

# MILP Model for Single Depot VRP with Customer and Locker Assignments

## 1 Mathematical Model

### 1.1 Decision Variables

$x_{i,j,v}$  Binary, 1 if vehicle  $v$  travels from node  $i$  to  $j$ , 0 otherwise  $\forall i, j \in N, v \in V, i \neq j$  (1)

$y_v$  Binary, 1 if vehicle  $v$  is used, 0 otherwise  $\forall v \in V$  (2)

$home_i$  Binary, 1 if customer  $i$  is served at home, 0 otherwise  $\forall i \in M$  (3)

$z_{i,lk}$  Binary, 1 if customer  $i$  is assigned to locker  $lk$ , 0 otherwise  $\forall i \in M, lk \in L$  (4)

$w_{lk}$  Binary, 1 if locker  $lk$  is used, 0 otherwise  $\forall lk \in L$  (5)

$d_{lk}$  Continuous, demand assigned to locker  $lk$   $\forall lk \in L$  (6)

$u_{i,v}$  Continuous, load of vehicle  $v$  at node  $i$   $\forall i \in M \cup L, v \in V$  (7)

$\delta_{v,i}$  Binary, 1 if vehicle  $v$  visits node  $i$ , 0 otherwise  $\forall i \in M \cup L, v \in V$  (8)

Additional variables for load constraints:  $load\_aux_{v,i}$   $\forall v \in V, i \in L$  (9)

### 1.2 Objective Function

Minimize the total cost, including routing and assignment costs:

$$\min \sum_{v \in V} \left[ \text{fixed cost of } v \cdot y_v + \sum_{i \in N} \sum_{j \in N} c_{i,j} \cdot x_{i,j,v} \right] + \sum_{i \in M} \sum_{lk \in L} c_{i,lk} \cdot z_{i,lk} \quad (10)$$

### 1.3 Constraints

- Each customer is either served at home or assigned to a locker:

$$home_i + \sum_{lk \in L} z_{i,lk} = 1, \quad \forall i \in M \quad (11)$$

- If a customer cannot be assigned to a locker, they must receive home delivery:

$$\text{If no available locker: } home_i = 1, \quad \forall i \in M \quad (12)$$

- Locker demand constraints:

$$d_{lk} = \sum_{i \in M} demand_i \cdot z_{i,lk}, \quad \forall lk \in L \quad (13)$$

- If a locker is not used, its assigned demand must be zero:

$$d_{lk} = 0, \quad \text{if no customer is assigned to locker } lk \quad (14)$$

- Locker capacity constraints:

$$d_{lk} \leq capacity_{lk} \cdot w_{lk}, \quad \forall lk \in L \quad (15)$$

- Flow conservation:

$$\sum_{i \in N} x_{i,k,v} = \sum_{j \in N} x_{k,j,v}, \quad \forall k \in M \cup L, v \in V \quad (16)$$

- MTZ subtour elimination:

$$u_{i,v} + demand_j \leq u_{j,v} + C(1x_{i,j,v}), \quad \forall i, j \in M \cup L, v \in V, i \neq j \quad (17)$$

- Customers assigned home must be visited:

$$\sum_{v \in V} \delta_{v,i} = home_i, \quad \forall i \in M \quad (18)$$

- If a locker is used, it must be visited by at least one vehicle:

$$\sum_{i \in N} x_{i,lk,v} = w_{lk}, \quad \forall lk \in L, v \in V \quad (19)$$

- Vehicle capacity constraint:

$$\sum_{i \in M} demand_i \cdot \delta_{v,i} + \sum_{i \in L} load_{aux_{v,i}} \leq capacity_v, \quad \forall v \in V \quad (20)$$

- Time limit constraint:

$$\sum_{i \in N} \sum_{j \in N} \frac{c_{i,j}}{speed} \cdot x_{i,j,v} \leq max\_duration_v, \quad \forall v \in V \quad (21)$$

- Forbidden group constraints:

$$\sum_{i \in G} \sum_{j \in G} x_{i,j,v} \leq 1, \quad \forall G \in \text{forbidden groups}, v \in V \quad (22)$$