialization step



initialization:

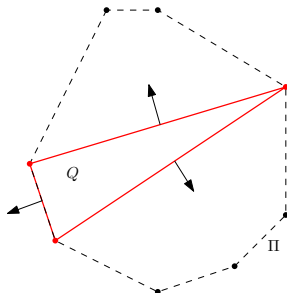
- ▶ $Q \subset \Pi$
- ▶ $\dim(Q) = \dim(\Pi)$

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2 kinds of hyperplanes of Q_H :

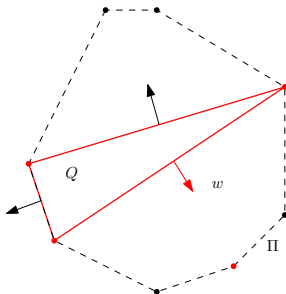
- **legal** if it supports facet $\subset \Pi$
- **illegal** otherwise

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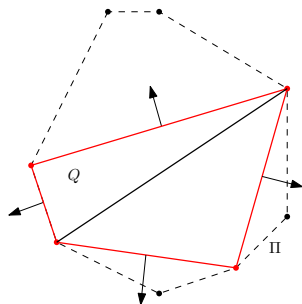
Extending an **illegal** facet

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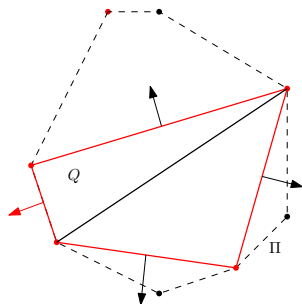
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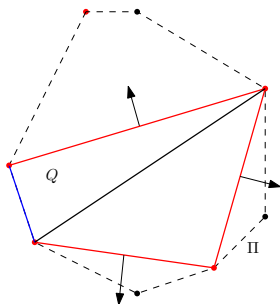
Validating a **legal** facet

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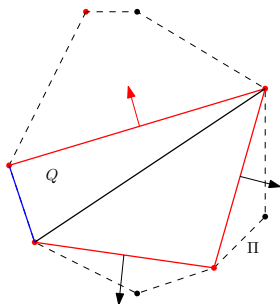
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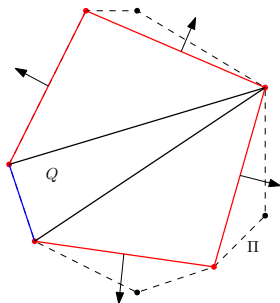


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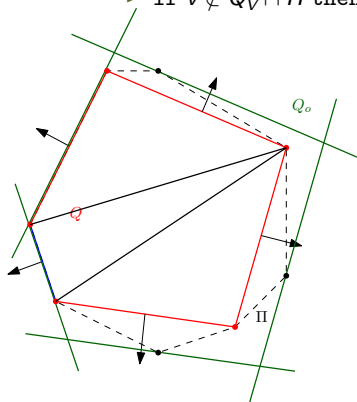
At any step, Q is an inner approximation ...

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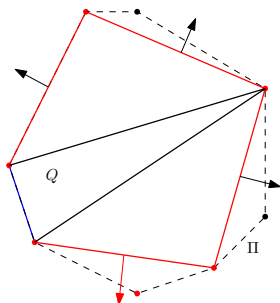
At any step, Q is an inner approximation ... from which we can compute an outer approximation Q_o .

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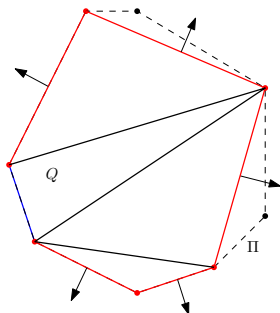


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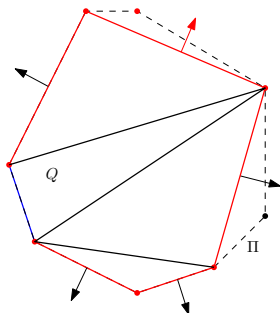


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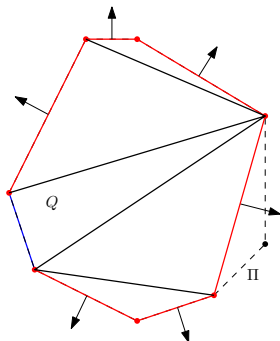


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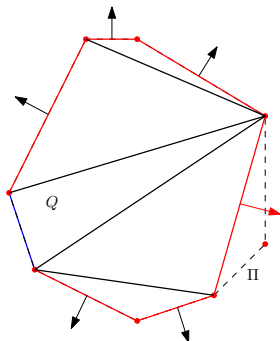


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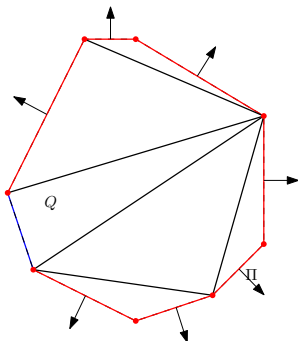


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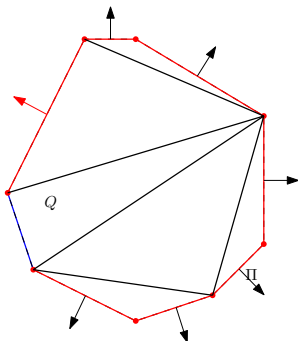


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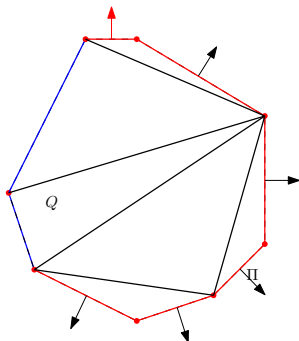


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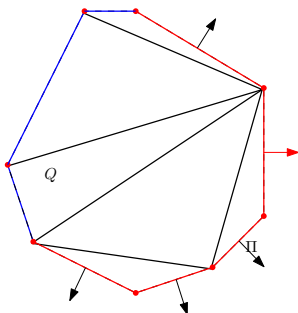


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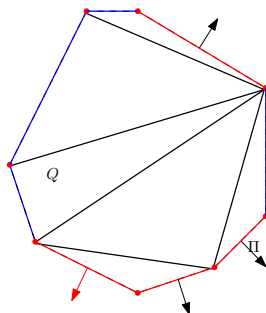


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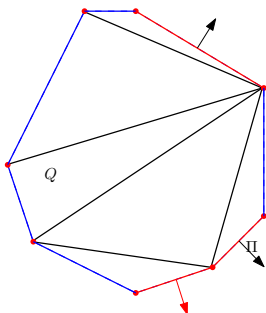


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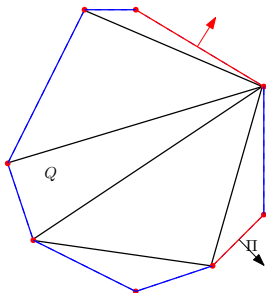


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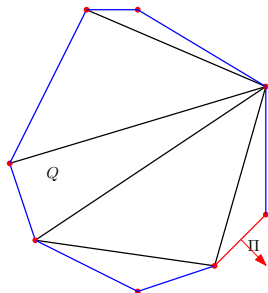


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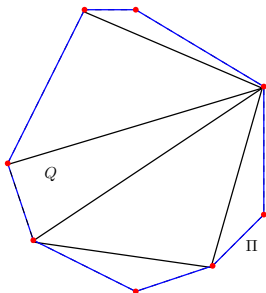


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Complexity

Theorem

We compute the Vertex- and Halfspace-representations of Π , as well as a triangulation T of Π , in

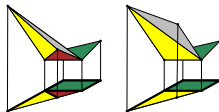
$$O^*(m^5 |\text{vtx}(\Pi)| \cdot |T|^2),$$

where $m = \dim \Pi$, and $|T|$ the number of full-dim faces of T .

Elements of proof

- ▶ Most computation is done in dimension $\leq m$.
- ▶ At most $\leq \text{vtx}(\Pi) + \text{fct}(\Pi)$ oracle calls
- ▶ Beneath-Beyond algorithm for converting V-rep. to H-rep. (bottleneck)

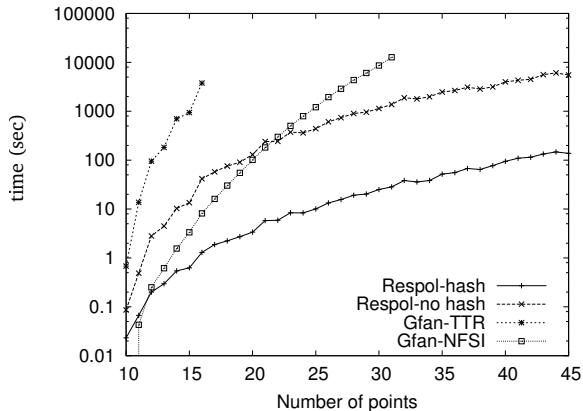
ResPol Implementation



Tools

C++, CGAL, triangulation [Boissonnat, Devillers, Hornus],
extreme_points_d [Gärtner]

Experiments ($\dim(\Pi) = 4$)



Future work

- ▶ approximate resultant polytopes ($\dim(\Pi) \geq 7$)
- ▶ preliminary results:

input	m $ \mathcal{A} $	3 200	3 490	4 20	4 30	5 17	5 20
exact	$\#\text{vtx}(\Pi)$	98	133	416	1296	1674	5093
	time	2.03	5.87	3.72	25.97	51.54	239.96
approx.	$\#\text{vtx}(Q_{in})$	15	11	63	121	—	—
	$\text{vol}(Q_{in})/\text{vol}(\Pi)$	0.96	0.95	0.93	0.94	—	—
	$\text{vol}(Q_{out})/\text{vol}(\Pi)$	1.02	1.03	1.04	1.03	—	—
	time	0.15	0.22	0.37	1.42	> 10hr	> 10hr

- ▶ approximate volume computation [Lovász-Vempala06]

References

The code

- ▶ <http://respol.sourceforge.net>

The full version of the paper

- ▶ <http://arxiv.org/abs/1108.5985v2>

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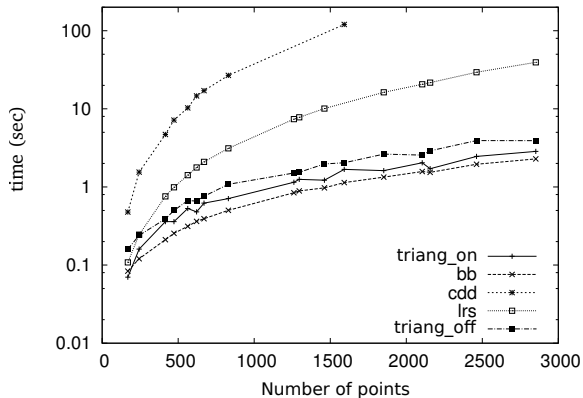
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Thank You !

Convex hull implementations

- From V- to H-rep. of Π .
- triangulation (on/off-line), polymake beneath-beyond, cdd, lrs



$$\dim(\Pi) = 4$$