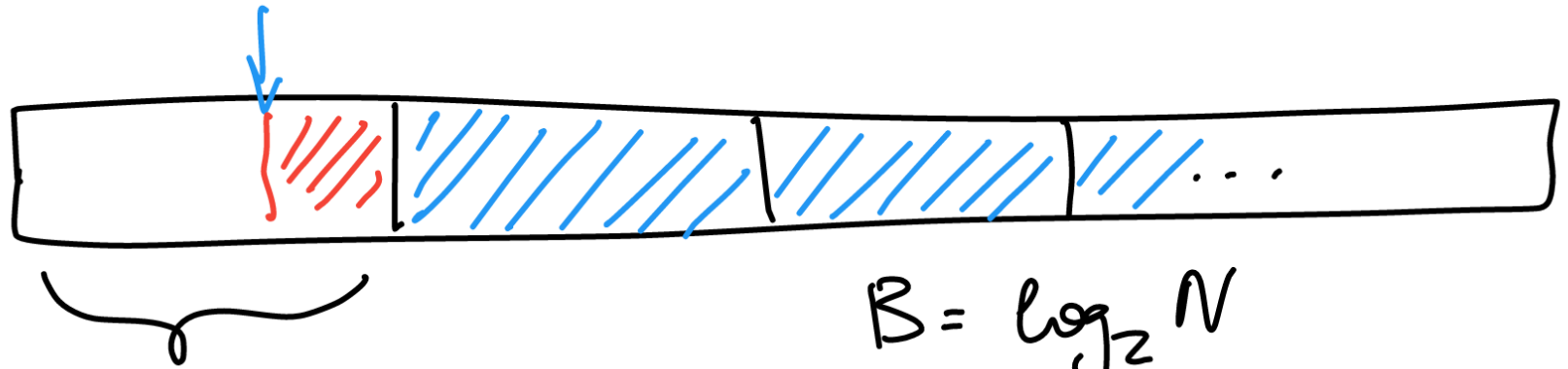


①



Sparse Table (block)

Global Sparse Table keeps block minimums

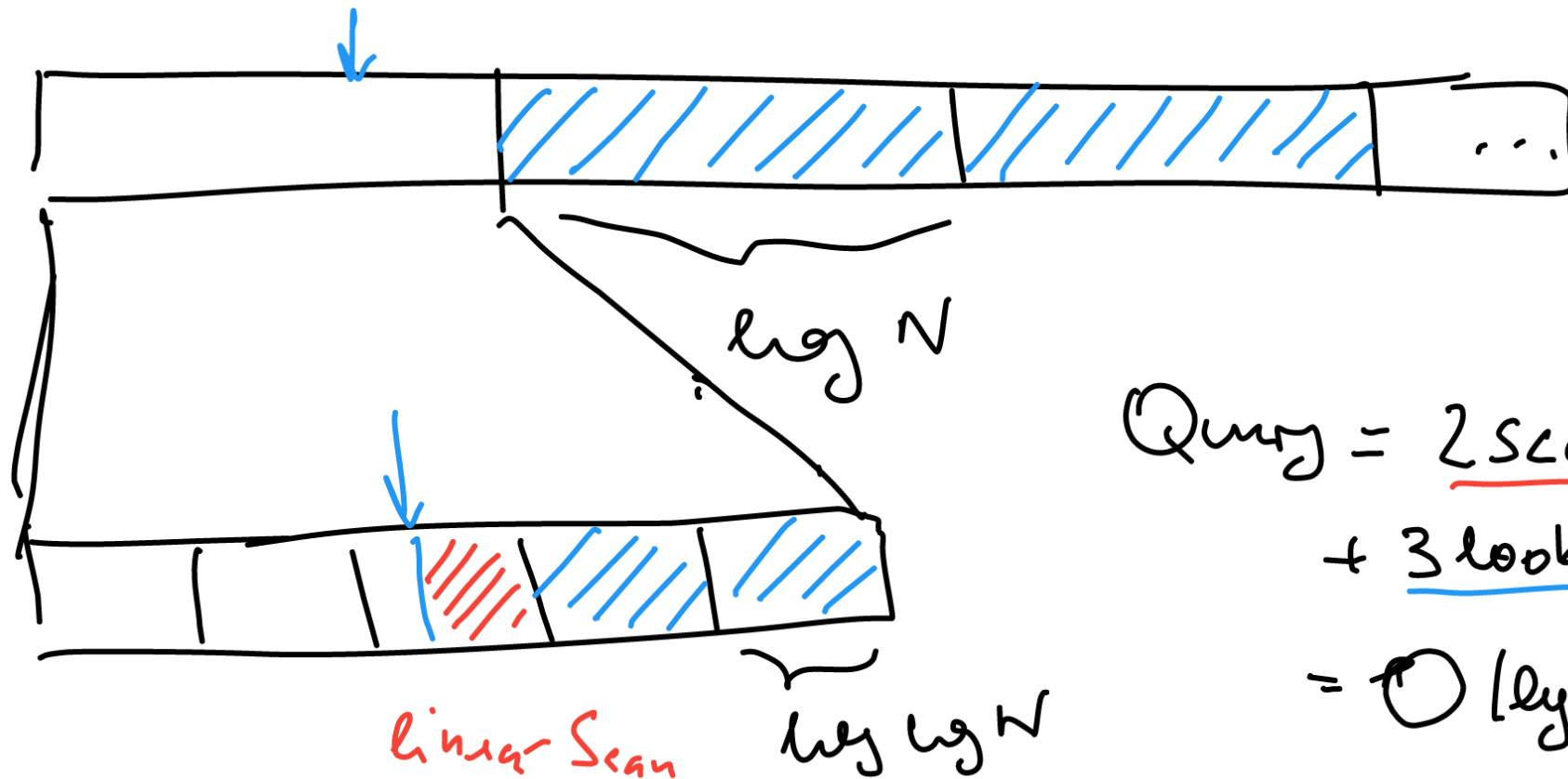
$$O\left(\frac{N}{\log N} \cdot \log\left(\frac{N}{\log N}\right)\right) = O(N)$$

$$+ \text{ST per each block} : \frac{N}{\log N} \cdot O(\log N \log \log N)$$

$$= O(N \log \log N)$$

$$\text{Query} = \underline{1 \text{ global}} + \underline{2 \text{ local}} = O(1)$$

Optimization: Store 1 ST with depth $O(\log \log N)$

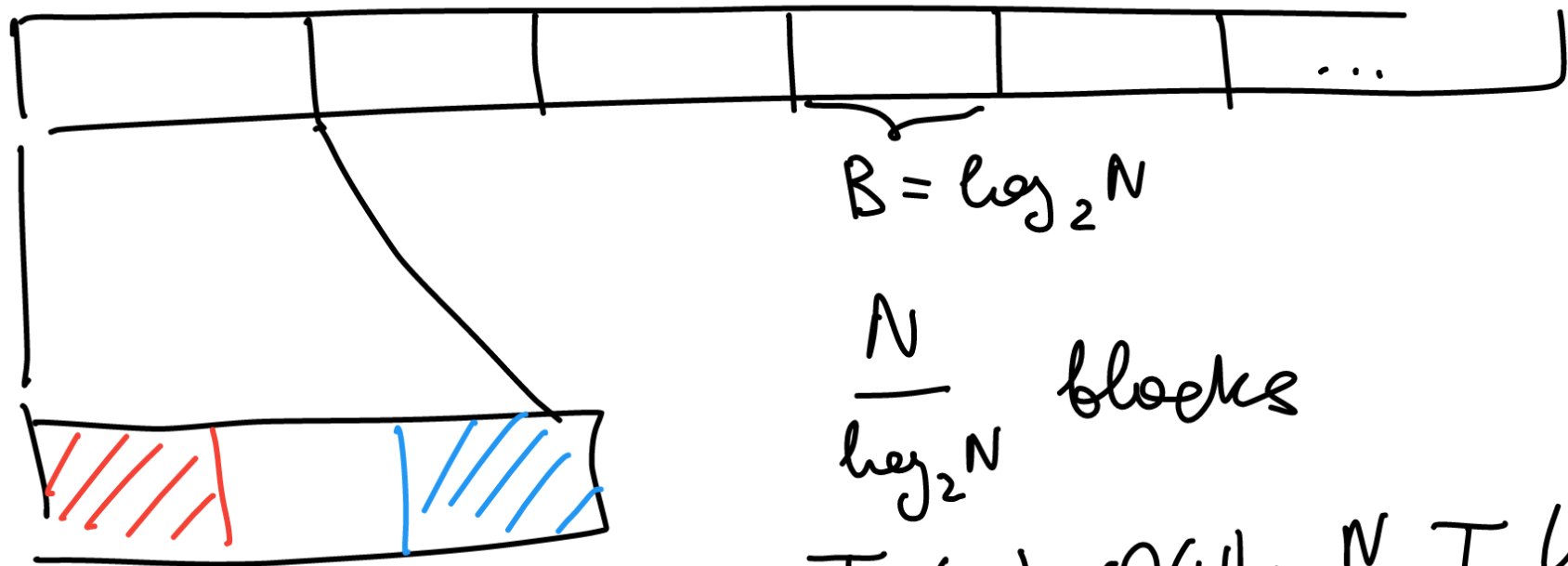


$$\begin{aligned} \text{Query} &= \underline{2 \text{ Scans}} \\ &+ \underline{3 \text{ lookups}} \\ &= \mathcal{O}(\log \log N) \end{aligned}$$

$$\mathcal{O}\left(\frac{N}{\log N} \cdot \log\left(\frac{N}{\log N}\right)\right) = \mathcal{O}(N)$$

$$\frac{N}{\log N} \cdot \mathcal{O}\left(\frac{\log N}{\log \log N} \cdot \log\left(\frac{\log N}{\log \log N}\right)\right) = \mathcal{O}(N)$$

Optimization: store 2nd lvl. in 1 table



Build recursively!

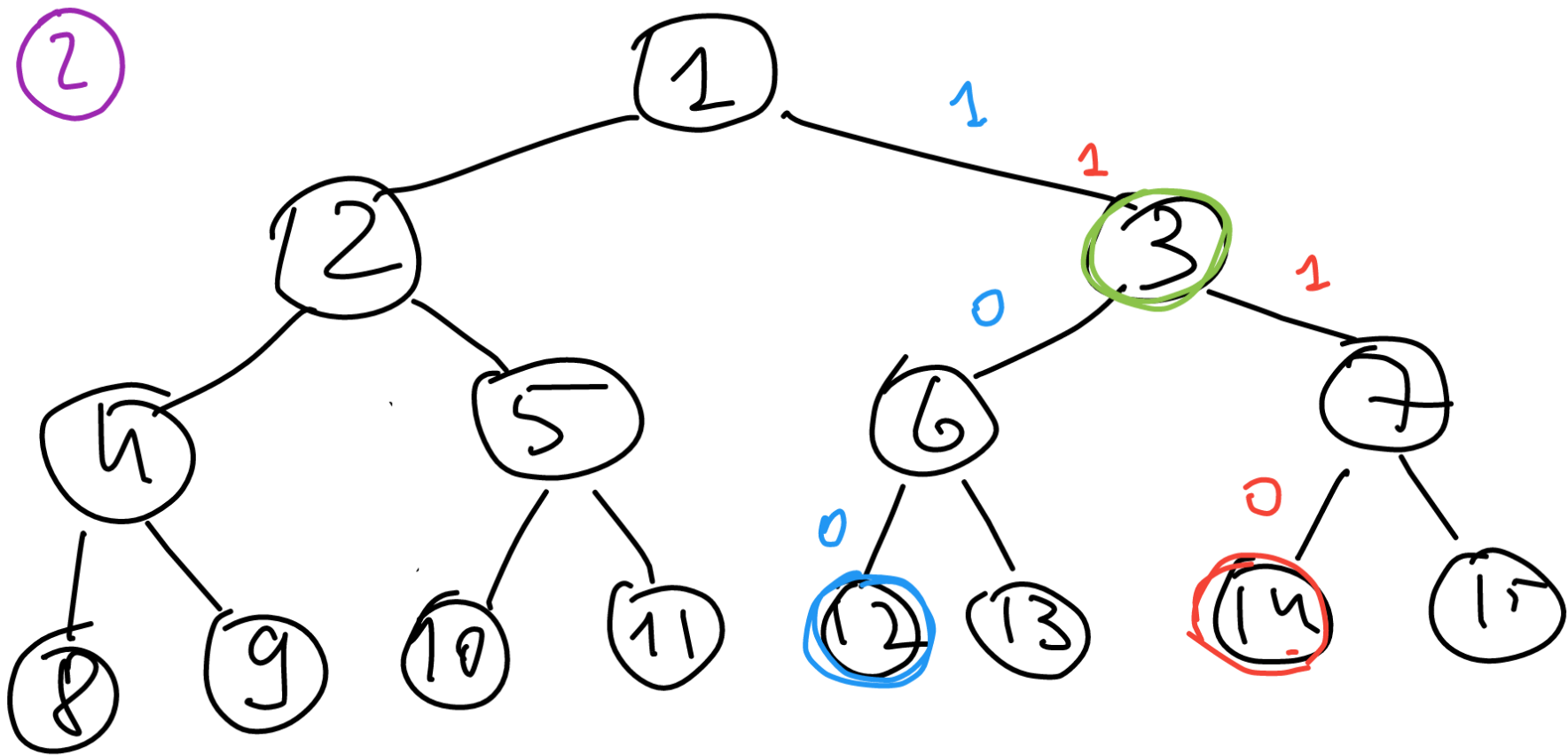
+ suffix min
+ prefix min

$$T_B(N) = O(N) + \frac{N}{\log N} T_B(\log N)$$

$$T_Q(N) = O(1) + T_Q(\log N)$$

without suffix / prefix:

$$T'_Q(N) = O(1) + 2 T'_Q(\log N) = O(2^{\log^2 N})$$



$$\begin{aligned}
 12 &= \boxed{11}00_2 \\
 14 &= \boxed{11}10_2 \\
 3 &= 11_2
 \end{aligned}$$

$$u \text{ xor } v = \underbrace{00 \dots 01}_{\text{LCA bits}} \dots$$

LCA bits

position of
highest bit

③ $b = \frac{1}{2} \log_2 N$. Store answer for all masks of length $1, 2, \dots, b$. Total $2^b + 2^{b-1} + \dots + 2 + 1 = 2^{b+1} - 1 = \mathcal{O}(\sqrt{N})$ masks.

mask = 1 1 0 0 0 1

$a = 0, 1, 2, 1, 0, -1, 0$

mask' = 1 0 0 0 1

$a' = 0, 1, 0, -1, -2, -1$

$$F(\text{mask}) = \min(0, F(\text{mask}') \pm 1)$$

$$\text{mask}' = \text{mask} \& (2^{b-1} - 1)$$

We need to extract leftmost bit:

$$\text{mask} \geq 2^{k-1} = \underbrace{100 \dots 0}_k 2$$

k bits

$+0-?$

Or just reverse mask \rightarrow

$$\begin{array}{l} +1 \swarrow \frac{\text{mask}}{\text{mask}'} \\ 0 \swarrow \\ -1 \swarrow \frac{\text{mask}}{\text{mask}'} \end{array}$$

mask mod 2 : rightmost bit

$$\text{mask}' = \text{mask} / 2$$

④ Calculate jump pointers
after adding a vertex u :

$$\text{jump}[u, 0] = \text{par}[u]$$

$$\text{jump}[u, k] = \text{jump}[\text{jump}[u, k-1], k-1]$$

$\mathcal{O}(\log N)$ per vertex

⑤ Jump pointers. Store

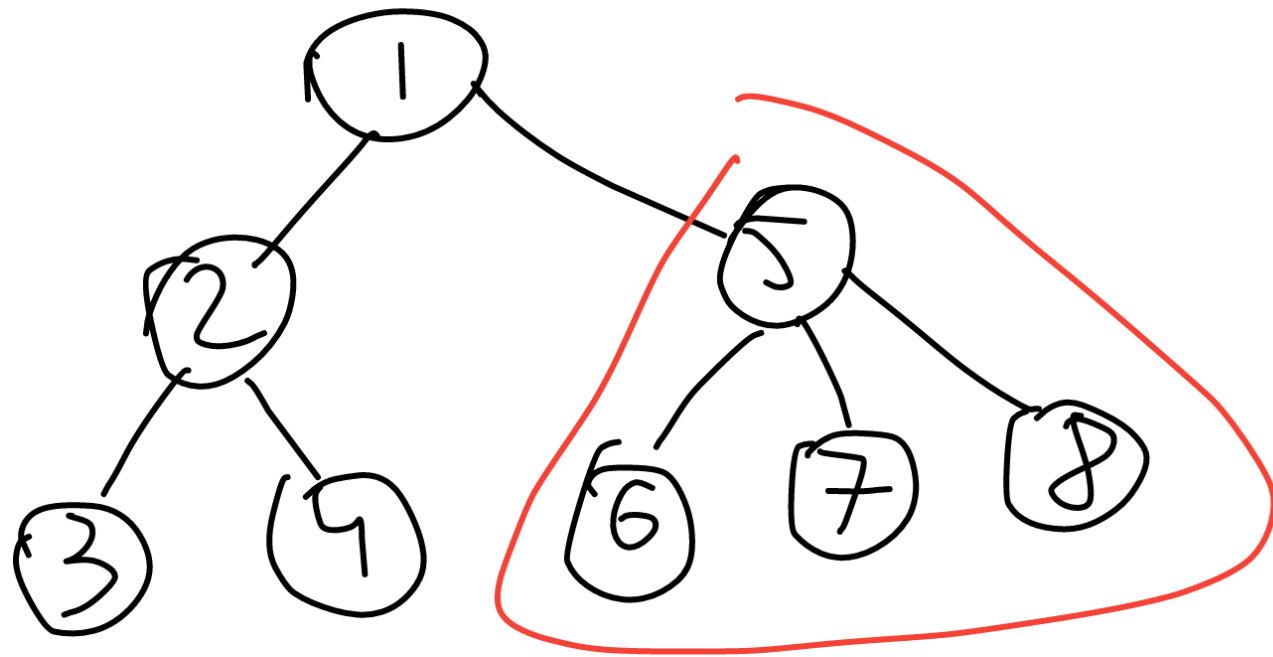
$$\min(a[u], a[\text{par}[u]], \dots, a[\text{par}^{k-1}[u]])$$

with jump pointer k

$$\text{par}^{k-1}[u] = \underbrace{\text{par}[\text{par}[\dots, \text{par}[u] \dots]]}_{k-1 \text{ times}}$$

$k-1$ times

⑥ Euler Tour. Store vertex once

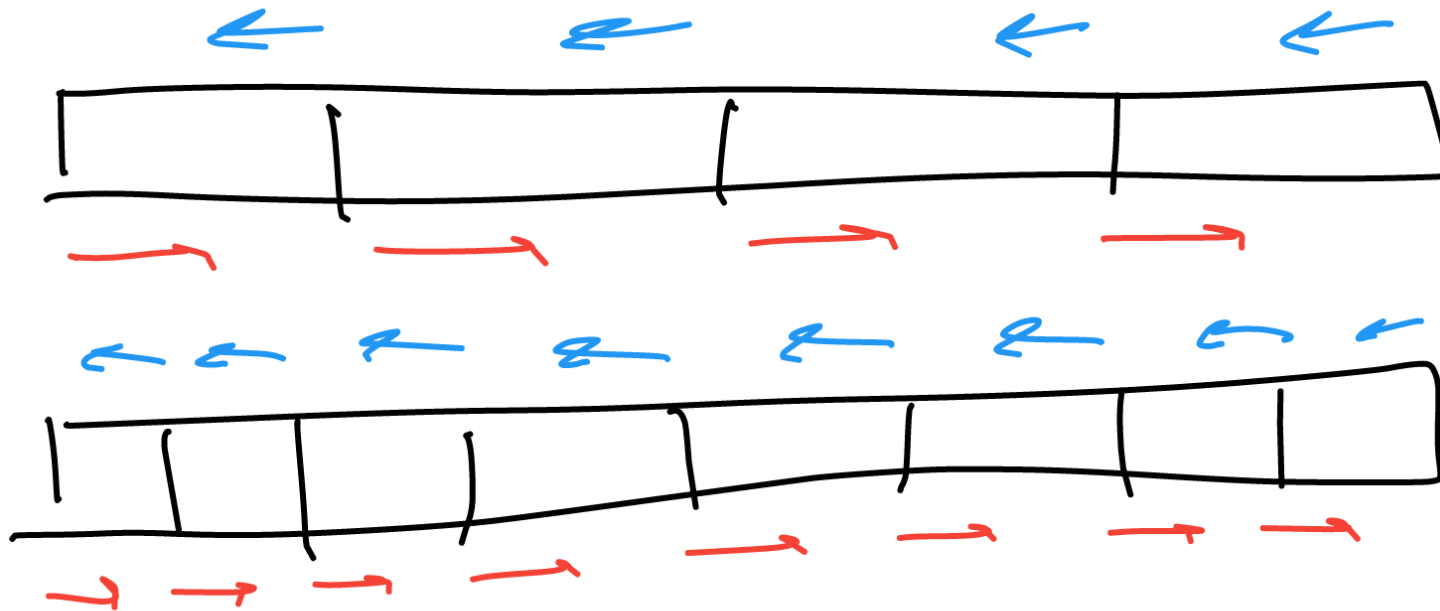
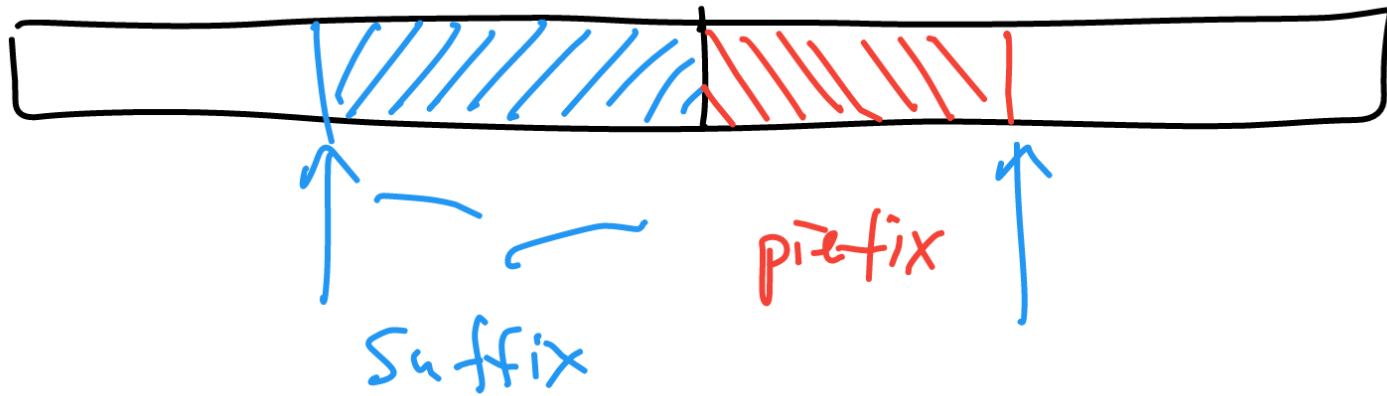


1 2 3 4 5 6 7 8

Then Segment Tree or BST to
store set of colored vert

⑦

$$R-L \geq N/2 :$$



Level ~ highest bit of $R-L$

level = 2

01011₂

$K=3$