#### Speeding up Association rules

Dynamic Hashing and Pruning technique

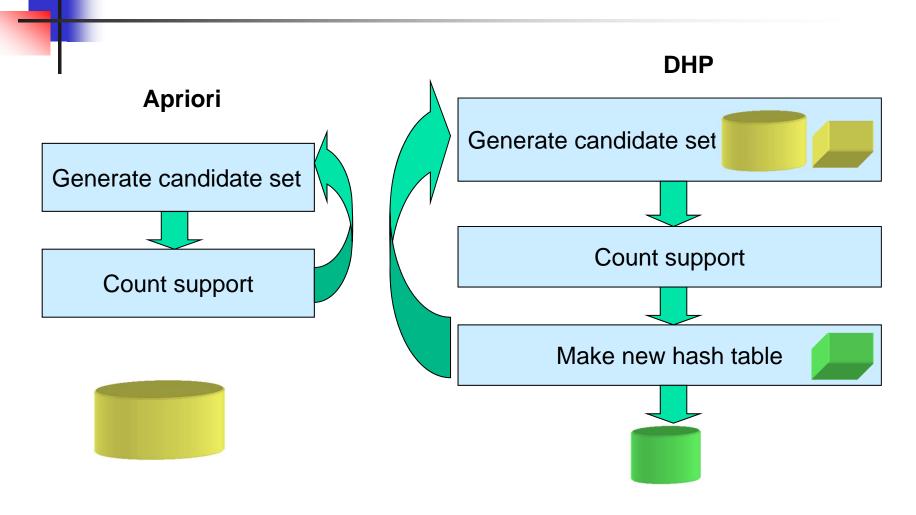
#### DHP: Reduce the Number of Candidates

- A k-itemset whose corresponding hashing bucket count is below the threshold cannot be frequent
  - Candidates: a, b, c, d, e
  - Hash entries: {ab, ad, ae} {bd, be, de} ...
  - Frequent 1-itemset: a, b, d, e
  - ab is not a candidate 2-itemset if the sum of count of {ab, ad, ae} is below support threshold
- J. Park, M. Chen, and P. Yu. An effective hash-based algorithm for mining association rules. In SIGMOD'95

## Still challenging, the niche for DHP

- DHP ( Park '95 ): Dynamic Hashing and Pruning
- Candidate large 2-itemsets are huge.
  - DHP: trim them using hashing
- Transaction database is huge that one scan per iteration is costly
  - DHP: prune both number of transactions and number of items in each transaction after each iteration

#### How does it look like?



#### Reducing Number of Comparisons

#### Candidate counting:

- Scan the database of transactions to determine the support of each candidate itemset
- To reduce the number of comparisons, store the candidates in a hash structure
  - Instead of matching each transaction against every candidate, match it against candidates contained in the hashed buckets

**Transactions** 

TID	Items	
1	Bread, Milk	
2	Bread, Diaper, Beer, Eggs	
3	Milk, Diaper, Beer, Coke	
4	Bread, Milk, Diaper, Beer	
5	Bread, Milk, Diaper, Coke	

Buckets

Hash Structure

# 4

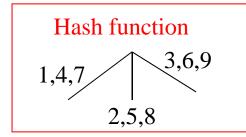
#### Generate Hash Tree

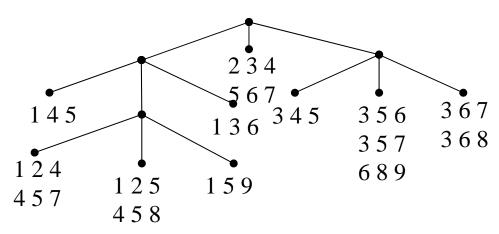
Suppose you have 15 candidate itemsets of length 3:

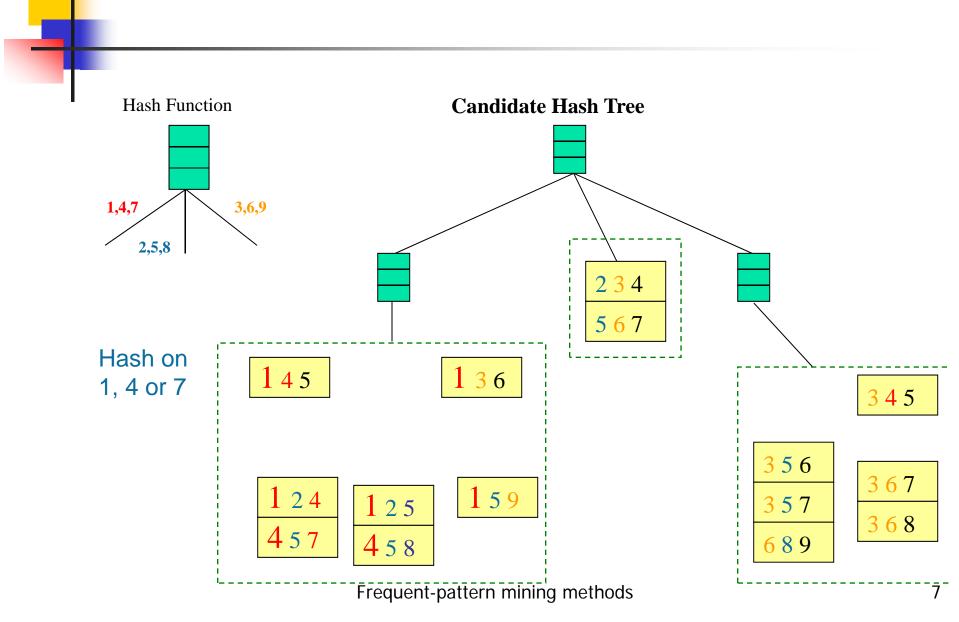
{1 4 5}, {1 2 4}, {4 5 7}, {1 2 5}, {4 5 8}, {1 5 9}, {1 3 6}, {2 3 4}, {5 6 7}, {3 4 5}, {3 5 6}, {3 5 6}, {3 5 6}, {6 8 9}, {3 6 7}, {3 6 8}

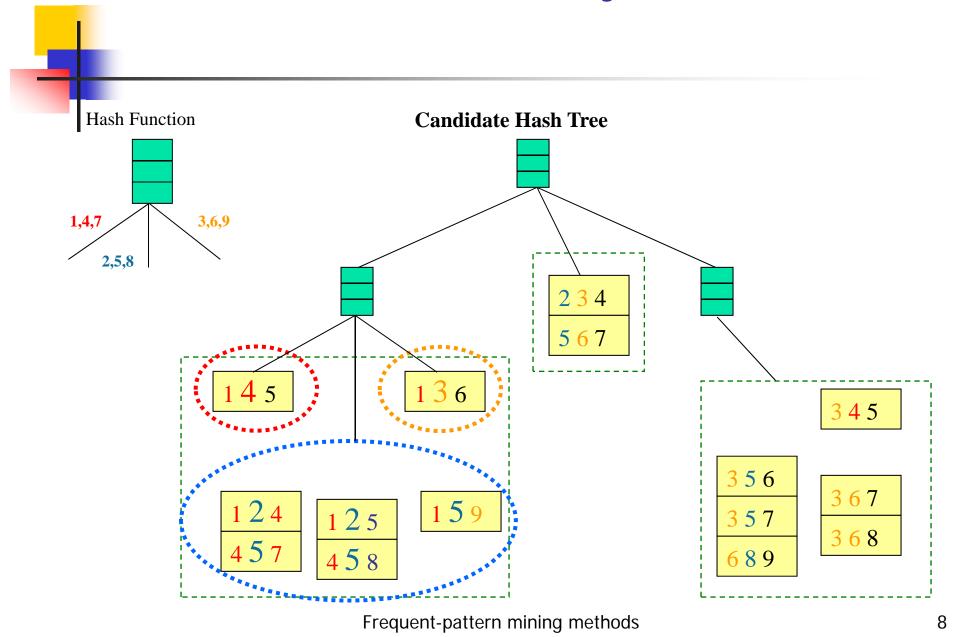
#### You need:

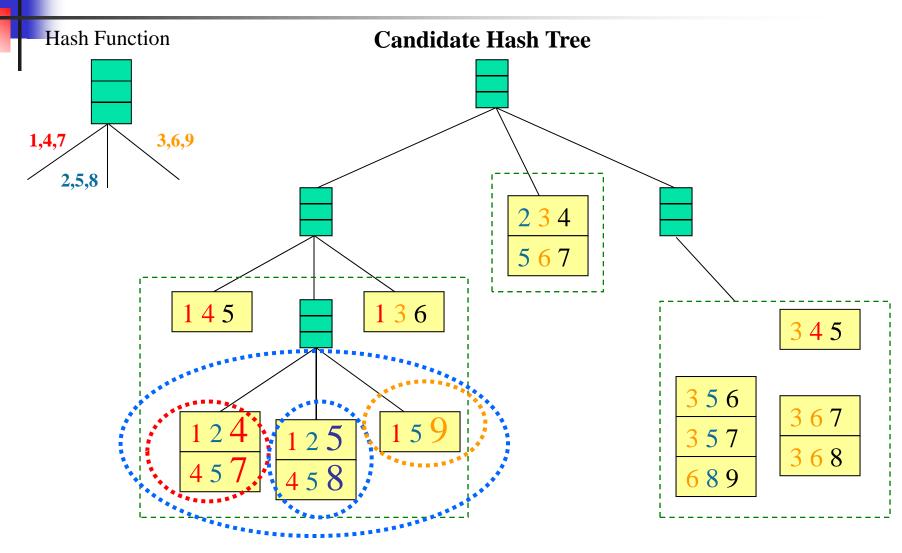
- Hash function
- Max leaf size: max number of itemsets stored in a leaf node (if number of candidate itemsets exceeds max leaf size, split the node)

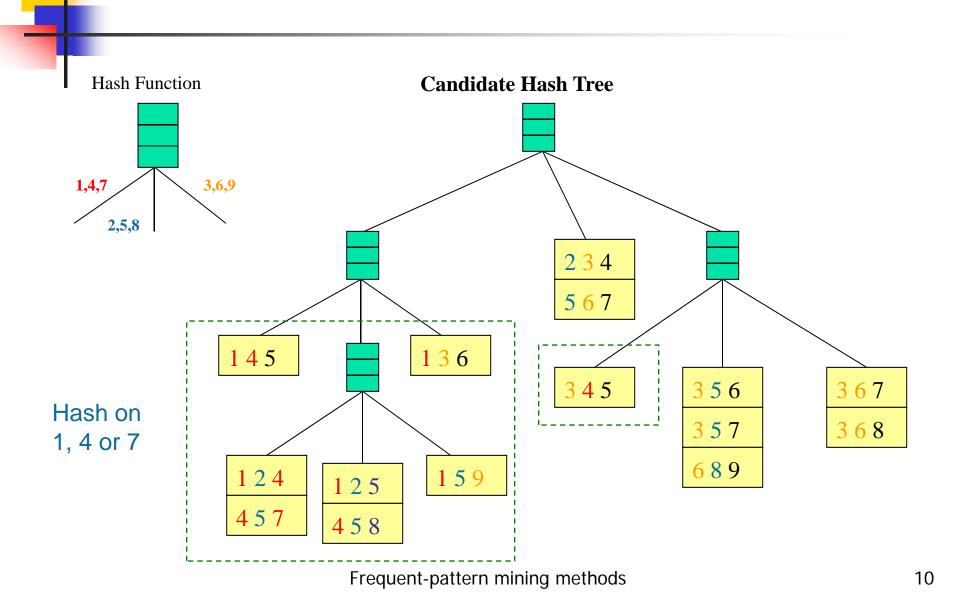


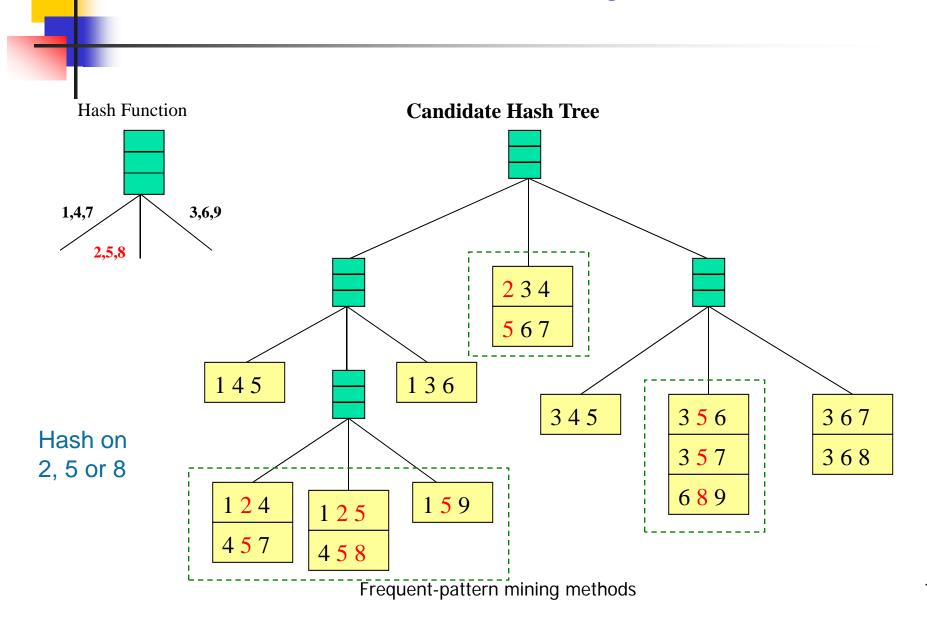


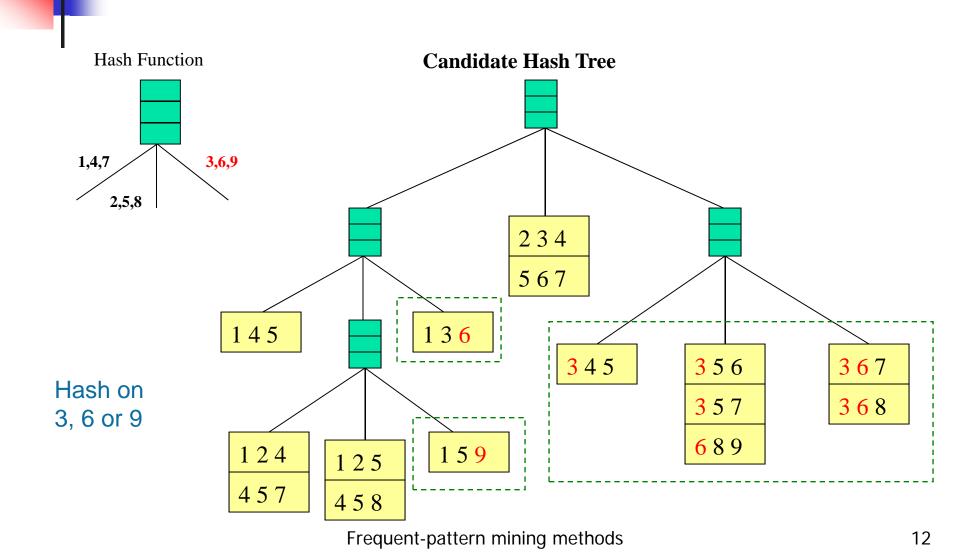












#### How to trim candidate itemsets

In k-iteration, hash all "appearing" k+1 itemsets in a hashtable, count all the occurrences of an itemset in the correspondent bucket.



 In k+1 iteration, examine each of the candidate itemset to see if its correspondent bucket value is above the support (necessary condition)

#### Hash Table Construction

- Consider two items sets, all itesms are numbered as i1, i2, ...in. For any any pair (x, y), has according to
  - Hash function bucket #= h({x y}) = ((order of x)\*10+(order of y)) % 7
- Example:
  - Items = A, B, C, D, E,
  - Order = 1, 2, 3 4, 5,
  - $H({C, E}) = (3*10 + 5) \% 7 = 0$
  - Thus, {C, E} belong to bucket 0.

# Example

TID	Items
100	ACD
200	BCE
300	ABCE
400	B E

Figure 1. An example transaction database

#### Generation of C1 & L1(1st iteration)

Itemset	Sup	
{A}	2	
{B}	3	
{C}	3	
{D}	1	
{E}	3	

C1

Itemset	Sup
{A}	2
{B}	3
{C}	3
{E}	3

L1

#### Hash Table Construction

Find all 2-itemset of each transaction

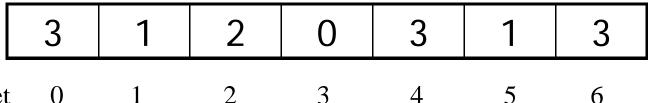
TID	2-itemset		
100	{A C} {A D} {C D}		
200	{B C} {B E} {C E}		
300	{A B} {A C} {A E} {B C} {B E} {C E}		
400 {B E}			

#### Hash Table Construction (2)

Hash function

$$h(\{x y\}) = ((order of x)*10+(order of y)) % 7$$

Hash table



bucket

1

4

5

## C2 Generation (2nd iteration)

L1*L1	# in the bucket	
{A B}	1	
{A C}	3	
{A E}	1	
{B C}	2	
{B E}	3	
{C E}	3	

Resulted C2		
{A C}		
{B C}		
{B E}		
{C E}		

C2 of Apriori
{A B}
{A C}
{A E}
{B C}
{B E}
{C E}
<del>_</del>

#### Effective Database Pruning

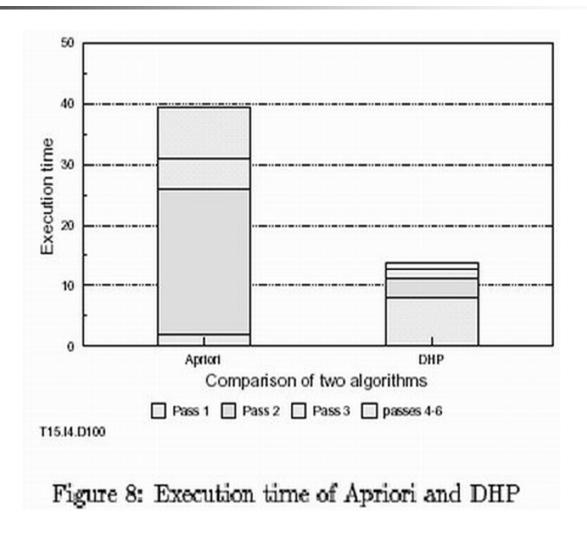
- 'Apriori
- Don't prune database.
- Prune C<sub>k</sub> by support counting on the original database.

- DHP
  - More efficient support counting can be achieved on pruned database.

### Performance Comparison

	Apriori		DHP
	number	number	$D_k$ , $D_k$
$L_1$	820	820	6,700KB, 100,000
$C_2$	335,790	338	6,700KB, 100,000
$L_2$	207	207	
C <sub>3</sub>	618	618	659KB, 20,602
$L_3$	201	201	10 - 00
$C_4$	184	184	546KB, 17,417
$L_4$	98	98	
$C_5$	30	30	332KB, 10,149
$L_5$	23	23	
C <sub>6</sub>	1	1	24KB, 756
$L_6$	1	1	
total time	39.39	13.91	

#### Performance Comparison (2)



#### Conclusion

- Effective hash-based algorithm for the candidate itemset generation
- Two phase transaction database pruning
- Much more efficient (time & space) than Apriori algorithm