1.

| Customer N. | Model | Shipping Address | Producer | Phone | Price (x100\$) |
|-------------|-------------|-----------------------|------------------------|----------------|-----------------------|
| Alan Smith | Golf | 35 Palm St, Miami | Volkswagen AG | (090) 555 6688 | 250 |
| Roger Banks | Fiesta | 47 Camp. Road, Boston | Ford MC | (090) 600 9090 | 300 |
| Evan Wilson | Golf, Focus | 28 Rock Av, Denver | Volkswagen AG, Ford MC | Both | 450 |
| Alan Smith | Fiesta | 47 Camp. Road, Boston | Ford MC | (090) 600 9090 | 300 |

The relation above records the info of car producer company and the sale operations.

Considering that relation perform:

a) 1NF decomposition

- Each cell should be atomic
- Entries in an attribute should be same type
- Rows should at least one attribute that uniquely identifies it.

| Cust ID* | Customer N. | Model | Shipping Address | Producer | Phone | Price(x100\$) |
|----------|-------------|--------|-----------------------|---------------|----------------|---------------|
| 1 | Alan Smith | Golf | 35 Palm St, Miami | Volkswagen AG | (090) 555 6688 | 250 |
| 2 | Roger Banks | Fiesta | 47 Camp. Road, Boston | Ford MC | (090) 600 9090 | 300 |
| 3 | Evan Wilson | Golf | 28 Rock Av, Denver | Volkswagen AG | (090) 555 6688 | 250 |
| 3 | Evan Wilson | Focus | 28 Rock Av, Denver | Ford MC | (090) 600 9090 | 200 |
| 4 | Alan Smith | Fiesta | 47 Camp. Road, Boston | Ford MC | (090) 600 9090 | 300 |

b) 2NF (but not 3NF) decomposition

To be in 2NF, each attribute (except non-key attributes) in the relation has to be dependent on only keys not any subset of keys.

Here attributes, *Cust ID* and *Model* can be a super key. In this case, attribute *Producer* is dependent on *Model*, which breaks the 2NF rule.

| Cust ID* | Customer N. | Model | Shipping Address |
|----------|-------------|--------|-----------------------|
| 1 | Alan Smith | Golf | 35 Palm St, Miami |
| 2 | Roger Banks | Fiesta | 47 Camp. Road, Boston |
| 3 | Evan Wilson | Golf | 28 Rock Av, Denver |
| 3 | Evan Wilson | Focus | 28 Rock Av, Denver |
| 4 | Alan Smith | Fiesta | 47 Camp. Road, Boston |

| Model* | Producer | Phone | Price(x100\$) |
|--------|---------------|----------------|---------------|
| Golf | Volkswagen AG | (090) 555 6688 | 250 |
| Fiesta | Ford MC | (090) 600 9090 | 300 |
| Focus | Ford MC | (090) 600 9090 | 200 |

| Cust ID* | Model* |
|----------|--------|
| 1 | Golf |
| 2 | Fiesta |
| 3 | Golf |
| 3 | Focus |
| 4 | Fiesta |

Quiz5 – Solutions

c) 3NF decomposition

All attributes can be dependent on only key not other attribute.

Here there is a dependency between *Producer* and *Phone* so those should be placed in another table and referred in the table of interest.

| Cust ID* | Customer N. | Model | Shipping Address |
|----------|-------------|--------|-----------------------|
| 1 | Alan Smith | Golf | 35 Palm St, Miami |
| 2 | Roger Banks | Fiesta | 47 Camp. Road, Boston |
| 3 | Evan Wilson | Golf | 28 Rock Av, Denver |
| 3 | Evan Wilson | Focus | 28 Rock Av, Denver |
| 4 | Alan Smith | Fiesta | 47 Camp. Road, Boston |

| Model* | Producer -FK | Price(x100\$) |
|--------|---------------|---------------|
| Golf | Volkswagen AG | 250 |
| Fiesta | Ford MC | 300 |
| Focus | Ford MC | 200 |

| Producer* | Phone |
|---------------|----------------|
| Volkswagen AG | (090) 555 6688 |
| Ford MC | (090) 600 9090 |

| Cust ID* | Model* |
|----------|--------|
| 1 | Golf |
| 2 | Fiesta |
| 3 | Golf |
| 3 | Focus |
| 4 | Fiesta |

Quiz5 – Solutions

- 2. Suppose that we decompose the schema <u>r (A, B, C, D, E)</u> into <u>r1 (D, A, E)</u> and <u>r2 (D, B, C)</u>
- a) Show that this decomposition is a lossless decomposition if the following set *F* of functional dependencies holds:

F:
$$D \rightarrow AE$$

 $EB \rightarrow C$
 $A \rightarrow B$
 $C \rightarrow D$

A decomposition $\{R1, R2\}$ is a lossless-join decomposition if $R1 \cap R2 \to R1$ or $R1 \cap R2 \to R2$. Let R1 = (D, A, E), R2 = (D, B, C), and $R1 \cap R2 = D$. Since D is a candidate key, therefore $R1 \cap R2 \to R1$.

b) Give a lossless-join decomposition into BCNF of schema R.

result : = {R}; $F^{+} = \{D \rightarrow ABCDE, A \rightarrow B, AE \rightarrow ABCDE, EB \rightarrow ABCDE, C \rightarrow ABCDE, ...\}.$ R is not in BCNF. A \rightarrow B is a non-trivial f.d. that holds on R, A \cap B = \emptyset , and A \rightarrow ABCDE is not in F^{+} .

Therefore, result := $(result - R) \cup (R - B) \cup (A, B)$, i.e.

- $(D, A, B, C) \cup (A, B).$
- (D, A, E, C) and (A,B) are in BCNF. So this is a decomposition of R into BCNF.