

# Robert Merton week 7. Formalisation in predicate.

7.1 i)  $\exists$ -elim done incorrectly on second line of both branches.

$$\begin{array}{c}
 \frac{[Pa \wedge Qa]^1}{Pa} \\
 \frac{\frac{\frac{[Pa \wedge Qa]^1}{Qa}}{\exists x Qx}}{\exists x Px \wedge \exists x Qx} \quad \frac{\frac{[Pa \wedge Qa]^1}{Qa}}{\exists x (Px \wedge Qx)}_1 \\
 \exists x Px \wedge \exists x Qx \quad \square
 \end{array}$$

(a' appears in undischarged assumptions)

ii)  $\exists$ -elim done incorrectly on second line (the constant 'a' can't appear in the <sup>initial</sup> assumption and the formula you conclude after  $\exists$ -elim)

Counter-example:  $D_A = \{0, 1\}$   
 $R_A^2 = \{ \langle 0, 0 \rangle, \langle 1, 1 \rangle \}$

~~Then note that~~  
 $\models \exists y \forall x Rxy \models T$  iff  $\models \forall x Rxy \models F$  under some assignment  $\alpha$  over  $A$

Case 1:  $\models y \models 0$ . But  $\langle 1, 0 \rangle \notin R_A^2$  so  $\models Rxy \models F$  so  $\models \forall x Rxy \models F$

Case 2:  $\models y \models 1$ .  $\langle 0, 1 \rangle \notin R_A^2$  so  $\models Rxy \models F$  so  $\models \forall x Rxy \models F$

So there is no assignment  $\alpha$  s.t.  $\models \forall x Rxy \models T$  so  $\models \exists y \forall x Rxy \models F$   
 $\forall x \exists y Rxy \not\models \exists y \forall x Rxy$ , since  $\models \forall x \exists y Rxy \models T$  (by inspection:  $\langle 0, 0 \rangle$  and  $\langle 1, 1 \rangle$ )

iii) There's an assumption of  $Pa$  which has no valid reason to be discharged before  $\forall$ -I in line 3, so the constant 'a' still appears undischarged in the proof of  $Pa \Rightarrow Qa$ .

Counterexample:  $D_A = \{0, 1\}$   $I/A = \{0, 1\}$   $Q/A = \{1\}$

$\models \exists y (Py \Rightarrow Qy) \models T$  since in the assignment  $\alpha$  where  $\models y \models 1$ ,  $\models Qy \models T$  as  $1 \in Q/A$  so  $\models Py \Rightarrow Qy \models T$ .

But  $\models \forall x (Px \Rightarrow Qx) \models F$  since in the assignment  $\alpha$  where  $\models x \models 0$ ,  $\models Px \models T$  and  $\models Qx \models F$  so  $\models Px \Rightarrow Qx \models F$ .

So  $\exists y (Py \Rightarrow Qy) \not\models \forall x (Px \Rightarrow Qx)$ .



7.2i)  $\neg \forall x (Px \rightarrow Qx)$   $P: \dots$  is a book author  $Q: \dots$  is famous

or equally  $\exists x (Px \wedge \neg Qx)$

ii)  $\exists x (Px \wedge Qx)$   $P: \dots$  is a book  $Q: \dots$  is famous

iii)  ~~$\forall x (Px \rightarrow Qx \leftrightarrow Rx)$~~   $\forall x (Px \rightarrow (Qx \leftrightarrow Rx))$  as above,  $m R: \dots$  is well-written

~~$\forall x (Px \wedge Qx \leftrightarrow Rx)$~~

iv)  $\neg P_1 a$   $P_1: \dots$  believes that <sup>not</sup> every book author is famous  $a: \text{Tom}$   
(extensional ambiguity here otherwise?)

7.3i)  $\forall x (Px \rightarrow Qax)$  or  $\exists x (Px \wedge Qax)$   $P: \dots$  is a logician

$Q: \dots$  despises  $a: \text{Ben}$

ii)  $Pab \wedge Pac$   $P^2: \dots$  slanders  $a: \text{Harry}$   $b: \text{Ron}$   
 $c: \text{Harry's parents OR Ron's parents}$

iii)  $\exists x \exists y (Px \wedge Qy \rightarrow Rxy)$  or  $\forall x \forall y (Px \wedge Qy \rightarrow Rxy)$   $P: \dots$  is a student  
 $Q: \dots$  is a tutor  $R: \dots$  is better than  $\dots$

iv)  $\forall x (Px \wedge Qx \rightarrow Rx)$  or  $\forall x (Px \rightarrow Rx \wedge Qx)$

v)  $\forall x (Px \rightarrow Rax)$   $P: \dots$  is an <sup>ast</sup> anchor  
or  $\exists x (Px \wedge Rax)$   ~~$Q: \dots$  is a ~~member~~~~

$R: \dots$  is rich  $Q: \dots$  is German  
 $P: \dots$  buys a house in Munich

$R: \dots$  likes  $a = \text{James}$

vi)  ~~$\forall x \exists y (Py \wedge Qxy)$~~   $\forall x \exists y (Py \wedge Qxy)$   $P: \dots$  is a mistake  
 $Q: \dots$  made  $\dots$

or  $\exists y (Py \wedge \forall x Qxy)$

assume that is domain of discourse

7.4i) (There exists a set such that for all sets, they are <sup>an element</sup> ~~members~~ of that set iff they're not an element of themselves.)

i.e. There is a set containing all sets which don't contain themselves.

ii) There's a set which all sets have iff they don't have themselves.

iii) From 6.5, we've established that  $\exists x \forall y (Ryx \leftrightarrow \neg Ryx) \vdash P$ .

So any sentence  $P$  can be proved from it, so any set  $\Gamma$  containing that as a premiss can be used to prove any sentence  $P$ . So there is no  $P$  s.t.  $\Gamma \not\vdash P$ ,  
so any such set  $\Gamma$  must be syntactically inconsistent.



- ii) Some such sets, like 'the set of all  $x$  such that  $x$  does not contain itself', may entail a contradiction (and thus not exist?)

7.5  $\forall x \exists y P_{xy} \vdash \exists y \forall x P_{xy}$   $P: \dots$  has  $\dots$  as its cause.

• This isn't valid in predicate logic - see 7.1 ii) for a counterexample.

• If we take 'a cause of everything' to mean 'some thing from which a forwards-only causal chain can be drawn to every other thing' then this argument might be logically valid if there is a finite number of objects (otherwise object  $n+1$  could be the cause of  $n$ , etc.). But here what we define terms as matters a lot: what does it mean for  $x$  to 'cause'  $y$ ?

• In English the argument seems not to be logically valid either. If there is a cause of everything, then that must itself also have a cause, ~~th~~ since everything does (as a premiss) and so it can't be the cause of everything.

- Also note that in the reflexive case, it's perfectly possible for everything to have a cause (itself) and there to be no single cause of everything.