

# NaCl (for the IoT)

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# A bit about me

2001-2006: Studies of computer science at RWTH Aachen



# A bit about me

2008-2011: Ph.D. student at TU Eindhoven



## A bit about me

2011-2012: Postdoc at Academia Sinica and NTU (Taiwan)



A bit about me

Since Jan. 2013: Assistant professor at Radboud University



## Space shuttles and elevators

*“OpenSSL is the space shuttle of crypto libraries. It will get you to space, provided you have a team of people to push the ten thousand buttons required to do so. NaCl is more like an elevator – you just press a button and it takes you there. No frills or options.*

*I like elevators.”*

Matthew Green in his blog entry *The anatomy of a bad idea*

# NaCl: Networking and Cryptography library

- ▶ Core development team: Daniel J. Bernstein, Tanja Lange, Peter Schwabe
- ▶ Acknowledgment: Contributions by
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  - ▶ Niels Duif (TU Eindhoven)
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  - ▶ Bo-Yin Yang (Academia Sinica)
- ▶ User's perspective: Bundle of functionalities rather than bundle of algorithms
- ▶ Focus on protecting Internet communication



# Protecting Internet communication

- ▶ Alice wants to send a message  $m$  to Bob
- ▶ Uses Bob's public key and her own private key to compute authenticated ciphertext  $c$ , sends  $c$  to Bob
- ▶ Bob uses his private key and Alice's public key to verify and recover  $m$

## Alice using a typical crypto library

- ▶ First choose algorithms and parameters, e.g. AES-128, RSA-2048, SHA-256
- ▶ Generate random AES key
- ▶ Use AES to encrypt packet
- ▶ Hash encrypted packet
- ▶ Read RSA private key from wire format
- ▶ Use key to sign hash
- ▶ Read Bob's RSA public key from wire format
- ▶ Use key to encrypt AES key and signature
- ▶ ...

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- ▶ ...
- ▶ Plus more code to allocate storage, handle errors etc.

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- ▶ `sk`: Alice's 32-byte private key
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- ▶ `c`: authenticated ciphertext, 16 bytes longer than plaintext `m`

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- ▶ Initial keypair generation for Alice and Bob:

```
pk = crypto_box_keypair(&sk)
```

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- ▶ NaCl also contains signatures with an easy-to-use interface:

```
pk = crypto_sign_keypair(&sk)
```

generates a 64-byte private key and a 32-byte public key

```
sm = crypto_sign(m, sk)
```

signs `m` under `sk`; `sm` is 64 bytes longer than `m`

```
m = crypto_sign_open(sm, pk)
```

verifies the signature and recovers `m`

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  - ▶ Poly1305 for one-time secret-key authentication
- ▶ At least 128 bits of security
- ▶ Easy to implement **efficiently and securely**

## NaCl Speed

# NaCl Speed

- ▶ **NaCl offers exceptionally high speeds, keeps up with the network**
- ▶ NaCl operations per second on AMD Phenom II X6 1100T for any reasonable packet size:
  - ▶ > 80000 `crypto_box`
  - ▶ > 80000 `crypto_box_open`
  - ▶ > 70000 `crypto_sign_open`
  - ▶ > 180000 `crypto_sign`
- ▶ Handles arbitrary packet floods up to  $\approx 30$  Mbps per CPU, depending on protocol
- ▶ Much faster than, e.g., TLS from OpenSSL

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- ▶ Very fast rejection of forged packets under known public keys
- ▶ Fast batch signature verification: doubling verification speed

# Introducing $\mu$ NaCl

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- ▶ Addresses have only 16 bits, so restrict message length to  $2^{16} - 1$  (avoid expensive arithmetic on 64-bit integers)



# AVR NaCl speeds and sizes

## High-speed configuration

- ▶ Secret-key authenticated encryption:  $\approx 500$  cycles/byte (268 bytes of RAM)
- ▶ Variable-basepoint scalar multiplication: 22 791 580 cycles (677 bytes of RAM)
- ▶ `crypto_sign`: 23 216 241 cycles (1 642 bytes of RAM)
- ▶ `crypto_sign_open`: 32 634 713 cycles (1 315 bytes of RAM)
- ▶ 27 962 bytes of ROM for NaCl

# AVR NaCl speeds and sizes

## Small-size configuration

- ▶ Secret-key authenticated encryption:  $\approx 520$  cycles/byte (273 bytes of RAM)
- ▶ Variable-basepoint scalar multiplication: 27 926 288 cycles (917 bytes of RAM)
- ▶ `crypto_sign`: 34 303 972 cycles (1 289 bytes of RAM)
- ▶ `crypto_sign_open`: 40 083 281 cycles (1 346 bytes of RAM)
- ▶ 17 373 bytes of ROM for NaCl

## How about hardware?

- ▶ Joint work with Michael Hutter, Jürgen Schilling, and Wolfgang Wieser: Tiny ASIC for `crypto_box`
- ▶ Target application: WISP nodes
- ▶ Optimize for low power, not low energy

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- ▶ 32-bit ASIP
- ▶ Results:
  - ▶ 14.6–18.0kGE
  - ▶ 40–70  $\mu$ W (half of power is spent for RAM)
  - ▶ 53.4–82.6ns (12–18MHz)
  - ▶ 811 170–3 455 394 cycles for DH

## More around NaCl

- ▶ Bernstein, Schwabe, 2012:

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<https://cryptojedi.org/crypto/#mu25519>
- ▶ Ongoing work together with Bernstein:  
**Full verification of Curve25519 (and generally ECC)**  
<https://cryptojedi.org/crypto/#gfverif> (online soon)

# NaCl online

<http://nacl.cr.yp.to>

<http://munacl.cryptojedi.org>

<http://cryptojedi.org/crypto/#naclhw>

- ▶ No license: NaCl is in the public domain
- ▶ No patents that we are aware of