

Load Factor

- ▶ Loading factor (LF), $\alpha = n / m$
n: number of keys
m: number of slots
- ▶ If uniform distribution ($1/m$) to get mapped to a slot, a slot will have an expectation of α elements.
- ▶ If m increases
 - ▶ Collision decreases
 - ▶ LF decreases
 - ▶ $0.5 > LF > 0.8$ is unacceptable
 - ▶ Storage requirements increases.
- ▶ Reduce collisions while keeping storage requirements low.

Linear Probing

$$h(k, i) = (h'(k) + i) \bmod m$$

- ▶ Always check the next index
- ▶ Increments index linearly with respect to i.
- ▶ Clustering problem

hash(10) = 2
hash(5) = 5
hash(15) = 7

0	72	72	72	72
1				15
2	18	18	18	18
3	43	43	43	43
4	36	36	36	36
5		10	10	10
6	6	6	6	6
7			5	5

Open Addressing – Quadratic Probing

$$h(k, i) = (h'(k) + c_1 i + c_2 i^2) \bmod m$$

- ▶ Instead of moving by one, move i^2

0			49	49	49
1					
2					
3					
4				58	58
5					69
6					
7					
8		18	18	18	18
9	89	89	89	89	89

$c_1=0, c_2=1$
hash(89)=9
hash(18)=8
hash(49)=9
hash(49, 1) = 0
hash(58) = 8
hash(58, 1) = 9
hash(58, 2) = 2
hash(69) = 9
hash(69, 1) = 0
hash(69, 2) = 3

Example 2 – Double Hashing

insert(76) 76%7 = 6	insert(93) 93%7 = 2	insert(40) 40%7 = 5	insert(47) 47%7 = 5 5 - (47%5) = 3	insert(10) 10%7 = 3	insert(55) 55%7 = 6 5 - (55%5) = 5
0 1 2 3 4 5 6	0 1 2 3 4 5 6	0 1 2 3 4 5 6	0 1 2 3 4 5 6	0 1 2 3 4 5 6	0 1 2 3 4 5 6
			47	47	47
	93	93	93	93	93
				10	10
					55
		40	40	40	40
76	76	76	76	76	76
probes: 1	1	1	2	1	2

Closed Hashing >>> <https://www.cs.usfca.edu/~galles/visualization/ClosedHash.html>

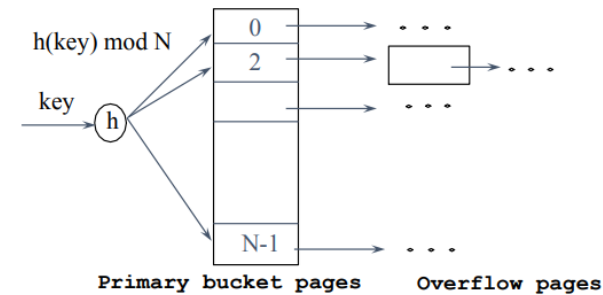
Open Hashing >>> <https://www.cs.usfca.edu/~galles/visualization/OpenHash.html>

Dynamic Hashing Methods

- ▶ As for any index, 2 alternatives for data entries k^* :
 - $\langle k, \text{rid of data record with search key value } k \rangle$
 - $\langle k, \text{list of rids of data records with search key } k \rangle$
- ▶ Choice orthogonal to the indexing technique
- ▶ Hash-based indexes are best for equality selections. **Cannot** support range searches.

Static Hashing

- ▶ # primary pages fixed, allocated sequentially, never de-allocated; overflow pages if needed.
- ▶ $h(k) \bmod M = \text{bucket to which data entry with key } k \text{ belongs. (} M = \# \text{ of buckets)}$



Extendible Hashing >>> <https://www.youtube.com/watch?v=TtkN2xRAgv4&t=519s>

Extendible Hashing Example

