

Exercise 21: (Not a law)

a

$$U_F = \{z\}, P^F = \{(x|x \in U_F)\}$$

$$U_G = \{y\}, P^G = \{(x|x \in U_G)\}$$

The only case in which we found this true was if the U's of F and G differ and $\forall x(F \vee G)$ uses a union of both U_F and U_G .

b

$$U_F = \mathbb{N}, P^F = \{(x|x \in \mathbb{N}, x = 3)\}$$

$$U_G = \mathbb{N}, P^G = \{(x|x \in \mathbb{N}, x = 4)\}$$

Both structures hold individually, but not when combined.

Exercise 22: (Equivalence vs consequence)

No Solution

Exercise 23: (Blue-eyed vampires)

We model our this Problem as follows: As convention, we will give persons truth-values, just like atomic formulae, where elves will be true (1) and vampires will be false (0). Similarly, we will give eye colors truth values, where blue eyes will be true (1) and brown eyes will be false (0).

Every person C has an eye color E and a parent P. Thus:

$$F_1 = (\neg E \wedge \neg P) \Rightarrow \neg C$$

Furthermore:

$$F_2 = P \Rightarrow E$$

And lastly:

$$F_3 = \neg E \Rightarrow \neg C$$

We will now bring $F_1 \wedge F_2 \wedge \neg F_3$ in CNF to show unsatisfiability by resolution:

$$\begin{aligned} F_1 \wedge F_2 \wedge \neg F_3 &= ((\neg E \wedge \neg P) \Rightarrow \neg C) \wedge (P \Rightarrow E) \wedge (\neg E \Rightarrow \neg C) \\ &= (\neg(\neg E \wedge \neg P) \vee \neg C) \wedge (\neg P \vee E) \wedge \neg(E \vee \neg C) \\ &= (E \vee P \vee \neg C) \wedge (\neg P \vee E) \wedge \neg E \wedge C \end{aligned}$$

Resolution:

$$\begin{aligned} Res^1(F_1 \wedge F_2 \wedge F_3) &= \{\{E, P, \neg C\}, \{\neg P, E\}, \{\neg E\}, \{C\}\} \\ Res^2(F_1 \wedge F_2 \wedge F_3) &= Res^1(F_1 \wedge F_2 \wedge F_3) \cup \{\{E, \neg C\}, \{\neg P\}, \{E, P\}\} \\ Res^3(F_1 \wedge F_2 \wedge F_3) &= Res^2(F_1 \wedge F_2 \wedge F_3) \cup \{\{E\}, \{\square\}\} \end{aligned}$$

$Res^3(F_1 \wedge F_2 \wedge F_3)$ contains \square and by Theorem 1.9 we now can conclude that $F_1 \wedge F_2 \wedge F_3$ is unsatisfiable.

Exercise 24: (Skolem normal form)

1)

$$\begin{aligned}
 F &= \neg \forall w P(a, w) \vee \forall x \exists y \forall z Q(f(x), y) \\
 \Leftrightarrow F &= \exists w \neg P(a, w) \vee \forall x \exists y \forall z Q(f(x), y) \\
 \Leftrightarrow F &= \exists w \forall x \exists y \forall z (\neg P(a, w) \vee Q(f(x), y)) \\
 \Leftrightarrow F &= \forall x \forall z (\neg P(a, w') \vee Q(f(x), y'))
 \end{aligned}$$

2)

$$\begin{aligned}
 G &= (\neg \exists x (P(x, z) \vee \forall y Q(x, f(y)))) \vee \forall y P(g(x, y), z) \\
 \Leftrightarrow G &= (\forall x \neg (P(x, z) \vee \forall y Q(x, f(y)))) \vee \forall q P(g(p, q), z) \\
 \Leftrightarrow G &= (\forall x (\neg P(x, z) \wedge \neg \forall y Q(x, f(y)))) \vee \forall q P(g(p, q), z) \\
 \Leftrightarrow G &= (\forall x (\neg P(x, z) \wedge \exists y \neg Q(x, f(y)))) \vee \forall q P(g(p, q), z) \\
 \Leftrightarrow G &= \forall x \exists y \forall q ((\neg P(x, z) \wedge \neg Q(x, f(y))) \vee P(g(p, q), z)) \\
 \Leftrightarrow G &= \forall x \forall q ((\neg P(x, z) \wedge \neg Q(x, f(y'))) \vee P(g(p, q), z))
 \end{aligned}$$

3)

$$\begin{aligned}
 H &= \forall z \exists y (P(x, g(y), z) \vee \neg \forall x Q(x)) \wedge \neg \forall z \exists x \neg R(f(x, z), z) \\
 \Leftrightarrow H &= \forall z \exists y (P(x, g(y), z) \vee \neg \forall s Q(s)) \wedge \neg \forall v \exists w \neg R(f(w, v), v) \\
 \Leftrightarrow H &= \forall z \exists y \exists s (P(x, g(y), z) \vee \neg Q(s)) \wedge \exists v \forall w R(f(w, v), v) \\
 \Leftrightarrow H &= \forall z \exists y \exists s \exists v \forall w ((P(x, g(y), z) \vee \neg Q(s)) \wedge R(f(w, v), v)) \\
 \Leftrightarrow H &= \forall z \forall w ((P(x, g(y'), z) \vee \neg Q(s')) \wedge R(f(w, v'), v'))
 \end{aligned}$$