# Lecture 11 Semester 1 Revision

COMP26120

Giles Reger and Lucas Cordeiro

December 2019

## Headlines

January Exam is 1.5 hours (90 minutes)

#### The Exam is Online Only via Blackboard

C programming is not examinable

The two main topics this semester have been

- Complexity Analysis and Divide-and-Conquer [3 lectures]
- Data Structures [3 lectures]

## Reminder: Course Assessment

## Coursework (50%):

- 25% each semester
- In each semester
  - 20% from coding exercises
  - 5% from 3 online quizzes (see website for when)
  - Quizzes provide early feedback on conceptual issues

## Exam (50%):

- 15% first semester, 35% second semester
- Material (general concepts) covered in labs is examinable
- Note: new staff means exams likely to be a little different in style

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# Reminder: Health Warning 1

## Warning 1

This course unit is rather different from others. It is important that you take the following into account:

- You will be expected to read a considerable amount of material outside the lectures - this includes the course textbook and other material too.
- Lectures will give some guidance about course content and topics, laboratory exercises, etc., but much of the material of the course unit is not lectured!

## **Examinable Topics**

C Programming is assessed in labs only.

## Complexity Analysis and Divide-and-Conquer

- Divide and Conquer paradigm
- Sorting Algorithms
- Algorithmic Complexity
- Sums and Integrals
- Proof by Induction

#### Data Structures

- Abstract Datatypes
- Binary trees
- AVL trees
- Hashing structures
- Skip Lists



# Relation to Previous Papers

Computing the complexity of recursive programs (and the required mathematics) is included in this year's syllabus and first appeared last year. .

With respect to past papers:

- We moved data structure material from semester 2 to semester 1 last year
- Relevant to this semester would be (e.g.) question 3 of paper 2 from 2015-2017 and question 4 from paper 2 in 2015
- Some parts of other questions may be relevant you can use your knowledge of the relevant topics from this semester to make a judgement

What is the most accurate (e.g. smallest correct) time complexity of this (C-like) pseudocode?

```
for (i=1; i < n; i *= 2){
 for (j = n; j > 0; j /= 2) {
   for (k = j; k < n; k += 2) {
     sum += (i + j * k);
```

#### Options:

- $O(n (\log n)^2)$
- $O(n^2)$
- $O(n (\log n)$
- None of the above.

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    for ( k = j; k < n; k += 2 ) {
      sum += (i + j * k );
    }
}</pre>
```

#### Options:

- $O(n (\log n)^2)$  Three nested loops. Inner is O(n), others are  $O(\log n)$
- $O(n^2)$
- $O(n (\log n)$
- None of the above.

#### Consider the scenario

John is working as a software developer at BoostCode UK Ltd. He needs to compare two implementations of sorting algorithms on the same machine, which will be later integrated into a software product. In particular, for inputs defined by size n, the first algorithm denoted by A runs in  $4n^2$ , while the second algorithm denoted by B runs in  $128n \log_2 n$ . John needs to compute values of n where algorithm A beats algorithm B so that he can make the best choice of which one will be integrated into the product. For which values of n will algorithm A run faster than algorithm B?

## Options:

• 256

- 64
- 128
   None of the Above

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Solve the recurrence equation

$$T(n) = T\left(\frac{2n}{3}\right) + 1$$

using the master method to give tight asymptotic bounds.

Options:

- $T(n) = \theta(\log n)$
- $T(n) = \theta(n)$
- $T(n) = \theta(1)$
- None of the Above

Other questions might include:

- Which Case applies when applying the master method?
- Identify the parts a, b, f(n)



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Solve the recurrence equation

$$T(n) = T\left(\frac{2n}{3}\right) + 1$$

using the master method to give tight asymptotic bounds.

#### Options:

- $T(n) = \theta(\log n)$  Case 2 applies
- $T(n) = \theta(n)$
- $T(n) = \theta(1)$
- None of the Above

Other questions might include:

- Which Case applies when applying the master method?
- Identify the parts a, b, f(n)



Which sentences about sums are correct?

- The perturbation method typically works for sums that do not contain exponentials
- ② The sum  $\sum_{i=1}^{n} a$  can be approximated by integrals to O(n)
- **1** The sum  $\sum_{i=1}^{n} i^3$  can be approximated by integrals to  $O(n^4)$
- The sum  $\sum_{i=1}^{n} \frac{1}{i}$  can be approximated by integrals to  $O(\ln n)$

Which sentences about sums are correct?

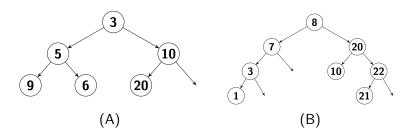
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b. 
$$\sum_{i=1}^{n} a = \int_{1}^{n} a = O(n)$$
.  
c.  $\sum_{i=1}^{n} i^{3} = O(n^{4})$ .  

$$\lim_{i=1}^{n} 1/i = \int_{1}^{n} 1/i = O(\ln n)$$
.

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## Consider the following two trees

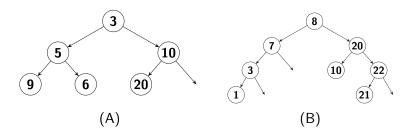


#### Which statements are TRUE:

- A is ordered
- B is ordered
- A is complete
- B is complete

- A has the Heap-Order property
- B has the Height-Balance property
- The balance of node 7 in A is -2
- The balance of node 8 in A is -2

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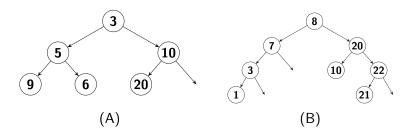


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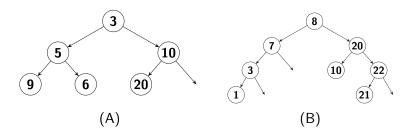


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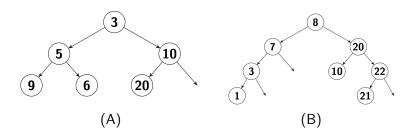


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## Consider the following scenario

I am creating an application that generates a Word of the Day. I want the words to get more interesting over time so I've associated an interestingness score with each word. I have the list in alphabetical order. At the start of the year I insert the words. Each day I return a new word.

#### Which ADT do I need?

- Set
  - Stack
  - Dictionary
  - Priority Queue

#### Which Data structure should I use?

- Dynamic Array
- AVL tree
- Skip List
- Binary Heap

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**Assertion:** Inserting 1, 10, 13, 23, 29 into a hash table of size 100 using the hash function  $h(n) = n \mod 9$  and linear probing would lead to 2 collisions.

#### **BECAUSE**

**Reason:** Quadratic Probing uses the formula  $A[(i+j^2) \mod N]$  to avoid clustering effects

- The assertion and reason are both correct and the reason is valid.
- The assertion and reason are both correct but the reason is invalid.
- The assertion is correct, but the reason is incorrect.
- The assertion is incorrect, but the reason is correct.
- Both the answer and the reason are incorrect.

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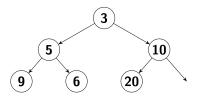
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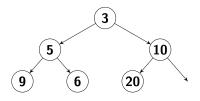


How many popMin operations are required before the tree has height 1?

How many popMin operations are required before 20 is the left of the root node?

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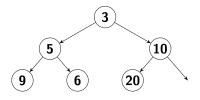
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How many popMin operations are required before the tree has height 1? 3

How many popMin operations are required before 20 is the left of the root node?

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How many popMin operations are required before the tree has height 1? 3

How many popMin operations are required before 20 is the left of the root node? 1

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## Data Structures Questions in General

The kinds of questions I might ask (similar to quiz):

- To identify the correct definition for a term
- To identify where a data structure has a property
- To identify whether a data structure is the result of certain operations e.g. inserting certain numbers into an hash table, AVL tree or binary heap, or removing a node from an AVL tree and re-balancing

# Summary

Good Luck!

Any Questions?

