

FIT2004

Algorithms and Data Structures

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Referencing materials by
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Ready?

Agenda

- Sorting Algorithms
 - Comparison based
 - Selection
 - Insertion
 - Non-comparison based (the IMBA ones)
 - Counting
 - Radix

Let us begin...

- We are back to sorting!
 - Bubble
 - Insertion
 - Selection
 - Merge
 - Quick

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- Bubble
- Insertion
- Selection
- Merge
- Quick



Janelle Shane @JanelleCShane · 14 Apr ✓

For example, there was an algorithm that was supposed to sort a list of numbers. Instead, it learned to delete the list, so that it was no longer technically unsorted.

💬 10

↻ 143

❤️ 635



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- All of these are known as comparison based sorting.
Why?

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Why? Because we compare between items to know if $a < b$ or $b > a$

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- All of these are known as comparison based sorting.
Why? Because we compare between items to know if $a < b$ or $b > a$

- Now let us analyze them based on what we have learnt!

Questions?

Sorting

Selection Sort

- Correctness
- Complexity

Sorting

Selection Sort

- **Correctness**
 - Loop invariant
 - Termination
- **Complexity**
 - Time
 - Space

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Sorting

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 - $\text{my_list}[0 \dots i-1]$ is sorted
 - $\text{my_list}[0 \dots i-1] \leq \text{my_list}[i \dots N]$
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 - `my_list[0...i-1]` is sorted
 - `my_list[0...i-1] ≤ my_list[i...N]`
- Termination
 - `i` and `j` always increment and both reach the end of the list
- So why is it working then?
 - `i` keep increment till `n` and we know from invariant `0...i-1` is sorted, thus we will sort the entire list!

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 - Space
 - $O(N)$ for the input list
 - Auxiliary? $O(1)$ **in place**

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 - Like comparing between words, you need to compare the alphabets

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 - But what if I tell you comparing the items have a cost of $O(k)$
 - Like comparing between words, you need to compare the alphabets
 - We know complexity is based on comparison $O(N^2)$ comparisons...
 - So our final complexity is $O(kN^2)$

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- Stable?

Sorting

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 - Relative ordering doesn't change

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 - Minimum is 1, so we swap

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- Complexity
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 - Is it stable? **No!** but **why?**
 - [4a, 2, 3, 4b, 1]
 - Minimum is 1, so we swap
 - [1, 2, 3, 4b, 4a]

Sorting

Selection Sort

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- Complexity
- Stable?
 - Relative ordering doesn't change
 - Is it stable? **No!** but **why?**
 - [4a, 2, 3, 4b, 1]
 - Minimum is 1, so we swap
 - [1, 2, 3, 4b, 4a]
 - Now we see that 4a is behind 4b!

Questions?

Sorting

Insertion Sort

- Correctness
- Complexity

- Correctness
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Problem 1. Write pseudocode for insertion sort, except instead of sorting the elements into non-decreasing order, sort them into non-increasing order. Identify a useful invariant of this algorithm.

- Correctness
- Complexity

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def insertion_sort(my_list):  
    for i in range(1, len(my_list)):  
        key = my_list[i]  
        j = i - 1  
        # keep shifting to left if left is greater  
        while j >= 0 and key < my_list[j]:  
            my_list[j+1] = my_list[j]  
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```

- Correctness
 - Loop invariant
 - Termination
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 - `my_list[0...i-1]` sorted
- Termination
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 - Each loop only look and compare with left item once
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- What about space?

- $O(N)$ for the input list
 - $O(1)$ auxiliary cause it is in-place

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- Stability
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 - Don't swap if value is the same
 - Most **shifting** will ensure stability

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Questions?


Summary

Sorting

	Best	Worst	Average	Stable?	In-place?
Selection Sort	$O(N^2)$	$O(N^2)$	$O(N^2)$	No	Yes
Insertion Sort	$O(N)$	$O(N^2)$	$O(N^2)$	Yes	Yes
Heap Sort	$O(N \log N)$	$O(N \log N)$	$O(N \log N)$	No	Yes
Merge Sort	$O(N \log N)$	$O(N \log N)$	$O(N \log N)$	Yes	No
Quick Sort	$O(N \log N)$	$O(N^2)$ – can be made $O(N \log N)$	$O(N \log N)$	Depends	No

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Auxiliary for Recursion

- The recursion stack takes up memory!!!

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 - If I have recursion $\log N$ times, then I take $O(\log N)$ space for the recursion alone!
 - If each recursion is k , then my total space is $O(k \log N)$!!!

- The recursion stack takes up memory!!!
 - So that is why it isn't in-place!
 - Iterative is easier to get in-place
 - If I have recursion $\log N$ times, then I take $O(\log N)$ space for the recursion alone!
 - If each recursion is k , then my total space is $O(k \log N)$!!!

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- So... what is the lower bound for the sorting algorithms that we have learnt?
 - Bubble
 - Insertion
 - Selection
 - Quick
 - Merge
- These are all comparison based
- $\Omega(N \log N)$
- We will see more of this **later**

Questions?

Thank you