Function Combinators with Ramda

Original Idea comes from Avaq/combinators.js.

Name	#	<u>Haskell</u>	<u>Ramda</u>	<u>Crocks</u>	Functional Signature	Functor m ⇒	Function f, g, h Evaluation
identity	ı	id	identity	identity	a → a		
constant	K	const	always	constant	$a \rightarrow b \rightarrow a$		
eager application ¹	A	(\$)	call		$(() \rightarrow b) \rightarrow b$		
lazy application			thunkify,		$(a \rightarrow b) \rightarrow a \rightarrow (() \rightarrow b)$		
thrush	Т	(&)	аррјуто	аррјуто	$a \rightarrow (a \rightarrow b) \rightarrow b$		
tap			tap	tap	$(a \rightarrow b) \rightarrow a \rightarrow a$		
flip	С	flip	flip	flip	$(a \rightarrow b \rightarrow c) \rightarrow b \rightarrow a \rightarrow c$		f(b, a)
compose	В	(·), fmap ²	map ² , o	composeB	$(b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c$	$(b \rightarrow c) \rightarrow m \ b \rightarrow m \ c$	f(g(a))
substitution	S	ap 2	ap ²	substitution	$(a \rightarrow b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow c$	$m (b \rightarrow c) \rightarrow m b \rightarrow m c$	f(a, g(a))
chain			chain 2		$(a \rightarrow b \rightarrow c) \rightarrow (b \rightarrow a) \rightarrow b \rightarrow c$	$(a \rightarrow m c) \rightarrow m a \rightarrow m c$	f(g(b), b)
duplication	W	join 2	unnest 2		$(a \rightarrow a \rightarrow b) \rightarrow a \rightarrow b$	$m (m b) \rightarrow m b$	f(a, a)
lift			converge,	converge	$(b \rightarrow c \rightarrow d) \rightarrow (a \rightarrow b) \rightarrow (a \rightarrow c) \rightarrow a \rightarrow d$	$(b \rightarrow c \rightarrow d) \rightarrow m \ b \rightarrow m \ c \rightarrow m \ d$	f(g(a), h(a))
useWith			useWith	compose2	$(c \rightarrow d \rightarrow e) \rightarrow (a \rightarrow c) \rightarrow (b \rightarrow d) \rightarrow a \rightarrow b \rightarrow e$		f(g(a), h(b))
psi	Р	on	on	psi	$(b \rightarrow b \rightarrow c) \rightarrow (a \rightarrow b) \rightarrow a \rightarrow a \rightarrow c$		f(g(a1), g(a2))
"compose inner"			4		$(a \rightarrow d \rightarrow c) \rightarrow (b \rightarrow d) \rightarrow (a \rightarrow b \rightarrow c)$	Functor $n \Rightarrow m (d \rightarrow c) \rightarrow n d \rightarrow m n c$	f(a, g(b))
"lift partial"			map,		$(d \rightarrow b \rightarrow c) \rightarrow (a \rightarrow d) \rightarrow (a \rightarrow b \rightarrow c)$	Functor $n \Rightarrow (d \rightarrow n c) \rightarrow m d \rightarrow m n c$	f(g(a), b)
unit				unit	() → undefined		

¹⁾ The A-combinator can be implemented as an alias of the I-combinator. Its implementation in Haskell exists because the infix nature gives it some utility. Its implementation in Ramda exists because it is overloaded with additional functionality.

²⁾ Algebras like ap have different implementations for different types. They work like Function combinators only for Function inputs.

^{3) @}visisoft/staticland/function/map

⁴⁾ Using map 3 can be implemented as (mfd2c, nd) => map(fd2c)(nd))(mfd2c)