Problem Description

You are given an array of integers, and you need to perform three types of queries on it. The allowed operations are updating any particular element based on some mathematical formula, querying for the sum of the whole array, and finally querying the number of different values you can come up with by flipping some elements in your array.

Solutions There are several ways to solve this problem; we have to do some optimization to pass such big constraints on time complexity. Now, let's discuss the major approaches:

Brute Force Approach: For every update or query operation, you directly compute the result by iterating over the whole array. It would be easy to implement this but quite inefficient, as the upper limits of input size would then be as high as 10⁶. Every operation would take a time complexity of O(n); thus, this method is not really practical for large-sized inputs.

Using Hashmap and Modifications: We could avoid recalculating the entire array at every operation by using a dictionary to track the distinct absolute values in the array. For update operations, instead of going for re-computation of the overall sum or distinct values, we only need to update the target element and adjust the total sum and count of distinct values. This cuts down the time complexity drastically for sum and distinct queries to O(1) and update operations to O(log n) because of the optimized handling of the data structure.

Solution Explanation The code implements the described approach with efficiency. Here's how the solution works:

Initialization: The array is read and processed, and the initial count of distinct values, ignoring sign flips, is calculated. A dictionary, counts, maintains the number of occurrences of each absolute value in the array. Special handling is done for zero values since they contribute differently to the distinct count.

Update Operation: In the case of any update operation, the old value at the index is removed, and the new value is supposed to be calculated based on the formula provided. Then, the sum and count of distinct elements are updated accordingly, which avoids the need for recomputing these values until an actual build for the entire array is done.

Query Operations: For queries regarding the sum, the present total sum is printed out. Resulting counts for different queries, modulo possible sign flips, are printed.

Q2

You are given a permutation of integers from 1 to n along with a series of operations. In each operation, some continuous subarray is cut into two parts, and the leftmost and rightmost elements of these two parts are recorded. The task is to check if a given permutation along with the recorded sequences of leftmost li and rightmost ri elements can be validly obtained using these operations.

Possible Solutions

Several approaches can be explored for this problem, particularly under the constraints that n can be as large as 10⁶. Care should be taken with the key challenge: how the given sequences of leftmost and rightmost elements could be obtained by this sequence of operations:

- 1. Brute Force Simulation: One might simulate the whole process of splitting arrays and check if at some point they can produce the given sequences of leftmost and rightmost elements. But this is inefficient because, for large n, it would not be feasible to simulate every operation step by step.
- 2. Efficient Stack-Based or Greedy Approach: The problem can be optimized by reversal: instead of directly simulating the process of splits, we check if the sequence of leftmost and rightmost elements fits the structure of a valid split. We can use a stack to ensure that every split gives the correct li and ri to efficiently validate the permutation and the operations conducted.

Explanation of Solution Approach

The provided solution employs an efficient approach involving the use of stacks. Here's an explanation of the steps:

- Parsing the Input: The input consists of a permutation of size n with two sequences: I and r of length q, representing the leftmost and rightmost elements of the subarrays after each split, respectively. The task is to verify if a given permutation can be split validly based on the provided sequences I and r.
- 2. Stack-Based Validation: The main idea is to emulate the process of top-down array division, performing its reverse operation. We use a stack to ensure every split operation conforms with the given leftmost and rightmost sequences by iterating over the operations and tracking which subarrays result from each split. If, after all splits, the entire array is processed correctly and all leftmost and rightmost elements match the given sequences, then the permutation is valid.
- 3. Handling Large Input Efficiently: The solution handles the input size by performing push and pop operations in constant time. Hence, the overall algorithm runs in linear time O(n), suitable for input sizes as large as 10^6.

Why This Solution is Better

- Time Complexity: This solution runs in O(n), optimal for this problem, as it
 processes the input permutation and operations in a single pass. Any
 brute-force approach would simulate each split, which could easily involve
 O(n^2) complexity, making it infeasible for large n.
- 2. Space Complexity: The space complexity is O(n) because a stack is used. No other large data structures are employed, ensuring efficient memory usage.
- 3. Correctness: This solution works by emulating an array split in reverse and ensuring each operation aligns with the given leftmost and rightmost sequences. Thus, it directly checks whether the given input values of li and ri can be validly produced by the sequence of splits.

Conclusion

The solution presented for the list problem is efficient and suited to the input size constraints. It uses a stack-based approach to validate the sequence of operations in linear time, making it much better than brute-force simulation. By processing the input in one pass and checking the validity of the operations, this solution is

optimal.