

# Advanced Counting Techniques

## Chapter 8

With Question/Answer Animations

# Chapter Summary

- Applications of Recurrence Relations
- Solving Linear Recurrence Relations
  - Homogeneous Recurrence Relations
  - Nonhomogeneous Recurrence Relations
- Divide-and-Conquer Algorithms and Recurrence Relations
- Generating Functions
- Inclusion-Exclusion
- Applications of Inclusion-Exclusion

# Applications of Recurrence Relations

Section 8.1

# Section Summary

- Applications of Recurrence Relations
  - Fibonacci Numbers
  - The Tower of Hanoi
  - Counting Problems
- Algorithms and Recurrence Relations (*not currently included in overheads*)

# Recurrence Relations

(recalling definitions from Chapter 2)

**Definition:** A *recurrence relation* for the sequence  $\{a_n\}$  is an equation that expresses  $a_n$  in terms of one or more of the previous terms of the sequence, namely,  $a_0, a_1, \dots, a_{n-1}$ , for all integers  $n$  with  $n \geq n_0$ , where  $n_0$  is a nonnegative integer.

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- A sequence is called a *solution* of a recurrence relation if its terms satisfy the recurrence relation.
- The *initial conditions* for a sequence specify *the terms that precede the first term where the recurrence relation takes effect*.

# Rabbits and the Fibonacci Numbers

**Example:** A young **pair of rabbits** (one of each gender) is placed on an island:












- *A pair of rabbits does not breed until they are 2 months old.*
- After they are 2 months old, *each pair of rabbits produces another pair each month.*

Find a recurrence relation for *the number of pairs of rabbits on the island after  $n$  months*, assuming that rabbits never die.

*This is the original problem considered by Leonardo Pisano (Fibonacci) in the thirteenth century.*



# Rabbits and the Fibonacci Numbers (*cont.*)

Reproducing pairs (at least two months old)	Young pairs (less than two months old)	Month	Reproducing pairs	Young pairs	Total pairs
		1	0	1	1
		2	0	1	1
		3	1	1	2
		4	1	2	3
		5	2	3	5
	 	6	3	5	8

**Modeling the Population Growth of Rabbits on an Island**

# Rabbits and the Fibonacci Numbers (*cont.*)

**Solution:** Let  $f_n$  be the *the number of pairs of rabbits after  $n$  months*.

- There are  $f_1 = 1$  pairs of rabbits on the island *at the end of the first month*.
- We also have  $f_2 = 1$  because *the pair does not also breed during the second month*.
- To find the number of pairs on the island after  $n$  months, *add the number on the island after the previous month,  $f_{n-1}$ , and the number of newborn pairs*, which equals  $f_{n-2}$ , *because each newborn pair comes from a pair at least two months old*.

# Rabbits and the Fibonacci Numbers (*cont.*)

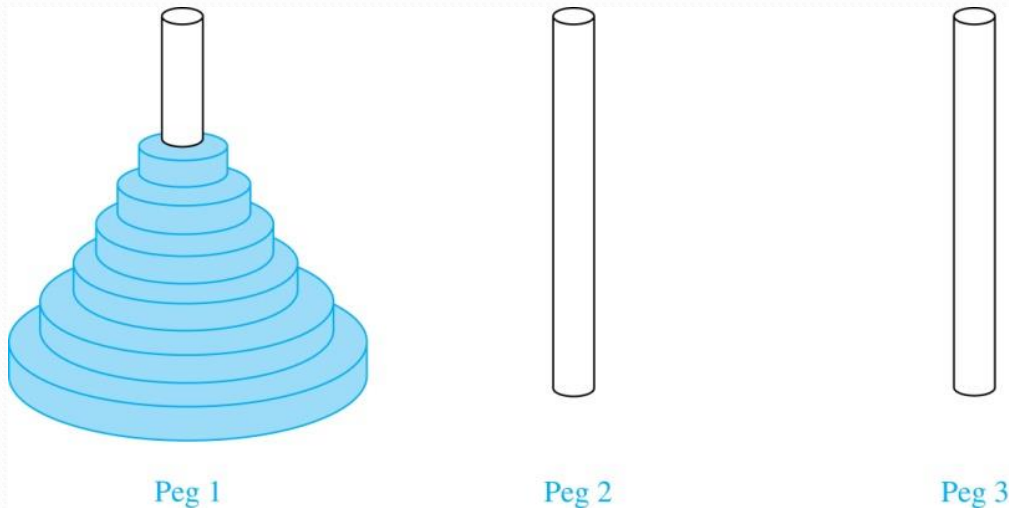
**Solution:** Let  $f_n$  be the the number of pairs of rabbits after  $n$  months.

- There are is  $f_1 = 1$  pairs of rabbits on the island at the end of the first month.
- We also have  $f_2 = 1$  because the pair does not breed during the first month.
- To find the number of pairs on the island after  $n$  months, add the number on the island after the previous month,  $f_{n-1}$ , and the number of newborn pairs, which equals  $f_{n-2}$ , because each newborn pair comes from a pair at least two months old.

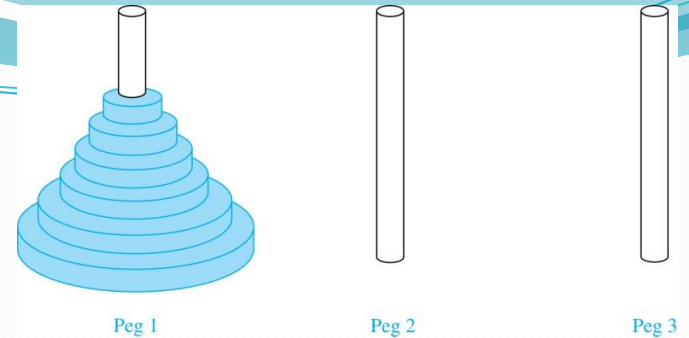
**Consequently** the sequence  $\{f_n\}$  satisfies the recurrence relation  $f_n = f_{n-1} + f_{n-2}$  for  $n \geq 3$  with the initial conditions  $f_1 = 1$  and  $f_2 = 1$ .  
The number of pairs of rabbits on the island after  $n$  months is given by the  *$n$ th Fibonacci number*.

# The Tower of Hanoi

- Three pegs on a board with **disks of different sizes**.
- Initially all of the disks are on the first peg *in order of size*, with **the largest on the bottom**.



# The Tower of Hanoi



## Rules:

You are allowed to move the disks **one at a time** from one peg to another as long as **a larger disk is never placed on a smaller.**

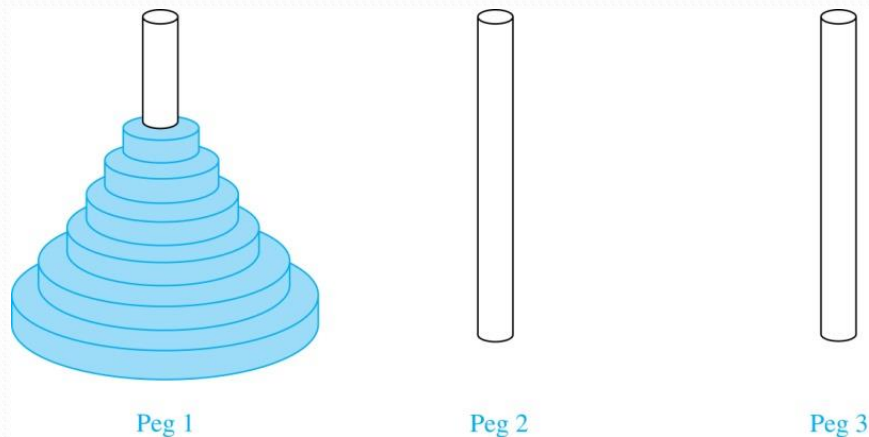
## Goal:

Using allowable moves, end up with all the disks **on the second peg** in order of size with **largest on the bottom.**

# The Tower of Hanoi (*continued*)

$\{H_n\}$ : the number of moves needed to solve the Tower of Hanoi Puzzle with  $n$  disks.

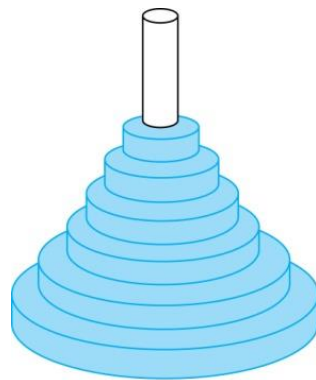
**Solution:** Set up a recurrence relation for the sequence  $\{H_n\}$ .



**The Initial Position in the Tower of Hanoi Puzzle**

# The Tower of Hanoi (*continued*)

- Begin with  $n$  disks on peg 1.



Peg 1



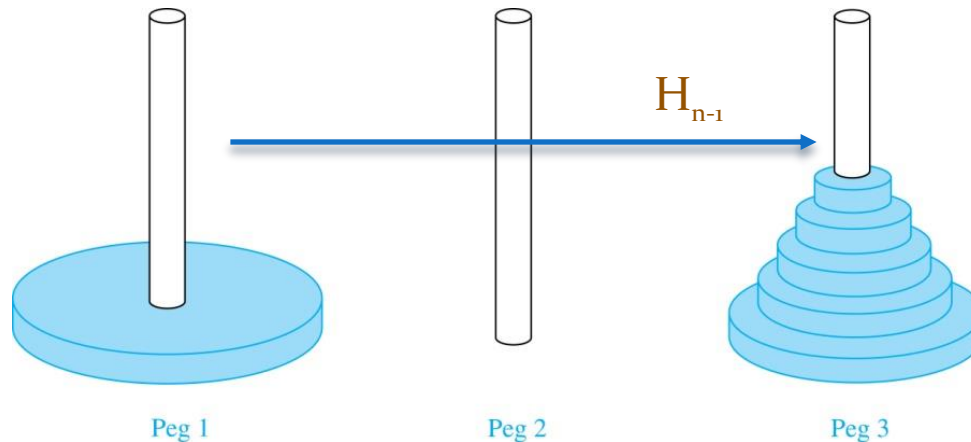
Peg 2



Peg 3

# The Tower of Hanoi (*continued*)

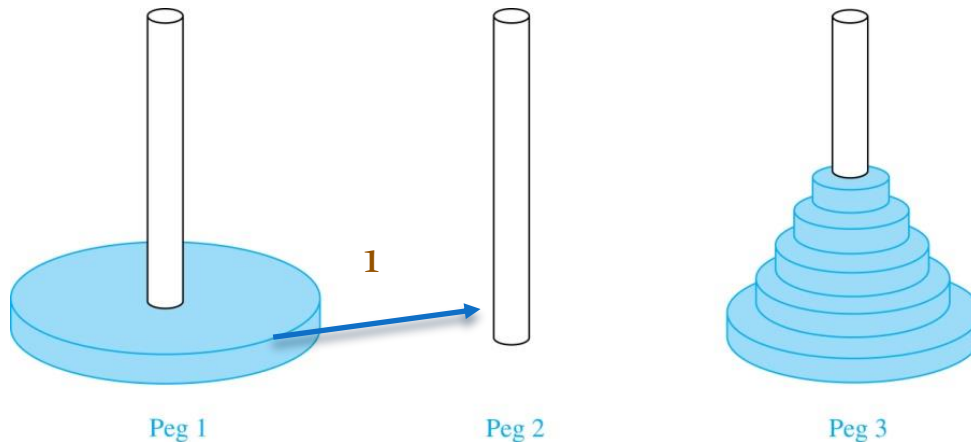
- Begin with  $n$  disks on peg 1.
- We can transfer the **top  $n - 1$  disks**, following the rules of the puzzle, to **peg 3** using  $H_{n-1}$  moves.





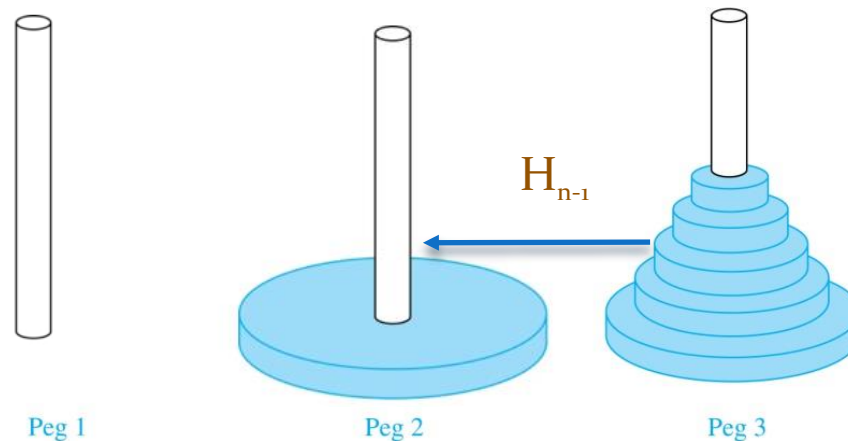
# The Tower of Hanoi (*continued*)

- First, we use **1 move** to transfer the largest disk to the second peg.



# The Tower of Hanoi (*continued*)

- Then we transfer the  $n - 1$  disks from peg 3 to peg 2 using  $H_{n-1}$  additional moves.

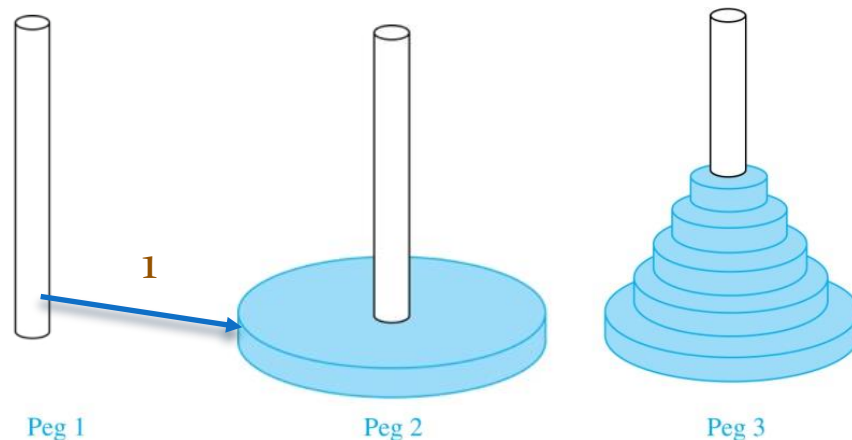


# The Tower of Hanoi (*continued*)

- Then we transfer the  $n - 1$  disks from peg 3 to peg 2 using  $H_{n-1}$  additional moves.

$$H_n = 2H_{n-1} + 1.$$

- The initial condition is  $H_1 = 1$ 
  - a single disk can be transferred from peg 1 to peg 2 in one move.



# The Tower of Hanoi (*continued*)

- Iterative approach to solve this recurrence relation
  - by repeatedly expressing  $H_n$  in terms of the previous terms of the sequence.

$$\begin{aligned}H_n &= 2H_{n-1} + 1 \\&= 2(2H_{n-2} + 1) + 1 = 2^2 H_{n-2} + 2 + 1 \\&= 2^2(2H_{n-3} + 1) + 2 + 1 = 2^3 H_{n-3} + 2^2 + 2 + 1 \\&\vdots \\&= 2^{n-1}H_1 + 2^{n-2} + 2^{n-3} + \dots + 2 + 1 \\&= 2^{n-1} + 2^{n-2} + 2^{n-3} + \dots + 2 + 1 \quad \text{because } H_1 = 1 \\&= 2^n - 1 \quad \text{using the formula for the sum of the terms of a geometric series}\end{aligned}$$

# Counting Bit Strings

- Find the number of bit strings of length  $n$  without two consecutive 0s.
  - Find a recurrence relation and give initial conditions
  - How many such bit strings are there of length five?

# Counting Bit Strings

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- $a_n$  denote the number of bit strings of length  $n$  without two consecutive 0s.
- The number of bit strings of length  $n$  that do not have two consecutive 0s:
  - the number of bit strings ending with a 0
  - the number of such bit strings ending with a 1

# Counting Bit Strings

Number of bit strings of length  $n$  without two consecutive 0s.

- $a_n$  denote the number of bit strings of length  $n$  without two consecutive 0s.
- Now assume that  $n \geq 3$ .
  - The bit strings of length  $n$  ending with 1 without two consecutive 0s are the bit strings of length  $n-1$  with no two consecutive 0s with a 1 at the end.
    - Hence, there are  $a_{n-1}$  such bit strings.
  - The bit strings of length  $n$  ending with 0 without two consecutive 0s are the bit strings of length  $n-2$  with no two consecutive 0s with 10 at the end.
    - Hence, there are  $a_{n-2}$  such bit strings.



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  - Hence, there are  $a_{n-1}$  such bit strings.
- The bit strings of length  $n$  ending with 0 without two consecutive 0s are the bit strings of length  $n - 2$  with no two consecutive 0s with 10 at the end.
  - Hence, there are  $a_{n-2}$  such bit strings.

$$a_n = a_{n-1} + a_{n-2} \text{ for } n \geq 3.$$

		Number of bit strings of length $n$ with no two consecutive 0s:	
End with a 1:	Any bit string of length $n - 1$ with no two consecutive 0s	1	$a_{n-1}$
End with a 0:	Any bit string of length $n - 2$ with no two consecutive 0s	1 0	$a_{n-2}$
		Total:	$a_n = a_{n-1} + a_{n-2}$

# Bit Strings (*continued*)

$$a_n = a_{n-1} + a_{n-2} \text{ for } n \geq 3.$$

The initial conditions are:

- $a_1 = 2$ , since both the bit strings **0** and **1** do not have consecutive 0s.
- $a_2 = 3$ , since the bit strings **01**, **10**, and **11** do not have consecutive 0s, while 00 does.

To obtain  $a_5$ , we use the recurrence relation three times to find that:

- $a_3 = a_2 + a_1 = 3 + 2 = 5$
- $a_4 = a_3 + a_2 = 5 + 3 = 8$
- $a_5 = a_4 + a_3 = 8 + 5 = 13$

Note that  $\{a_n\}$  satisfies the same recurrence relation as the Fibonacci sequence. Since  $a_1 = 2 = f_3$  and  $a_2 = 3 = f_4$ , we conclude that  $a_n = f_{n+2}$ .

# Counting the Ways to Parenthesize a Product

- Find a recurrence relation for  $C_n$ :
  - The number of ways to parenthesize the product of  $n + 1$  numbers,  $x_0 \cdot x_1 \cdot x_2 \cdot \cdots \cdot x_n$ , to specify the order of multiplication.
  - For example,  $C_3 = 5$ , since all the possible ways to parenthesize 4 numbers are
$$\begin{aligned} & ((x_0 \cdot x_1) \cdot x_2) \cdot x_3, \\ & (x_0 \cdot (x_1 \cdot x_2)) \cdot x_3, \\ & (x_0 \cdot x_1) \cdot (x_2 \cdot x_3), \\ & x_0 \cdot ((x_1 \cdot x_2) \cdot x_3), \\ & x_0 \cdot (x_1 \cdot (x_2 \cdot x_3)) \end{aligned}$$

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**Solution:** Note that however parentheses are inserted in  $x_0 \cdot x_1 \cdot x_2 \cdot \dots \cdot x_n$ ,

- one “ $\cdot$ ” operator remains outside all parentheses.
- This final operator appears between two of the  $n + 1$  numbers, say  $x_k$  and  $x_{k+1}$ .
- Since there are  $C_k$  ways to insert parentheses in the product  $x_0 \cdot x_1 \cdot x_2 \cdot \dots \cdot x_k$  and
- $C_{n-k-1}$  ways to insert parentheses in the product  $x_{k+1} \cdot x_{k+2} \cdot \dots \cdot x_n$ , we have

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- $C_{n-k-1}$  ways to insert parentheses in the product  $x_{k+1} \cdot x_{k+2} \cdot \dots \cdot x_n$ , we have

$$C_n = C_0 C_{n-1} + C_1 C_{n-2} + \dots + C_{n-2} C_1 + C_{n-1} C_0$$

$$= \sum_{k=0}^{n-1} C_k C_{n-k-1}$$

The initial conditions are  $C_0 = 1$  and  $C_1 = 1$ .

# Counting the Ways to Parenthesize a Product

$$\begin{aligned} & ((x_0 \cdot x_1) \cdot x_2) \cdot x_3, \\ & (x_0 \cdot (x_1 \cdot x_2)) \cdot x_3, \\ & (x_0 \cdot x_1) \cdot (x_2 \cdot x_3), \\ & x_0 \cdot ((x_1 \cdot x_2) \cdot x_3), \\ & x_0 \cdot (x_1 \cdot (x_2 \cdot x_3)) \end{aligned}$$

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- $C_{n-k-1}$  ways to insert parentheses in the product  $x_{k+1} \cdot x_{k+2} \cdot \cdots \cdot x_n$ , we have

$$C_n = C_0 C_{n-1} + C_1 C_{n-2} + \cdots + C_{n-2} C_1 + C_{n-1} C_0$$

$$= \sum_{k=0}^{n-1} C_k C_{n-k-1}$$

The sequence  $\{C_n\}$  is the sequence of **Catalan Numbers**.

# Solving Linear Recurrence Relations

Section 8.2

# Section Summary

- Linear Homogeneous Recurrence Relations
- Solving Linear Homogeneous Recurrence Relations with Constant Coefficients.
- Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients.



# Linear Homogeneous Recurrence Relations

**Definition:** A *linear homogeneous recurrence relation of degree  $k$  with constant coefficients* is a recurrence relation of the form:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k},$$

where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$

- *linear* : the right-hand side is a sum of the previous terms of the sequence each multiplied by a function of  $n$ .
- *homogeneous* because no terms occur that are not multiples of the  $a_j$ 's. Each coefficient is a constant.
- *degree is  $k$*  because  $a_n$  is expressed in terms of the previous  $k$  terms of the sequence.

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where  $c_1, c_2, \dots, c_k$  are real numbers, and  $c_k \neq 0$

By strong induction, a sequence satisfying such a recurrence relation is uniquely determined by the recurrence relation and the  *$k$  initial conditions*

$$a_0 = C_0, a_1 = C_1, \dots, a_{k-1} = C_{k-1}.$$

# Examples of Linear Homogeneous Recurrence Relations

- $P_n = (1.11)P_{n-1}$  linear homogeneous recurrence relation of degree one
- $f_n = f_{n-1} + f_{n-2}$  linear homogeneous recurrence relation of degree two
- $a_n = a_{n-1} + a_{n-2}^2$  not linear
- $H_n = 2H_{n-1} + 1$  not homogeneous
- $B_n = nB_{n-1}$  coefficients are not constants

# Solving Linear Homogeneous Recurrence Relations

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- Algebraic manipulation yields the *characteristic equation*:
$$r^k - c_1 r^{k-1} - c_2 r^{k-2} - \cdots - c_{k-1} r - c_k = 0$$

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- The solutions to the characteristic equation are called the *characteristic roots* of the recurrence relation.
- The roots are used to give an explicit formula for all the solutions of the recurrence relation.



# Solving Linear Homogeneous Recurrence Relations of Degree Two

**Theorem 1:** Let  $c_1$  and  $c_2$  be real numbers.

Suppose that  $r^2 - c_1r - c_2 = 0$  has two distinct roots  $r_1$  and  $r_2$ .

Then the sequence  $\{a_n\}$  is a solution to the recurrence relation  $a_n = c_1a_{n-1} + c_2a_{n-2}$  if and only if

for  $n = 0, 1, 2, \dots$ , where  $a_n = \alpha_1 r_1^n + \alpha_2 r_2^n$  and  $\alpha_1$  and  $\alpha_2$  are constants.

# Using Theorem 1

**Example:** What is the solution to the recurrence relation

$$a_n = a_{n-1} + 2a_{n-2} \text{ with } a_0 = 2 \text{ and } a_1 = 7?$$

**Solution:**

- The characteristic equation is  $r^2 - r - 2 = 0$ .
- Its roots are  $r = 2$  and  $r = -1$ .
- $\{a_n\}$  is a solution to the recurrence relation if and only if
  - $a_n = \alpha_1 2^n + \alpha_2 (-1)^n$ , for some constants  $\alpha_1$  and  $\alpha_2$ .

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- $\{a_n\}$  is a solution to the recurrence relation if and only if
  - $a_n = \alpha_1 2^n + \alpha_2 (-1)^n$ , for some constants  $\alpha_1$  and  $\alpha_2$ .

To find the constants  $\alpha_1$  and  $\alpha_2$ :

$$a_0 = 2 = \alpha_1 + \alpha_2$$

$$a_1 = 7 = \alpha_1 2 + \alpha_2 (-1).$$

Solving these equations, we find that  $\alpha_1 = 3$  and  $\alpha_2 = -1$ .

- Hence, the solution is the sequence  $\{a_n\}$  with  $a_n = 3 \cdot 2^n - (-1)^n$ .

# An Explicit Formula for the Fibonacci Numbers

We can use Theorem 1 to find an **explicit formula** for the Fibonacci numbers.

The sequence of Fibonacci numbers satisfies the recurrence relation  $f_n = f_{n-1} + f_{n-2}$  with the initial conditions:  $f_0 = 0$  and  $f_1 = 1$ .

**Solution:** The roots of the characteristic equation  $r^2 - r - 1 = 0$  are

$$r_1 = \frac{1+\sqrt{5}}{2}$$

$$r_2 = \frac{1-\sqrt{5}}{2}$$

# Fibonacci Numbers (*continued*)

Therefore by Theorem 1

$$f_n = \alpha_1 \left( \frac{1+\sqrt{5}}{2} \right)^n + \alpha_2 \left( \frac{1-\sqrt{5}}{2} \right)^n$$

for some constants  $\alpha_1$  and  $\alpha_2$ .

Using the initial conditions  $f_0 = 0$  and  $f_1 = 1$ , we have

$$f_0 = \alpha_1 + \alpha_2 = 0$$

Solving, we obtain  $f_1 = \alpha_1 \left( \frac{1+\sqrt{5}}{2} \right) + \alpha_2 \left( \frac{1-\sqrt{5}}{2} \right) = 1$ .

Hence,

$$\alpha_1 = \frac{1}{\sqrt{5}} \quad , \quad \alpha_2 = -\frac{1}{\sqrt{5}}$$

$$f_n = \frac{1}{\sqrt{5}} \left( \frac{1+\sqrt{5}}{2} \right)^n - \frac{1}{\sqrt{5}} \left( \frac{1-\sqrt{5}}{2} \right)^n$$

## The Solution when there is a Repeated Root

**Theorem 2:** Let  $c_1$  and  $c_2$  be real numbers with  $c_2 \neq 0$ . Suppose that  $r^2 - c_1r - c_2 = 0$  has one *repeated root*  $r_0$ . Then the sequence  $\{a_n\}$  is a solution to the recurrence relation  $a_n = c_1a_{n-1} + c_2a_{n-2}$  if and only if

$$a_n = \alpha_1 r_0^n + \alpha_2 n r_0^n$$

for  $n = 0, 1, 2, \dots$ , where  $\alpha_1$  and  $\alpha_2$  are constants.

# Using Theorem 2

**Example:** What is the solution to the recurrence relation  $a_n = 6a_{n-1} - 9a_{n-2}$  with  $a_0 = 1$  and  $a_1 = 6$ ?

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**Solution:** The characteristic equation is  $r^2 - 6r + 9 = 0$ .

The only root is  $r = 3$ . Therefore,  $\{a_n\}$  is a solution to the recurrence relation if and only if

$$a_n = \alpha_1 3^n + \alpha_2 n(3)^n$$

where  $\alpha_1$  and  $\alpha_2$  are constants.

To find the constants  $\alpha_1$  and  $\alpha_2$ , note that

$$\begin{aligned} a_0 = 1 &= \alpha_1 \quad \text{and} \\ a_1 = 6 &= \alpha_1 \cdot 3 + \alpha_2 \cdot 3. \end{aligned}$$

Solving, we find that  $\alpha_1 = 1$  and  $\alpha_2 = 1$ .

Hence,  $a_n = 3^n + n3^n$ .



# Solving Linear Homogeneous Recurrence Relations of Arbitrary Degree

This theorem can be used to solve linear homogeneous recurrence relations with constant coefficients.

**Theorem 3:** Let  $c_1, c_2, \dots, c_k$  be real numbers. Suppose that the characteristic equation

$$r^k - c_1 r^{k-1} - \dots - c_k = 0$$

has  $k$  distinct roots  $r_1, r_2, \dots, r_k$ . Then a sequence  $\{a_n\}$  is a solution of the recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$

if and only if

for  $n = 0, 1, 2, \dots$ , where  $\alpha_1, \alpha_2, \dots, \alpha_k$  are constants.

$$a_n = \alpha_1 r_1^n + \alpha_2 r_2^n + \dots + \alpha_k r_k^n$$

# The General Case with Repeated Roots Allowed

**Theorem 4:** Let  $c_1, c_2, \dots, c_k$  be real numbers. Suppose that the characteristic equation

$$r^k - c_1 r^{k-1} - \dots - c_k = 0$$

has  $t$  distinct roots  $r_1, r_2, \dots, r_t$  with multiplicities  $m_1, m_2, \dots, m_t$ , respectively so that  $m_i \geq 1$  for  $i = 1, 2, \dots, t$  and  $m_1 + m_2 + \dots + m_t = k$ . Then a sequence  $\{a_n\}$  is a solution of the recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k}$$

if and only if

$$a_n = (\alpha_{1,0} + \alpha_{1,1}n + \dots + \alpha_{1,m_1-1}n^{m_1-1})r_1^n + (\alpha_{2,0} + \alpha_{2,1}n + \dots + \alpha_{2,m_2-1}n^{m_2-1})r_2^n + \dots + (\alpha_{t,0} + \alpha_{t,1}n + \dots + \alpha_{t,m_t-1}n^{m_t-1})r_t^n$$

for  $n = 0, 1, 2, \dots$ , where  $\alpha_{i,j}$  are constants for  $1 \leq i \leq t$  and  $0 \leq j \leq m_i - 1$ .

# Linear Nonhomogeneous Recurrence Relations with Constant Coefficients

**Definition:** A linear *nonhomogeneous* recurrence relation with constant coefficients is a recurrence relation of the form:

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k} + F(n),$$

where  $c_1, c_2, \dots, c_k$  are real numbers, and  $F(n)$  is a function not identically zero *depending only on n*.

The recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \dots + c_k a_{n-k},$$

is called the *associated homogeneous recurrence relation*.

# Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*cont.*)

The following are linear nonhomogeneous recurrence relations with constant coefficients:

$$a_n = a_{n-1} + 2^n,$$

$$a_n = a_{n-1} + a_{n-2} + n^2 + n + 1,$$

$$a_n = 3a_{n-1} + n3^n,$$

$$a_n = a_{n-1} + a_{n-2} + a_{n-3} + n!$$

where the following are the **associated** linear homogeneous recurrence relations, respectively:

$$a_n = a_{n-1},$$

$$a_n = a_{n-1} + a_{n-2},$$

$$a_n = 3a_{n-1},$$

$$a_n = a_{n-1} + a_{n-2} + a_{n-3}$$

## Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients

**Theorem 5:** If  $\{a_n^{(p)}\}$  is a particular solution of the nonhomogeneous linear recurrence relation with constant coefficients

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k} + F(n),$$

then every solution is of the form  $\{a_n^{(p)} + a_n^{(h)}\}$ , where  $\{a_n^{(h)}\}$ : solution of the associated homogeneous recurrence relation

$$a_n = c_1 a_{n-1} + c_2 a_{n-2} + \cdots + c_k a_{n-k}.$$

# Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*continued*)

**Example:** Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ .  
What is the solution with  $a_1 = 3$ ?

**Solution:** The associated linear homogeneous equation is  $a_n = 3a_{n-1}$ .  
Its solutions are  $a_n^{(h)} = \alpha 3^n$ , where  $\alpha$  is a constant.

# Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*continued*)

**Example:** Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ .  
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**Solution:** The associated linear homogeneous equation is  $a_n = 3a_{n-1}$ .  
Its solutions are  $a_n^{(h)} = \alpha 3^n$ , where  $\alpha$  is a constant.

Because  $F(n) = 2n$  is a polynomial in  $n$  of **degree one**, to find a particular solution **we might try** a linear function in  $n$ , say  $p_n = cn + d$ , where  $c$  and  $d$  are constants.

# Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*continued*)

**Example:** Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ .

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Suppose that  $p_n = cn + d$  is such a solution.

Then  $a_n = 3a_{n-1} + 2n$  becomes  $cn + d = 3(c(n-1) + d) + 2n$ .



# Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*continued*)

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Because  $F(n) = 2n$  is a polynomial in  $n$  of degree one, to find a particular solution we might try a linear function in  $n$ , say  $p_n = cn + d$ , where  $c$  and  $d$  are constants.

Then  $a_n = 3a_{n-1} + 2n$  becomes  $cn + d = 3(c(n-1) + d) + 2n$ .

Simplifying yields  $(2 + 2c)n + (2d - 3c) = 0$ . It follows that  $cn + d$  is a solution if and only if  $2 + 2c = 0$  and  $2d - 3c = 0$ . Therefore,  $cn + d$  is a solution if and only if  $c = -1$  and  $d = -3/2$ .

Consequently,  $a_n^{(p)} = -n - 3/2$  is a particular solution.

# Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*continued*)

**Example:** Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ .  
What is the solution with  $a_1 = 3$ ?

**Solution:** The associated linear homogeneous equation is  $a_n = 3a_{n-1}$ .  
 $a_n^{(h)} = \alpha 3^n$ , homogeneous solution.

$a_n^{(p)} = -n - 3/2$  is a particular solution.

By Theorem 5, all solutions are of the form

$a_n = a_n^{(p)} + a_n^{(h)} = -n - 3/2 + \alpha 3^n$ , where  $\alpha$  is a constant.

# Solving Linear Nonhomogeneous Recurrence Relations with Constant Coefficients (*continued*)

**Example:** Find all solutions of the recurrence relation  $a_n = 3a_{n-1} + 2n$ .  
What is the solution with  $a_1 = 3$ ?

**Solution:** The associated linear homogeneous equation is  $a_n = 3a_{n-1}$ .  
 $a_n^{(h)} = \alpha 3^n$ , homogeneous solution.

$a_n^{(p)} = -n - 3/2$  is a particular solution.

By Theorem 5, all solutions are of the form

$$a_n = a_n^{(p)} + a_n^{(h)} = -n - 3/2 + \alpha 3^n, \text{ where } \alpha \text{ is a constant.}$$

To find the solution with  $a_1 = 3$ ,  
let  $n = 1$  in the above formula for the general solution.

Then  $3 = -1 - 3/2 + 3\alpha$ , and  $\alpha = 11/6$ .

Hence, the solution is  $a_n = -n - 3/2 + (11/6)3^n$ .