

计算语言学 Computational Linguistics

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第二章 自然语言的特点及其计算复杂性

- 结构性
- 无限性(递归性)

This is the cat.

This is the cat that caught the rat.

This is the cat that caught the rat that ate the cheese.

- This is the house
- This is the house that Jack built
- This is the grain that lay in the house that Jack built
- This is the rat that ate the grain that lay in the house that Jack built
- This is the cat that killed the rat that ate the grain that lay in the house that Jack built
- This is the dog that chased the cat that killed the rat that ate the grain that lay in the house that Jack built

Recursive Structures

NP → NP PP The flight to Boston

VP → VP PP departed Miami at noon

Flights to Miami

Flights to Miami from Boston

Flights to Miami from Boston in April

Flights to Miami from Boston in April on Friday

Flights to Miami from Boston in April on Friday under \$300.

Hights to Miami from Boston in April on Hriday under \$300 with lunch.

Conjunctions

 $S \rightarrow S$ and S

NP → NP and NP

VP → VP and VP

歧义性(ambiguity)

Lexical ambiguity
多音字(词) (polyphone)

朝辞白帝彩云间, 千里江陵一日还。 两岸猿声啼不住, 轻舟已过万重山。



TTS for Mandarin

多义词 (polysemy)

同形异义字(词) (homograph)

- "Minute": (1) a unit for measuring time(noun); (2) to make a written record of what is said or decided. during a meeting(verb); (3) tiny(adj)
 - 1a. One minute has sixty seconds.
 - 1b. Part of the job of a secretary is to minute meetings.
 - 1c. There is only minute difference between these pictures.

"编辑"

Structural ambiguity

亚洲语言学会

彩色铅笔盒子

关于鲁迅的书

他讲不清楚。

(句法结构歧义) n+n+n

n+n+n

(句法结构歧义)

prep+n+的+n(句法结构歧义)

v+不+adi (句法结构歧义)

漂亮的姑娘和小伙子 adj+的+n+的+n(句法结构歧义)

小张的处理意见

他在看病。

他借我一本书。

(语义结构歧义)

(语义结构歧义)

(语义结构歧义)

中国队打败了。 中国队被打败了。 中国队打败了对手。

热爱人民的总理 v+n+的+n 咬死 猎人的鸡 咬死 猎人的鸡 咬死 | 猎人的鸡 咬死鸡的 狗 咬死鸡的 | 狗 咬死 猎人的狗 咬死 猎人的 | 狗 咬死 | 猎人的狗

统计性(Markov链)

我爱吃红	

A brute force solution. Shannon proposed an interesting scheme to generate text according to a Markov model of order 1.

To construct [an order 1 model] for example, one opens a book at random and selects a letter at random on the page. This letter is recorded. The book is then opened to another page and one reads until this letter is encountered. The succeeding letter is then recorded. Turning to another page this second letter is searched for and the succeeding letter recorded, etc. It would be interesting if further approximations could be constructed, but the labor involved becomes enormous at the next stage.

- 模糊性 "下半旗"
- 文化因素

"lying on top of a bed in English and Chinese": "in bed" (English) and "在床上" (Chinese)

In Chinese, a bed means the bed frame but certainly not the duvet nor the blanket; whereas in the English sense of bed, pillows and duvets are often considered as part of the bed (that is why the English phrase "to make the bed" means to arrange the sheets and covers neatly on the bed instead of really constructing the bed from wood and/or metal). In this sense, if the preposition meaning "in" in Chinese is used instead, the meaning of the phrase will be distorted. For instance, if "John is in bed." is translated to Chinese word by word, the resulting Chinese sentence will sound odd unless the bed is actually like a box where John hides himself in.

Grammar:
$$G = (V, T, S, P)$$

V: Set of (meta) symbols, or variables

T: Set of terminal symbols

S: Start symbols, from V

P: Set of Production rules

 $P = \{S \rightarrow aSb, S \rightarrow \lambda\}$

$$G \qquad S \to aSb$$

$$S \to \lambda$$

$$G = (V, T, S, P)$$

$$V = \{S\} \qquad T = \{a, b, \lambda\}$$

Sentential Form:

A sentence that contains variables and terminals

Example:

$$S \Rightarrow aSb \Rightarrow aaSbb \Rightarrow aaaSbbb \Rightarrow aaabbb$$

Sentential Forms

sentence

*

We write:

$$S \Rightarrow aaabbb$$

Derivations

Instead of:

$$S \Rightarrow aSb \Rightarrow aaSbb \Rightarrow aaaSbbb \Rightarrow aaabbb$$

In general we write:

$$w_1 \Rightarrow w_n$$

If:
$$w_1 \Rightarrow w_2 \Rightarrow w_3 \Rightarrow \cdots \Rightarrow w_n$$

Language of a Grammar

For a grammar G with start variable S:

$$L(G) = \{w: S \stackrel{*}{\Longrightarrow} w\}$$

String of terminals

Chomsky hierarchy

The Chomsky hierarchy is an ordering of types of grammar according to generality. The classification in fact only depends on the type of grammar rule (rewrite rule) used. The grammar types include:

unrestricted grammars(type 0): rules of the form $\alpha \to \beta$ with no restrictions on the sequence of symbols α and β . context sensitive grammars(type 1): rules of the form $\alpha X\beta \to \alpha \Psi\beta$ where X is a non-terminal symbol, α and β are (possibly empty) sequences of symbols, and Ψ is nonempty sequence of symbols.

context free grammars(type 2): rules of the form $X \rightarrow \alpha$ where X is a single non-terminal symbol, α are nonempty sequence of symbols.

(right) regular grammars(type 3): rules of the form $X \rightarrow a$ and $X \rightarrow aN$ where X and N are nonterminal symbols, and a is a terminal symbol.

(left) regular grammars(type 3): rules of the form $X \rightarrow a$ and $X \rightarrow Na$ where X and N are nonterminal symbols, and a is a terminal symbol.

The Chomsky Hierarchy

 unrestricted or type-0 grammars, generate the recursively enumerable languages, automata equals Turing machines

 context-sensitive grammars, generate the context-sensitive languages, automata equals Linear Bounded Automata

 context-free grammars, generate the context-free languages, automata equals Pushdown Automata

 regular grammars, generate the regular languages, automata equals Finite-State Automata

* A language is *recursively enumerable* if there exists a Turing machine that accepts every string of the language, and does not accept strings that are not in the language.

"Does not accept" is *not* the same as "reject" -- the Turing machine could go into an infinite loop instead, and never get around to either accepting *or* rejecting the string.

The languages generated by unrestricted grammars are precisely the recursively enumerable languages.

* A language is *recursive* if there exists a Turing machine that accepts every string of the language and rejects every string (over the same alphabet) that is not in the language.

Recursively enumerable languages

Recursive languages

Decidable language (definition)

Definition: A language for which membership can be decided by an *algorithm* that halts on all inputs in a finite number of steps --- equivalently, can be recognized by a *Turing machine* that halts for all inputs.

Generative capacity of grammars

Any grammar G that is a type n (> 0) grammar is also a type n-1 grammar.

Any language that is a type n > 0 language is also a type n-1 language.

L0: (right) regular grammar

$$S \rightarrow a S1$$
 $S \rightarrow d$ $S1 \rightarrow d$ $S3 \rightarrow d$ $S \rightarrow b S2$ $S1 \rightarrow b S2$ $S2 \rightarrow c S3$ $S1 \rightarrow c S3$ $S2 \rightarrow d$

$$L_1 = \{a^n b^n\}, n \ge 1$$

ab, aabb, aaabbb,.....

$$L_2 = \{\alpha\alpha^*\}$$

aa, bb, abba, aaaa, bbbb, aabbaa, abbbba,... 镜象语言 $L_3 = \{\alpha\alpha\}$

aa, bb, abab,aaaa, bbbb, aabaab, abbabb, ...

L1不能用RG生成,可用CFG生成:

 $S \rightarrow a b$ $S \rightarrow a S b$

ab, aabb, aaabbb,.....

L2:不能用RG生成,可用CFG生成:

 $S \rightarrow a S a S \rightarrow b S b$

 $S \rightarrow a a$ $S \rightarrow b b$

aa, bb, abba, aaaa, bbbb, aabbaa, abbbba,... 镜象语言

L3不能用CFG生成,可用CSG生成:

 $S \rightarrow a S$ $S \rightarrow b S$

 $\alpha S \rightarrow \alpha \alpha$

 α 是集合 $\{a,b\}$ 上的任意非空符号串

aa, bb, abab,aaaa, bbbb, aabaab, abbabb, ...

●自然语言不能用RG完全生成

```
The rat disappeared.

a a

The rat the cat caught disappeared.

a b b a

The rat the cat the dog chased caught disappeared.

a b c c b a
```

L2

Consider the following set of English sentences (strings)

$$-S = If S_1 then S_2$$

$$-S = Either S_3$$
, or S_4

-S = The man who said S_5 is arriving today

 Map If, then → a and either, or → b. This results in strings like abba or abaaba or abbaabba

- ●自然语言不能用CFG完全生成
- (Shieber, 1985) and (Huybregts, 1984) showed this using examples from Swiss-German:

mer	d'chind	em Hans	es huus lönd hälfed aastriiche
we	the children-ACC	Hans-dat	the house-Acc let helped paint
w	а	b	x c d y
	N_1	N_2	N_3 V_1 V_2 V_3

... we let the children help Hans paint the house

P. Postal (1964) 发现, 印第安的Mohawk语中: "我读书" 我书读书 "我喜欢读书" 我书读书喜欢书读书 b a b b a b 我尝到了读书的甜头 我书读书的甜头尝到了书读书的甜头 b a b c d babc d

L3

大姐、二姐、三姐分别是二十、十八和十六岁。 a b c a b c

Language	Automaton	Grammar	Recognition	Dependency
Recursively Enumerable Languages	Turing Machine	Unrestricted Baa → A	Undecidable	Arbitrary
Context- Sensitive Languages	Linear-Bounded	Context- Sensitive At → aA	NP-Complete	Crossing
Context- Free Languages	Pushdown (stack)	Context-Free S → gSc	Polynomial	Nested
Regular Languages	Finite-State Machine	Regular A → cA	Linear	Strictly Local

Natural language 属于CSG, 接近于CFG

CFGs are very **important** because:

- * powerful enough to describe most of the structure in natural languages;
- * restricted enough so that efficient parsers can be built to analyze sentences.

• Chomsky 范式

任何上下文无关语言都能由那样的文法产生,其中所有规则的形式或者是 $U \rightarrow XY$ 或者是 $U \rightarrow T$,这里X,Y,U属于VN, T属于VT.

- ●上下文无关语言的可判定性
- 文法的二义性问题是不可判定的(上下文无关文法)寻找充分条件

DFA vs. NFA

Noam Chomsky
Institute Professor; Professor of Linguistics
Linguistic Theory, Syntax, Semantics,
Philosophy of Language, MIT

http://web.mit.edu/linguistics/www/chomsky.ho me.html

The Chomsky hierarchy is a <u>containment</u> <u>hierarchy</u> of classes of <u>formal grammars</u> that generate <u>formal languages</u>. This hierarchy was described by <u>Noam Chomsky</u> in <u>1956</u>.



(December 7, 1928)

Chomsky has written and lectured widely on linguistics, philosophy, intellectual history, contemporary issues, international affairs and U.S. foreign policy. His works include: Aspects of the Theory of Syntax; Sound Pattern of English (with Morris Halle); Language and Mind; American Power and the New Mandarins; At War with Asia; For Reasons of State; Peace in the Middle East?; Reflections on Language; The Political Economy of Human Rights, Vol. I and II (with E.S. Herman); Rules and Representations; Lectures on Government and Binding; Towards a New Cold War; Radical Priorities; Fateful Triangle; Knowledge of Language; Turning the Tide; Pirates and Emperors; On Power and Ideology; Language and Problems of Knowledge; The Culture of Terrorism; Manufacturing Consent (with E.S. Herman); Necessary Illusions; Deterring Democracy; Year 501; Rethinking Camelot: JFK, the Vietnam War and US Political Culture; Letters from Lexington; World Orders, Old and New; The Minimalist Program; Powers and Prospects; The Common Good; Profit Over People; The New Military Humanism; New Horizons in the Study of Language and Mind; Rogue States; A New Generation Draws the Line; 9-11; and Understanding Power.

乔姆斯基"言语获得装置"(language acquisition device):认为儿童的大脑里有一种天生的"言语获得装置"。这是人类头脑中固有的内在的语法规则。儿童运用这种普遍语法,就很容易掌握这种语言。

1871年, 达尔文首先提出语言是一种本能的理论。"牙牙学语"...

2005年,英国的《展望》(Prospect)和美国的《外交政策》(Foreign Policy)两本杂志联合进行了一次跨大西洋两岸的读者投票,以期选出全球最著名的公众知识分子。共两万余名读者填写了选票,最后生成了一份百人大榜。 乔姆斯基位列头名。

目前人文领域被引次数最高的十位作家之一。超过黑格尔,紧跟马克思、列宁、莎士比亚、《圣经》、亚里士多德、柏拉图和弗洛伊德之后,唯一在世

2001年 "9·11"事件发生, 当月, 他的《9·11》一书 便告上市。此后在美国主流媒体上,很难再见到他 , 主流知识分子——无论左右, 往往也与他这样的 好斗者保持距离。他的声音更多是通过校园演讲、 "油印"小报、海外报刊、互联网,以及小出版社 的出版物达于外界。他在MIT(乔氏从1955年起任 教于此)的讲座,从来都是人满为患,其场面之热 烈, 堪与校园内的摇滚音乐会媲美。讲话时声调不 高,更像是在做学术报告,而非发表反战演说,但 极其雄辩,锋芒毕露又不失机巧