

# RMQ, LCA

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## Problems

1. Use Sparse Table idea to solve static RMQ problem...
  - (a) ...in  $\langle \mathcal{O}(n \log \log n), \mathcal{O}(1) \rangle$  time.
  - (b) ...in  $\langle \mathcal{O}(n), \mathcal{O}(\log \log n) \rangle$  time.
  - (c) ...in  $\langle \mathcal{O}(n \log^* n), \mathcal{O}(\log^* n) \rangle$  time. [What is  \$\log^\* n\$ ?](#)
2. You are given a full binary tree of size  $2^h - 1$  labelled in level order fashion. I.e. the root has label 1, and for each non-leaf vertex labelled  $i$  its left and right children has labels  $2i$  and  $2i + 1$ . Find a simple way to answer  $\text{LCA}(u, v)$  queries in this tree in  $\mathcal{O}(1)$  time.
  - (a) How the binary representation of label  $i$  and the path from root to  $i$  are related?
  - (b) Solve the problem if  $\text{depth}(u) = \text{depth}(v)$ .
  - (c) Solve the problem in general case.
3. Bender-Farach-Colton algorithm takes  $\mathcal{O}(\sqrt{n} \log^2 n)$  time and space to precompute answers for all short  $\pm 1$  sequences. Show how to do that in just  $\mathcal{O}(\sqrt{n})$  time and space, keeping the same  $\frac{1}{2} \log_2 n$  block size and  $\mathcal{O}(1)$  query time.
4. Create a data structure which starts with a single root vertex and can quickly answer the following queries: (1) add a new leaf vertex and (2) find  $\text{LCA}(u, v)$ , both in  $\mathcal{O}(\log n)$  time. You can use  $\mathcal{O}(n \log n)$  memory.
5. You are given a tree where each vertex has an assigned value. Create a data structure which finds minimum value on a shortest path between a pair of vertices  $u, v$  in  $\langle \mathcal{O}(n \log n), \mathcal{O}(\log n) \rangle$  time.
6. You are given a tree where each vertex can be colored or not. Create a data structure which can quickly answer the following queries: (1) color/uncolor any vertex and (2) find the number of colored vertices in a subtree, both in  $\mathcal{O}(\log n)$  time. You can use  $\mathcal{O}(n)$  memory.
7. You are given a binary operation  $\circ : X \times X \rightarrow X$  which is associative ( $(a \circ b) \circ c = a \circ (b \circ c)$ ), but not necessarily idempotent. You are asked to answer range queries  $a_L \circ a_{L+1} \circ \dots \circ a_{R-1}$  in  $\mathcal{O}(1)$  time. Construct a data structure which allows that after  $\mathcal{O}(n \log n)$  preprocessing.
  - (a) Solve the special case  $R - L \geq N/2$  with  $\mathcal{O}(n)$  preprocessing.
  - (b) Solve the general case for  $N = 2^k$ .