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Algorithms by Sedgewick and Wayne (4th edition) [SW11]

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2.5: Applications

Exercise 1. Consider the following implementation of the complareTo() method for String. How doe the third line help with efficiency?

```
public int compareTo(String that)
{
   if (this == that) return 0; // this line
   int n = Math.min(this.length(), that.length());
   for (int i = 0; i < n; i++)
   {
      if (this.charAt(i) < that.charAt(i)) return -1;
      else if (this.charAt(i) > that.charAt(i)) return +1;
   }
   return this.length() - that.length();
}
```

Solution. In general, the method is linear in the length of the shortest of the two strings. However, it may be that the strings are aliased, so that effectively a string is being compared to itself. The indicated line detects this condition and reduces the duration of the compare to constant time.

Exercise 2. Write a program that reads a list of words from standard input and prints all two-word compound words in the list. For example, if after, thought, and afterthought are in the list, then afterthought is a compound word.

Solution. See com.segarciat.algs4.ch2.sec5.ex02.TwoWordCompoundWords.

Exercise 3. Criticize the following implementation of a class intended to represent account balances. Why is compareTo() a flawed implementation of the Comparable interface?

```
public class Balance implements Comparable<Balance>
{
    // ...
    private double amount;
    public int compareTo(Balance that)
    {
        if (this.amount < that.amount - 0.005) return -1;
        if (this.amount > that.amount + 0.005) return +1;
        return 0;
    }
    // ...
}
```

Describe a way to fix this problem.

Solution. It appears that the implementation is attempting to assert that the two Balance instances compare equal when their amount is within 0.005. For example, this would certify that 0.10 and 0.104 are the same, presumably both 10 cents. However, numbers of type double are known to be subject to rounding errors. Moreover, such an implementation does not define a total ordering. For example, suppose we had objects a, b, and c of type Balance, such that

- (i) a.amount = 0.097
- (ii) b.amount = 0.10
- (iii) c.amount = 0.103

Assuming no rounding errors, we would have a.compareTo(b) == 0 and b.compareTo(c) == 0, but a.compareTo(c) == -1, so that we don't have transitivity.

To fix this, we can choose a different representation for the amount. We can use two instance variables: one for the amounts smaller than 1 (for example, the number of cents if we are speaking of dollars), and another for the amounts that are 1 or larger (like dollars bills). Then the compareTo() method can exactly compare these quantities.

Exercise 4. Implement a method String[] dedup(String[] a) that returns the objects in a[] in sorted order, with duplicates removed.

Solution. See com.segarciat.algs4.ch2.sec5.ex04.DeduplicatedStrings.

Exercise 5. Explain why selection sort is not stable.

Solution. [SW11] describes a sorting method as *stable* if "it preserves the relative order of equal keys in the array". The reason this is so is because at any point, the "next minimum" that the algorithm searches for could be anywhere in the array. If the two elements equal elements are adjacent to one another, and the "next minimum" is somewhere to the right of them, then they could end up not in relative order.

Considered, for example, the following array:

[2] 2 3 4 1

, and exchanges the first 2 to get:

1 [2] 3 4 2

Notice that the relative order of the 2's changed. On the next iteration, the 2 in the second place (which has not been subject to a swap) stays in place, because no other key in the array is smaller than it:

1 2 [3] 4 2

Next, the next smallest is the 2 at the end, which is swapped with the 3 to get:

1 2 2 [4] 3

The elements to the left of the scan pointers are not moved anymore, so the 2's do not end up in the same relative order they started with.

Exercise 6. Implement a recursive version of select().

Solution. See com.segarciat.algs4.ch2.sec5.ex06.RecursiveSelect.

Exercise 7. About how many compares are required, on average, to find the smallest of n items using select()?

Solution. By Proposition U, the average number of compares to find the kth smallest is $\sim 2n + 2 \cdot k \ln(n/k) + 2(n-k) \cdot \ln(n/(n-k))$. As $k \to 0$, this quantity approaches 2n, suggesting the average.

Exercise 8. Write a program Frequency that reads strings from standard input and prints the number of times each string occurs, in descending order of frequency.

Solution. See com.segarciat.algs4.ch2.sec5.ex08.Frequency. I have implemented this by using a minimum-oriented priority queue with String objects read from standard input, and a max-oriented priority queue with StringCountNode objects, a data type I defined that simply holds a String read from standard input and its frequency.

Exercise 9. Develop a data type that allows you to write a client that can sort a file such as the one shown on below:

```
# input (DJIA volumes for each day)
1-Oct-28 3500000
2-Oct-28 3850000
3-Oct-28 4060000
4-Oct-28 4330000
5-Oct-28 4360000
. . .
30-Dec-99 554680000
31-Dec-99 374049984
3-Jan-00 931800000
4-Jan-00 1009000000
5-Jan-00 1085500032
# output
 19-Aug-40 130000
 26-AUg-40 160000
 24-Jul-40 200000
 10-Aug-42 210000
 23-Jun-42 210000
 23-Jul-02 2441019904
 17-Jul-02 2566500096
 15-Jul-02 2574799872
 19-Jul-02 2654099968
 24-Jul-02 2775555936
```

Solution. See com.segarciat.algs4.ch2.sec5.ex09.DJIAVolume.

Exercise 10. Create a data type Version that represents a software version number, such as 115.1.1, 115.10.1, 115.10.2. Implement the Comparable interface so that 115.1.1 is less than 115.10.1, and so forth.

Exercise 12. Scheduling. Write a program SPT. java that reads job names and processing times from standard input and prints a schedule that minimizes average completion time using the shortest processing time first rule, as described on page 349.

Solution. See com.segarciat.algs4.ch2.sec5.ex12.

Exercise 13. Load balancing. Write a program LPT. java that takes an integer m as a command-line argument, reads job names and processing times from standard input and prints a schedule assigning the jobs to m processors that approximately minimizes the time when the last job completes using the longest processing time first rule, as described on page 349.

Solution. See com.segarciat.algs4.ch2.sec5.ex13.

Exercise 14. Sort by reverse domain. Write a data type Domain that represents domain names, including an appropriate compareTo() where the natural order is in order of the reverse domain name. For example, the reverse domain name of cs.princeton.edu is edu.princeton.cs. This is useful for web log analysis. Hint: Use s.split("\\.") to split the string s into tokens, delimited by dots. Write a client that reads domain names from standard input and prints the reverse domains in sorted order.

Solution. See com.segarciat.algs4.ch2.sec5.ex14.

Exercise 15. Spam campaign. To initiate an illegal spam campaign, you have a list of email addresses from various domains (the part of the email address that follows the @ symbol). To better forge the return addresses, you want to send the email from another user at the same domain. For example, you might want to forge an email from wayne@princeton.edu to rs@princeton.edu. How would you process the email list to make this an efficient task?

Solution. I would use a priority queue to sort the list by domain. Then I would begin by taking address from the queue, which is the first forge senders. I would take other address off and as long as the the it is in the same domain as the current sender, I would set its sender as the first sender. I would keep track of the last sender removed at each step. As soon as the domain of the current domain differs from that of the last sender removed, I would arrange for the first sender to receive an email from the last sender removed. Then I would update the first sender to be the email just removed, the one from the new domain.

Exercise 16. Unbiased election. In order to thwart bias against candidates whose names appear toward the end of the alphabet, California sorted the candidates appearing on its 2003 gubernational ballot by using the following order of characters:

RWQOJMVAHBSGZXNTCIEKUPDYFL

Create a data type where this is the natural order and write a client California with a single static method main() that sorts strings according to this ordering. Assume that each string is composed solely of uppercase letters.

Exercise 17. Check stability. Extend your check() method from Exercise 2.1.16 to call sort() for a given array and return true if sort() sorts the array in a stable manner, false otherwise. Do not assume that sort() is restricted to move data only with exchange().

Solution. See com.segarciat.algs4.ch2.sec5.ex17.Stability.

Exercise 18. Force stability. Write a wrapper method that makes any sort stable by creating a new key type that allows you to append each key's index to the key, call sort(), then restore the original key after the sort.

Solution. See com.segarciat.algs4.ch2.sec5.ex18.ForceStability.

Exercise 19. Kendall tau distance. Write a program KendallTau.java that computes the Kendall tau distance between two permutations in linearithmic time.

Solution. From page 345, we know the following:

- A permutation (rankings) is an array of n integers where each of the integers between 0 and n-1 appears exactly once.
- The Kendall tau distance between two permutations is the number of pairs that are in different order in the two rankings.
- The number of inversions in an array is the Kendall tau distance between the array and the identity permutation.

Given permutations a and b, we can have the order of the elements in b define the "sort order". For example, if we have arrays:

```
a: 0 3 1 6 2 5 4
b: 1 0 3 6 4 2 5
```

Then, according to array b, the sort order is:

```
s: [1, 0, 5, 2, 4, 6, 3]
```

Notice that if we consider **b** as a mapping from the integer indices to its array values, then **s** is the inverse.

Now we can use mergesort to sort a according to the order define by b (according to s), counting inversions along the way.

Exercise 25. Points in the plane. Write three static comparators for the Point2 data type of page 77, one that compares by their x coordinate, one that compares by their y-coordinate, and one that compares by their distance from the origin. Write two non-static comparators for the Point2D data type, one that compares them by their distance to a specified point and one that compares them by their polar angle with respect to a specified point.

Solution. See com.segarciat.algs4.ch2.sec5.ex25.

Exercise 26. Simple polygon. Given N points in the plane, draw a simple polygon with the N points as vertices. Hint: Find the point p with the smallest y coordinate, breaking ties with the smallest x coordinate. Connect the points in increasing order of the polar angle they make with p.

Solution. See com.segarciat.algs4.ch2.sec5.ex26.DrawPolygon.

Exercise 27. One-dimensional intervals. Write three comparators for the Interval1D data type of page 77, one that compares intervals by their minimum endpoint, one that compares intervals by their maximum endpoint, and one that compares intervals by their length.

Solution. See com.segarciat.algs4.ch2.sec5.ex27.

Exercise 28. Sort files by name. Write a program FileSorter that takes the name of a directory as a command-line argument and prints out all of the files in that directory sorted by file name. Hint: Use the File data type.

Solution. See com.segarciat.algs4.ch2.sec5.ex28.FileSorter.

Exercise 29. Sort files by size and date of last modification. Write comparators for the type File to order by increasing/decreasing order of file size, ascending/descending order of file name, and ascending/descending order of last modification date. Use these comparators in a program LS that takes a command-line argument and lists the file in the current directory according to a specified order, e.g., "-t" to sort by timestamp. Support multiple flags to break ties. Be sure to use a stable sort.

Solution. See com.segarciat.algs4.ch2.sec5.ex29.LS.

References

[SW11] Robert Sedgewick and Kevin Wayne. *Algorithms*. 4th ed. Addison-Wesley, 2011. ISBN: 9780321573513.