CS410 - Automata Theory and Formal Languages

Project 2 Design Report

Introduction

This project aims to create a program that converts a context-free grammar to its equivalent Chomsky normal form. The program will get CFG as an input with the help of a pre-designed .txt file and it will print an equivalent CNF of it with the same design.

Problem Definition

Context Free Grammar can be defined by four tuples which are:

1-N is a set of non-terminal symbols.

2-T is a set of terminals.

3-P is a set of rules.

4-S is the start symbol.

To check if a CFG is in CNF all production rules must satisfy one of the following conditions.

1-Only the start symbol generates epsilon. (ex: if S is start S->(epsilon))

2-A non terminal generates 2 non terminals. (ex: A->BC)

3-A non terminal generates a terminal. (ex: A->b)

Since all CFG's have an equivalent CNF of it, it is possible to convert every CFG input to equivalent CNF of it.

Tools

JAVA : Programming language for implementation Eclipse: Integrated development environment for JAVA

Usage of The program

To use the java application CFG to CNF Converter follow these steps:

Compile the file:

javac Selami Karakas S018705.java

Run the program with giving your path to input file as an argument:

java Selami Karakas S018705 <Path to Input File>

Example: java Selami_Karakas_S018705 /home/selami/CS410/G1.txt

Note: Please be sure that your path does not include any whitespaces.

Method

To create the program following structures and functions needed:

- A method that reads input files and contains its tuples in arraylists.
- A method that checks whether there is any start terminal on the right hand side of the production rules.
- A method that returns a free char that is not used in CFG from the alphabet.
- A method that returns the index of a production rule if it contains epsilon.

- . A method that returns the index of a production rule if it is Unit Production.
- A method that returns the index of a production rule if it has a terminal with some terminals or non terminals on the right hand side.
- A method that returns the index of a production rule if it has more than two terminals on the right hand side.

Implementation Steps Of Conversion

In this section the conversion function is explained step by step with an example. The input example is at the end of this report.

Input Output Format

Here is the format mentioned in a few parts.

NON-TERMINAL

S

F

TERMINAL

0

1

RULES

S:00S

S:11F

F:00F

F:e

START

S