

Maximilian W. Gotthardt

Evaluation of protocols for control stage lighting

Bachelor Thesis in Computer Engineering

3 March 2022

Please cite as:

Maximilian W. Gotthardt, "Evaluation of protocols for control stage lighting," Bachelor Thesis (Bachelorarbeit), Institute of Telecommunication Systems, Technische Universität Berlin, Germany, March 2022.

Evaluation of protocols for control stage lighting

Bachelor Thesis in Computer Engineering

vorgelegt von

Maximilian W. Gotthardt

geb. am 13. July 1993
in Berlin

angefertigt in der Fachgruppe

Fachgebiet Telekommunikationsnetze

**Institut für Telekommunikationsnetze
Technische Universität Berlin**

Betreuer: **Anatolij Zubow**
Gutachter: **Falko Dressler**
Thomas Sikora

Abgabe der Arbeit: **3. März 2022**

Erklärung

Ich versichere, dass ich die Arbeit ohne fremde Hilfe und ohne Benutzung anderer als der angegebenen Quellen angefertigt habe und dass die Arbeit in gleicher oder ähnlicher Form noch keiner anderen Prüfungsbehörde vorgelegen hat und von dieser als Teil einer Prüfungsleistung angenommen wurde.

Alle Ausführungen, die wörtlich oder sinngemäß übernommen wurden, sind als solche gekennzeichnet.

(Maximilian W. Gotthardt)

Berlin, den 3 March 2022

Abstract

about 1/2 page:

1. Motivation (Why do I care?)
2. Problem statement (What problem are I trying to solve?)
3. Approach (How did I go about it)
4. Results (What's the answer?)
5. Conclusion (What are the implications of the answer?)

In the field of lighting and stage technology, the challenge of controlling the individual installations, called 'fixtures', quickly and without complications is a recurring one. Established solutions are realized via cables.

However, due to the progress in radio technology, wireless solutions are becoming more and more common. Therefore is often expensive hardware needed. Parallel to this there is a fast growing market around creative and individually developed DIY projects, which have found their own niche. Durch niedrigpreisige While most commercial solutions still rely on expensive and complex wired control, it is particularly suitable for smaller projects to experiment with the new wireless technologies. In this thesis I try to implement a wireless solution, which does not need an IP-Layer using the popular platform ESP and the proprietary protocol ESP-NOW to distribute the light information to each fixture. ESP-Now instead works more like a direct radio communication.

Kurzfassung

- kabellose lösungen werden interessant
- chips werden günstiger
- für kleine projekte leider sehr teure Hardware
- 802.11 wird als standard benutzt
- protokolle wie art-net könnten optimiert werden
- esp plattform bietet interessante möglichkeiten, wegen der geringen kosten der chips und esp-now
- entwicklung einer plattform die esp-now nutzt
- verschiedene ansätze studiert, wie broadcast und unicasts
- mit jeweils unterschiedlichen modifikationen

Contents

Abstract	iii
Kurzfassung	iv
1 Introduction	1
1.1 Motivation/Requirements	1
1.2 Challenges	2
1.3 Problemstatement and Contribution WICHTIG	2
1.4 Thesis Outline	2
2 Related Work	3
2.1 References	3
3 Fundamentals	7
3.1 IEEE 802.11 Specification Family	7
3.1.1 Physical layer	8
3.1.2 Data Link Layer	9
3.1.3 Carrier Sense Multiple Access/Collision Avoidance	10
3.1.3.1 Unicast	11
3.1.3.2 Broadcast	12
3.1.3.3 Multicast	12
3.2 Light protocols	12
3.2.1 DMX-512A	13
3.2.2 Art-Net	14
3.3 ESP Platform	14
3.3.1 ESP32 Hardware	14
3.3.2 ESP-Now	15
3.3.3 ESP-Now vs Art-Net Baseline	17
4 Proposed Approach	18
4.1 Design	18

4.1.1	Slim Unicast	19
4.1.2	Slim Broadcast	22
4.1.3	Rapid Repetition	24
4.1.4	Delayed Repetition	25
4.2	Implementation	25
4.2.1	ESP Programming	25
4.2.2	Collecting measurment results	26
5	Evaluation	29
5.1	Methodic	29
5.2	Protocols under Study	29
5.3	Results	30
6	Conclusion & Discussion	31
	Bibliography	35

The table of contents should fit on one page. When in doubt, adjust the tocdepth counter.

Chapter 1

Introduction

- general motivation for your work, context and goals.
- context: make sure to link where your work fits in
- problem: gap in knowledge, too expensive, too slow, a deficiency, superseded technology
- strategy: the way you will address the problem
- recommended length: 1-2 pages.

In the subject area

====

In the field of lighting and stage technology, the challenge of controlling the individual installations, called 'fixtures', quickly and without complications is a recurring one. Established solutions are realized via cables.

1.1 Motivation/Requirements

- Reliability
- .. and why 100% Reliability is not important (pyrotechnics)
- Lower latency
- Synchronisation
- higher update frequency
- Range

1.2 Challenges

- low cost
- (ESP Platform)
- DIY community

1.3 Problemstatement and Contribution WICHTIG

- open source available on github [link]
- thought-provoking impulse for different approaches
- Protocol auf DL Layer/App Layer Ebene
- Art-Net baseline
- simulativ und experimentel untersucht

1.4 Thesis Outline

- A First Implementation and Evaluation of the IEEE 802.11aa Group Addressed Transmission Service
 - unsosliced Repetition
 - blockack
- Evaluation of Error Control Mechanisms for 802.11b Multicast Transmissions
 - packet loss rate
 - ARQ, FEC
- ESP-NOW communication protocol with ESP32
 - ESP-NOW details
- The Working Principles of ESP32 and Analytical Comparision of using Low-Cost Microcontroller Modules in Embedded Systems Design
 - why the ESP32 is superior over arduino
- Adaptive Cross-Layer Protection Strategies for Robust Scqalable Video Transmissions Over 802.11 WLANs
- Voice Capacity of IEEE 802.11b, 802.11a and 802.11g Wireless LANs

Chapter 2

Related Work

- Wie der und der in Paper so gezeigt hat
- Auch Ding et al haben versucht
- ...
- 10 Paper
- halbe seite

Foo and bar [1] are of equal value. Thus, any can be used.

According to [2]

Wireless solutions for stage lighting are growing fast.

2.1 References

What follows is just a very quick refresher on how to use references. It is not a guide on scientific writing in general, nor copyright and plagiarism in particular. Please refer to an actual guide on technical writing and scientific practices to make sure you understand how, where, and when to cite.

Simply speaking, proper scientific writing has to deal with two closely related (but not identical) concepts:

- a) Copyright
- b) Plagiarism

Do not confuse the need for properly citing your sources as something related to copyright. Questions of (a) copyright or the corresponding national equivalent deal with who has the right to reproduce a certain text excerpt, an image, or something similar. Questions of (b) plagiarism deal with who came up with a certain idea or

insight, e.g., a certain finding, a certain concept, or a certain way of illustrating a concept. By way of analogy, consider a car: after buying a car you have the right to (a) do whatever you want with it, but you still cannot claim that you (b) invented it. Conversely, properly (b) crediting who invented your neighbor's car does not give you the right to (a) use it. Put yet another way, problem (a) is a legal one: to be allowed to publish a scientific work you (or, rather, your publisher) needs to have permission to reproduce it – or suffer legal consequences like heavy fines. Problem (b) is an academic one: claiming someone else's ideas as one's own is plagiarism; similarly, re-selling old ideas as new ones is self-plagiarism. Both incur heavy penalties like exclusion from schools and professional associations or being blacklisted from publishing with scientific outlets for any number of years.

You will need to address both problems in writing your thesis. Problem (a) can be addressed in two ways: First, by creating original content (that is, text or figures) yourself, which is always preferable as this gives you the freedom to present the content your way. Second, by obtaining a license to reproduce content (e.g., by way of buying a license or adhering to the terms of an existing copyleft license). Problem (b) can be addressed in two ways: First, presenting original ideas and insights (as you will do when presenting own results). Second, by clearly pointing out the (primary) source of an idea. The latter is the topic of this section.

In brief, use references whenever you cite from related work (either directly or indirectly), or when you build on related work (this includes their way of illustrating a particular concepts, in text form as well as in the overall design of a figure). Also use references to point a reader to related work. Clearly distinguish between these uses. Make it very clear which part of a statement a reference belongs to. Compare the following three, vastly different uses (where the cited idea appears in **boldface**):

- “**Foo and bar are of equal value. Thus, any can be used.**” [2], [3]
- According to [2] and [3], **foo and bar are of equal value, and any can be used.**

versus

- “**Foo and bar are of equal value**” [2], [3]. Thus, any can be used.
- According to [2] and [3], **foo and bar are of equal value.** From this it follows that any can be used.

versus

- Foo and **bar** [2], [3] are of equal value. Thus, any can be used.
- Foo and **bar** (detailed in [2] and [3]) are of equal value. Thus, any can be used.

versus

- **Foo** [2] and **bar** [3] are of equal value. Thus, any can be used.
- **Foo** (detailed in [2]) and **bar** (detailed in [3]) are of equal value. Thus, any can be used.

Never typeset a reference after the final full stop of a paragraph (or sentence) and expect your reader to figure out which part of the paragraph is an indirect citation and which part is original (i.e., your own) work. When paraphrasing longer passages of text, use an indirect citation. Make sure to clearly point out when you are finished paraphrasing, like so: *According to Akyildiz et al. [2], Foo and Bar can be characterized as follows. They are big. They are bright. The authors further argue that one can be substituted for the other. In the following I will go on to prove that this is not true.*

When citing more than a few pages worth of text, point the reader to the specific part you are referring to in your citation, like so: *In recent years, an increasing number of cyclists are switching from air filled tires to cement filled ones [4, Table IV].*

If a figure or a table is closely based on another one, make sure to cite its source, preferably in its caption, like so: *Figure 1 – the relation of ravens and writing desks (based on [4, Figure 42]).* Be aware that, while there is a well-established convention on how to illustrate a verbatim quote of text (by using quotation marks), there is no well-established convention for indicating that an image was copied verbatim. Thus, when citing a figure or table, you must explicitly state whether it was copied verbatim, ed, or whether it served as inspiration for your own.

Do not cite URLs. Content found there is not peer reviewed and it is likely to change during the lifetime of your work. For pointing a reader to interesting websites, use footnotes – but trust your reader to know how to use a web search engine.

Your text reads nicer if you do not use citations as a substitute for nouns (like this section did). Instead of *The benefits of cement filled tires has been shown by [2]*, consider writing *Akyildiz et al. [2] have shown the benefits of cement filled tires*. The `textcite` command makes this straightforward.

Make sure to read your bibliography section (that is, the typeset list of references) after you are done adding all citations to your text. Does it contain all information needed to uniquely identify to references you used? Do not trust BibTeX files you find on the web: Digital libraries frequently have their contents wrong, are missing information, or are using different field names than your bibliography style expects (leading to missing information in the typeset bibliography). To give a few examples: Check the authors' list (making sure all authors are listed in the same order and in the same way they are listed in the publication). Check the conference location (it's most likely not "New York, New York"). Check the publisher name (many digital libraries use a field that is not typeset by your bibliography style; have a look at the demo

bibliography in this template for how to deal with that). Check the page numbers (many digital libraries put “1–5” here despite the paper starting at a later page – or despite it not having any page numbers to begin with). Check the conference name, put its parts in a logical order, and lose the “in proceedings of” (it’s not “Mobicom, in proceedings of, 1999 series MobiCom99” but “5th ACM International Conference on Mobile Computing and Networking (MobiCom 1999)”).

triple-check all references

Chapter 3

Fundamentals

In this chapter the fundamentals required for understanding the different approaches in this thesis using are explained. This contains basic knowledge of the physical- and data link layer, which are located in the first and second layer of the Open Systems Interconnection (OSI) Model..

reference to table below

In order to understand the upcoming ESP-Now protocol we have to take a look at the Data Link Layer (DL) layer in 802.11. It is the second layer of the OSI model of computer networking illustrated in Table 3.1.

Rewrite introduction in chapter Fundamentals!

3.1 IEEE 802.11 Specification Family

The Institut of Electrical and Electronics Engineers (IEEE) 802 is a family of standards dealing with area networks different kinds.

- 802.11 Wireless Local Area Network (WLAN)
- 802.15.1 Wireless Personal Area Network (WPAN)
- 802.15.4 Low-rate WPAN (LR-WPAN)

Application layer
Presentation layer
Session layer
Network layer
Data Link layer
Physical layer

Table 3.1 – OSI model

- 802.16 Wireless metropolitan area network (WMAN)

For this thesis is the focus set to the 802.11, because of the accessibility and wide functionality. There are two Basic Service Set (BSS) defined:

- Infrastructure BSS

A central element manages the network and all the traffic goes through. Every Station (STA) must always communicate via the Access Point (AP) and never directly - exceptional: Direct Link Mode. An initial association must take place to use this BSS. This is the most common mode a WLAN is used.

- Independent BSS

A network without a central station, where the network topology can flexible change over time. The communication happens directly between the Wireless Endsystems. Efficient routing can become a problem in more complex topologies.

The most common use in 802.11 is the Infrastructure mode, which is commonly used in office and home environments.

3.1.1 Physical layer

In this thesis we could take a brief look into the Physical Network Layer (PHY) of the IEEE 802.11 standard, which is the first layer of the OSI model 3.1. This layer provides mechanical, electrical and other functional tools to activate or deactivate physical connections, maintain them and transmit bits over them. These can be, for example, electrical signals, optical signals (fiber optics, lasers) or electromagnetic waves (wireless networks). There are several complements to the 802.11 standard:

- 802.11b

supports larger bitrates with Direct Sequence Spread Spectrum (DSSS) or Frequency Hopping Spread Spectrum (FHSS) as modulation from 1Mbit/s to 11Mbit/s. It uses the 2.4 GHz ISM band.

- 802.11a and 802.11g

with Orthogonal Frequency Division Multiplexing (OFDM) data rates are increased up to 54 Mbit/s. Where 802.11a is in the 5GHz ISM band 802.11g uses the 2.4GHz ISM band.

- 802.11n

It also uses OFDM and improves with additionally Multiple Input-Multiple Output (MIMO), channel bonding and frame aggregation to increase the bandwidth and decrease the overhead. Using 2.4 GHz and 5GHz ISM band.

- **802.11ac**
Support of wider channel and out of it higher bitrates. It also includes features like Multi-User MIMO. It only uses the 5 GHz ISM band.
- **802.11ax**
Like 802.11ac but with additional use of the 6GHz ISM band and better power control. Also called WiFi6.

In this thesis the rather basic 802.11b is used with a transmission rate of 1Mbit/s.

3.1.2 Data Link Layer

The DL Layer is the second lowest layer of the OSI Model 3.1 and is split in two sublayers. The Locig Link Control (LLC) sublayer which multiplex protocols over the MAC layer while transmitting and to de-multiplex the protocols while receiving. LLC provides the hop-to-hop flow and error control, allows multipoint communication over networks and it also adds frame sequence numbers. But in this thesis we focus on the other data link sublayer.

The Media Access Control (MAC) includes network protocols that regulate how multiple computers share the physical transmission medium they use. Without regulation, collisions and data loss would occur in the shared medium if several WES were to transmit simultaneously. The MAC Protocol Data Unit is additional added inside of the PHY Payload. It contains the MAC Header and encapsulated in it the MAC Service Data Unit (MSDU).

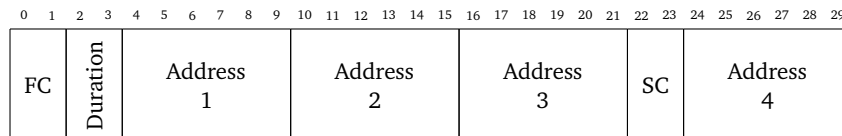


Figure 3.1 – MAC header of a WLAN frame

- **Frame Control Field:** Describes the Type of frame:
 - 00 Manegement Frame
 - 01 Control Frame
 - 10 Data Frame
- **Duration:** Contains the Network Allocation Vector (NAV) value, specifies the transmission time required for the frame. In order to save power to save energy, WES can defer access to the medium for this duration
- **Address fields:** Certain address fields are specified by the relative position of the address field. Not every address field is needed by certain frames.

Payload and FCS are missing

- Basic Service Set Identifier (BSSID)
- Source Address
- Destination Address
- Transmitting STA Address
- Receiving STA Address

- **Sequence Control:** Sequence number of the current frame modulo 4096.
- **MAC Payload:** The actual payload information of the MAC layer. The actual payload can differ, because the headers of the LLC and ip etc. has to be subtracted.
- **Frame Check Frequency:** The sender calculates the checksum for the entire data block and appends it to the end of the block.

802.11ac and later using frame aggregation in order to reduce overhead.

3.1.3 Carrier Sense Multiple Access/Collision Avoidance

shot explanation of CSMA/CD

Multiple Access/Multiplexing: When Signals to/from different users share a common channel using time division methods (TDM/TDMA, CSMA)

eigene Worte

DSSS: Usage of multiple antennas Direct Sequence Spread Spectrum. Spreading of the signal over a given bitsequence PN

Addressing

In a LAN environment, devices are logically separated using 48-bit globally unique MAC addresses: example In IPv4 networks (e.g. Internet), nodes are logically separated using 32-bit globally unique IP addresses: example

Routing

- Routing in a (W)LAN is based on MAC addresses, never IP addresses.
- A router (e.g. integrated with an access point) performs mapping between these two address types:

Address allocation

- MAC addresses are associated with the hardware devices.
- IP addresses can be allocated to (W)LAN devices either on a permanent basis or dynamically from an address pool using the Dynamic Host Configuration Protocol (DHCP).

Mesh networks

are able to relay frames from one device to another. • Provide coverage extension over multiple hops (e.g. Internet access) • Sufficient address information is required to be able to relay data from a source device to the ultimate destination (IP or MAC address). This can be used to extend the range from on Wireless Endsystem (WES) to another WES over some other WES. Since range isn't a critical parameter in this thesis, it hasn't to be further discussed.

Beacon Frames

contain the channel information found during passive scanning. Probe request are used in active scanning.

Backoff:

random time delay to avoid collisions

DIFS/SIFS:

Delay between transmissions used for Carrier-sense multiple access with collision avoidance (CSMA/CA)

3.1.3.1 Unicast

The link layer unicast is used to send data over an single hop to the target WES destination. The link layer of each WES checks the destination MAC address in the link layer header and discards the frame if the destination address does not match its own address.

Unicast is by default reliable.

E.g. the AP wants to transmit a packet to one specific WES

When the Unicast reaches the destination WES an acknowledgement frame is sent back after the Short Inter Frame Spaces (SIFS) + backoff.

If the acknowledgement is not successfully received by the sender, the sender will repeat the transmission for a given number. When the number is exceeded, the packet could not be delivered. If the number is set to zero, the unicast can be considered as non-reliable.

when to explain CS-MA/DC?

example

complete figure UC



Figure 3.2 – Unicast Transmission

3.1.3.2 Broadcast

If a packet should be received from all WES's it can be distributed as broadcast. The MAC address of the destination address in the link layer is set to the common broadcast address, which is ff:ff:ff:ff:ff:ff.

In contrast to unicast, broadcast is not reliable. This is mainly because the packet is addressed to all nodes at the same time, and if link layer acknowledgements would be used, the acknowledgements would be sent by all nodes at the same time, because there is no mechanism in which order acknowledgements should be answered. In addition, the sender of a broadcast does not know how many WESs he is addressing the packet to in the first place. Retransmitting acknowledgements would lead to massive collision and loss of acknowledgements. E.g. management information in a WLAN is sent in a broadcast mode, because it has to reach every WES and isn't worth to be acknowledged.

complete figure BC



Figure 3.3 – Broadcast Transmission

3.1.3.3 Multicast

Multicast explain multicast mac address

When the same packet should be transmitted to multiple WES's, but not to all, multicast can be used. Transmitting the same packet multiple times via unicast is wasteful. There are different approaches to realize acknowledgements for multicasts, they differ mainly by the respective field of application.

examples for multicast ack + related work

Level 2 multicast is often used for large files in audio or video streams, where a big amount of data is distributed and multiple clients listen simultaneously.

complete figure MC



Figure 3.4 – Multicast Transmission

3.2 Light protocols

There are several lighting protocols that are used. The field of application ranges from wired CAN buses over ethernet cables to wireless WLAN networks. To give a short insight, some of the most important protocols are explained below.

3.2.1 DMX-512A

Digital Multiplex (DMX) 512A, is the current industry standard for stage lighting. It is based on Controller Area Network (CAN), therefore it uses wires. Physically is the DMX protocol transmitted over a differential pair of lines using the RS-485 voltage levels. The bus signal is updated with 44Hz. According to the specification are XLR-5 type connectors are to be used.

image der connectoren

Show Hardware e.g. DMX Plug

The endsystems are called fixture because it's most likely a lighting installation which is mounted somewhere, this could be a moving-head, fresnel, spotlight, stroboscope or any other light installation. It could also be a fog machine that emits fog on an appropriate signal.

All devices are daisy chained together visualized in 3.5. The DMX controller is in the begin of each chain. The receiving endsystems, are chained behind each other from output to input. A terminator, specified in the DMX specification, is to be connected to the final output.

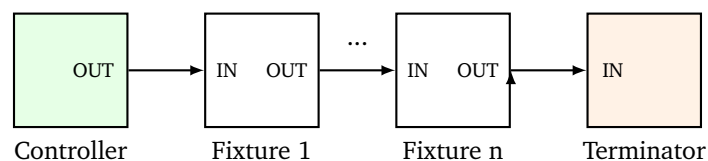


Figure 3.5 – Block Diagram of an DMX Universe

The whole chain is called DMX-Universe and can contain a set of 512 channels. If there is a need for more channels, one needs more DMX universes - each channel consists of one byte. Due to the fact that a DMX universe always has its own bus, starting from a new controller can lead to inconveniences.

A channel in the event technology is used to distinguish between e.g. a WiFi channel. Each endsystem is assigned at least one, but usually several channels. Every endsystem knows which channel is intended for it, this must be preset.

For example: An RGB-LED spotlight could have three channels, one for each color. Due to the resolution of one byte, the individual colors can (theoretically) be controlled in 256 different intensities. If the Channel 0-56 are already used, it could be set to channel 57, 58, 59. Any other free channel range would also be possible, provided it is connected. There exist hardware with automatic-address-assignment.

Image of fixtures distributed on channel

Since DMX is unidirectional it can be assumed that the endsystems generally only receive or forward (daisy chain) control signals sent from the control console.

This is a major limitation of DMX, beside of the rather small universe size. It is also not reliable, the use of fire installations is therefore considered too dangerous.

3.2.2 Art-Net

- 2.4 or 5GHz
- DMX-Like
- related work

Due the limitation of 512 channel for each universe there where protocols implemented using the Art-Net also called Art-Net DMX is

3.3 ESP Platform

Almost every 802.11 capable Microcontroller Unit (MCU) could be picked for this research. But there are several reasons why the ESP Platform from Espressif is a valid choice. There are several chips provided by Espressif with WiFi specifications, these chips are very affordable () and although the ongoing chip crisis (2021) there are easy to get, in contrast of the also very popular Chips from the manufacturer Arduino, which are also more expensive. Espressif supports an own development IDF to flash the chips, with minor tweaks it's also to use the Arduino IDE. However the proprietary protocol ESP-Now which, just supported in the ESP Ecosystem, is discussed below ?? and has promising properties for a solid and fast realisation of a low level protocol.

ESP32 Kosten in €
2021 aufführen? Link?
Datum?

paper about esp above arduino

3.3.1 ESP32 Hardware

The chip ESP32 is quite common in DIY projects around everything from home automation to light installations and can be bought on development boards, which are ready to use. The chip 3.6 is promoted with several features:

- Fast CPU (2 cores at 240 MHz)
- 802.11 b/g/n with up to 150 Mbps (2.4GHz)
- Wifi Multimedia (WMM)
- Immediate Block ACK
- Automatic Beacon monitoring (hardware TSF)
- Virtual Wi-Fi Interfaces

this is cited from: link
zum esp32 datasheet

- Simultaneous support for Infrastructure, SoftAP, and Promiscuous modes
- Bluetooth v4.2 BR/EDR and Bluetooth LE
- Advanced Peripheral Interfaces: GPIO, ADC, DAC, touch sensors, hall sensor, SPI, I2S, I2C, UART, CAN, RMT (TX/RX), Motor/LED PWM

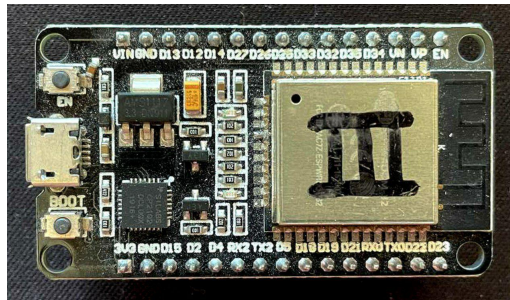


Figure 3.6 – ESP32 Devboard (Devkit V1)

3.3.2 ESP-Now

- How does ESP-Now match the requirements from the motivation
- 250 payload
- gaps in the documentation

ESP-NOW is a proprietary protocol developed by Espressif. ESP-NOW is widely used in smart light, remote controlling, sensor, etc. It is a connectionless protocol, so the WES's are in Ad-Hoc mode instead of STA. It is just supported on the ESP8266, ESP32 and ESP32s, all chipsets from Espressif, but they are compatible with each other. Because of this, an ESP-Chip as gateway is needed to interact from the outside to the ESP-NOW communication.

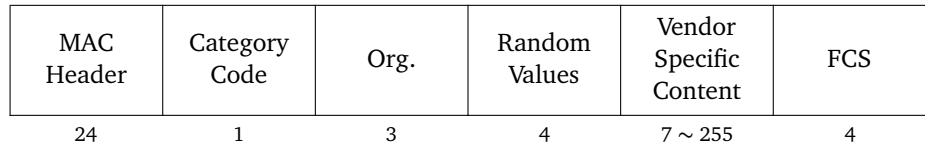
cite ESP documentation website

Through the hardware limitation of the boards it can just be used on the 2.4 GHz frequency band. ESP-NOW allows 10 ESPs for pairing with encryption and up to 20 without encryption. Espressif promises throughput of up to 30MBit/s with a possible range of up to 1km. However *Roberto Pasic [5]* measured a range of the unmodified onboard antenna of the ESP32 and just got a *Roberto Pasic [5] stable communication up to 190m in open field.*

The focus of the ESP-NOW protocol is on low power consumption. A connectionless communication between WES's not only saves energy during the authentication process, Additionally, is the communication the the properties of the ad-hoc mode, direct and not over a second access point. The protocol has a limitation of a limited

payload of 250 byte for each transmission. It also has a much less overhead, which results in shorter airtime, less disturbances and also less power consumption through the antenna (latter is not relevant for this thesis). There is no TCP/IP header to be transmitted. For very small payloads, this offset can become dispropotional.

The default ESP-NOW bit rate is 1 Mbps it uses a channelwidth of 20MHz, there is no double channel (40Mbit/s or higher) used. But e.g. the low energy, high range protocol Long Range Wide Area Network (LoRaWAN) suffers from a too slow throughput for this application.

To understand what ESP-NOW does it needs to take a look to the vendor-specific action frame transmitting ESP-NOW data.  visualized in 3.2.

MAC Header	Category Code	Org.	Random Values	Vendor Specific Content	FCS
24	1	3	4	7 ~ 255	4

Table 3.2 – ESP-NOW Frame Format

- **MAC Header:** As ESP-NOW is connectionless, the MAC header differs from that of standard frames.
- **Category Code:** The Category Code field is set to the value(127) indicating the vendor-specific category.
- **Organization Identifier:** The Organization Identifier contains a unique identifier (0x18fe34), which is the first three bytes of MAC address applied by Espressif.
- **Random Value:** The Random Value field is used to prevent relay attacks.
- **Vendor Specific Content:** The Vendor Specific Content contains vendor-specific fields (table 3.3)
- **Frame Check Sequence:** Used for error correction in layer 2.

Inside of the ESP-NOW frame 3.2 is the vendor specific content visualized in 3.3.

Element ID	Length	Org. Identifier	Type	Version	Body
1	1	3	1	4	7 ~ 250

Table 3.3 – Vendor Specific Action Frame

cite somehow the ESP-NOW documentation pdf: ESP-IDF Programming Guide: ESP-NOW, source: https://docs.espressif.com/projects/esp-idf/en/latest/esp32/api-reference/network/esp_now.html

- **Element ID:** The Element ID field is set to the value (221), indicating the vendor-specific element.
- **Length:** The length is the total length of Organization Identifier, Type, Version and Body.
- **Organization Identifier:** The Organization Identifier contains a unique identifier(0x18fe34), which is the first three bytes of MAC address applied by Espressif.
- **Type:** The Type field is set to the value (4) indicating ESP-NOW.
- **Version:** The Version field is set to the version of ESP-NOW.
- **Body:** The Body contains the ESP-NOW data.

this is cited from espressif manual!!

It is worth to mention, that the vendor specific content 3.2 is allowed to contain up to 255 byte, but the sum over all values in 3.3 if the body would contain the maximum of 250 bytes, leads to a total of 260 bytes. The values are from the documentation of ESP-NOW from Espressif. They also claim, that broadcast is not supported in ESP-NOW, but it is. It seems that the documentation isn't completely finished (or translated).

3.3.3 ESP-Now vs Art-Net Baseline

subsection can be on a wrong position!

remove newpage command

- Network stack diagram
- baseline

ESP-NOW Baseline Artnet should be moved to Design part??

Chapter 4

Proposed Approach

Different approaches are presented and discussed in this chapter using data link unicast or broadcast in different specifications, these were also empirically tested and evaluated in the Chapter 5. At the end of this chapter, the test setup was introduced and it is briefly explained how the chips are programmed and how the measured values were generated and collected.

4.1 Design

The use of high-layer protocols, such as Art-Net Section 3.2.2, in lighting technology involves a considerable overhead. Because the lighting console does not talk directly to the WES, communication must be controlled via an AP, which means that the Internet Protocol (IP) (layer 3) must be used for addressing and User Datagram Protocol (UDP) (layer 4) for transporting the data. They both come with additional headers. Such an overhead can lead to latency, channel congestion and packet loss.

In an ad-hoc network Section 3.1, on the other hand, packets can be sent directly on the MAC (layer 2) Table 4.1. With a payload of a few bytes to each WES, keeping overhead small can be quite important. Complexity problems as often typical in Ad-Hoc networks are not to be assumed, since the controller at the light desk must normally stand in line of sight to the individual WES, from there finally the lighting technician from there must have everything in the range of vision to be able to intervene. One can therefore assume a simple star topology. In the following the ESP-NOW Section 3.3.2 protocol was chosen to distribute the packets low-level, because even if it was not developed for this purpose the specification fits quite well to the requirements.

Fehlt hier nicht ESP-Now?

Art-Net	Slim Application
UDP	
IP	
802.11 DL/Unicast	802.11 DL/UC or BC
802.11* PHY	802.11b/g/n PHY

Table 4.1 – Art-Net Layer compared with Slim Data Link Layer

For the purpose of this analysis, the controller transmits 20 Byte (analogue to 20 DMX channels) to every WESs. Four different metrics are considered, which were discussed in the requirements.

ref to requirements

- **Latency**

What is the latency from commanding the controller to the estimated reaction at the WES e.g. lighting of a light? For the sake of simplicity, delays caused by the microcontroller instruction set, data distribution and control of the light installation are neglected and the focus is placed only on the airtime.

- **Update Frequency**

How often can we update all WESs per second? This results in how smooth movement of moving heads are moving or how smooth the transition of the color of an LED can be performed.

- **Reliability**

Wie sicher kommen die vom Controller gesendeten Daten bei den WESs an? Lack of reliability can result in two WESs positioned next to each other not behaving the same because one of them only receives half of the signals.

- **Synchronisation**

Are the signals sent to different WESs carried out at the same time? If two of the WESs are to be controlled simultaneously, but the signal was transmitted one after the other, they have to wait for each other so that the lights change at the same time.

4.1.1 Slim Unicast

The implementation that probably comes closest to Art-Net's is to replace the TCP packets sent by Art-Net to the respective IP-Address of the WESs with unicasts to the MAC-address of the WES. Of course, the MAC address of all WESs must be known and they must all be paired with the controller, but this process is just analogous to mapping the corresponding IP addresses after dialing the WESs into a WLAN.

Latency

Latency describes the time between a command and the expected response, here considered as airtime. The airtime using 1Mbit/s is rather easy calculated, every Byte (8 Bits) takes $8\mu\text{s}$.

For a full transmission the PHY and MAC preamble and header be transmitted twice, once for the data and once for the acknowledgement. The MAC body contains the payload, which depends of the needs of the addressed WES. In a perfect clean channel the sender hasn't to defer, but has to take a DCF Inter Frame Spaces (DIFS) plus a backoff. A perfect empty channel resets the CW to CW_{min} , which are 16 slots in 802.11b, so the average backoff should take $\frac{CW_{min}}{2}$ slots, with a slottime of $20\mu\text{s}$ follows an average backoff of $160\mu\text{s}$.

Frame segment	Byte	Airtime in μs
DIFS	-	50
Average Backoff	-	160
PHY header: PLCP preamble	18	144
PHY header: PLCP header	6	48
MAC headers	28	224
MAC body	20	160
= tx time data		746
SIFS	-	10
PHY header: PLCP preamble	18	144
PHY header: PLCP header	6	48
MAC headers, no MAC body	18	112
= tx time ack		314

Table 4.2 – Composition of the Total Airtime (tx + ack)

The total airtime of the transmission of data and ack, assuming the transmission arrived successfully, is $t_{tx} = 1100\mu\text{s}$. Strictly speaking, the light could also be changed before the acknowledgement is sent, i.e. after $746\mu\text{s}$. The latency scales linearly, the delay to the n-th WES is:

$$\text{Airtime} = N \cdot t_{tx} = N \cdot 1100\mu\text{s} \quad (4.1)$$

In order to avoid unnecessary load on the radio channel, Art-Net transmit only the changes. In the worst case, however, changes affect all WESs at the same time.

Update Frequency

Following the approach of DMX and updating the 'bus' every 44Hz, would made by sending the packets round robin via unicast. With a correspondingly high number of WESs, this could be challenge with a transmission speed of 1Mbit/s, it also scales

Discuss the inimportance of order of round robin in unicast

linearly with each additional WES. In an labor steril empty channel, denying all side latencys, there airtime could be for N WESs:

$$\text{Frequency} = \frac{1}{N \cdot 1100\mu s} = \frac{9090}{N} \text{Hz} \quad (4.2)$$

For 10 WESs, addressed with respectively 20 Byte, it would still be 909Hz. This is far above the update frequency of DMX, but also very unrealistic and just intended to show, that it could theoretically be within the realm of possibility.

Reliability

One benefit of the unicast is the support of acknowledgements. The acknowledgements trigger a retransmission if no packet has arrived, therefore a controlled light will receive its signal in any case. So the reliability should be very good.

Synchronisation

Synchronising the unicast transmissions costs a lot of latency. This is because not only does each WES have to wait until its own packet has arrived, but until the packets have arrived at all the others. The implementation is chosen in such a way that each WES knows at which position of the round-robin it is and delays the execution of the successfully received transmission until the last WES has also received its signal. The delay must then be calculated deterministically. In ESP-NOW, the default is set to 8 retransmissions, so in the worst case it is assumed that a packet is sent 8 times and that for each WES.

$$\text{Airtime}_{sync} = 8 \cdot N \cdot 1100\mu s = N \cdot 8800\mu s \quad (4.3)$$

$$\text{Frequency}_{sync} = \frac{1}{N \cdot 8800\mu s} = \frac{1136}{N} \text{Hz} \quad (4.4)$$

The idea of the slim unicast is, that a transmission to each device is very fast, because the transmitted payload small. However, since we are sending many small packets, it can be assumed that we will be sending a lot of overhead. So we playing off reliability against transmission speed. It can be said that synchronisation is a feature that should be dispensed with in the slim unicast for the sake of latency.

Unfortunately the ESP-Now protocol does not allow to control the number of retransmissions before the packet is discarded.

Is it tru, that retransmission can't be controlled in ESP-NOW?

4.1.2 Slim Broadcast

The ESP-Now protocol supports both unicast and broadcast. Instead of transmitting every unicast after each other, Slim Broadcast transmits a broadcast with the payload of all channels at the same time to all fixtures. If there is a need for more than 250 Byte (DMX channel) a second broadcast has to be sent to transmit containing the missing data. To achieve this, each WES must be told in advance at which position in the payload its data is located. The broadcast must also spend one byte of payload for the sequence number. Only in the application layer do the WESs discard the incorrect broadcasts and read out their area from the entire payload.

For the unicast the payload was assumed to be 20 byte for each WES, the same amount is assumed for each WES in the Slim Broadcast calculations. The maximum payload of the broadcast is also fixed to 200 byte, because when it is close to the 250 byte limit, reliability is supposed to collapse.

quelle 200 byte

Latency

The latency of the broadcast is easier to calculate than that of unicast, because the acknowledgments are no longer necessary. In return, the payload of a single transmission increases. Assuming 10 WESs are to be addressed, each with 20 byte payload 4.3.

Frame segment	Byte	Airtime in μs
DIFS	-	50
Average Backoff	-	160
PHY header: PLCP preamble	18	144
PHY header: PLCP header	6	48
MAC headers	28	224
MAC body	200	1600
= tx time data		2286

Table 4.3 – Composition the Broadcast Airtime

Due to the fact that a second broadcast is needed if the maximum payload of ESP-NOW is exceeded, the expected transmission time is not continuous, as shown in Figure 4.1. It is easy to see that the additional overhead caused by adding another package creates noticeable latency.

Update Frequency

However, while comparing this to the latency of unicast 4.2, it becomes clear that the low overhead and the missing acknowledgments lead to a significantly higher

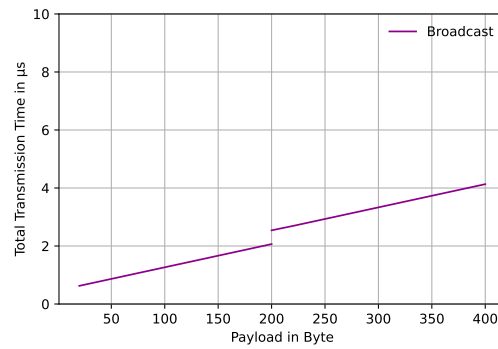


Figure 4.1 – Transmission Time of Broadcasts Depending on Payload

rate. A complete pass, i.e. addressing all WESs with 20 bytes, is already an order of magnitude faster from a number of 10 WESs.

$$\text{Frequency}_{UC} = \frac{1}{10 \cdot 1100\mu s} = 909Hz \quad (4.5)$$

$$\text{Frequency}_{BC} = \frac{1}{2286\mu s} = 4347Hz \quad (4.6)$$

Even if these values are only remotely comparable with real measurement data, it is clear, that the throughput is significantly higher with broadcast than with unicast. This effect should be even clearer under real conditions.

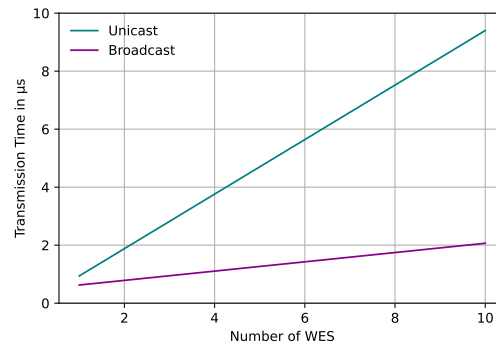


Figure 4.2 – Transmission Time of Unicast vs Broadcast

Reliability

The huge advantage that the Slim Broadcast has over the Slim Unicast in terms of update frequency, comes at the cost of lower reliability. Broadcasts can't do acknowledgements, a WES that has poor reception to the controller, will not receive

ref to fundamentals,
datalink, broadcast

packets and the controller cannot take countermeasures.

Synchronisation

Insead of transmitting to several fixtures after each other we just transmitt to all fixtures at the same time. This solves the problem of synchronization for less than 200 channel. For more than 200 each WES has to wait until the last broadcast is arrived, even if the broadcast must be discarded anyway because the required channel has already been arrived previous, deterministically calculated, the latency of all required broadcasts added up.

ist der Satz gramatisch falsch?

4.1.3 Rapid Repetition

To improve the reliability of the slim broadcast, the same transmission can simply be repeated unsolicited. The idea is not to wait for a missing acknolegedment, but to increase the probability that one of the packets got through. The reliability of the Slim Broadcast with Rapid Repetition is improved with every rapid repetition (RR). In the formula below Equation (4.7), with RR set to zero, there happens no repetition.

Is Rapid Repetition a appropriate name? Unsolicited Repetition is better siehe Paper?

Cite paper A First Implementation and Evaluation of the IEEE 802.11aa Group Addressed Transmission Service

$$SR_{RR}(RR) = 1 - (1 - SR)^{RR+1} \quad (4.7)$$

$$SR_{RR}(0) = SR \quad (4.8)$$

$$SR_{RR}(1) = 1 - (1 - SR)^2 \quad (4.9)$$

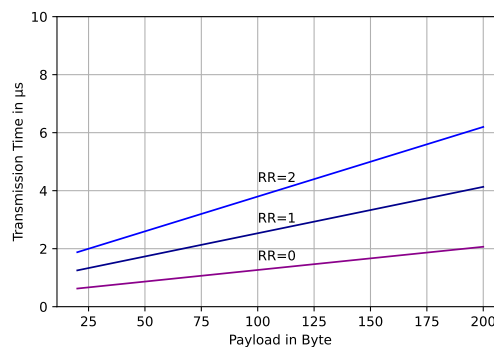


Figure 4.3 – Transmission Time with with different sets of Rapid Repetition (RR)

The RR makes it possible to reach WESs with poor reception much more reliably, However, WESs that have very good reception also receive the same packet redundantly. The update frequency of a BC with RR must be divided by the number of

repetitions, compared to one without repetitions, the same applies to the latency when synchronisation is required. However, in contrast to unicast, broadcast offers such shorter transmission times, that at least a few repetitions can be accepted. Figure 4.3

4.1.4 Delayed Repetition

- when to perform repetition
- buffering delay
- explain why its only relevant for BC?

To push the idea of rapid repetition even further, we should take a look to temporarily occurring disturbances.

Figure of bad channel time

4.2 Implementation

- ESP programming
- code examples
- python script

4.2.1 ESP Programming

- broadcast unicast
- IDF/Arduino
- IDE & ESP hardware flashing
- setup devices

for later use: ESP-NOW User Guide, V1,
source: <https://www.espressif.com/en/support/documents/>

Unfortunately in the documentation of the ESP-protocol is written, that broadcast is not supported but actually it is. Instead of adding the MAC Address of a fixture, we can use `ff:ff:ff:ff:ff:ff` to add a peer with the broadcast MAC.

has this a proper name, like broadcast address?

```
1 void addFixtureToPeerList(const uint8_t *mac_addr)
2 {
3     if (esp_now_is_peer_exist(mac_addr)) return;
4
5     peer_info.channel = 1;           // 1-14
6     peer_info.ifidx   = ESP_IF_WIFI_STA; // Station mode
```



```
7  peer_info.encrypt = false;           // not needed
8  memcpy(peer_info.peer_addr, mac_addr, 6);
9
10  esp_err_t status = esp_now_add_fixture(&peer_info);
11  if (ESP_OK != status)
12  {
13      Serial.println("[ERROR] Could not add fixture");
14  }
15  else
16  {
17      Serial.println("[OK] fixture added");
18  }
19 }
```

fix colorscheme omf code examples

4.2.2 Collecting measurment results

- Collecting values
- state machine
- python script
- saving values
- digital encoding e.g.: 777477472717

Visual Paradigm Online Free Edition

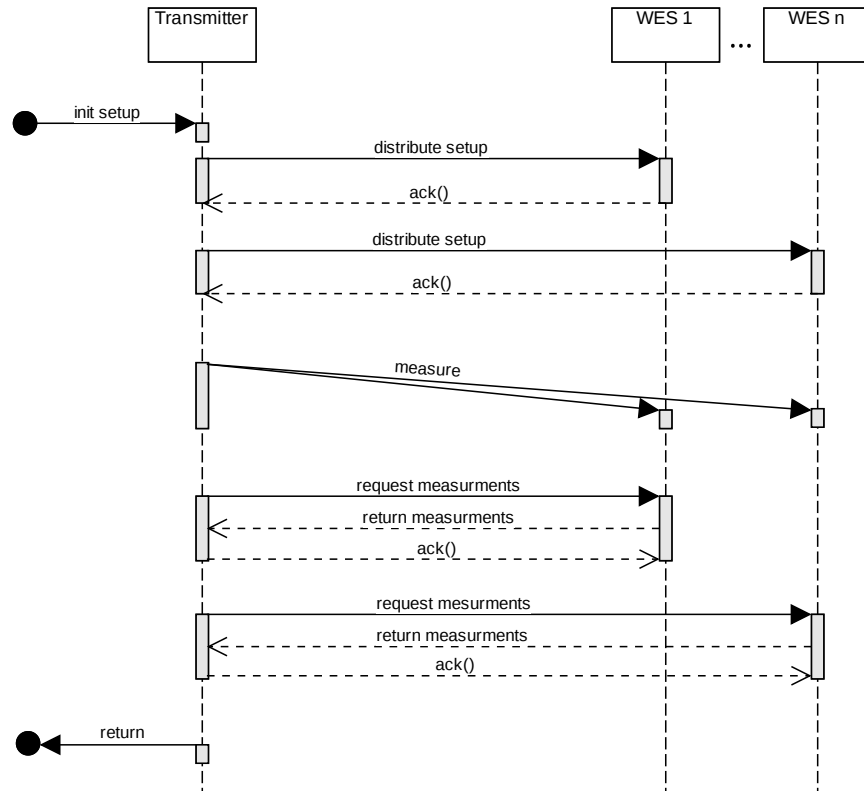
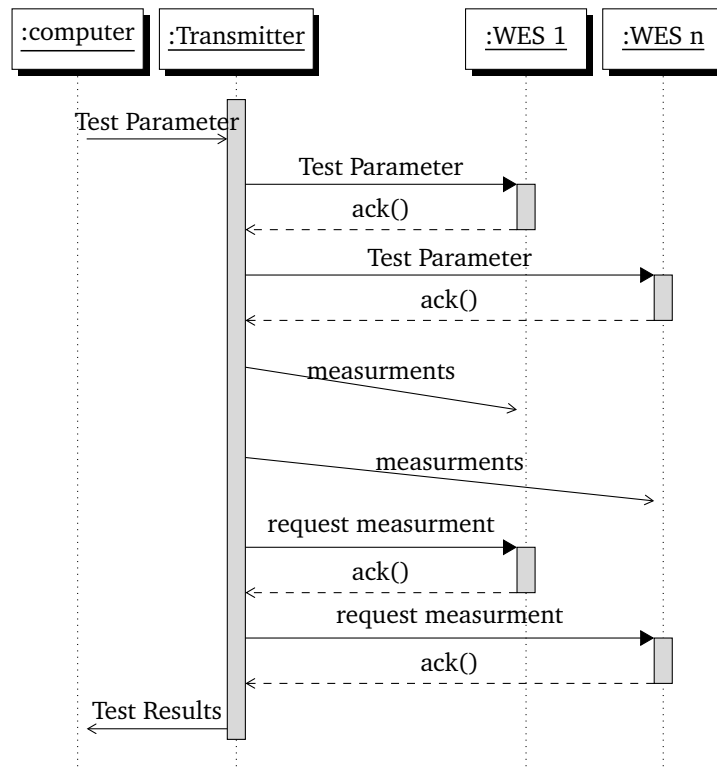


Figure 4.4 – Sequence diagram of the measurement



Chapter 5

Evaluation

- measurement setup / results / evaluation / discussion
- whatever you have done, you must comment it, compare it to other systems, evaluate it
- usually, adequate graphs help to show the benefits of your approach
- each result/graph must not only be described, but also discussed (What's the reason for this peak? Why have you observed this effect? What does this tell about your architecture/system/implementation?)
- recommended length: approximately one third of the thesis.

Keep in Mind

- metrics (SR, Latency, ...)
- compare with art-net all the time
- wireshark

5.1 Methodic

- Testbed
- Collect Data
- Sequence Diagram to explain

5.2 Protocols under Study

-

Unicast vs Broadcast

Rapid Repetition

Delayed Repetition

Grouping

5.3 Results

Difference between Results and Discussion?

- Grafen miteinander vergleichen?
- Which method had the best results?

Chapter 6

Conclusion & Discussion

- summarize again what your paper did, but now emphasize more the results, and comparisons
- write conclusions that can be drawn from the results found and the discussion presented in the paper
- future work (be very brief, explain what, but not much how, do not speculate about results or impact)
- recommended length: one page.

Why not 5GHz -> too expensive.

List of Abbreviations

AP	Access Point
BSS	Basic Service Set
BSSID	Basic Service Set Identifier
CAN	Controller Area Network, <i>when referring to the bus protocol</i>
CSMA/CA	Carrier-sense multiple access with collision avoidance
DIFS	DCF Inter Frame Spaces
DL	Data Link Layer
DMX	Digital Multiplex
DSSS	Direct Sequence Spread Spectrum
FHSS	Frequency Hopping Spread Spectrum
IEEE	Institut of Electrical and Electronics Engineers
IP	Internet Protocol
LLC	Locig Link Control
MAC	Media Access Control
MCU	Microcontroller Unit
MIMO	Multiple Input-Multiple Output)
MSDU	MAC Service Data Unit
NAV	Network Allocation Vector
OFDM	Orthogonal Frequency Division Multiplexing)
OSI	Open Systems Interconnection
PHY	Physical Network Layer
SIFS	Short Inter Frame Spaces
STA	Station
UDP	User Datagram Protocol
WES	Wireless Endsystem
WLAN	Wireless Local Area Network

List of Figures

3.1	MAC header of a WLAN frame	9
3.2	Unicast Transmission	11
3.3	Broadcast Transmission	12
3.4	Multicast Transmission	12
3.5	Block Diagram of an DMX Universe	13
3.6	ESP32 Devboard (Devkit V1)	15
4.1	Transmission Time of Broadcasts Depending on Payload	23
4.2	Transmission Time of Unicast vs Broadcast	23
4.3	Transmission Time with with different sets of Rapid Repetition (RR)	24
4.4	Sequence diagram of the measurment	27

List of Tables

3.1	OSI model	7
3.2	ESP-NOW Frame Format	16
3.3	Vendor Specific Action Frame	16
4.1	Art-Net Layer compared with Slim Data Link Layer	19
4.2	Composition of the Total Airtime (tx + ack)	20
4.3	Composition the Broadcast Airtime	22

Bibliography

- [1] D. D. Husam Kareem, “The Working Principles of ESP32 and aAnalytical Comparison of using Low-Cost Microcontroller Modules in Embedded Systems Design,” *4th International Conference on Circuits, Systems and Simulation*, vol. 40, no. 8, pp. 130–135, Aug. 2021. DOI: 10.1109/ICCSS51193.2021.9464217.
- [2] I. F. Akyildiz, W. Su, Y. Sankarasubramaniam, and E. Cayirci, “A Survey on Sensor Networks,” *IEEE Communications Magazine*, vol. 40, no. 8, pp. 102–114, Aug. 2002. DOI: 10.1109/MCOM.2002.1024422.
- [3] T. Arampatzis, J. Lygeros, and S. Manesis, “A Survey of Applications of Wireless Sensors and Wireless Sensor Networks,” in *13th Mediterrean Conference on Control and Automation*, Limassol, Cyprus, Jun. 2005, pp. 719–724. DOI: 10.1109/.2005.1467103.
- [4] I. Dietrich and F. Dressler, “On the Lifetime of Wireless Sensor Networks,” *ACM Transactions on Sensor Networks (TOSN)*, vol. 5, no. 1, pp. 1–39, Feb. 2009. DOI: 10.1145/1464420.1464425.
- [5] K. A. Roberto Pasic Ivo Kuzmanov, “ESP-NOW communication protocol with ESP32,” *Journal of Universal Excellence*, vol. 40, no. 8, pp. 53–60, Feb. 2021. DOI: 10.37886/ip.2021.019.

Todo list

■ The table of contents should fit on one page. When in doubt, adjust the tocdepth counter.	vi
■ triple-check all references	6
■ refrence to table below	7
■ Rewrite introduction in chapter Fundamentals!	7
■ Payload and FCS are missing	9
■ shot explanation of CSMA/CD	10
■ eigene Worte	10
■ when to explain CSMA/DC?	11
■ example	11
■ complete figure UC	11
■ complete figure BC	12
■ Mutlicast explain multicast mac address	12
■ examples for multicast ack + related work	12
■ complete figure MC	12
■ image der connectoren	13
■ Show Hardware e.g. DMX Plug	13
■ Image of fixtures distributed on channel	13
■ ESP32 Kosten in € 2021 aufführen? Link? Datum?	14
■ paper about esp above arduino	14
■ this is cited from: link zum esp32 datasheet	14
■ cite ESP documentation website	15
■ cite somehow the ESP-NOW documentation pdf: ESP-IDF Programming Guide: ESP-NOW, source: https://docs.espressif.com/projects/ esp-idf/en/latest/esp32/api-reference/network/esp_now.html	16
■ this is cited from espressif manual!!	17
■ subsection can be on a wrong position!	17
■ remove newpage command	17
■ ESP-NOW Baseline Artnet should be moved to Design part??	17
■ Fehlt hier nicht ESP-Now?	18

ref to requirements	19
Discuss the inimportance of order of round robin in unicast	20
Is it tru, that retransmission can't be controlled in ESP-NOW?	21
quelle 200 byte	22
ref to fundamentionals, datalink, broadcast	23
ist der Satz gramatisch flasch?	24
Is Rapid Repetition a appropriate name? Unsosliced Repetition is better siehe Paper?	24
Cite paper A First Implementation and Evaluation of the IEEE 802.11aa Group Addressed Transmission Service	24
Figure of bad channel time	25
for later use: ESP-NOW User Guide, V1, source: https://www.espressif.com/en/support/documents/	25
has this a proper name, like broadcast address?	25
fix colorscheme omf code examples	26
Difference between Results and Discussion?	30