Types

TABLE 7-1

Category Name	Categorial Definition of Name	Nearest Transformational Equivalent	Basic Expressions
e		None	None
t		Sentence	None
IV	t/e	Verb Phrase and Intransitive Verb (IV is mnemonic for "Intransitive Verb Phrase")	run, walk, talk, rise, change
T	t/IV	Noun Phrase and Proper Name (T is mnemonic for "Term Phrase")	John, Mary, Bill, ninety, he ₀ , he ₁ , he ₂ ,
TV	IV/T	Transitive Verb	find, lose, eat, love, date, be, seek, conceive
IAV	IV/IV	Verb Phrase Adverb (IA V is mnemonic for "Intransitive Adverb")	rapidly, slowly, voluntarily, allegedly
CN	t//e	Common Noun	man, woman, pari fish, pen, unicorn, price, temperature
t/t		Sentence Adverb	necessarily
IAV/T		Preposition (one that forms a VP-modifying prepositional phrase)	in, about
IV/t		Sentence-complement Verb	believe, assert
IV/ IV		Infinitive-complement Verb	try, wish
DET ²	T/CN	Determiner	every, the, a(n)

Figure 1: English Types

Syntactic Rules

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S2. \delta \in Det, \zeta \in CN \implies F_2(\delta,\zeta) \in Tm. [a/an/the/every] F2 :: Det -> CN -> Tm
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S3. $\zeta \in CN, \phi \in T \implies F_{3,n}(\zeta,\phi) \in CN$. [such that]

F3 :: Int -> CN -> T -> CN

TABLE 7-2

Category Name	Categorial Definition of Name	Corresponding Type by Rule (7-2)	Name of Semantical Object Denoted by this Type
e	e	e	individual
t	t	t	truth value
IV	t/e	((s,e),t)	set of individual concepts
T	$t/IV \\ (= t/(t/e))$	$\langle\langle s, \langle\langle s, e \rangle, t \rangle\rangle, t \rangle$	set of properties of individual concepts
CN	t//e	$\langle\langle s,e \rangle,t \rangle$	set of individual concepts
t/t	t/t	$\langle (s,t),t \rangle$	set of propositions
etc.	etc.	etc.	etc.

Figure 2: Intensional Logic Types

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S4. \alpha \in \text{Tm}, \delta \in \text{IV} \implies F_4(\alpha, \delta) \in \text{T. [subj-IV]}
F4 :: Tm -> IV -> T

S5. \delta \in \text{TV}, \beta \in \text{Tm} \implies F_5(\delta, \beta) \in \text{IV. [TV-obj]}
F5 :: TV -> Tm -> IV

S6. \delta \in \text{Prep}, \alpha \in \text{Tm} \implies F_5(\delta, \alpha) \in \text{IAV. [prep-obj]}
F5 :: Prep -> Tm -> IAV

S7. \alpha \in \text{StV}, \phi \in \text{T} \implies F_{11}(\alpha, \phi) \in \text{IV. [that]}
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$$SI. \ \alpha \in StV, \ \phi \in I \implies F_{11}(\alpha, \phi) \in IV. \text{ [that]}$$

$$F11 :: StV \rightarrow T \rightarrow IV$$

S8.
$$\delta \in \text{ItV}, \ \beta \in \text{IV} \implies F_{17}(\delta, \beta) \in \text{IV}. \ [\text{to}]$$

F17 :: ItV -> IV -> IV

S9.
$$\delta \in \text{SmA}, \ \phi \in T \implies F_6(\delta, \phi) \in T.$$
 [necessarily]

TABLE 7-3

		TABLE 7-3	
Category Name	Categorial Definition of Name	Corresponding Type by Bennett's Rule (7-3)	Name of Semantical Object Denoted by this Type
t	t	1	truth value
CN	CN	(e, t)	set of individuals
IV	IV	(e, t)	set of individuals
T	t/IV	$\langle\langle s, \langle e, t \rangle\rangle, t \rangle$	set of properties of individuals
IAV	IV/IV	$\langle\langle s, \langle e, t \rangle\rangle, \langle e, t \rangle\rangle$	function from pro- perties of individuals to sets of individuals
TV	IV/T = IV/(t/IV)	$\langle\langle s, \langle\langle s, \langle e, t \rangle\rangle, t \rangle\rangle, \langle e, t \rangle\rangle$	function from pro- perties of properties of individuals to sets of individuals
T/CN	(t/IV)/CN	$((s, \langle e, t \rangle), ((s, \langle e, t \rangle), t \rangle)$	function from pro- perties of individuals to sets of properties of individuals
t/t	t/t	$\langle (s,t),t \rangle$	set of propositions
IV/t	IV/t	$\langle \langle s, t \rangle, \langle e, t \rangle \rangle$	function from propositions to sets of individuals
IV IV	IV IV	$\langle\langle s, \langle e, t \rangle\rangle, \langle e, t \rangle\rangle$	function from pro- perties of individuals to sets of individuals
IAV/T	(IV IV) T	((s, ((s, (e, t)), t)), ((s, (e, t)), (e, t))	function from pro- perties of properties of individuals to functions from pro- perties of individuals to sets of individuals

Figure 3: Mapping

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F6 :: SmA -> T -> T
S10. \delta \in IAV, \beta \in IV \implies F_7(\delta, \beta) \in IV. [adverb]
F7 :: IAV -> IV -> IV
S11. \phi, \psi \in T \implies F_8(\phi, \psi), F_9(\phi, \psi) \in T. [and, or]
F8 :: T -> T -> T
F9 :: T -> T -> T
S12. \delta, \gamma \in IV \implies F_8(\delta, \gamma), F_9(\delta, \gamma) \in IV. [and, or]
F8 :: IV -> IV -> IV
F9 :: IV -> IV -> IV
S13. \alpha, \beta \in \text{Tm} \implies F_9(\alpha, \beta) \in \text{Tm. [or]}
F9 :: Tm -> Tm -> Tm
S14. \alpha \in \text{Tm}, \phi \in \text{T} \implies F_{10,n}(\alpha,\phi) \in \text{T.} [replacement]
F10 :: Int -> Tm -> T -> T
S17. \alpha \in \text{Tm}, \delta \in \text{IV} \implies F_{12}(\alpha, \delta), F_{13}(\alpha, \delta), F_{14}(\alpha, \delta), F_{15}(\alpha, \delta), F_{16}(\alpha, \delta) \in
T. [neg., fut., neg. fut., pres. perf., neg. pres. perf.]
F12 :: Tm -> IV -> T
F13 :: Tm -> IV -> T
F14 :: Tm -> IV -> T
F15 :: Tm -> IV -> T
F16 :: Tm -> IV -> T
Translation Rules
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T2. $F_2(\delta,\zeta) \to \delta'(^{\wedge}\zeta')$.

T3. $F_{3,n}(\zeta,\phi) \to \lambda x_n(\zeta'(x_n) \wedge \phi')$.

T4. $F_4(\alpha, \delta) \to \alpha'(^{\wedge}\delta')$.

T5. $F_5(\delta, \beta) \to \delta'(^{\wedge}\beta')$.

T6. $F_5(\delta, \alpha) \to \delta'(^{\wedge}\alpha')$.

T7. $F_{11}(\alpha, \phi) \to \alpha'(^{\wedge}\phi')$.

T8. $F_{17}(\delta, \beta) \rightarrow \delta'(^{\wedge}\beta')$.

T9. $F_6(\delta, \phi) \to \delta'(^{\wedge}\phi')$.

T10. $F_7(\delta, \beta) \to \delta'(^{\wedge}\beta')$.

T11. $F_8(\phi, \psi) \rightarrow [\phi' \wedge \psi]$ $F_9(\phi, \psi) \rightarrow [\phi' \vee \psi'].$

T12. $F_8(\delta, \gamma) \to \lambda x [\delta'(x) \land \gamma'(x)]$ $F_9(\delta, \gamma) \to \lambda x [\delta'(x) \lor \gamma'(x)].$

T13. $F_9(\alpha, \beta) \to \lambda P[\alpha'(P) \land \beta'(P)].$

T14. $F_{10,n}(\alpha,\phi) \to \alpha'(^{\wedge}\lambda x_n \phi')$.

T17. $F_{12}(\alpha, \delta) \rightarrow \neg \alpha'(^{\wedge}\delta')$

 $F_{13}(\alpha,\delta) \to \mathbf{F}\alpha'(^{\wedge}\delta')$

 $F_{14}(\alpha, \delta) \rightarrow \neg \mathbf{F} \alpha'(^{\wedge} \delta')$

 $F_{15}(\alpha, \delta) \to \mathbf{P}\alpha'(^{\wedge}\delta')$

 $F_{16}(\alpha, \delta) \to \neg \mathbf{P}\alpha'(\hat{\delta}').$