

A Guide to Approximations

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Floating Point Approximations

Most embedded processors don't know how to compute trig and other complex functions. Programming in C we're content to call a library routine that does all of the work for us. Unhappily this optimistic approach often fails in real time systems where size, speed and accuracy are all important issues.

The compiler's runtime package is a one-size-fits-all proposition. It gives a reasonable trade-off of speed and precision. But every embedded system is different, with different requirements. In some cases it makes sense to write our own approximation routines. Why?

Speed – Many compilers have very slow runtime packages. *A clever approximation may eliminate the need to use a faster CPU.*

Predictability – Compiler functions vary greatly in execution time depending on the input argument. Real time systems must be *predictable* in the time domain. The alternative is to always assume worst case execution time, which again may mean your CPU is too slow, too loaded, for the application.

Accuracy – Read the compiler's manuals carefully! Some explicitly do not support the ASNI C standard, which requires all trig to be double precision. (8051 compilers are notorious for this). Alternatively, why pay the cost (in time) for double precision math when you only need 5 digits of accuracy?

Size – When memory is scarce, using one of these approximations may save much code space. If you only need a simple cosine, why include the entire floating point trig library?

This collection is not an encyclopedia of all possible approximations; rather, it's the most practical ones distilled from the bible of the subject, *Computer Approximations* by John Hart (ISBN 0-88275-642-7). Unfortunately this work is now out of print. It's also very difficult to use without a rigorous mathematical background.

All of the approximations here are polynomials, or ratios of polynomials. All use very cryptic coefficients derived from Chebyshev series and Bessel functions. Rest assured that these numbers give minimum errors for the indicated ranges. Each approximation (with a few exceptions) has an error chart so you can see exactly where the errors occur. In some cases, if you're using a limited range of input data, your accuracies can far exceed the indicated values. For instance, `cos_73` is accurate to 7.3 decimal digits over the 0 to 360 degree range. But as the graph shows, in the range 0 to 30 degrees you can get almost an order of magnitude improvement.

Do be wary of the precision of your compiler's floating point package. Some treat doubles as floats. Others, especially for tiny CPUs like the PIC, cheat and offer less than full 32 bit floating point precision.

All of the code for the following approximations was compiled with Microsoft's Visual C++ 6.0. The source is available at www.ganssle.com/approx/sincos.cpp. It includes test code that writes a text file of results and errors; imported to a spreadsheet we can see just how accurate they are.

General Trig Notes

We generally work in radians rather than degrees. The 360 degrees in a circle are equivalent to 2π radians; thus, one radian is $360/(2\pi)$, or about 57.3 degrees. This may seem a bit odd till you think of the circle's circumference, which is $2\pi r$; if r (the circle's radius) is one, the circumference is indeed 2π .

The conversions between radians and degrees are:

```
Angle in radians= angle in degrees * 2 p /360
Angle in degrees= angle in radians * 360/(2 p)
```

Degrees	Radians	Sine	Cosine	Tangent
0	0	0	1	0
45	$\pi/4$	$\sqrt{2}/2$	$\sqrt{2}/2$	1
90	$\pi/2$	1	0	infinity
135	$3\pi/4$	$\sqrt{2}/2$	$-\sqrt{2}/2$	-1
180	π	0	-1	Infinity
225	$5\pi/4$	$-\sqrt{2}/2$	$-\sqrt{2}/2$	1
270	$3\pi/2$	-1	0	Infinity
315	$7\pi/4$	$-\sqrt{2}/2$	$\sqrt{2}/2$	-1
360	2π	0	1	0

Cosine and Sine

The following examples all approximate the cosine function; sine is derived from cosine via the relationship:

$$\sin(x) = \cos(\pi/2 - x)$$

In other words, the sine and cosine are the same function, merely shifted 90° in phase. The sine code is (assuming we're calling `cos_32`, the lowest accuracy cosine approximation):

```
// The sine is just cosine shifted a half-pi, so
// we'll adjust the argument and call the cosine approximation.
//
float sin_32(float x){
    return cos_32(halfpi-x);
}
```

All of the cosine approximations in this chapter compute the cosine accurately over the range of 0 to $\pi/2$ (0 to 90°). That surely denies us of most of the circle! Approximations in general work best over rather limited ranges; it's up to us to reduce the input range to something the approximation can handle accurately.

Therefore, before calling any of the following cosine approximations we assume the range has been reduced to 0 to $\pi/2$ using the following code:

```
// Math constants
double const pi=3.1415926535897932384626433;// pi
double const twopi=2.0*pi;           // pi times 2
double const halfpi=pi/2.0;          // pi divided by 2
//
// This is the main cosine approximation "driver"
// It reduces the input argument's range to [0, pi/2],
// and then calls the approximator.
//
float cos_32(float x){
    int quad;                          // what quadrant are we in?

    x=fmod(x, twopi);                  // Get rid of values > 2* pi
    if(x<0)x=-x;                       // cos(-x) = cos(x)
    quad=int(x/halfpi);                // Get quadrant # (0 to 3)    switch (quad){
    case 0: return cos_32s(x);
    case 1: return -cos_32s(pi-x);
    case 2: return -cos_32s(x-pi);
    case 3: return cos_32s(twopi-x);
    }
}
```

This code is configured to call `cos_32s`, which is the approximation (detailed shortly) for computing the cosine to 3.2 digits accuracy. Use this same code, though, for all cosine approximations; change `cos_32s` to `cos_52s`, `cos_73s` or `cos_121s`, depending on which level of accuracy you need. See the complete listing for a comprehensive example.

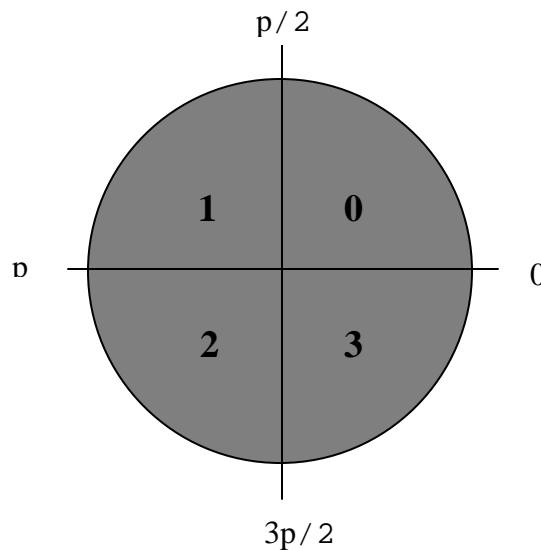
If you can guarantee that the input argument will be greater than zero and less than 2π , delete the two red lines in the listing above to get even faster execution.

Be clever about declaring variables and constants. Clearly, working with the `cos_32` approximation nothing must be declared “double”. Use `float` for more efficient code.

Reading the complete listing you'll notice that for `cos_32` and `cos_52` we used floats everywhere; the more accurate approximations declare things as doubles.

One trick that will speed up the approximations is to compute x^2 by incrementing the characteristic of the floating point representation of x . You'll have to know exactly how the numbers are stored, but can save hundreds of microseconds over performing the much clearer " $x*x$ " operation.

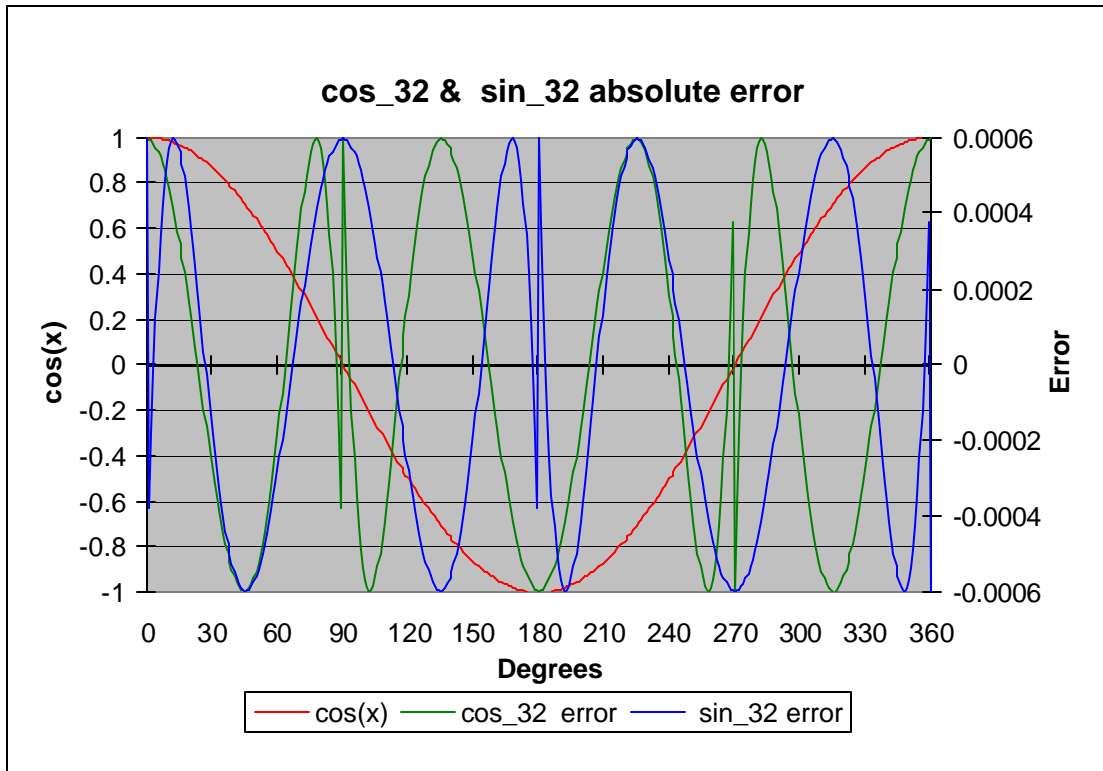
How does the range reduction work? Note that the code divides the input argument into one of four "quadrants" – the very same quadrants of the circle shown below:



Quadrants 0 to 3 of the circle

- For the first quadrant (0 to $p/2$) there's nothing to do since the cosine approximations are valid over this range.
- In quadrant 1 the cosine is symmetrical with quadrant 0, if we reduce it's range by subtracting the argument from p . The cosine, though, is negative for quadrants 1 and 2 so we compute $-\cos(p-x)$.
- Quadrant 2 is similar to 1.
- Finally, in 3 the cosine goes positive again; if we subtract the argument from $2p$ it translates back to something between 0 and $p/2$.

The approximations do convert the basic polynomial to a simpler, much less computationally expensive form, as described in the comments. All floating point operations take appreciable amounts of time, so it's important to optimize the design.

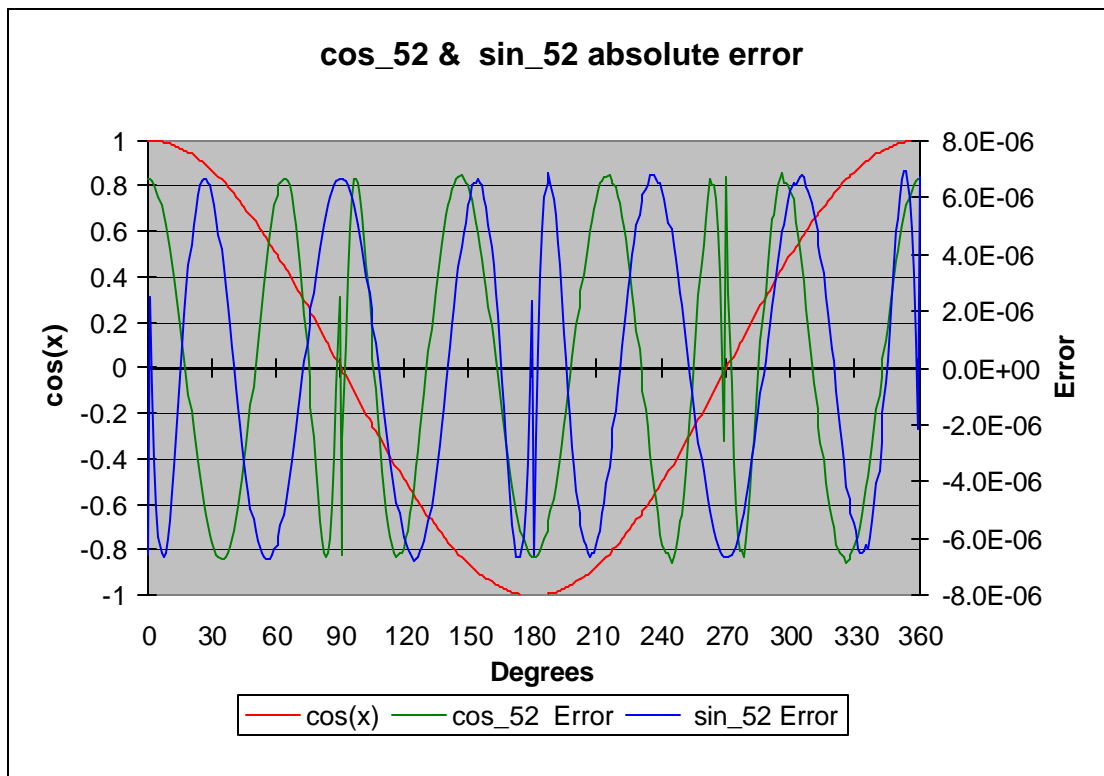


```
//      cos_32s computes cosine (x)
//
//  Accurate to about 3.2 decimal digits over the range [0, pi/2].
//  The input argument is in radians.
//
//  Algorithm:
//      cos(x)= c1 + c2*x**2 + c3*x**4
//  which is the same as:
//      cos(x)= c1 + x**2(c2 + c3*x**2)
//
float cos_32s(float x)
{
    const float c1= 0.99940307;
    const float c2=-0.49558072;
    const float c3= 0.03679168;

    float x2;                                // The input argument squared

    x2=x * x;
    return (c1 + x2*(c2 + c3 * x2));
}
```

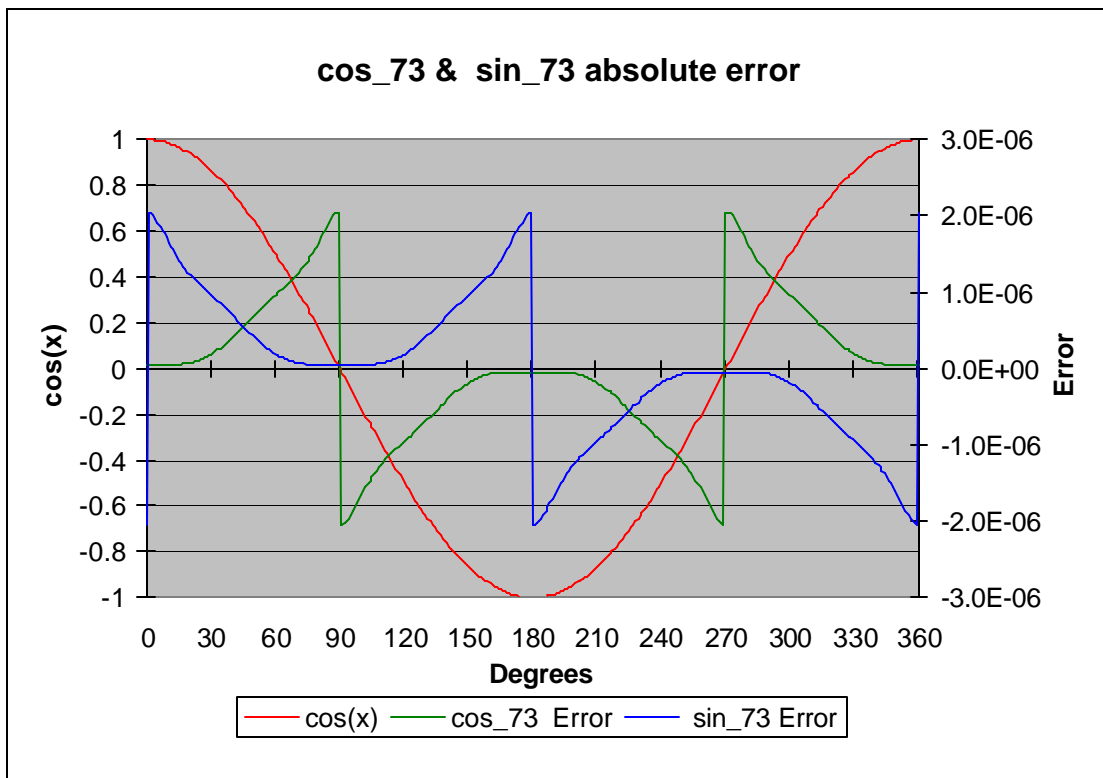
cos_32 computes a cosine to about 3.2 decimal digits of accuracy. Use the range reduction code (listed earlier) if the range exceeds 0 to $\pi/2$. The plotted errors are absolute (not percent error).



```
//      cos_52s computes cosine (x)
//
//  Accurate to about 5.2 decimal digits over the range [0, pi/2].
//  The input argument is in radians.
//
//  Algorithm:
//      cos(x)= c1 + c2*x**2 + c3*x**4 + c4*x**6
//  which is the same as:
//      cos(x)= c1 + x**2(c2 + c3*x**2 + c4*x**4)
//      cos(x)= c1 + x**2(c2 + x**2(c3 + c4*x**2))
//
float cos_52s(float x)
{
    const float c1= 0.9999932946;
    const float c2=-0.4999124376;
    const float c3= 0.0414877472;
    const float c4=-0.0012712095;
    float x2;                                // The input argument squared

    x2=x * x;
    return (c1 + x2*(c2 + x2*(c3 + c4*x2)));
}
```

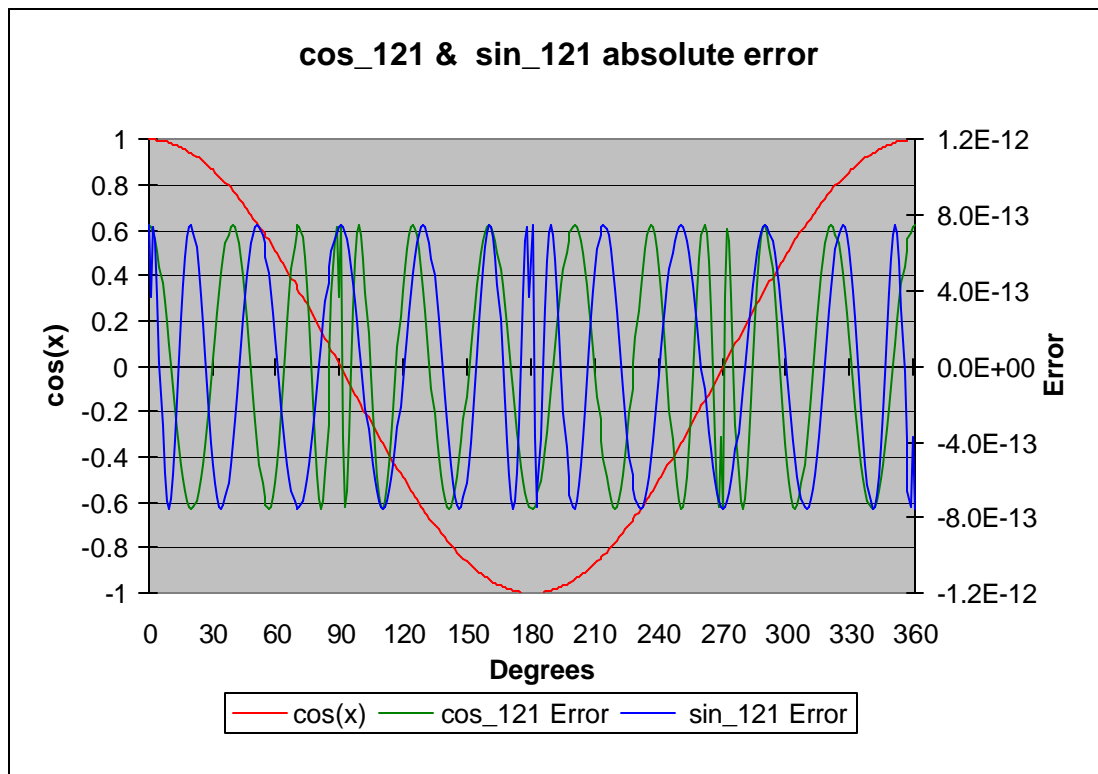
cos_52 computes a cosine to about 5.2 decimal digits of accuracy. Use the range reduction code (listed earlier) if the range exceeds 0 to $\pi/2$. The plotted errors are absolute (not percent error).



```
//      cos_73s computes cosine (x)
//
// Accurate to about 7.3 decimal digits over the range [0, pi/2].
// The input argument is in radians.
//
// Algorithm:
//      cos(x)= c1 + c2*x**2 + c3*x**4 + c4*x**6 + c5*x**8
//      which is the same as:
//      cos(x)= c1 + x**2(c2 + c3*x**2 + c4*x**4 + c5*x**6)
//      cos(x)= c1 + x**2(c2 + x**2(c3 + c4*x**2 + c5*x**4))
//      cos(x)= c1 + x**2(c2 + x**2(c3 + x**2(c4 + c5*x**2)))
//
double cos_73s(double x)
{
    const double c1= 0.999999953464;
    const double c2=-0.4999999053455;
    const double c3= 0.0416635846769;
    const double c4=-0.0013853704264;
    const double c5= 0.00002315393167;
    double x2;
    // The input argument squared

    x2=x * x;
    return (c1 + x2*(c2 + x2*(c3 + x2*(c4 + c5*x2))));
}
```

cos_73 computes a cosine to about 7.3 decimal digits of accuracy. Use the range reduction code (listed earlier) if the range exceeds 0 to $\pi/2$. Also plan on using double precision math for the range reduction code to avoid losing accuracy. The plotted errors are absolute (not percent error).



```
//      cos_121s computes cosine (x)
//
//  Accurate to about 12.1 decimal digits over the range [0, pi/2].
//  The input argument is in radians.
//
//  Algorithm:
//      cos(x)= c1+c2*x**2+c3*x**4+c4*x**6+c5*x**8+c6*x**10+c7*x**12
//  which is the same as:
//      cos(x)= c1+x**2(c2+c3*x**2+c4*x**4+c5*x**6+c6*x**8+c7*x**10)
//      cos(x)= c1+x**2(c2+x**2(c3+c4*x**2+c5*x**4+c6*x**6+c7*x**8 ))
//      cos(x)= c1+x**2(c2+x**2(c3+x**2(c4+c5*x**2+c6*x**4+c7*x**6 )))
//      cos(x)= c1+x**2(c2+x**2(c3+x**2(c4+x**2(c5+c6*x**2+c7*x**4 ))))
//      cos(x)= c1+x**2(c2+x**2(c3+x**2(c4+x**2(c5+x**2(c6+c7*x**2 )))))
//
double cos_121s(double x)
{
    const double c1= 0.99999999999925182;
    const double c2=-0.49999999997024012;
    const double c3= 0.0416666666473384543;
    const double c4=-0.001388888418000423;
    const double c5= 0.0000248010406484558;
    const double c6=-0.0000002752469638432;
    const double c7= 0.0000000019907856854;
    double x2;                                // The input argument squared

    x2=x * x;
    return (c1 + x2*(c2 + x2*(c3 + x2*(c4 + x2*(c5 + x2*(c6 + c7*x2))))));
}
```

cos_121 computes a cosine to about 12.1 decimal digits of accuracy. Use the range reduction code (listed earlier) if the range exceeds 0 to $\pi/2$. Also plan on using double

precision math for the range reduction code to avoid losing accuracy. The plotted errors are absolute (not percent error).

Higher Precision Cosines

Given a large enough polynomial there's no limit to the possible accuracy. A few more algorithms are listed here. These are all valid for the range of 0 to $\pi/2$, and all can use the previous range reduction algorithm to change any angle into one within this range. All take an input argument in radians.

No graphs are included because these exceed the accuracy of the typical compiler's built-in cosine function... so there's nothing to plot the data against.

Note that C's `double` type on most computers carries about 15 digits of precision. So for these algorithms, especially for the 20.2 and 23.1 digit versions, you'll need to use a data type that offers more bits. Some C's support a `long double`. But check the manual carefully! Microsoft's Visual C++, for instance, while it does support the `long double` keyword, converts all of these to `double`.

Accurate to about 14.7 decimal digits over the range $[0, \pi/2]$:

```
c1= 0.9999999999999806767
c2=-0.4999999999998996568
c3= 0.041666666666581174292
c4=-0.001388888886113613522
c5= 0.000024801582876042427
c6=-0.0000002755693576863181
c7= 0.0000000020858327958707
c8=-0.00000000011080716368
cos(x)= c1 + x^2(c2 + x^2(c3 + x^2(c4 + x^2(c5 +
x^2(c6 + x^2(c7 + x^2*c8))))))
```

Accurate to about 20.2 decimal digits over the range $[0, \pi/2]$:

```
c1 = 0.9999999999999999999936329
c2 = -0.4999999999999999999948362843
c3 = 0.0416666666666666665975670054
c4 = -0.0013888888888888885302082298
c5 = 0.000024801587301492746422297
c6 = -0.00000027557319209666748555
c7 = 0.0000000020876755667423458605
c8 = -0.000000000114706701991777771
c9 = 0.00000000000000477687298095717
c10=-0.00000000000000015119893746887
cos(x)= c1 + x^2(c2 + x^2(c3 + x^2(c4 + x^2(c5 + x^2(c6 +
x^2(c7 + x^2(c8 + x^2(c9 + x^2*c10))))))))
```

Accurate to about 23.1 decimal digits over the range $[0, \pi/2]$:

```
c1 = 0.999999999999999999999914771
c2 = -0.49999999999999999999991637437
c3 = 0.041666666666666666665319411988
c4 = -0.001388888888888888880310186415
c5 = 0.00002480158730158702330045157
c6 = -0.000000275573192239332256421489
```

```

c7 = 0.000000002087675698165412591559
c8 = -0.0000000000114707451267755432394
c9 = 0.0000000000000477945439406649917
c10 = -0.0000000000000015612263428827781
c11 = 0.000000000000000039912654507924
cos(x) = c1 + x^2(c2 + x^2(c3 + x^2(c4 + x^2(c5 + x^2(c6 +
      x^2(c7 + x^2(c8 + x^2(c9 + x^2(c10 + x^2*c11))))))))))

```

Tangent

The tangent of an angle is defined as $\tan(x) = \sin(x) / \cos(x)$. Unhappily this is not the best choice, though, for doing an approximation. As $\cos(x)$ approaches zero the errors propagate rapidly. Further, at some points like $\pi/4$ (see the previous graphs of sine and cosine errors) the errors of sine and cosine reinforce each other; both are large and have the same sign.

So we're best off using a separate approximation for the tangent. All of the approximations we'll use generate a valid tangent for angles in the range of 0 to $\pi/4$ (0 to 45 degrees), so once again a range reduction function will translate angles to this set of values.

```
//
// This is the main tangent approximation "driver"
// It reduces the input argument's range to [0, pi/4],
// and then calls the approximator.
// Enter with positive angles only.
//
// WARNING: We do not test for the tangent approaching infinity,
// which it will at x=pi/2 and x=3*pi/2. If this is a problem
// in your application, take appropriate action.
//
float tan_32(float x){
    int octant; // what octant are we in?

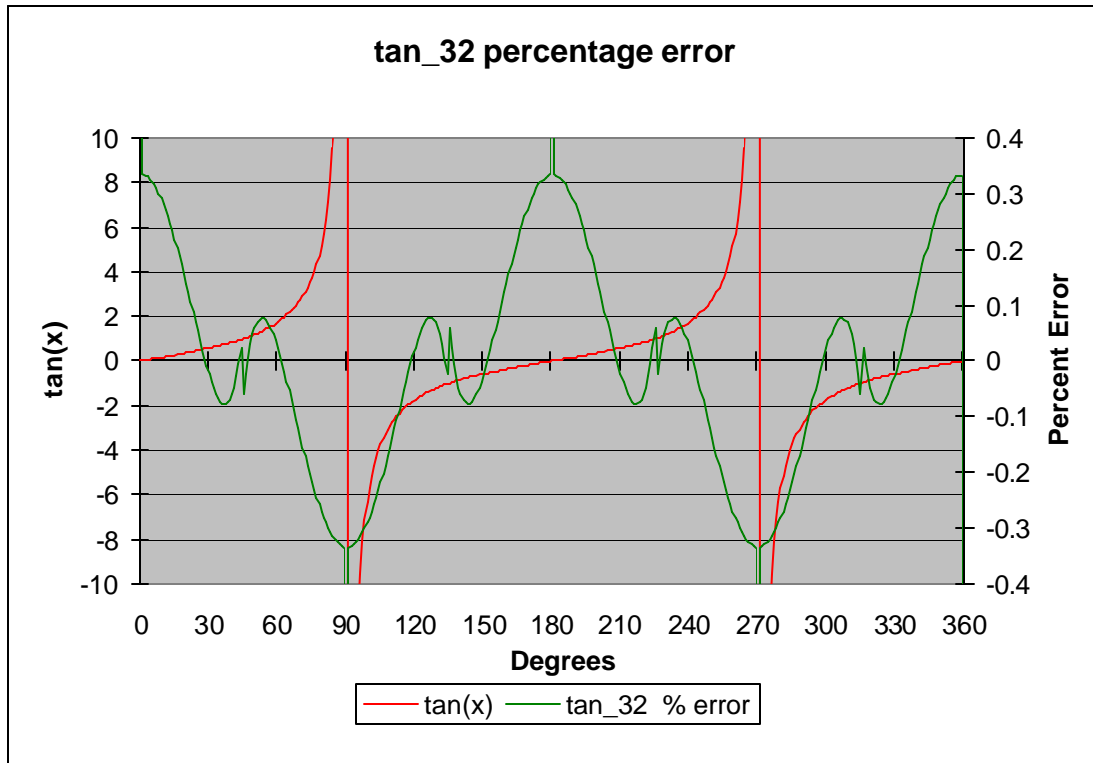
    x=fmod(x, twopi); // Get rid of values >2 *pi
    octant=int(x/qtrpi); // Get octant # (0 to 7)
    switch (octant){
    case 0: return tan_32s(x) *four_over_pi;
    case 1: return 1.0/tan_32s(halfpi-x) *four_over_pi;
    case 2: return -1.0/tan_32s((x-halfpi) *four_over_pi);
    case 3: return -tan_32s(pi-x) *four_over_pi;
    case 4: return tan_32s((x-pi) *four_over_pi);
    case 5: return 1.0/tan_32s(threehalfpi-x) *four_over_pi;
    case 6: return -1.0/tan_32s((x-threehalfpi) *four_over_pi);
    case 7: return -tan_32s(twopi-x) *four_over_pi;
    }
}
```

The code above does the range reduction and then calls `tan_32`. When using the higher precision approximations substitute the appropriate function name for `tan_32`.

The reduction works much like that for cosine, except that it divides the circle into octants and proceeds from there. One quirk is that the argument is multiplied by $4/\pi$. This is because the approximations themselves actually solve for $\tan((\pi/4)x)$.

The listings that follow give the algorithms needed.

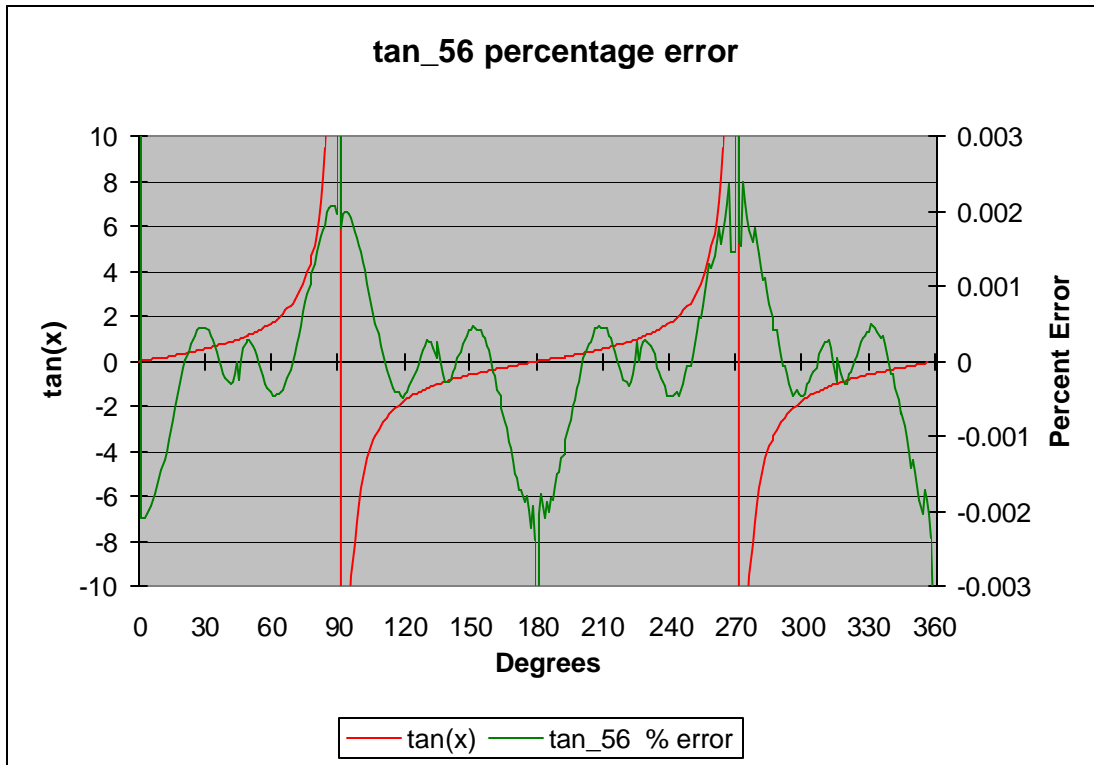
Remember that $\tan(90)$ and $\tan(270)$ both equal infinity. As the input argument gets close to 90 or 270 the value of the tangent skyrockets, as illustrated on the following error charts. *Never take a tangent close to 90 or 270 degrees!*



```
// *****
// ***
// ***   Routines to compute tangent to 3.2 digits
// ***   of accuracy.
// ***
// *****
//
//           tan_32s computes tan(pi*x/4)
//
// Accurate to about 3.2 decimal digits over the range [0, pi/4].
// The input argument is in radians. Note that the function
// computes tan(pi*x/4), NOT tan(x); it's up to the range
// reduction algorithm that calls this to scale things properly.
//
// Algorithm:
//           tan(x)= x*c1/(c2 + x**2)
//
float tan_32s(float x)
{
    const float c1=-3.6112171;
    const float c2=-4.6133253;
    float x2;                                // The input argument squared

    x2=x * x;
    return (x*c1/(c2 + x2));
}
```

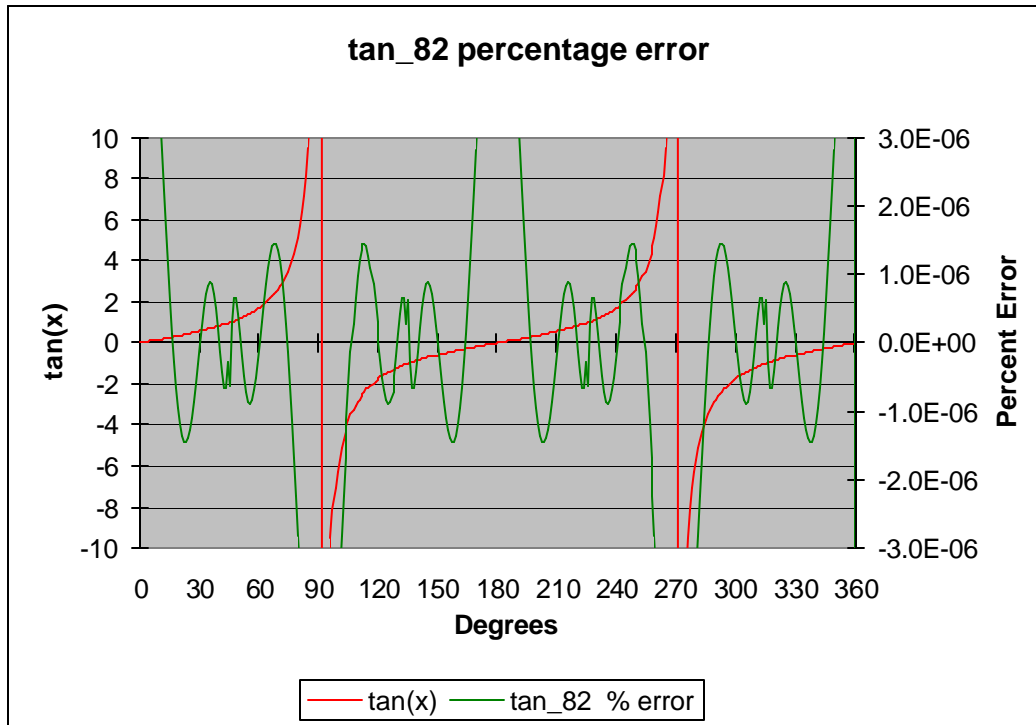
tan_32 computes the tangent of $p/4 \cdot x$ to about 3.2 digits of accuracy. Use the range reduction code to translate the argument to 0 to $p/4$, and of course to compensate for the peculiar “ $p/4$ ” bias required by this routine. Note that the graphed errors are percentage error, not absolute.



```
// *****
// ***
// ***   Routines to compute tangent to 5.6 digits
// ***   of accuracy.
// ***
// *****
//
//           tan_56s computes tan(pi*x/4)
//
// Accurate to about 5.6 decimal digits over the range [0, pi/4].
// The input argument is in radians. Note that the function
// computes tan(pi*x/4), NOT tan(x); it's up to the range
// reduction algorithm that calls this to scale things properly.
//
// Algorithm:
//           tan(x)= x(c1 + c2*x**2)/(c3 + x**2)
//
float tan_56s(float x)
{
    const float c1=-3.16783027;
    const float c2= 0.134516124;
    const float c3=-4.033321984;
    float x2;                                // The input argument squared

    x2=x * x;
    return (x*(c1 + c2 * x2)/(c3 + x2));
}
```

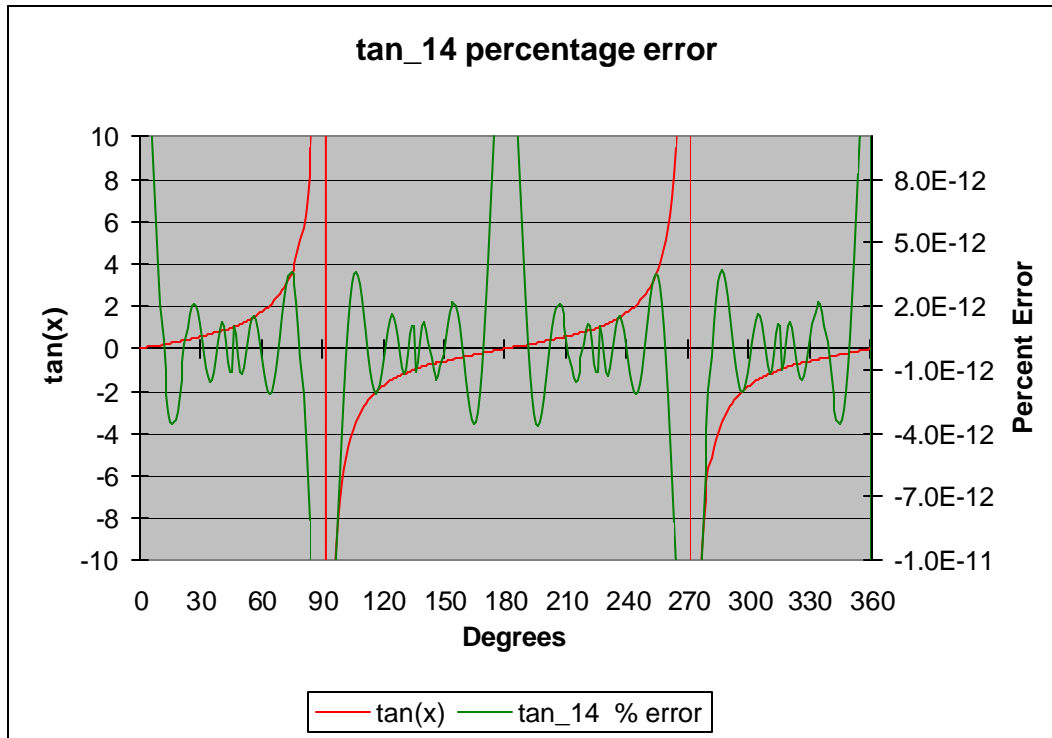
tan_56 computes the tangent of $\pi/4 \cdot x$ to about 5.6 digits of accuracy. Use the range reduction code to translate the argument to 0 to $\pi/4$, and of course to compensate for the peculiar “ $\pi/4$ ” bias required by this routine. Note that the graphed errors are percentage error, not absolute.



```
// *****
// ***
// ***  Routines to compute tangent to 8.2 digits
// ***  of accuracy.
// ***
// *****
//
//          tan_82s computes tan(pi*x/4)
//
// Accurate to about 8.2 decimal digits over the range [0, pi/4].
// The input argument is in radians. Note that the function
// computes tan(pi*x/4), NOT tan(x); it's up to the range
// reduction algorithm that calls this to scale things properly.
//
// Algorithm:
//          tan(x)= x(c1 + c2*x**2)/(c3 + c4*x**2 + x**4)
//
double tan_82s(double x)
{
    const double c1= 211.849369664121;
    const double c2=- 12.5288887278448 ;
    const double c3= 269.7350131214121;
    const double c4=- 71.4145309347748;
    double x2;                                // The input argument squared

    x2=x * x;
    return (x*(c1 + c2 * x2)/(c3 + x2*(c4 + x2)));
}
```

tan_82 computes the tangent of $p/4 \cdot x$ to about 8.2 digits of accuracy. Use the range reduction code to translate the argument to 0 to $p/4$, and of course to compensate for the peculiar “ $p/4$ ” bias required by this routine. Note that variables are declared as “double”. The graphed errors are percentage error, not absolute.



```
// *****
// ***
// ***   Routines to compute tangent to 14 digits
// ***   of accuracy.
// ***
// *****
//
//           tan_14s computes tan(pi*x/4)
//
// Accurate to about 14 decimal digits over the range [0, pi/4].
// The input argument is in radians. Note that the function
// computes tan(pi*x/4), NOT tan(x); it's up to the range
// reduction algorithm that calls this to scale things properly.
//
// Algorithm:
//           tan(x)= x(c1 + c2*x**2 + c3*x**4)/(c4 + c5*x**2 + c6*x**4 + x**6)
//
double tan_14s(double x)
{
    const double c1=-34287.4662577359568109624;
    const double c2= 2566.7175462315050423295;
    const double c3=- 26.5366371951731325438;
    const double c4=-43656.1579281292375769579;
    const double c5= 12244.4839556747426927793;
    const double c6=- 336.611376245464339493;
    double x2;
    // The input argument squared

    x2=x * x;
    return (x*(c1 + x2*(c2 + x2*c3))/(c4 + x2*(c5 + x2*(c6 + x2))));
}
```

tan_14l computes the tangent of $p/4 \cdot x$ to about 14.1 digits of accuracy. Use the range reduction code to translate the argument to 0 to $p/4$, and of course to compensate for the peculiar “ $p/4$ ” bias required by this routine. Note that variables are declared as “double”. The graphed errors are percentage error, not absolute.

Higher Precision Tangents

Given a large enough polynomial there's no limit to the possible accuracy. A few more algorithms are listed here. These are all valid for the range of 0 to $\pi/4$, and all should use the previous range reduction algorithm to change any angle into one within this range. All take an input argument in radians, though it is expected to be mangled by the $\pi/4$ factor. The prior range reducer will correct for this.

No graphs are included because these exceed the accuracy of the typical compiler's built-in cosine function... so there's nothing to plot the data against.

Note that C's `double` type on most computers carries about 15 digits of precision. So for these algorithms, especially for the 20.2 and 23.1 digit versions, you'll need to use a data type that offers more bits. Some C's support a `long double`. But check the manual carefully! Microsoft's Visual C++, for instance, while it does support the `long double` keyword, converts all of these to `double`.

Accurate to about 20.3 digits over the range of 0 to $\pi/4$:

```
c1= 10881241.46289544215469695742
c2=- 895306.0870564145957447087575
c2= 14181.99563014366386894487566
c3=- 45.63638305432707847378129653
c4= 13854426.92637036839270054048
c5=- 3988641.468163077300701338784
c6= 135299.4744550023680867559195
c7=- 1014.19757617656429288596025
tan(xp/4)=x(c1 + x2(c2 + x2(c3 + x2*c4)))
/(c5 + x2(c6 + x2(c7 + x2)))
```

Accurate to about 23.6 digits over the range of 0 to $\pi/4$:

```
c1= 4130240.558996024013440146267
c2=- 349781.8562517381616631012487
c3= 6170.317758142494245331944348
c4=- 27.94920941480194872760036319
c5= 0.0175143807040383602666563058
c6= 5258785.647179987798541780825
c7=-1526650.549072940686776259893
c8= 54962.51616062905361152230566
c9=- 497.495460280917265024506937
tan(xp/4)=x(c1 + x2(c2 + x2(c3 + x2(c4 + x2*c5))))
/(c6 + x2(c7 + x2(c8 + x2(c9 + x2))))
```

Arctangent, arcsine and arccosine

The arctangent is the same as the inverse tangent, so $\arctan(\tan(x))=x$. It's often denoted as “ $\text{atan}(x)$ ” or “ $\tan^{-1}(x)$ ”.

In practice the approximations for inverse sine and cosine aren't too useful; mostly we derive these from the arctangent as follows:

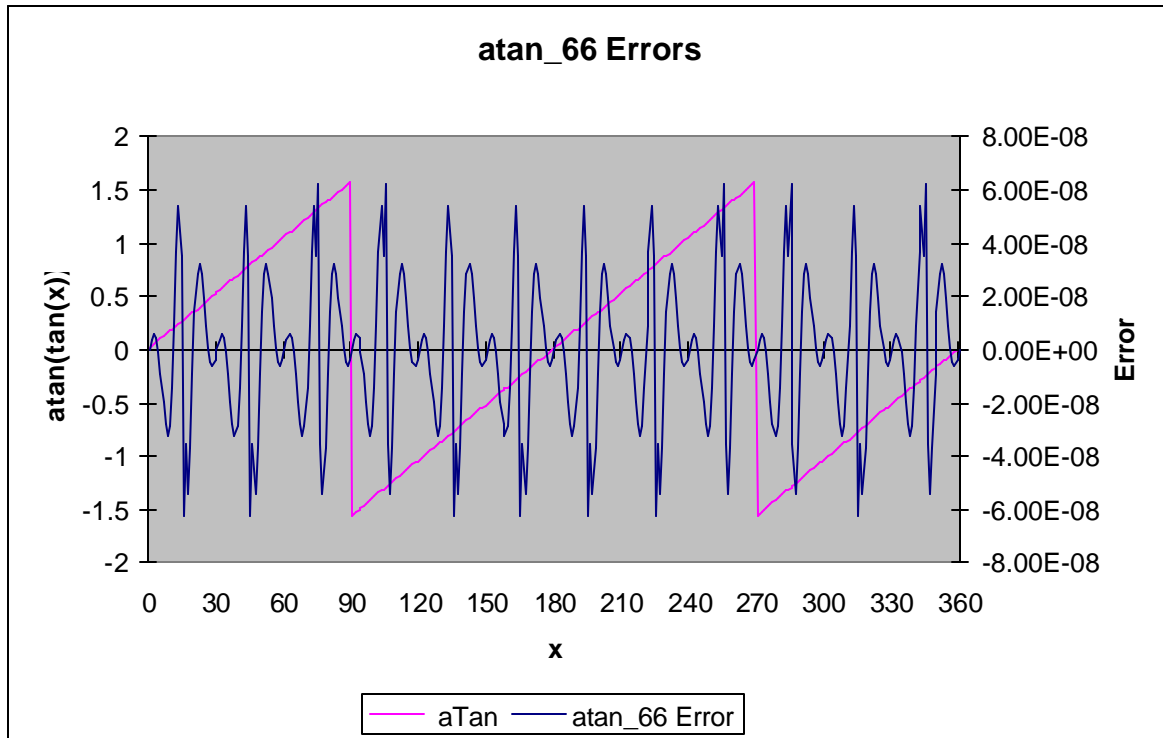
$$\begin{aligned}\text{Arcsine}(x) &= \text{atan}(x/\sqrt{1-x^2}) \\ \text{Arccosine}(x) &= \pi/2 - \text{arcsine}(x) \\ &= \pi/2 - \text{atan}(x/\sqrt{1-x^2})\end{aligned}$$

The approximations are valid for the range of 0 to $\pi/12$. The following code, based on that by Jack Crenshaw in his *Math Toolkit for Real-Time Programming*, reduces the range appropriately:

```
//
// This is the main arctangent approximation "driver"
// It reduces the input argument's range to [0, pi/12],
// and then calls the approximator.
//
//
double atan_66(double x){
double y;                                // return from atan__s function
int complement= FALSE;                  // true if arg was >1
int region= FALSE;                      // true depending on region arg is in
int sign= FALSE;                        // true if arg was < 0

if (x < 0 ){
    x=-x;
    sign=TRUE;                          // arctan(-x)=-arctan(x)
}
if (x > 1.0){
    x=1.0/x;                            // keep arg between 0 and 1
    complement=TRUE;
}
if (x > tantwelfthpi){
    x = (x-tansixthpi)/(1+tansixthpi*x); // reduce arg to under tan(pi/12)
    region=TRUE;
}

y=atan_66s(x);                          // run the approximation
if (region) y+=sixthpi;                  // correct for region we're in
if (complement)y=halfpi-y;              // correct for 1/x if we did that
if (sign)y=-y;                           // correct for negative arg
return (y);
}
```



```
// *****
// ***
// ***  Routines to compute arctangent to 6.6 digits
// ***  of accuracy.
// ***
// *****
//
//          atan_66s computes atan(x)
//
// Accurate to about 6.6 decimal digits over the range [0, pi/12].
//
// Algorithm:
//          atan(x)= x(c1 + c2*x**2)/(c3 + x**2)
//
double atan_66s(double x)
{
    const double c1=1.6867629106;
    const double c2=0.4378497304;
    const double c3=1.6867633134;

    double x2;                                // The input argument squared

    x2=x * x;
    return (x*(c1 + x2*c2)/(c3 + x2));
}
```

atan_66 computes the arctangent to about 6.6 decimal digits of accuracy using a simple rational polynomial. It's input range is 0 to $\pi/12$; use the previous range reduction code.

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Jack Ganssle has written over 300 articles in Embedded Systems Programming, EDN, and other magazines. His three books, *The Art of Programming Embedded Systems*, *The Art of Developing Embedded Systems*, and his most recent, *The Embedded Systems Dictionary* are the industry's standard reference works

Jack lectures internationally at conferences and to businesses, and was this year's keynote speaker at the Embedded Systems Conference. He founded three companies, including one of the largest embedded tool providers. His extensive product development experience forged his unique approach to building better firmware faster.

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- *Manage* features... or miss the schedule!
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Stamp Out Bugs!

- Unhappy truths of ICEs, BDMs, and debuggers.
- *Managing* bugs to get good code fast.
- *Quick* code inspections that keep the schedule on-track.
- Cool ways to find hardware/software glitches.

Managing Real-Time Code

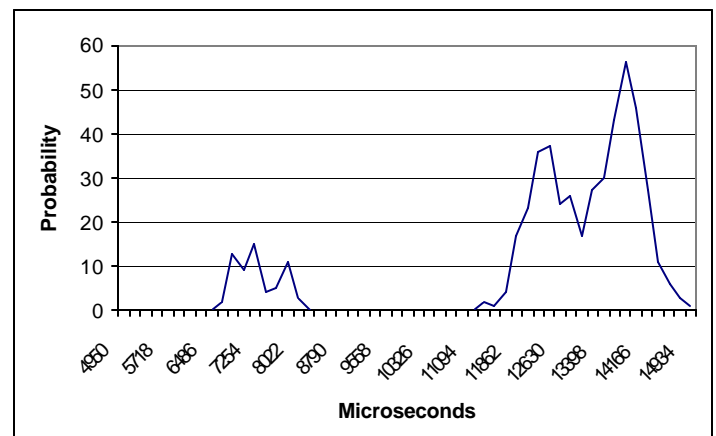
- Design *predictable* real-time code.
- Preventing system performance debacles.
- Troubleshooting and eliminating *erratic crashes*.
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- Understanding high-speed signal problems.
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Thanks for the terrific seminar here at ALSTROM yesterday!
It got rave reviews from a pretty tough crowd.

Cheryl Saks, ALSTROM

Thanks for a valuable, pragmatic, and informative lesson in embedded systems design.
All the attendees thought it was well worth their time.

Craig DeFilippo, Pitney Bowes

I just wanted to thank you again for the great class last week. With no exceptions, all of the feedback from the participants was extremely positive. We look forward to incorporating many of the suggestions and observations into making our work here more efficient and higher quality.

Carol Bateman, INDesign LLC

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