

## 3.1: Symbol Tables

**Exercise 1.** Write a client that creates a symbol table mapping letter grades to numerical scores, as in the table below, then reads from standard input a list of letter grades and computes and prints the GPA (the average of the numbers corresponding to the grades).

A+	A	A-	B+	B	B-	C+	C	C-	D	F
4.33	4.00	3.67	3.33	3.00	2.67	2.33	2.00	1.67	1.00	0.00

**Solution.** See `com.segarciat.algs4.ch3.sec1.ex01.GPA`

**Exercise 2.** Develop a symbol-table implementation `ArrayST` that uses an (unordered) array as the underlying data structure to implement our basic symbol-table API.

**Solution.** See `com.segarciat.algs4.ch3.sec1.ex02.ArrayST`.

**Exercise 3.** Develop a symbol-table implementation of `OrderedSequentialSearchST` that uses an ordered linked list as the underlying data structure to implement our ordered symbol-table API.

**Solution.** See `com.segarciat.algs4.ch3.sec1.ex03.OrderedSequentialSearchST`.

**Exercise 4.** Develop `Time` and `Event` ADTs that allow processing of data as in the example illustrated on page 367.

**Solution.** See `com.segarciat.algs4.ch3.sec1.ex04.Time`. The class is immutable, and it implements `Comparable<Time>`, so that it has a natural order. I was unclear about what an `Event` ADT would include, so I did not provide an implementation for this ADT.

**Exercise 5.** Implement `size()`, `delete()`, and `keys()` for `SequentialSearchST`.

**Solution.** See `com.segarciat.algs4.ch3.sec1.ex05.SequentialSearchST`.

**Exercise 6.** Give the number of calls to `put()` and `get()` issued by `FrequencyCounter`, as a function of the number  $W$  of words and the number  $D$  of distinct words in the input.

**Solution.** The following assumes that the minimum length accepted for a word is 1.

During the first phase, the program builds the symbol tables by processing all  $W$  words. For each word, there is a call to `contains()`, for a total of  $W$  such calls. Since a call to `put()` is made regardless of the result, there are  $W$  calls to `put()` during this phase. Each result of `false` from the call to `contains()` corresponds to a distinct word, so there are  $D$  such outcomes. Thus, there are  $W - D$  direct calls to `get()` in the branch

of the `if-else`, where `get()` is used to retrieve the count of a previously-seen word. Note also that each call to `contains()` leads to a call to `get()`, accounting for  $W$  more calls.

In the second phase, one additional call to `put()` is made, which enters the empty string, so that there are now  $D + 1$  keys in the symbol table. In the loop, 2 calls to `get()` are made in each iteration, for a total of  $2(D + 1)$  calls. A final call to `get()` is made after the loop.

Thus, if  $f$  is the number of calls made to `put()`, and  $g$  is the number of calls made to `get()`, then

$$\begin{aligned} f(W, D) &= W + 1 \\ g(W, D) &= W + (W - D) + 2(D + 1) + 1 \\ &= 2W + D + 3 \end{aligned}$$

**Exercise 7.** What is the average number of distinct keys that `FrequencyCounter` will find among  $N$  random nonnegative integers less than 1,000, for  $N = 10, 10^2, 10^3, 10^4, 10^5$ , and  $10^6$ ?

**Solution.** Consider the random experiment of picking  $N$  integers at random, where each integer is between 1 and 1,000, and is chosen independently of the other. Then each outcome is an  $N$ -tuple, where each component is an integer between 1 and 1,000. Let  $X$  be a random variable that counts the number of distinct keys in an  $N$ -tuple. Then  $X$  is a discrete random variable, whose values range from 1 through  $\min\{N, 1000\}$ .

Consider the number of outcomes with  $X = k$  distinct integers. If  $k$  is not an integer, or  $k > \min\{N, 1000\}$ , or  $k \leq 0$ , then there are 0 such outcomes. Otherwise, we can count in two steps:

- (i) Choose  $k$  distinct integers. There are  $\binom{1000}{k}$  ways of doing this, for  $1 \leq k \leq 1000$ , and 0 for  $k > 1000$ .
- (ii) Of the  $k$  integers chosen, choose  $N - k$  integers to fill in the remaining spots. There are  $k^{N-k}$  ways to do this.
- (iii) Permute the  $N$  chosen integers, in  $N!$  ways.

By the multiplication principle of counting, we find that there are  $\binom{1000}{k} \cdot k^{N-k}$  ways to choose  $k$  distinct integers when choosing a total of  $N$  integers. There are a total of  $1000^N$  outcomes. Since every outcome is equally likely,

## References

- [SW11] Robert Sedgewick and Kevin Wayne. *Algorithms*. 4th ed. Addison-Wesley, 2011.  
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