

Week 3

1a) $P \rightarrow (P \rightarrow P)$ \square PL1 ✓

b) ~~$P \rightarrow (P \rightarrow P)$~~

1. $P \rightarrow ((P \rightarrow P) \rightarrow P)$ PL2 ✓

2. $(P \rightarrow ((P \rightarrow P) \rightarrow P)) \rightarrow ((P \rightarrow (P \rightarrow P)) \rightarrow (P \rightarrow P))$ PL2 ✓

3. $(P \rightarrow (P \rightarrow P)) \rightarrow (P \rightarrow P)$ 1, 2, MP ✓

4. $P \rightarrow (P \rightarrow P)$ PL2 ✓

5. $P \rightarrow P$ \square

3, 4, MP ✓

c) 1. $\sim P \rightarrow \sim P$ by using the proof above but with everywhere $\sim P$
 2. $(\sim P \rightarrow \sim P) \rightarrow ((\sim P \rightarrow P) \rightarrow P)$ PL3
 3. $(\sim P \rightarrow P) \rightarrow P$ \square 1, 2, MP ✓

2a) 1. ϕ Ass. ✓
 2. $\phi \rightarrow \psi$ Ass. ✓
 3. ψ 1, 2, MP

So $\phi, \phi \rightarrow \psi \vdash \psi$

$\therefore \phi \vdash (\phi \rightarrow \psi) \rightarrow \psi$ DT

$\vdash \phi \rightarrow ((\phi \rightarrow \psi) \rightarrow \psi)$ \square DT ✓

b) 1. $\phi \rightarrow (\psi \rightarrow \phi)$ PL7
 2. ϕ Ass.
 3. $\psi \rightarrow \phi$ \square 1, 2, MP

or

1. ϕ Ass.
 2. ψ Ass.

So $\phi, \psi \vdash \phi$ ✓

$\phi \vdash \psi \rightarrow \phi$ DT

c) 1. ϕ Ass. ✓

2. $\sim \phi$ Ass. ✓

3. $(\sim \psi \rightarrow \sim \phi) \rightarrow ((\sim \psi \rightarrow \phi) \rightarrow \psi)$ PL3 ✓

4. $\sim \phi \rightarrow (\sim \psi \rightarrow \sim \phi)$ PL2

5. $\sim \psi \rightarrow \sim \phi$ 2, 4, MP

6. $(\sim \psi \rightarrow \phi) \rightarrow \psi$ 3, 5, MP

7. $\phi \rightarrow (\sim \psi \rightarrow \phi)$ PL7

8. $\sim \psi \rightarrow \phi$ 1, 7, MP

9. ψ 6, 8, MP ✓

So $\phi, \sim \phi \vdash \psi$ ✓

$\sim \phi \vdash \phi \rightarrow \psi$ DT

$\vdash \sim \phi \rightarrow (\phi \rightarrow \psi)$ DT

\square

Good stuff

- d) 1. $\neg\neg\phi$ Ass. ✓
 2. $\neg\neg\phi \rightarrow (\neg\phi \rightarrow \neg\neg\phi)$ ✓ PL7
 3. $(\neg\phi \rightarrow \neg\neg\phi) \rightarrow ((\neg\neg\phi \rightarrow \neg\phi) \rightarrow \phi)$ ✓
 4. ~~$(\neg\phi \rightarrow \neg\neg\phi) \rightarrow \neg\neg\phi$~~ 1, 2, MP ✓
 5. $(\neg\phi \rightarrow \neg\phi) \rightarrow \phi$ 3, 4, MP ✓
 6. ϕ 5, MP with 1b) solution
 So $\neg\neg\phi \vdash \phi$
 $\therefore \vdash \neg\neg\phi \rightarrow \phi$ DT ✓

- e) 1. $(\neg\neg\neg\phi \rightarrow \neg\phi) \rightarrow ((\neg\neg\neg\phi \rightarrow \phi) \rightarrow \neg\neg\phi)$ PL3
 2. $\neg\neg\neg\phi \rightarrow \neg\phi$ from 2(d)
 3. $(\neg\neg\neg\phi \rightarrow \phi) \rightarrow \neg\neg\phi$ 1, 2, MP so $\neg\neg\neg\phi \vdash \neg\neg\phi$ DT
 4. $\phi \rightarrow (\neg\neg\neg\phi \rightarrow \phi)$ PL7 so $\phi \vdash \neg\neg\neg\phi \rightarrow \phi$ DT
 5. ~~$\neg\neg\neg\phi$~~ ~~so~~ $\phi \rightarrow \neg\neg\phi$ by 3, 4, cut
 ~~$\neg\neg\neg\phi$~~ $\therefore \phi \vdash \neg\neg\phi$ DT ✓

- 4a) 1. $\Box P \rightarrow P$ T ✓
 2. $\Box\Box P \rightarrow \Box P$ T
 3. $(\Box P \rightarrow P) \rightarrow (\Box\Box P \rightarrow (\Box P \rightarrow P))$ PL7 ✓
 4. $\Box\Box P \rightarrow (\Box P \rightarrow P)$ 2, 3, MP ✓
 5. $(\Box\Box P \rightarrow (\Box P \rightarrow P)) \rightarrow ((\Box\Box P \rightarrow \Box P) \rightarrow (\Box\Box P \rightarrow P))$ PL2 ✓
 6. $(\Box\Box P \rightarrow \Box P) \rightarrow (\Box\Box P \rightarrow P)$ 4, 5, MP
 7. $\Box\Box P \rightarrow P$ 1, 6, MP ✓
 Or just use the PL tautology "syllogism"

- b) 1. $\Box\neg P \rightarrow \neg P$ T
 2. $(\psi \rightarrow \neg\phi) \rightarrow (\phi \rightarrow \neg\psi)$ PL i.e. contrapositive.
 3. $P \rightarrow \neg\Box\neg P$ 1, 2, MP
 You don't want to write the schema (but rather the relevant instance of it) in a proof

- c) Goal: show $\vdash_{S4} \neg\Box\neg(\neg\Box\neg\neg\Box\neg P \rightarrow \neg\Box\neg P)$
 1. $\Box\neg P \rightarrow \Box\Box\neg P$ S4 ✓
 2. $\Box\neg P \rightarrow \neg\neg\Box\neg P$ PL ✓
 3. $\Box\Box\neg P \rightarrow \Box\neg\neg\Box\neg P$ 2, K, NEC, MP ✓
 4. $\Box\neg P \rightarrow \Box\neg\neg\Box\neg P$ 2, 3, MP ✓
 5. $\Box\neg\neg\Box\neg P \rightarrow \Box\neg\neg\Box\neg\neg\Box\neg P$ (repeat 1-4) ✓
 Yup, or just say "Becker Box"

6. $\Box \sim P \rightarrow \Box \sim \neg \Box \sim \sim \Box \sim P$ 4, 5, PL
 7. $\sim \Box \sim \sim \Box \sim \sim \Box \sim P \rightarrow \sim \Box \sim P$ 6, PL

Good

d) Goal: show $\Box P \rightarrow \Box \sim \Box \sim \Box P$

1. $\Box \sim \Box P \rightarrow \sim \Box P$ T ✓
 2. $\Box P \rightarrow \sim \Box \sim \Box P$ 1, PL ✓
~~3. $\Box \Box P \rightarrow \Box \sim \Box \sim \Box P$~~
~~4. $\Box \Box P \rightarrow \Box \sim \Box \sim \Box P$~~
 3. $\Box \Box P \rightarrow \Box \sim \Box \sim \Box P$ 2, K, NEC, MP ✓
 4. $\Box P \rightarrow \Box \Box P$ S4 ✓
 5. $\Box P \rightarrow \Box \sim \Box \sim \Box P$ 3, 4, PL ✓

Contraposition (btw it helps an exam marker a lot to mention the PL tautology you're appealing to!)

e) ~~$\Box(P \rightarrow Q) \rightarrow \Box(Q \rightarrow P)$~~ PL

1. $(\Box P \rightarrow \Box Q) \vee (\Box Q \rightarrow \Box P)$ PL
 2. $(\Box P \rightarrow \Box Q) \rightarrow \Box(\Box P \rightarrow \Box Q)$
 3. $(\Box Q \rightarrow \Box P) \rightarrow \Box(\Box Q \rightarrow \Box P)$
 4. $\Box(\Box Q \rightarrow \Box P) \vee \Box(\Box P \rightarrow \Box Q)$ 2, 3, PL

not totally sure how to justify this but I think it's true.

~~is this 2012? It is strictly weaker than S5.~~

Yeah it's a tough one, isn't it! Hope you caught the solution we covered in class -- drop me an email if you missed it

5a) ~~Goal: show $\Box(\phi \wedge \psi) \rightarrow \Box \phi$~~

1. $\Box(\phi \wedge \psi) \rightarrow \Box \phi$ PL, Becker
 2. $\Box(\phi \wedge \psi) \rightarrow \Box \psi$ "
 3. $\Box(\phi \wedge \psi) \rightarrow \Box \phi \wedge \Box \psi$ 1, 2, PL

~~5b) $\sim(\Box \phi \wedge \Box \psi) \rightarrow \sim \Box(\phi \wedge \psi)$ 3, PL~~

1. $\Box((\phi \wedge \psi) \rightarrow \phi)$ PL, NEC ✓
 2. $\Box((\phi \wedge \psi) \rightarrow \psi)$ "
 3. $\Diamond(\phi \wedge \psi) \rightarrow \Diamond \phi$ 2, K ✓
 4. $\Diamond(\phi \wedge \psi) \rightarrow \Diamond \psi$ 2, K ✓
 5. $\Diamond(\phi \wedge \psi) \rightarrow \Diamond \phi \wedge \Diamond \psi$ 3, 4, PL ✓

Yup, that'll do!

Nice

- ~~1. $\Box(\phi \rightarrow \psi) \rightarrow (\sim \psi \rightarrow \sim \Box \phi)$ NEC, PL
 2. $\Diamond(\phi \rightarrow \psi) \rightarrow \Diamond(\sim \psi \rightarrow \sim \Box \phi)$ K, MP, 1
 3. $\sim \sim \Box \sim(\sim \psi \rightarrow \sim \Box \phi) \rightarrow \sim \sim \Box \sim(\phi \rightarrow \psi)$ 2, PL~~

1. $\Box(\neg\psi \rightarrow \neg\phi) \rightarrow (\Diamond\neg\psi \rightarrow \Diamond\neg\phi)$ $K\Diamond$ ✓
2. $\Box(\phi \rightarrow \psi) \rightarrow \Box(\neg\Diamond\neg\phi \rightarrow \neg\Diamond\neg\psi)$ PL
3. $\Box(\phi \rightarrow \psi) \rightarrow (\neg\neg\Box\neg\neg\phi \rightarrow \neg\neg\Box\neg\neg\psi)$ $\Diamond df$
4. $\Box(\phi \rightarrow \psi) \rightarrow (\Box\phi \rightarrow \Box\psi)$ PL

Are we sure about that?
It looks like you're using the distribution of Box over \neg here, together with the definition of Box. But you can't help yourself to K in this question!

b) ~~(\Rightarrow) If $\vdash_K \phi$ then $\vdash_{K\Diamond} \phi$~~

Goal: show that $\vdash_K \phi$ iff $\vdash_{K\Diamond} \phi$.

(\Rightarrow) If some wff is provable in K, then take any parts of that proof using the (K) axiom, insert a subproof as in ii, then continue with the proofs in $K\Diamond$.

(\Leftarrow) If some wff is provable in $K\Diamond$, then take any parts of that proof using $(K\Diamond)$ or $(\Diamond df)$, inserting above a subproof of the axiom in K. Then continue with the proof in K.

You have roughly the right idea! In an exam I might be a bit more explicit and write out things like "assume there exists a proof of ϕ from [system], then every line of this proof is either an axiom or derived from prior lines via MP or NEC,.... you know how it goes.