

Efficient Globally Optimal Consensus Maximisation with Tree Search

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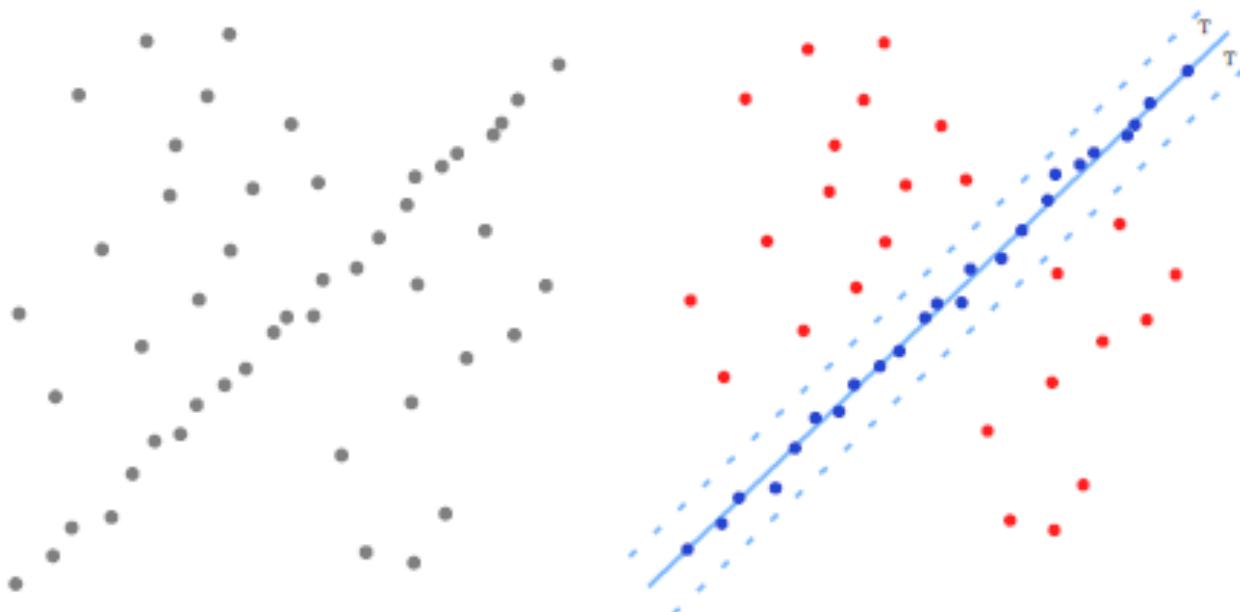
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Maximum consensus

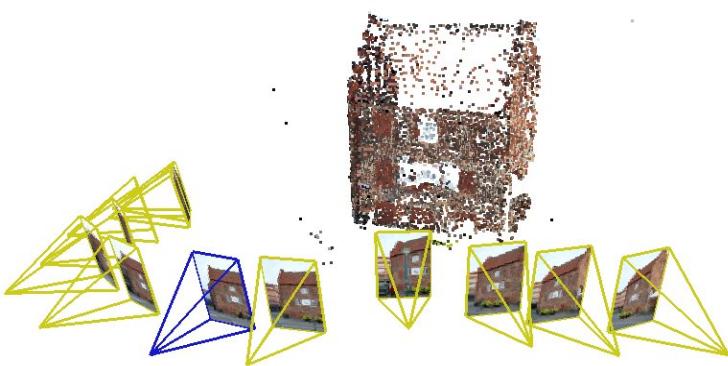
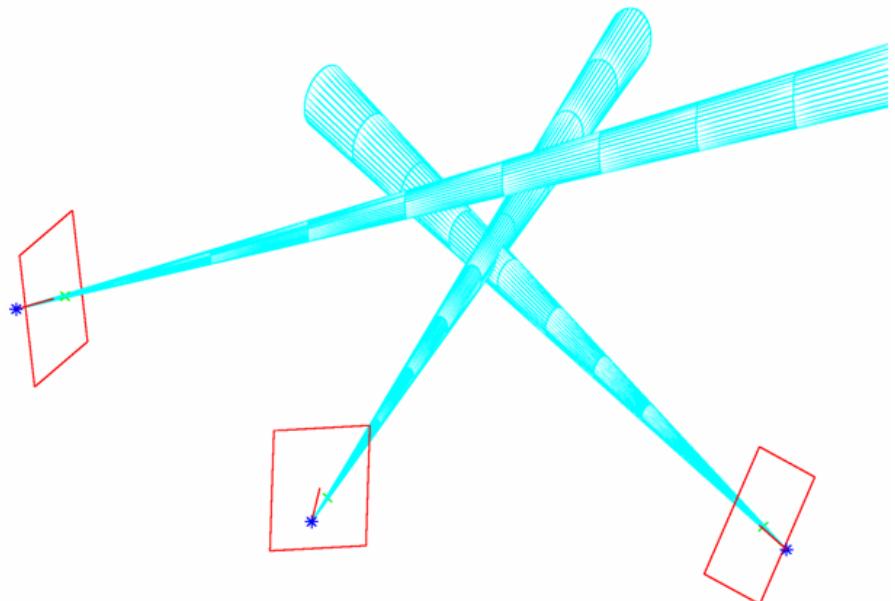
$$\max_{\theta, \mathcal{I} \subseteq \mathcal{X}} \quad |\mathcal{I}|$$

subject to $r_i(\theta) \leq \epsilon \quad \forall \mathbf{x}_i \in \mathcal{I}$

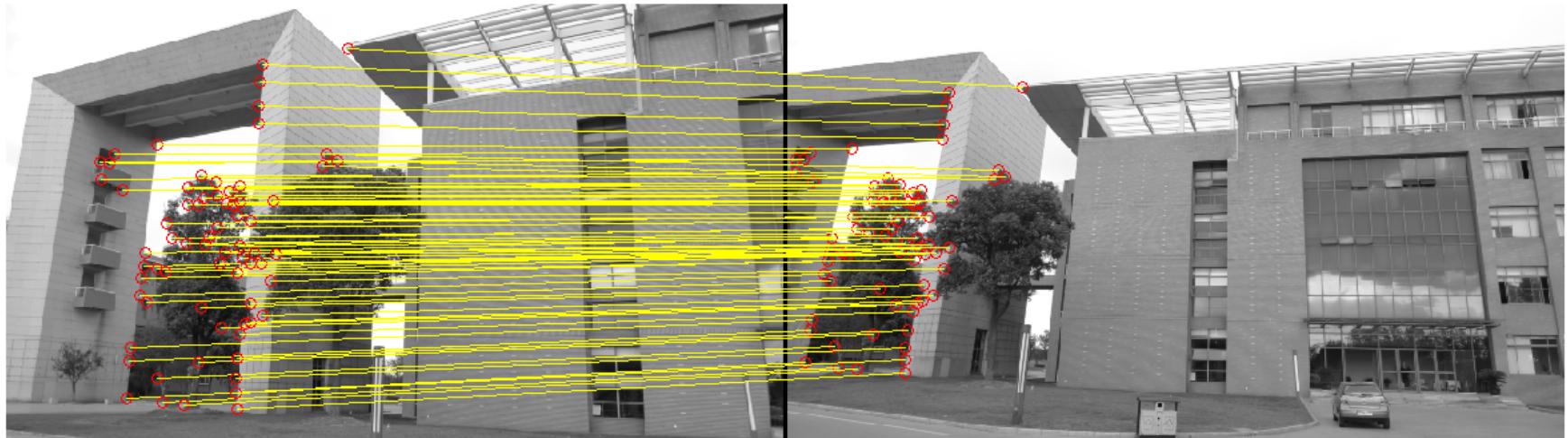
Example 1: Line fitting



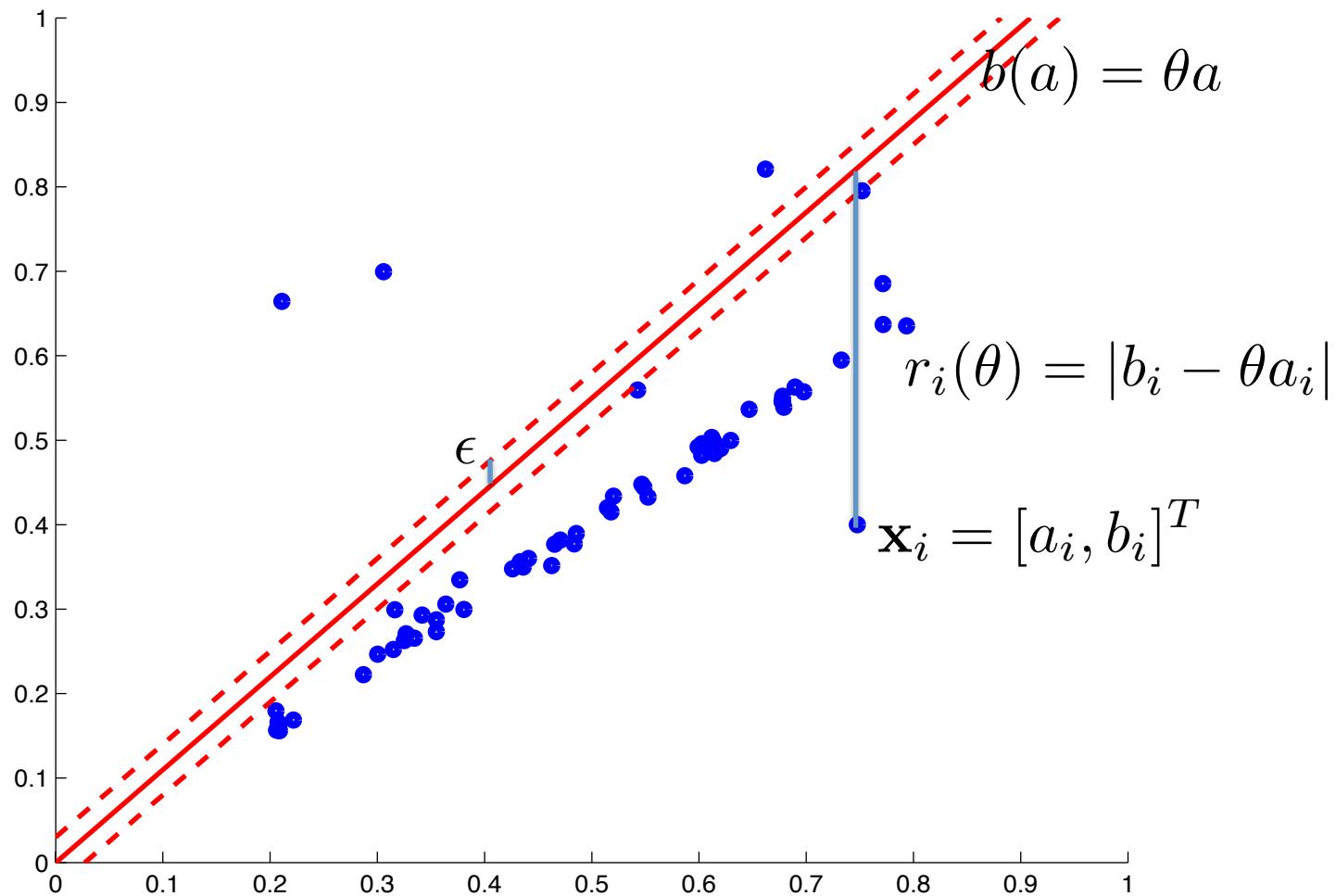
Example 2: Triangulation



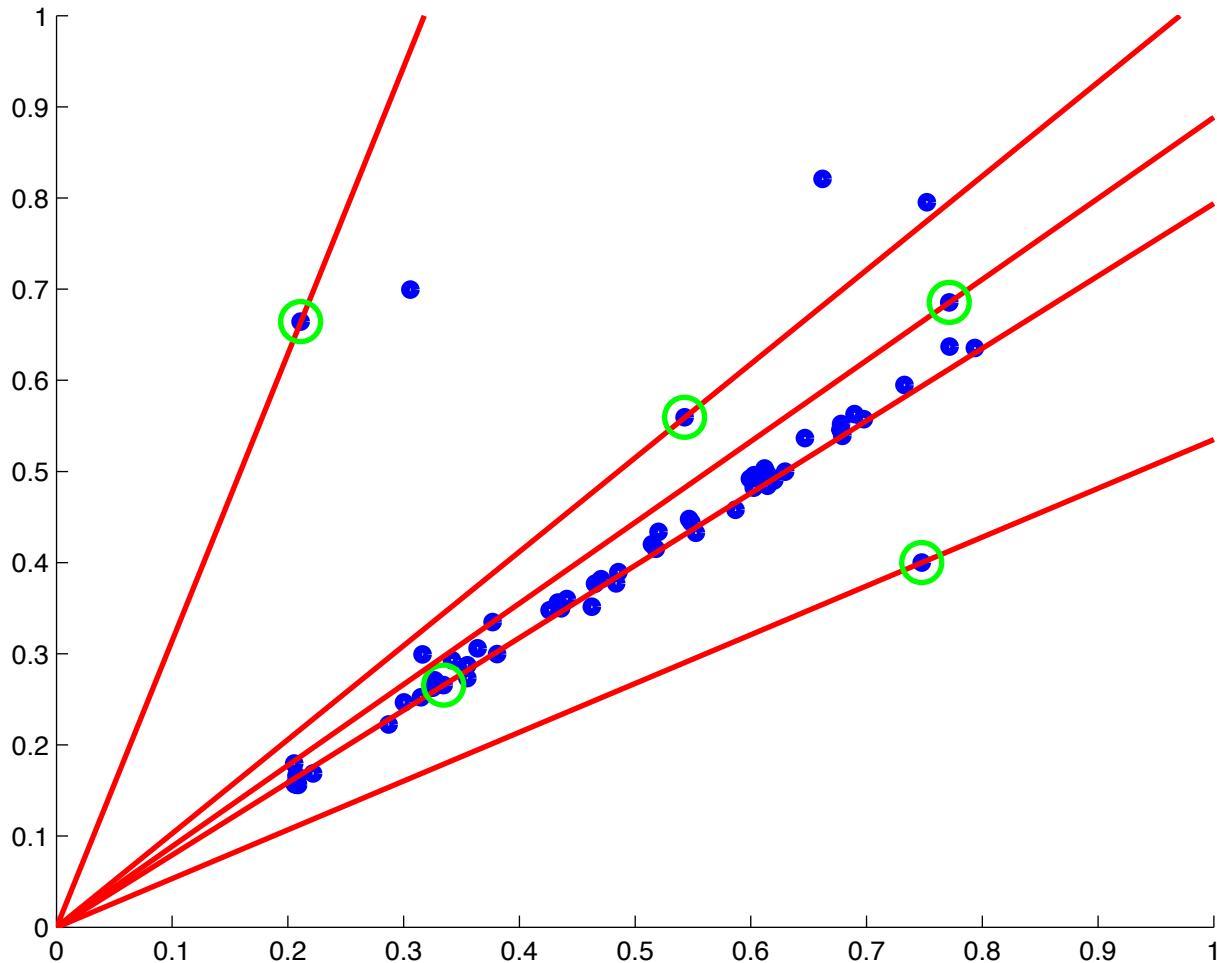
Example 3: Homography fitting



Running example: Linear regression



RANSAC, minimal subset size = p



Minmax problem

- Minimise the maximum residual:

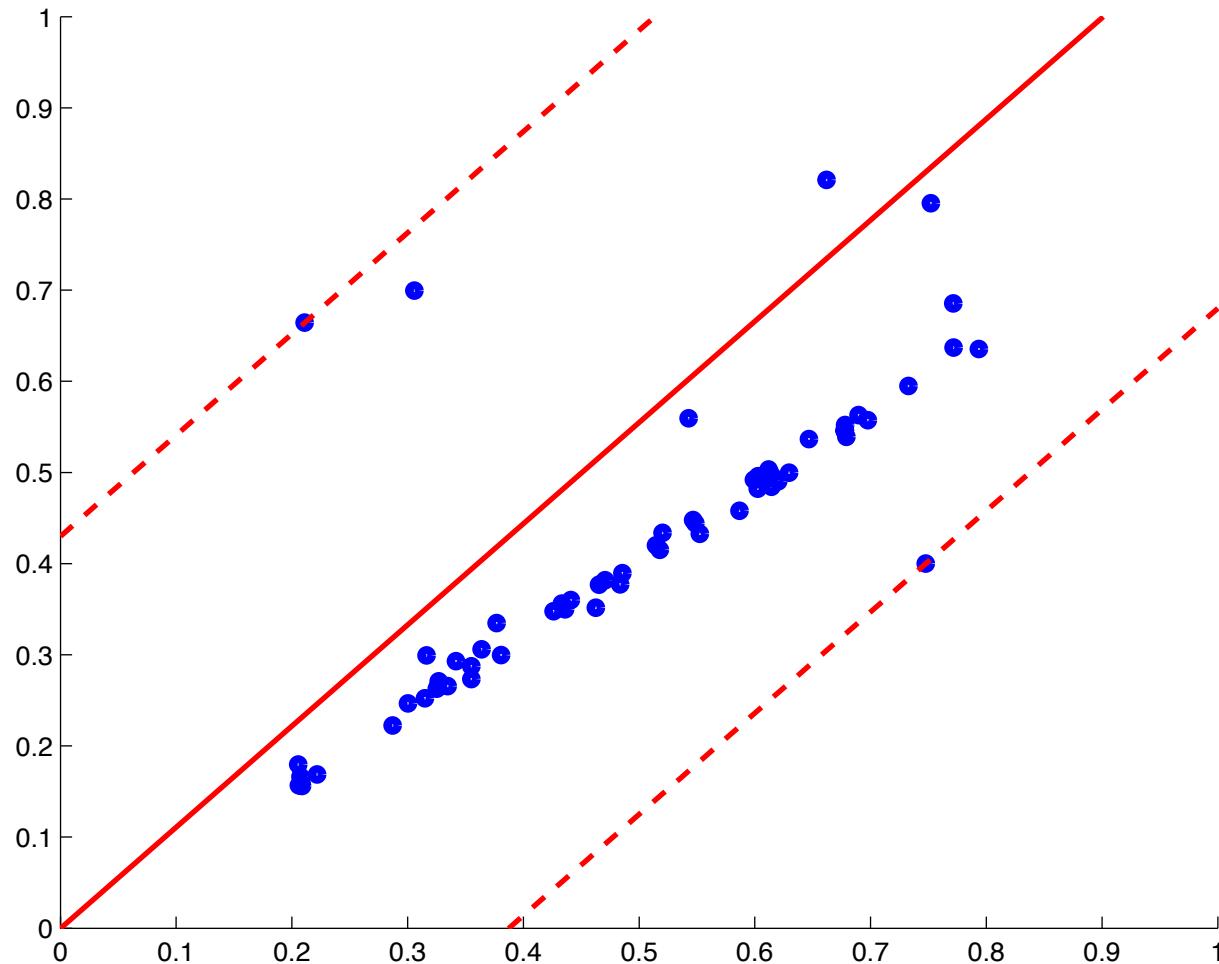
$$\min_{\theta} \max_i |b_i - \theta a_i|$$

- Same as L-infinity minimisation^{1,2}:

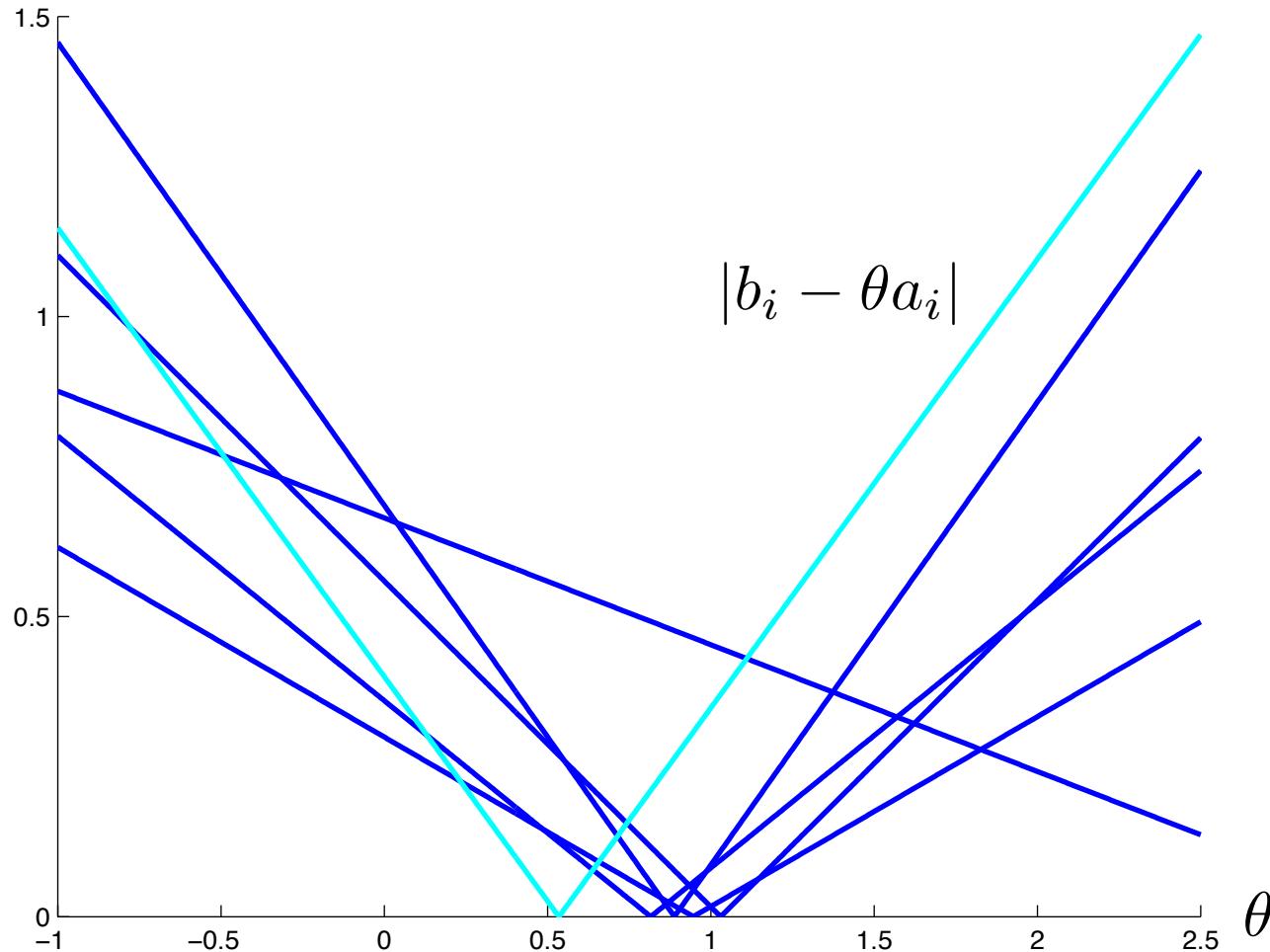
$$\min_{\theta} \left\| \begin{bmatrix} |b_1 - \theta a_1| \\ |b_2 - \theta a_2| \\ \vdots \\ |b_N - \theta a_N| \end{bmatrix} \right\|_{\infty}$$

- A.k.a. Chebyshev approximation/regression.

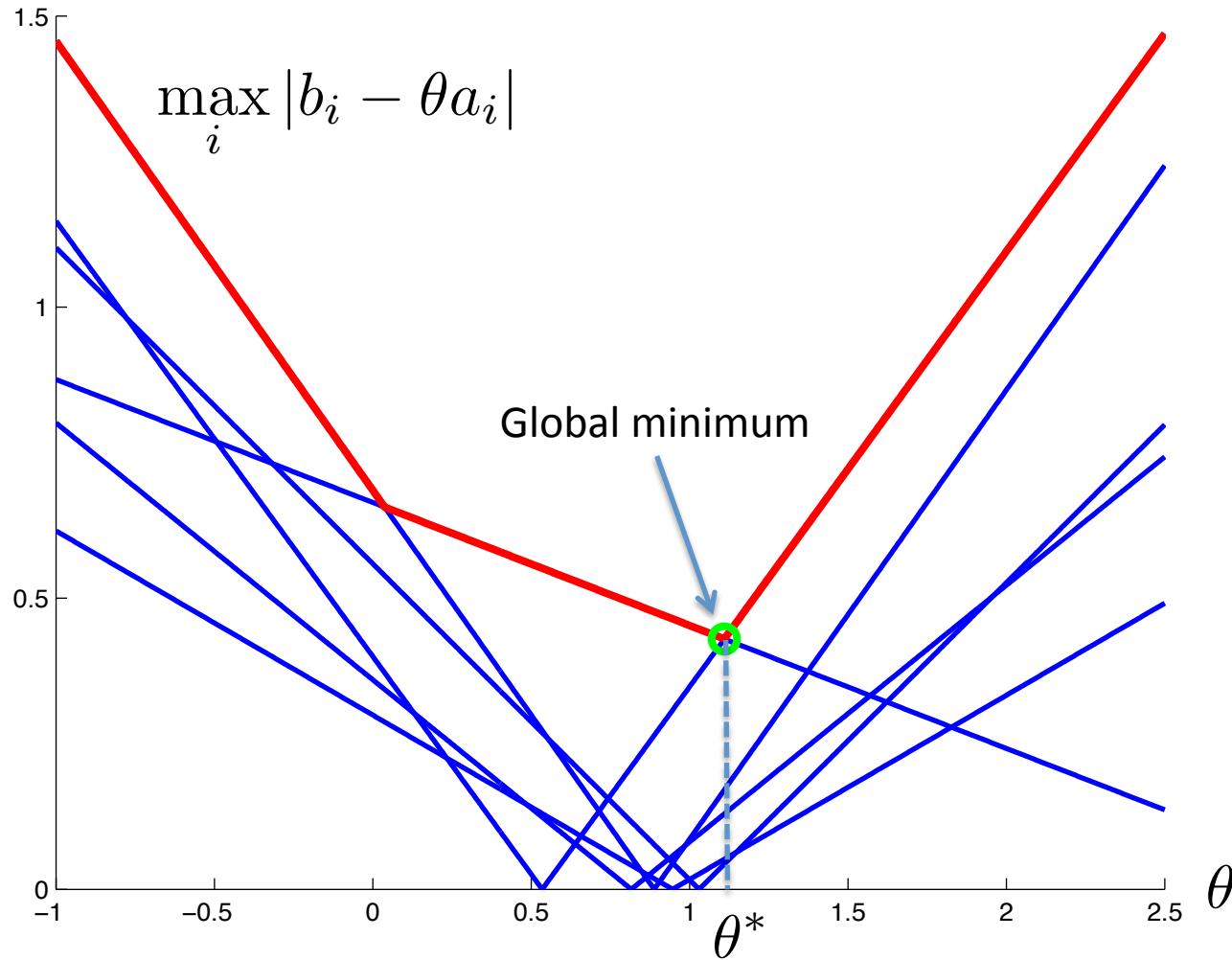
Minmax problem



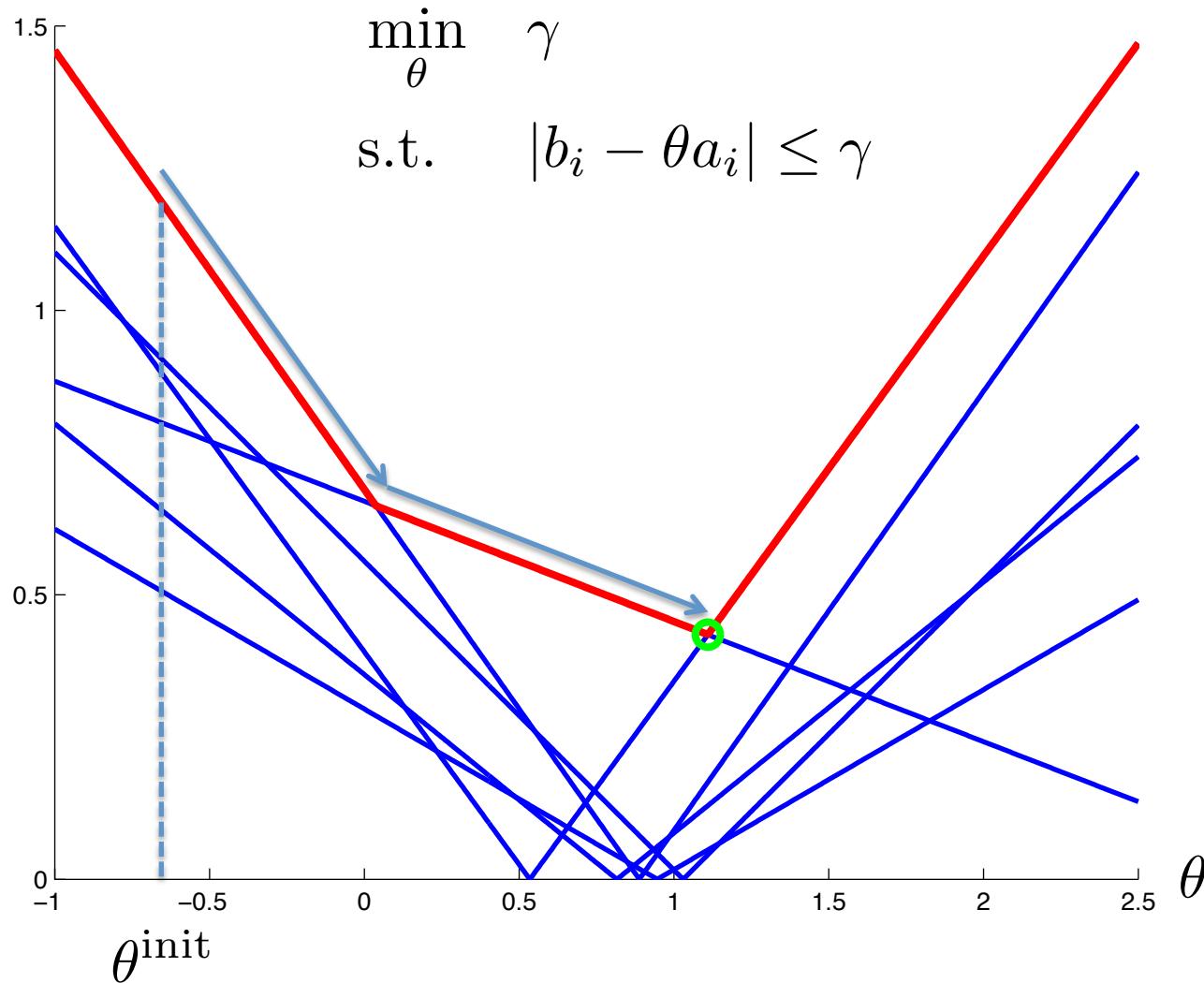
Minmax problem



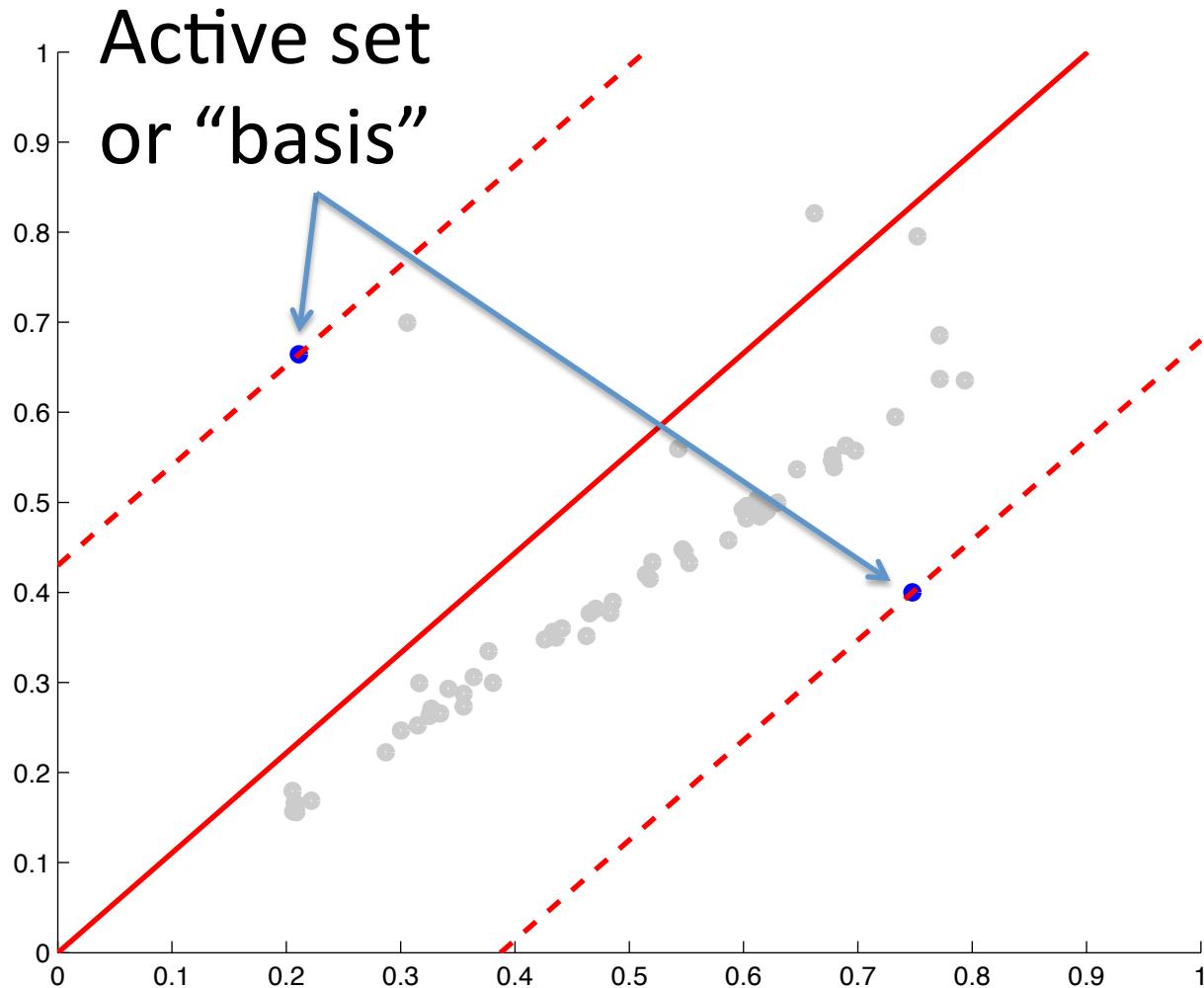
Minmax problem



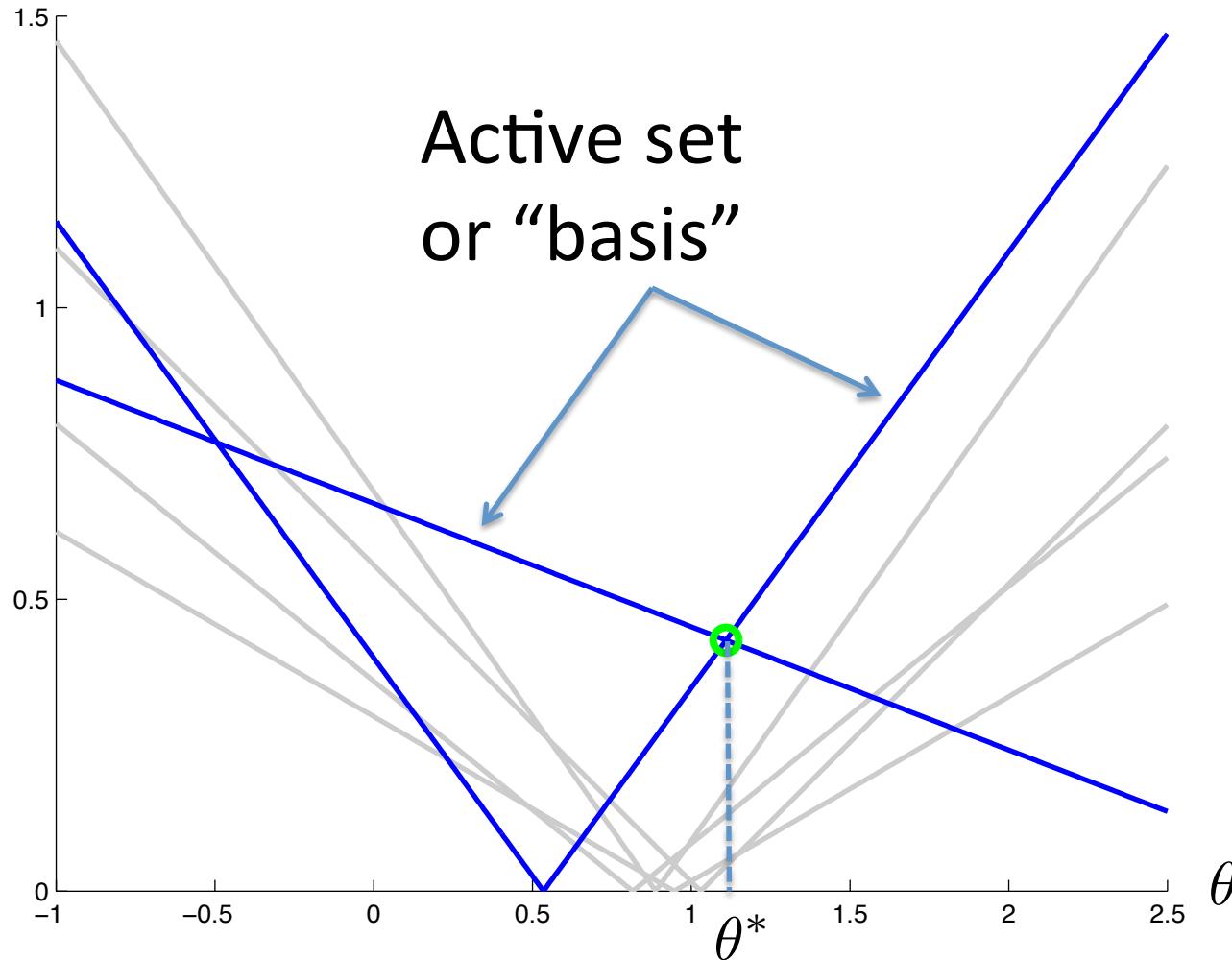
Simplex algorithm



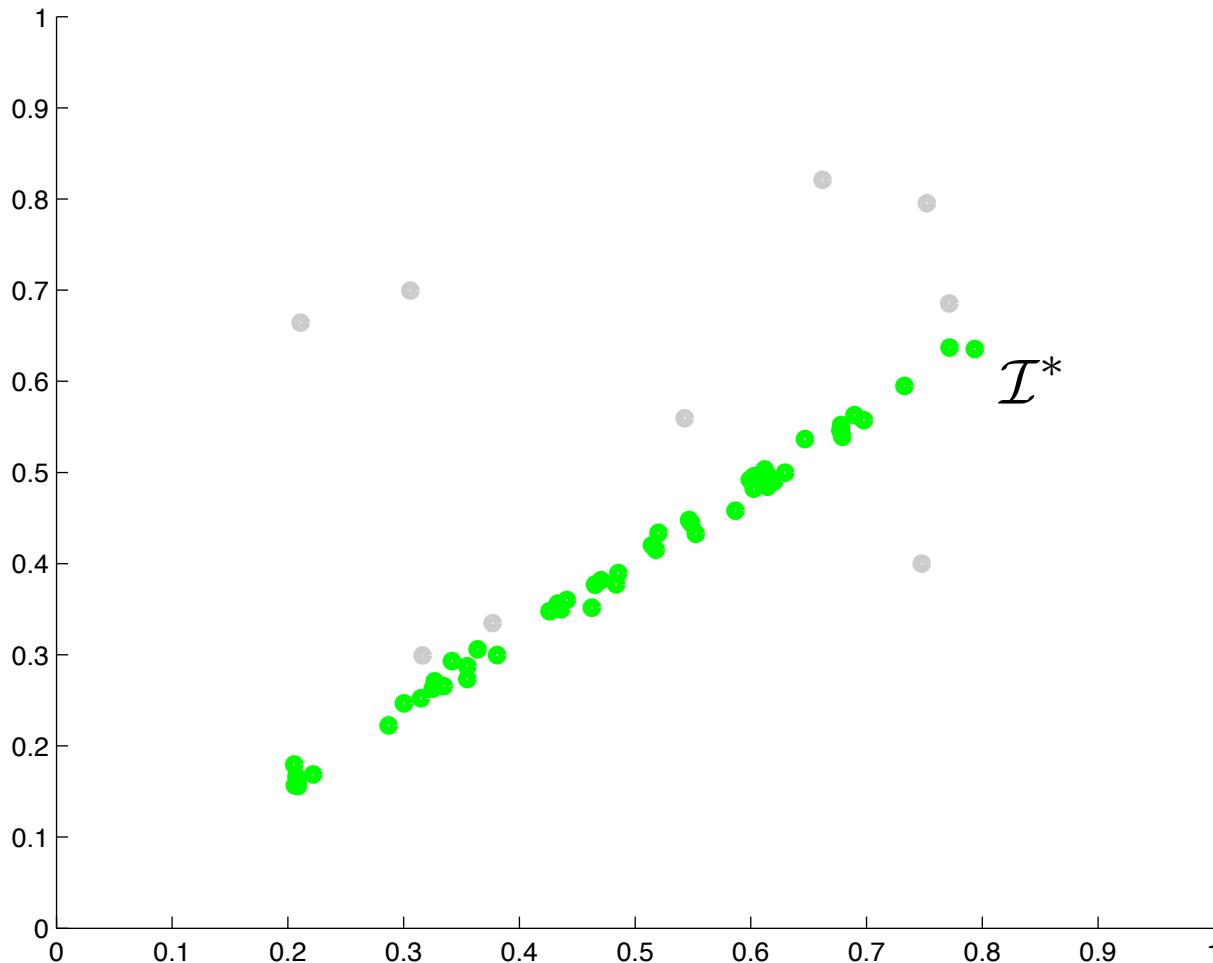
Combinatorial dimension = $p+1$



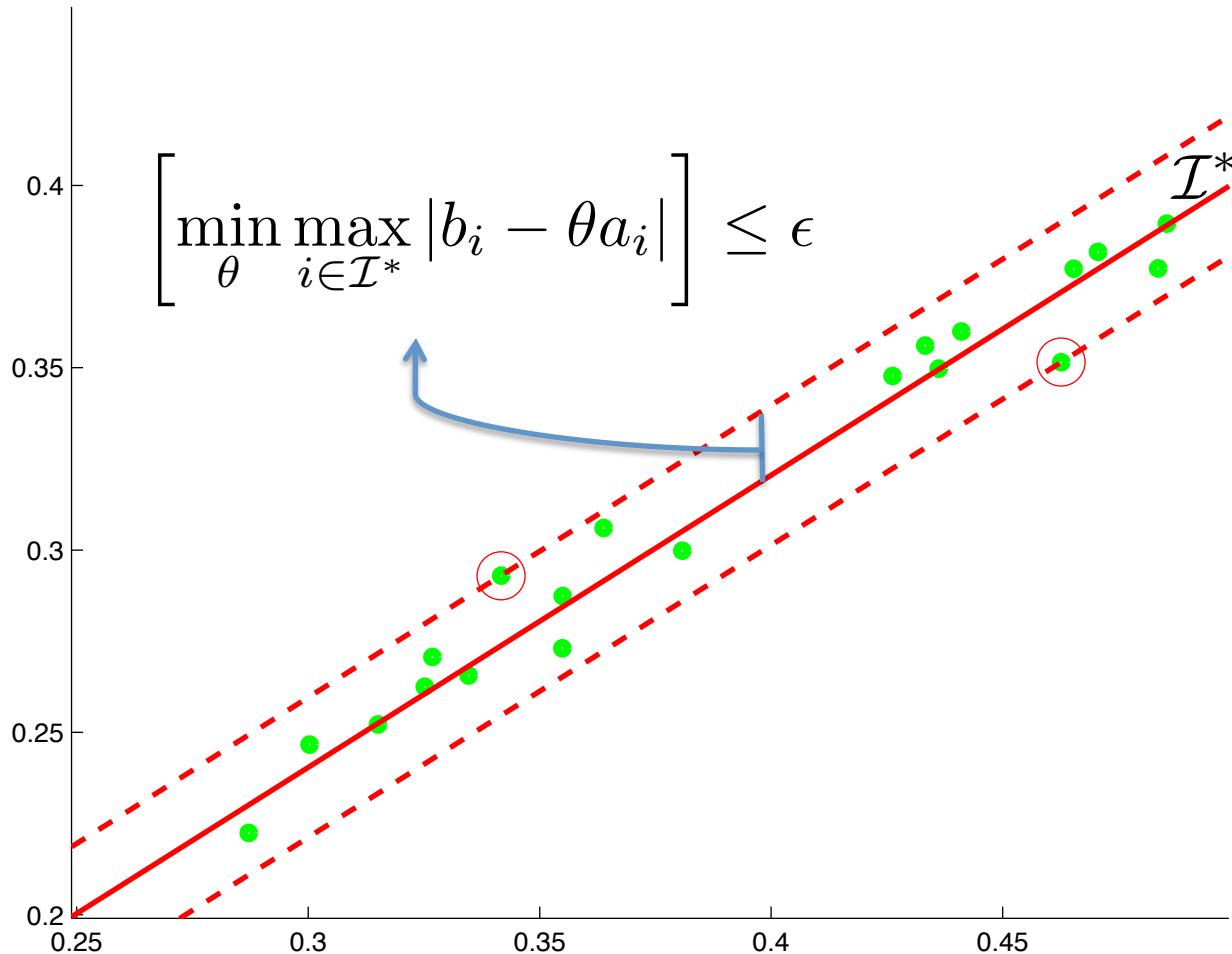
Combinatorial dimension = p+1



Maximum consensus



Maximum consensus

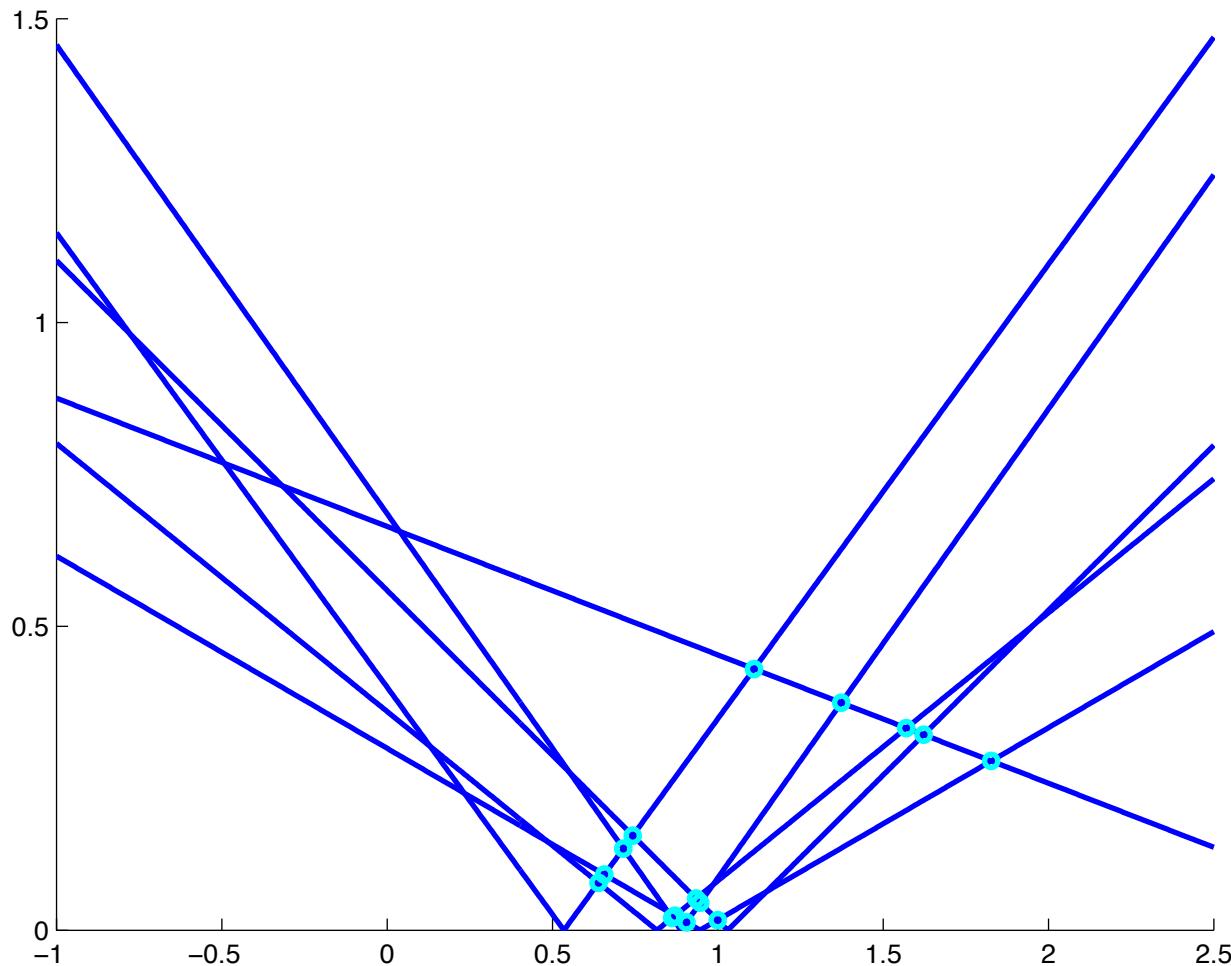


An algorithm

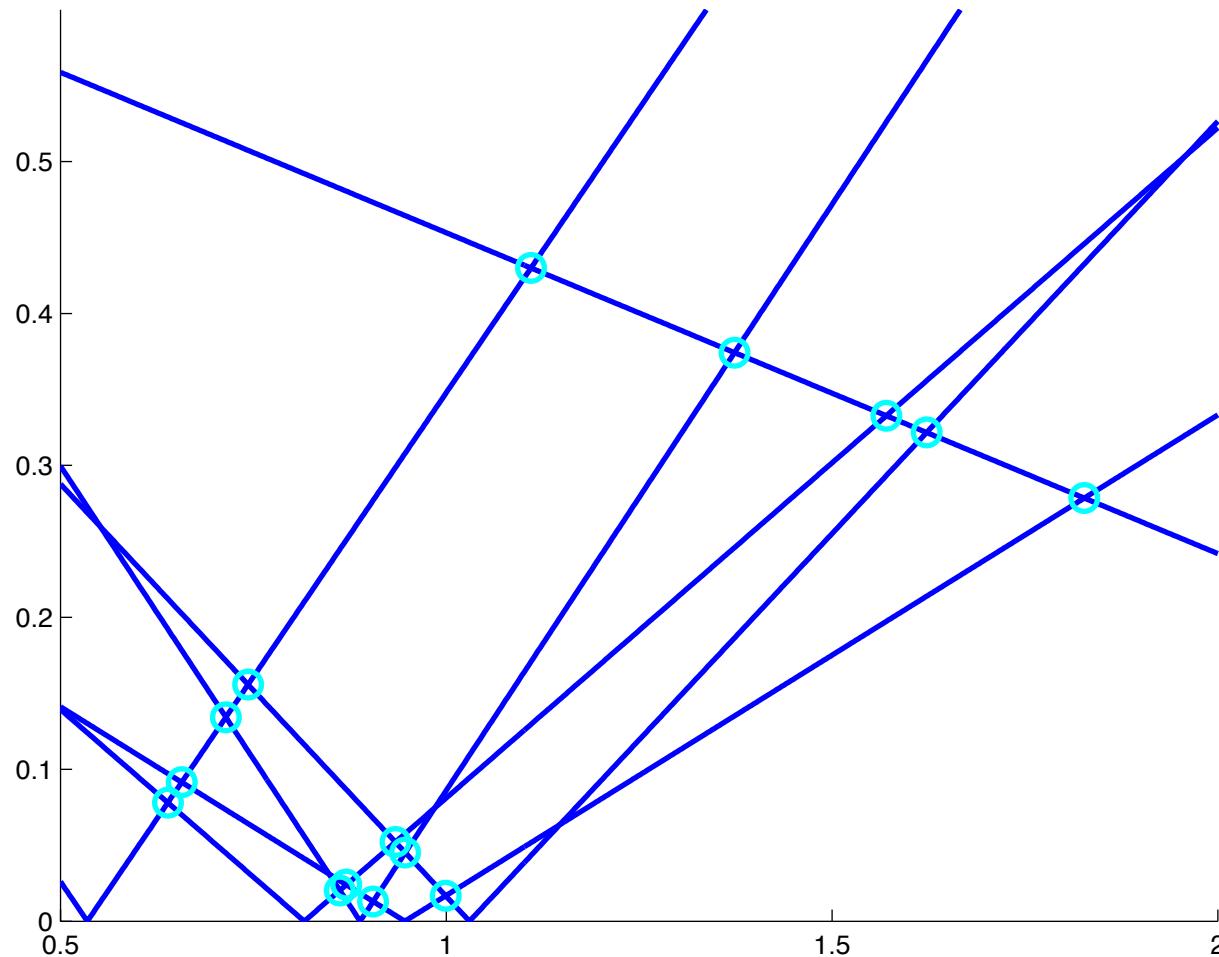
- For all subsets $\mathcal{B} \subset \mathcal{X}$ of size $(p+1)$;
 - Solve minmax problem on \mathcal{B} .
 - If maximum residual of \mathcal{B} is $\leq \epsilon$;
 - If the coverage of \mathcal{B} is greater than the current largest;
 - Set \mathcal{I}^* as the coverage of \mathcal{B} .

$$\begin{aligned}\binom{N}{p+1} &= \frac{N!}{(p+1)!(N-p-1)!} \\ &= \frac{1}{(p+1)!} N(N-1)\dots(N-p) \equiv \mathcal{O}(N^{p+1})\end{aligned}$$

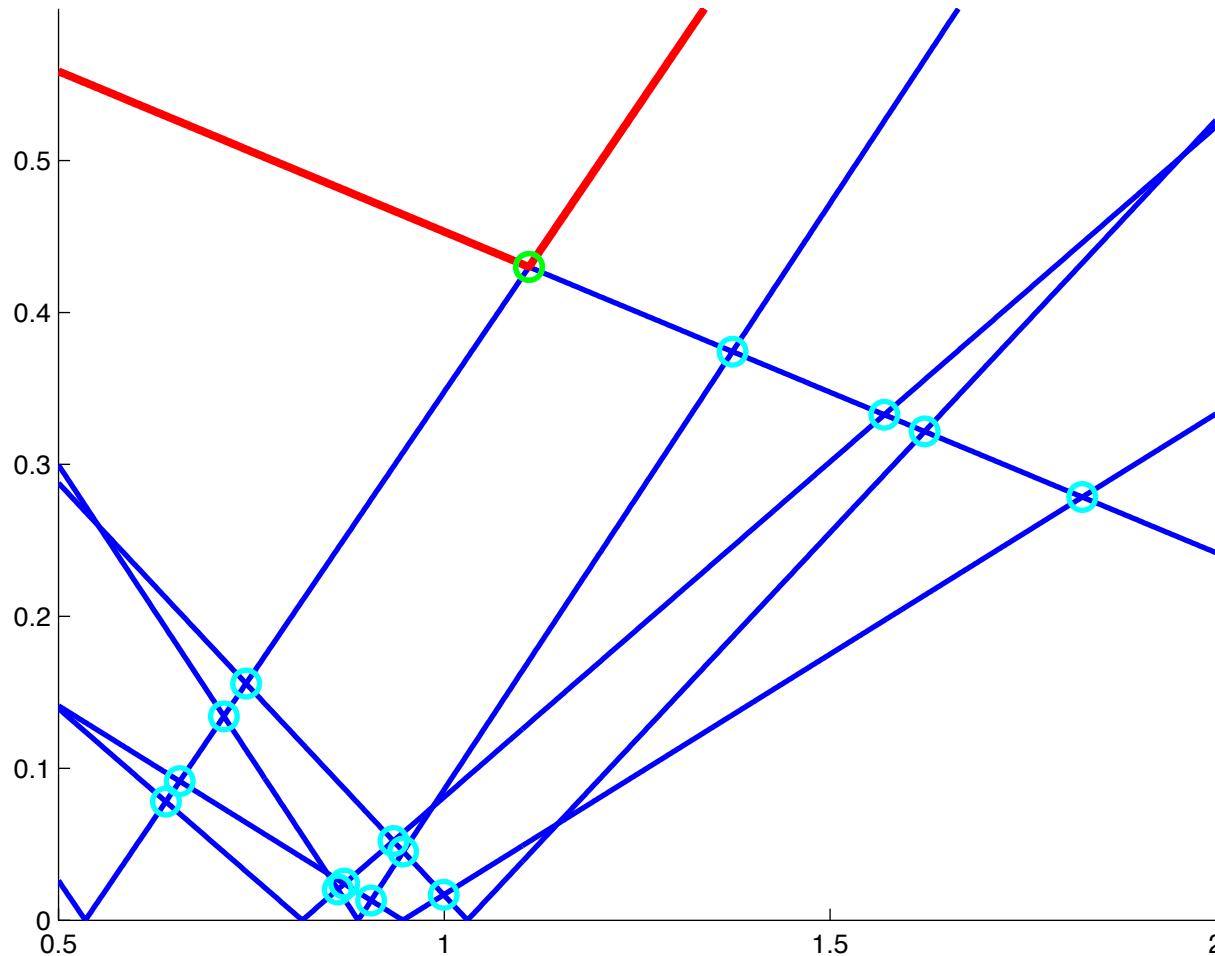
Minmax problem



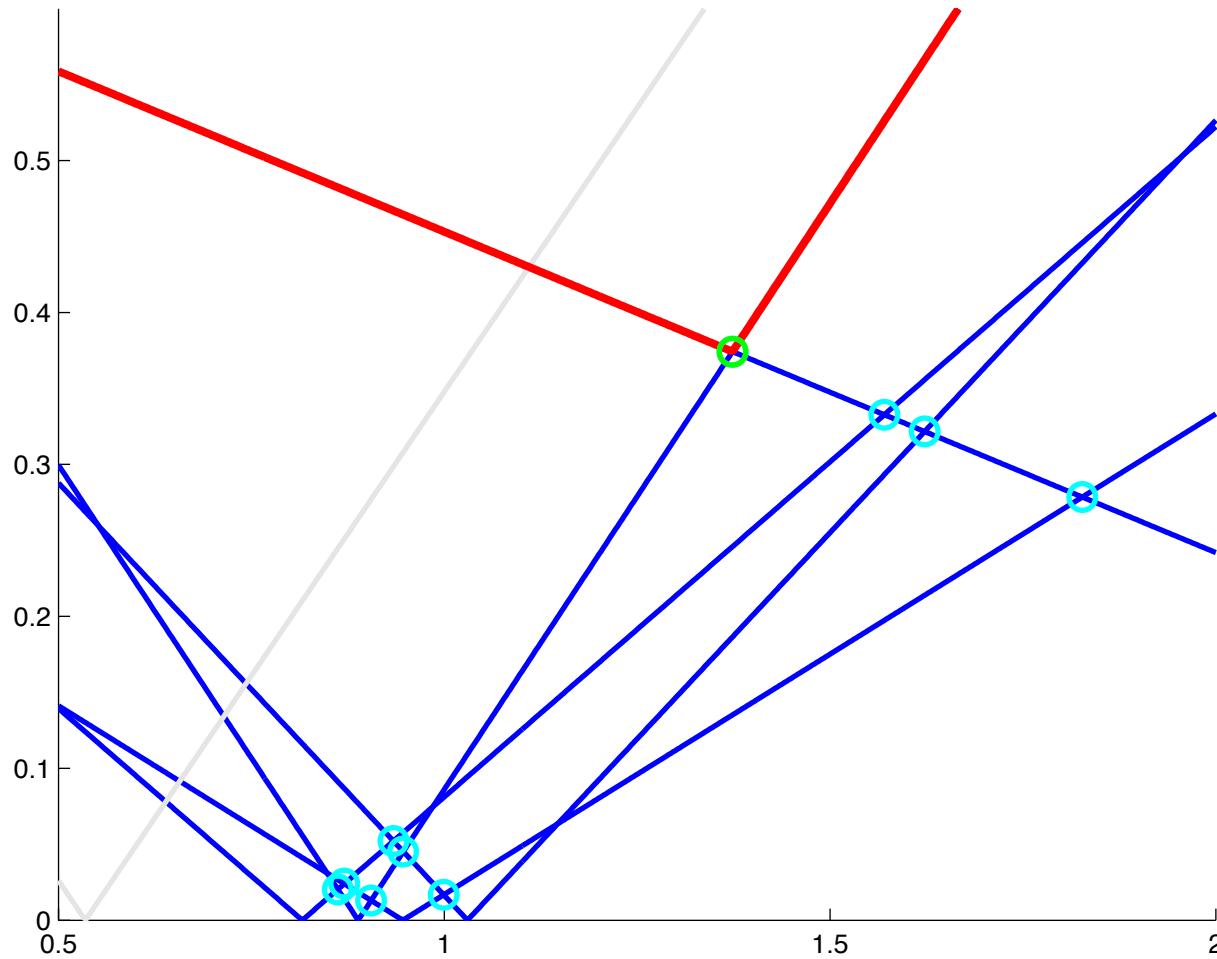
Minmax problem



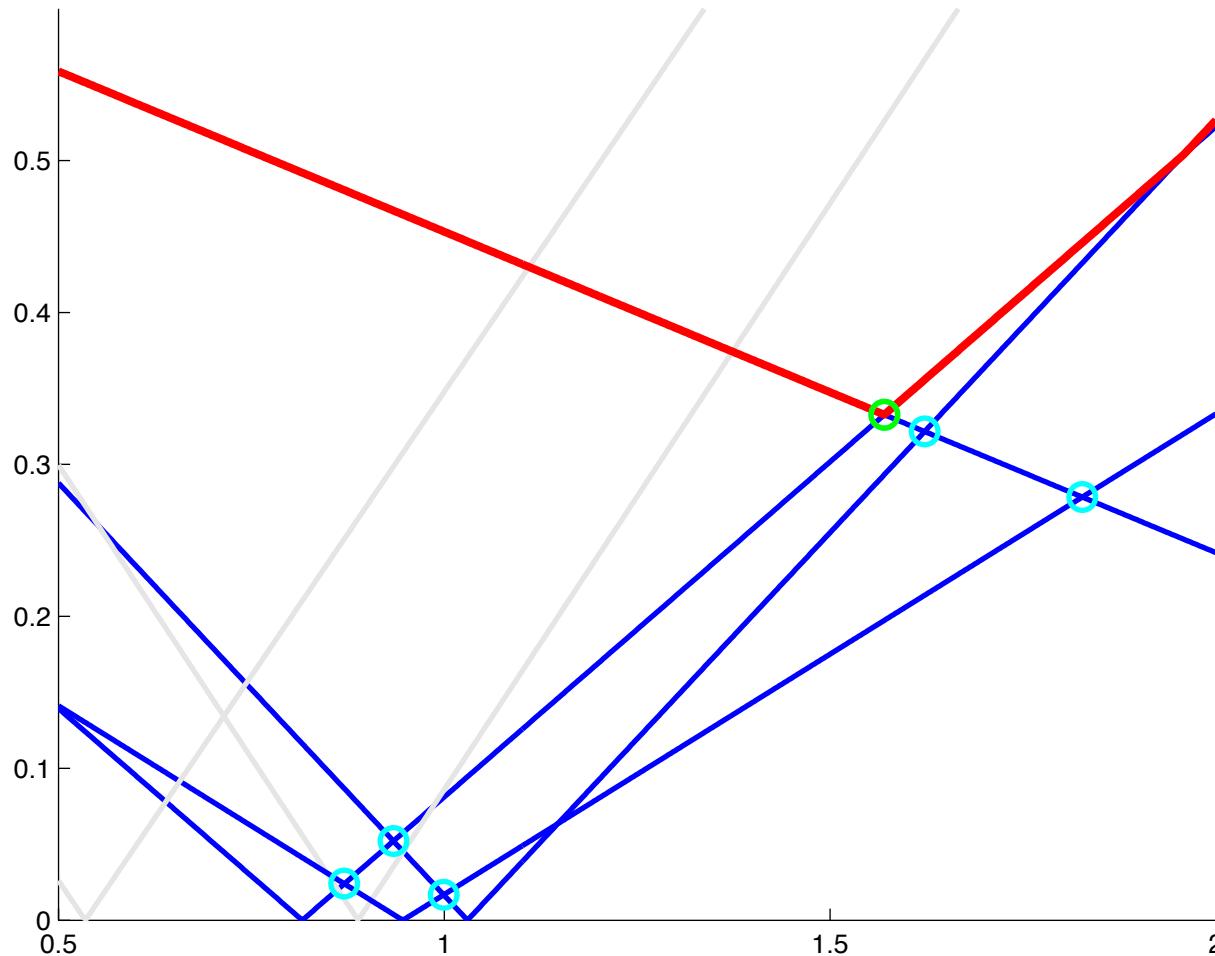
Recursive minmax



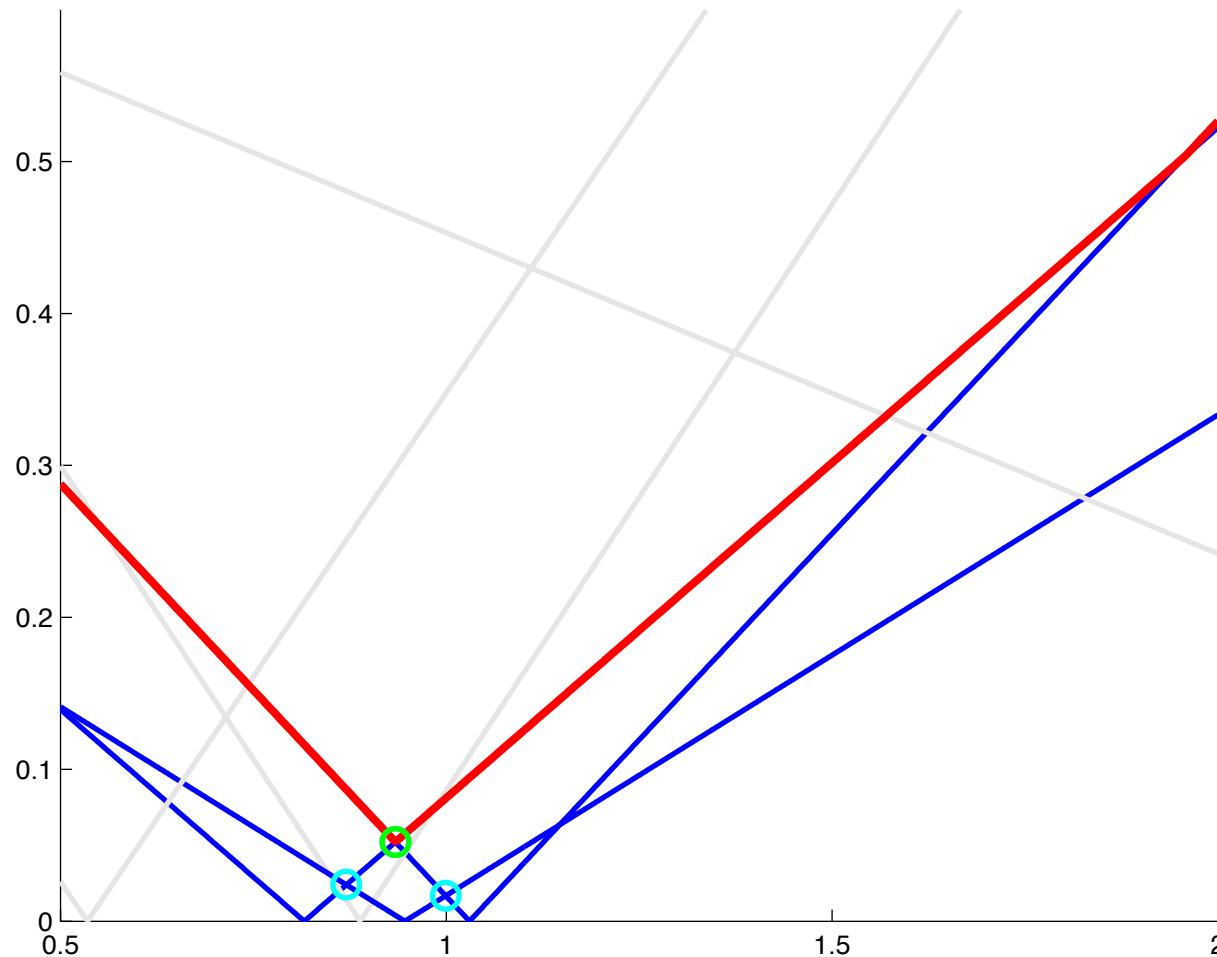
Recursive minmax



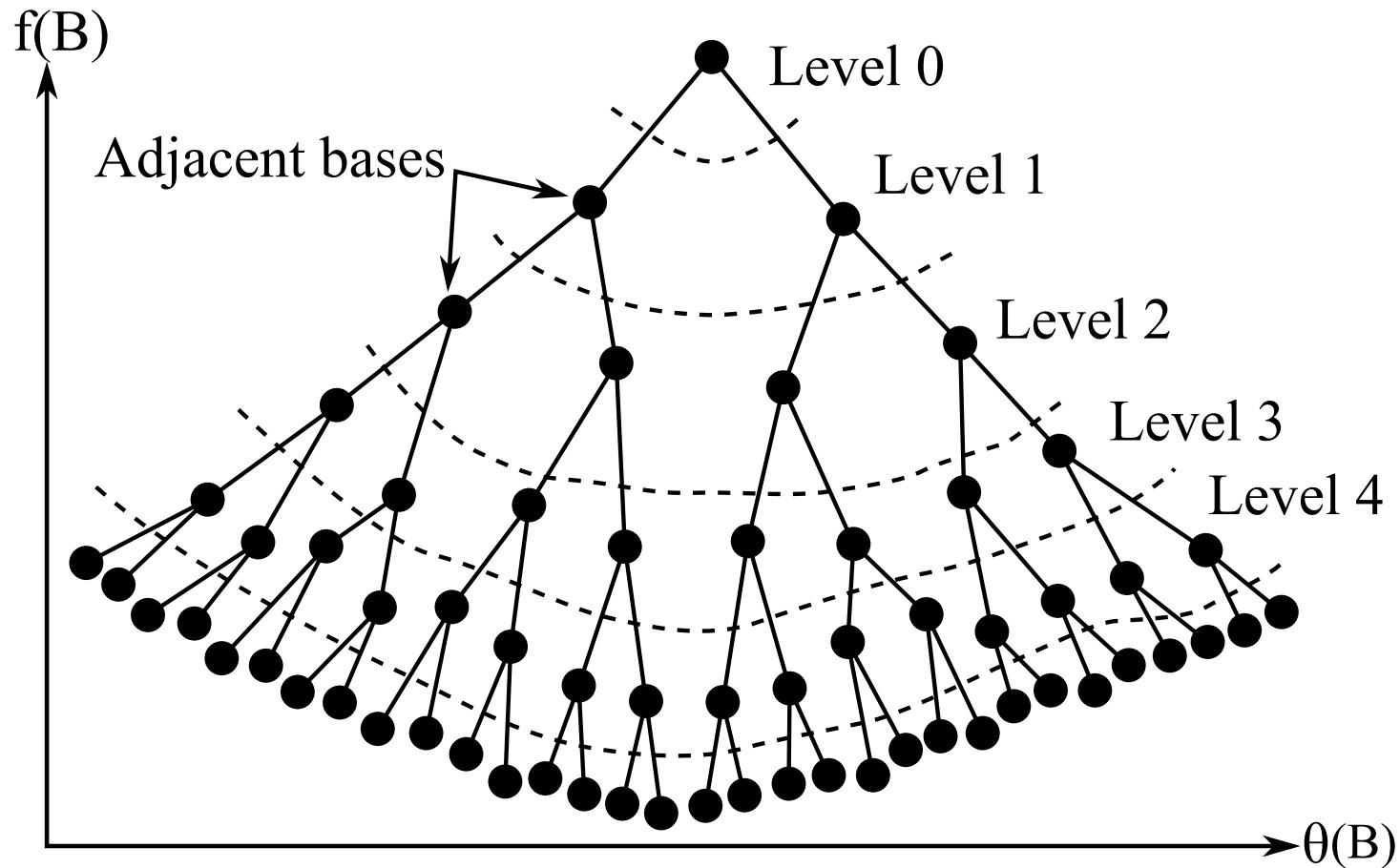
Recursive minmax



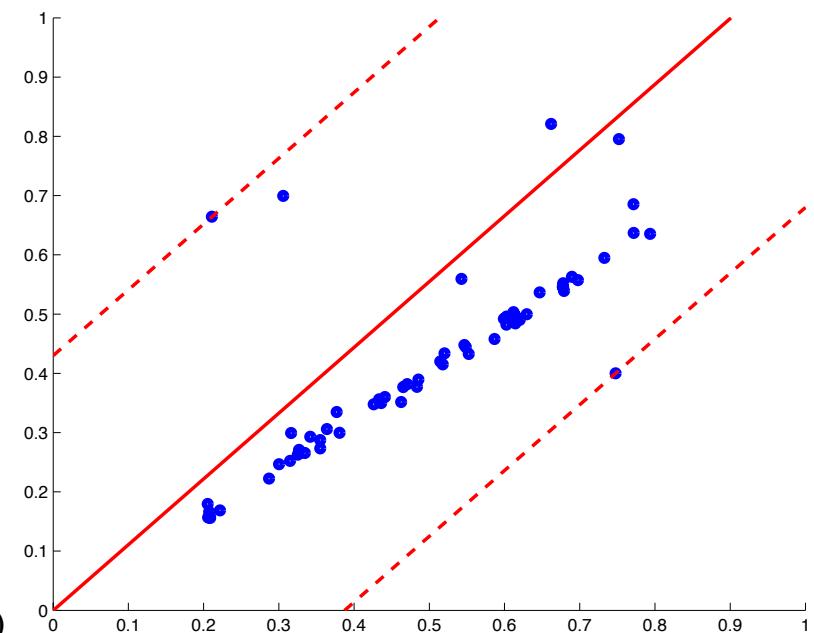
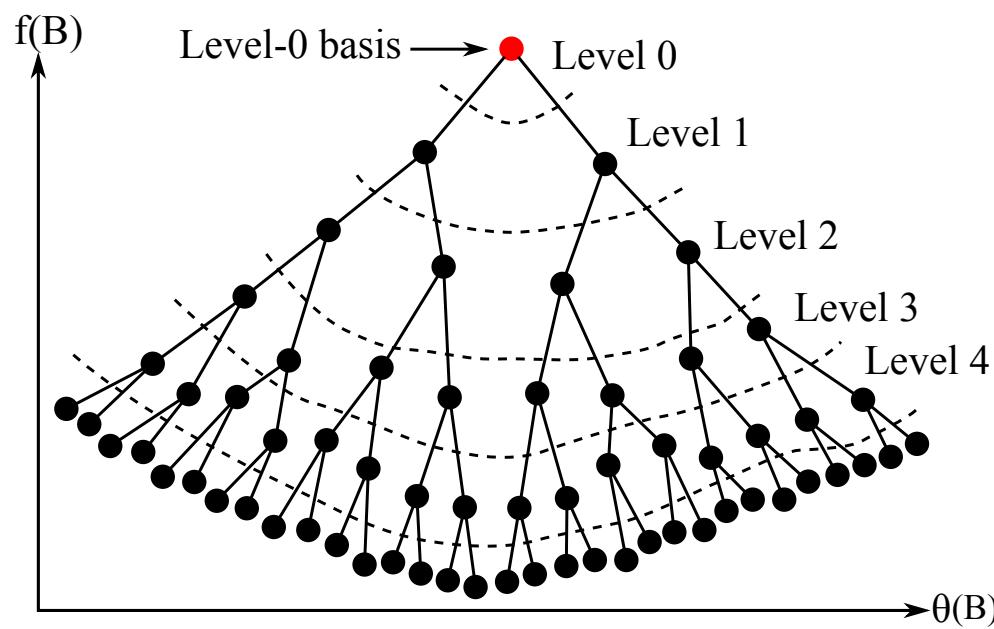
Recursive minmax



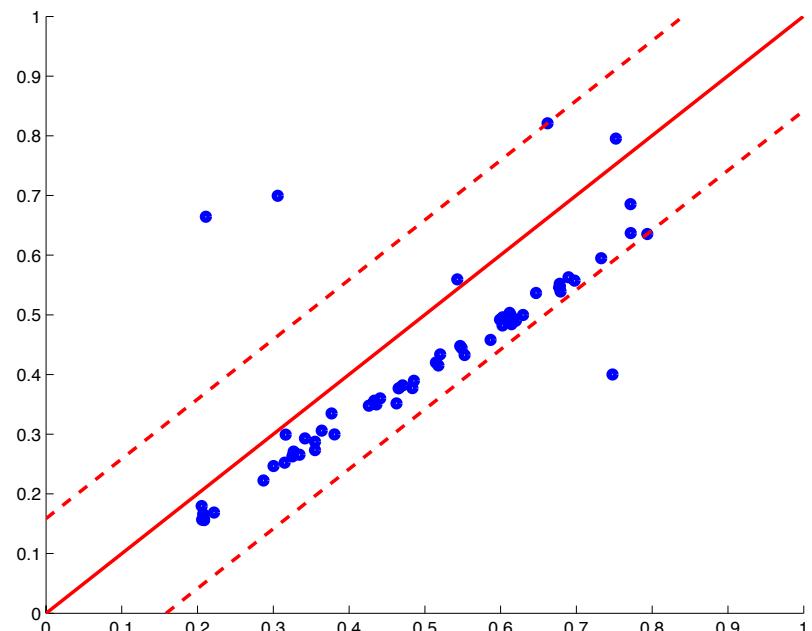
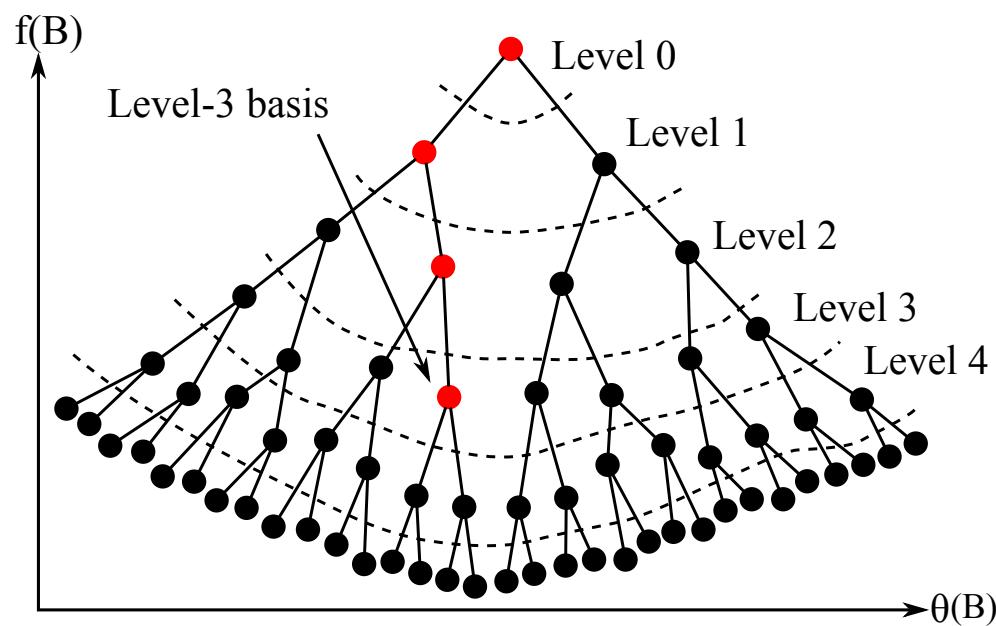
It's a tree!



Tree depth or level

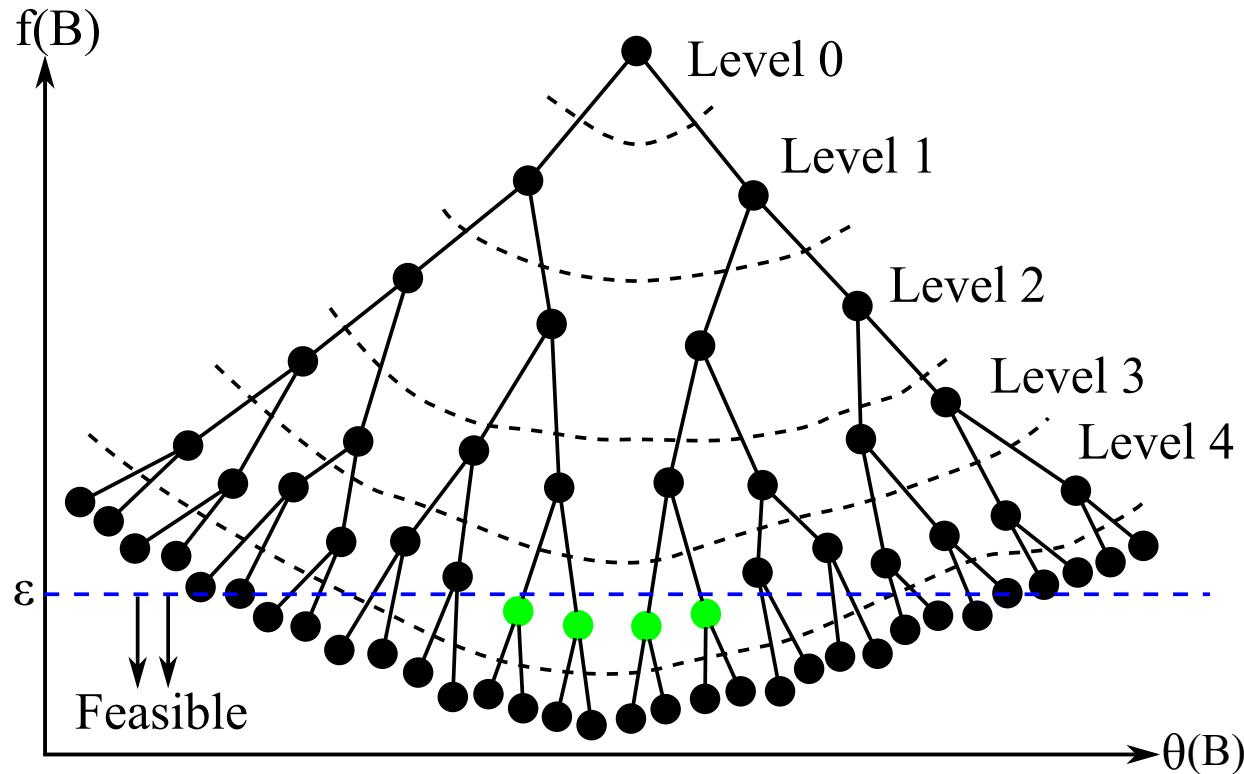


Tree depth or level

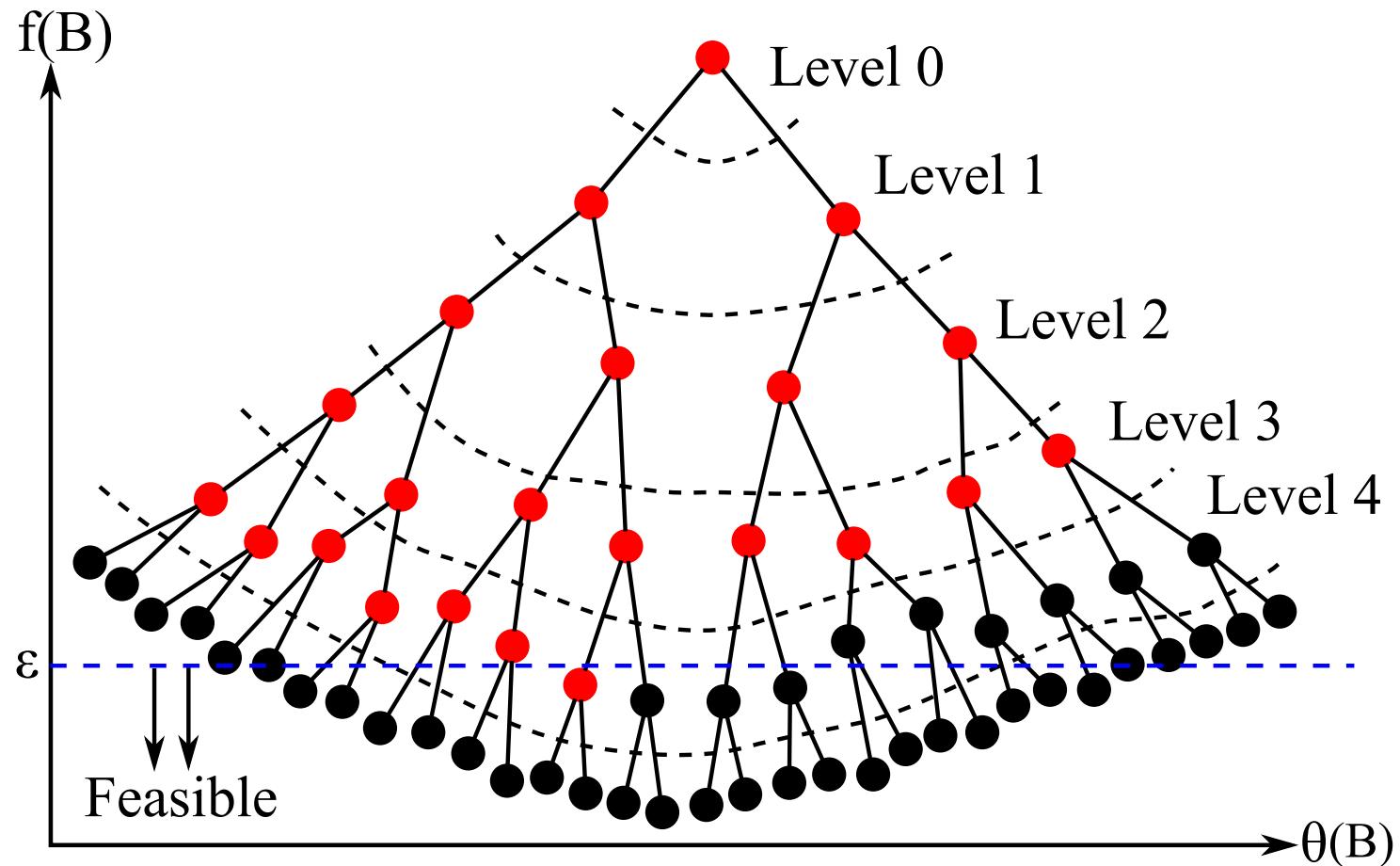


Maximum consensus II

$$\min_{\mathcal{B}} l(\mathcal{B}) \quad \text{s.t.} \quad f(\mathcal{B}) \leq \epsilon.$$



Breadth-first search (BFS)



A* algorithm

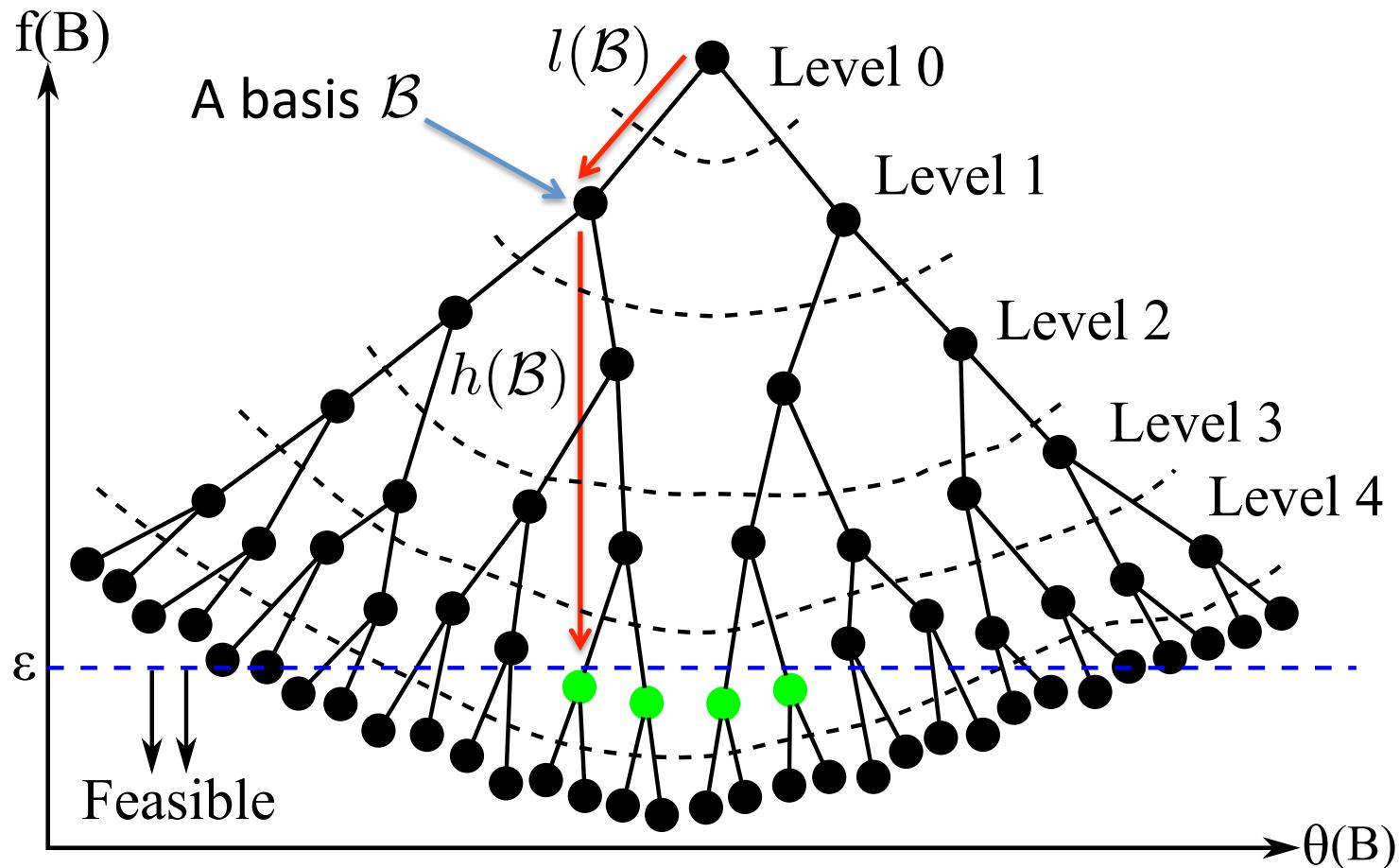
- Basis expansion is prioritised by

$$e(\mathcal{B}) = l(\mathcal{B}) + h(\mathcal{B})$$

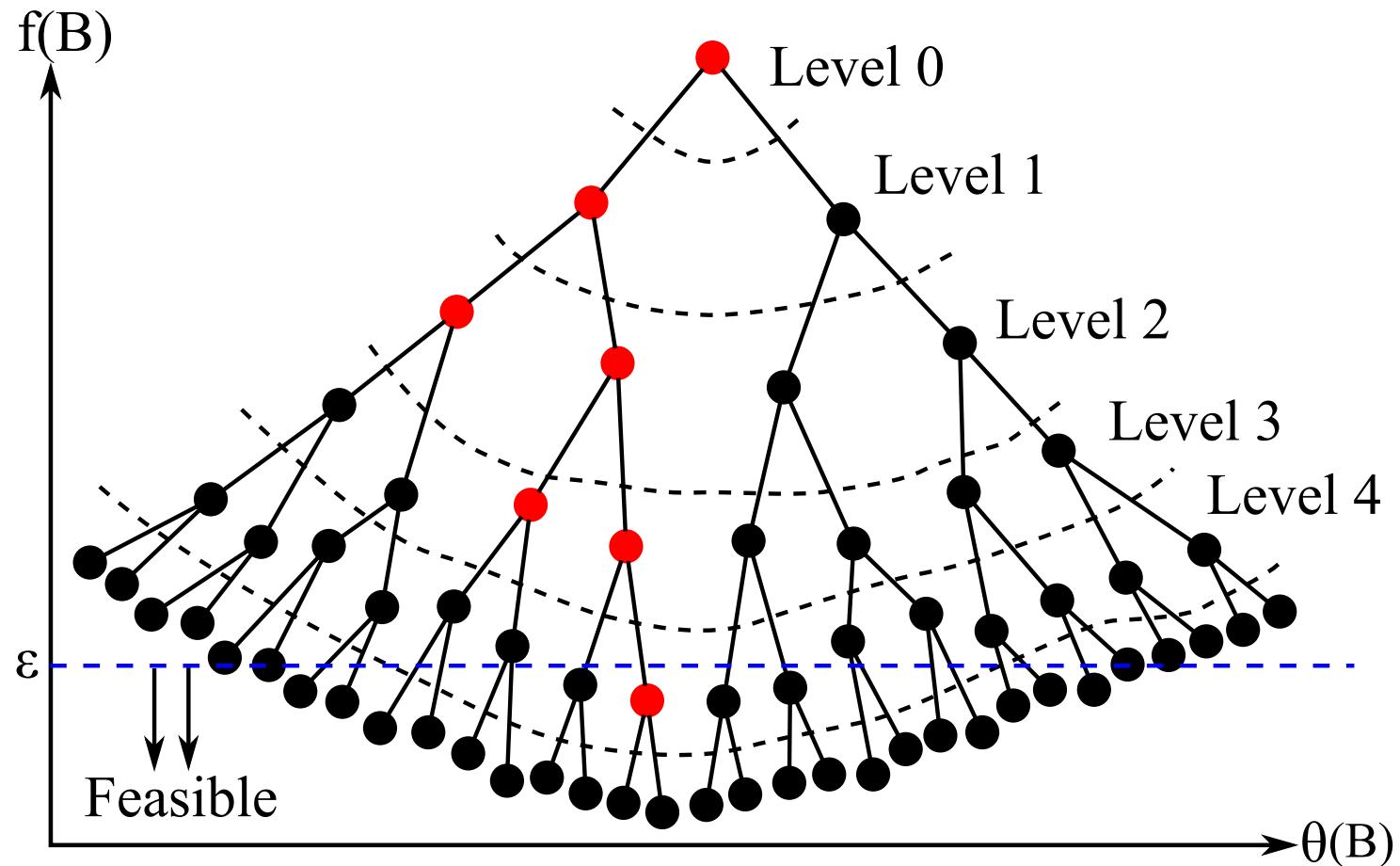
$l(\mathcal{B})$: Level of basis \mathcal{B} .

$h(\mathcal{B})$: An estimate of the number of steps remaining to feasibility.

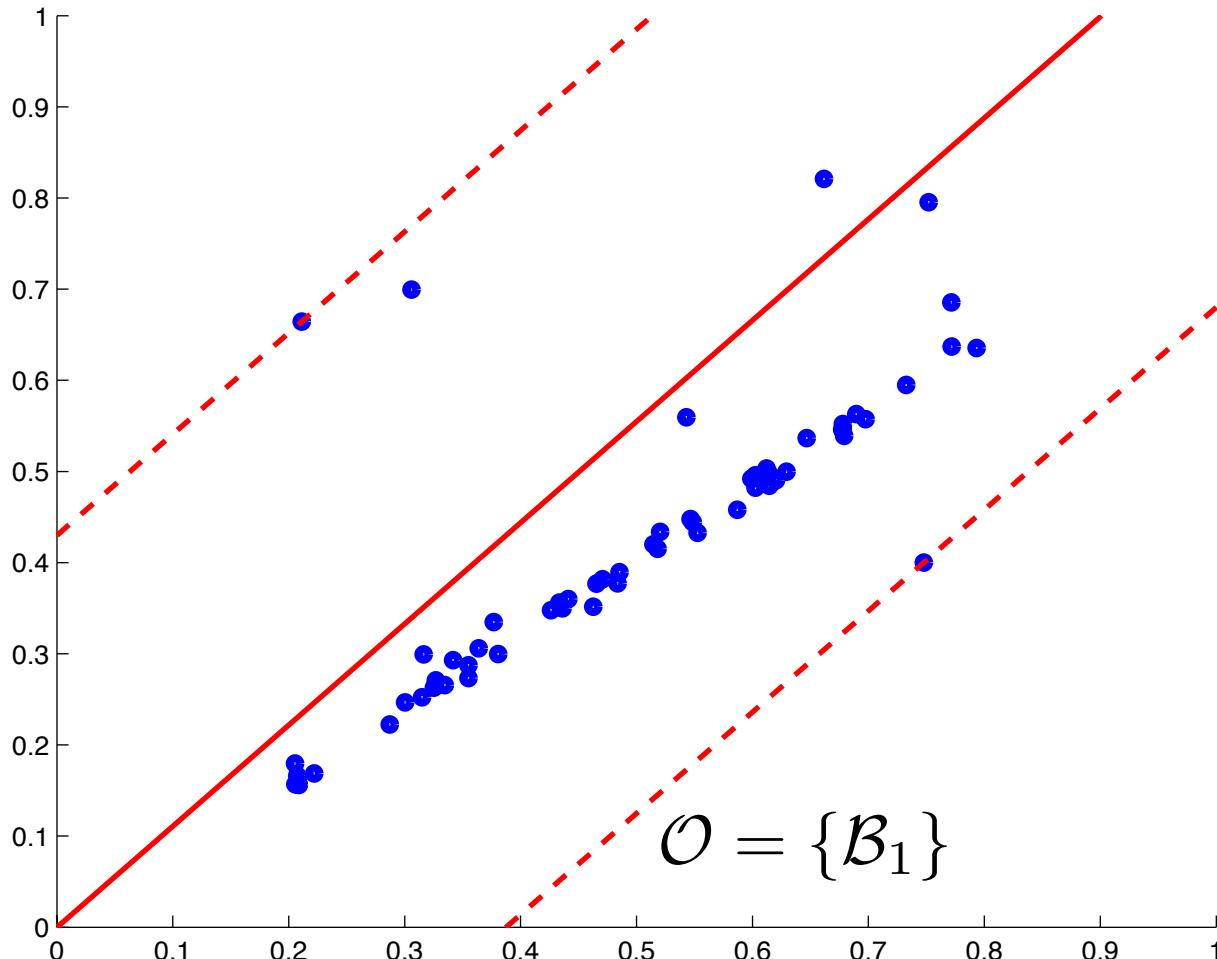
A* algorithm



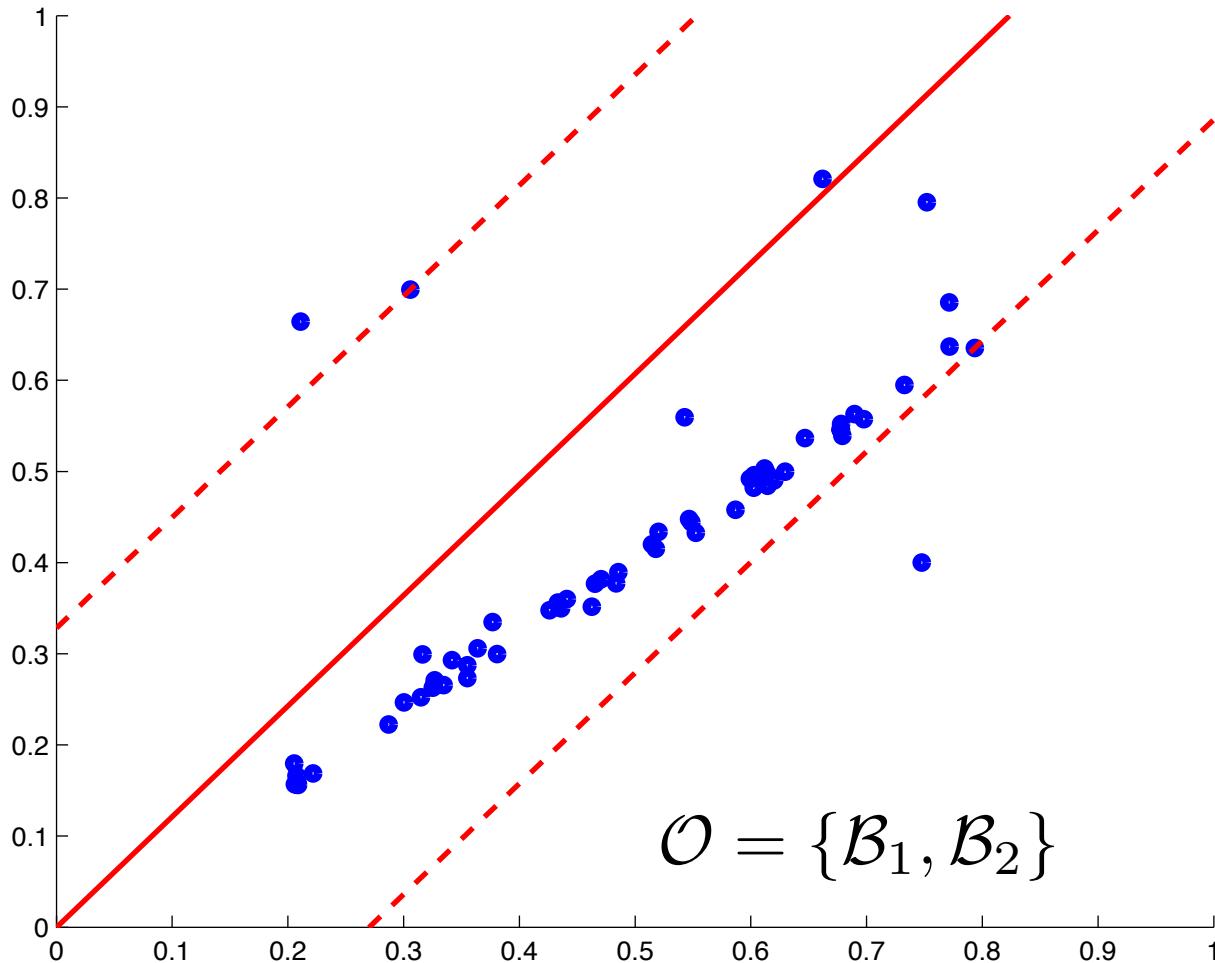
A* algorithm



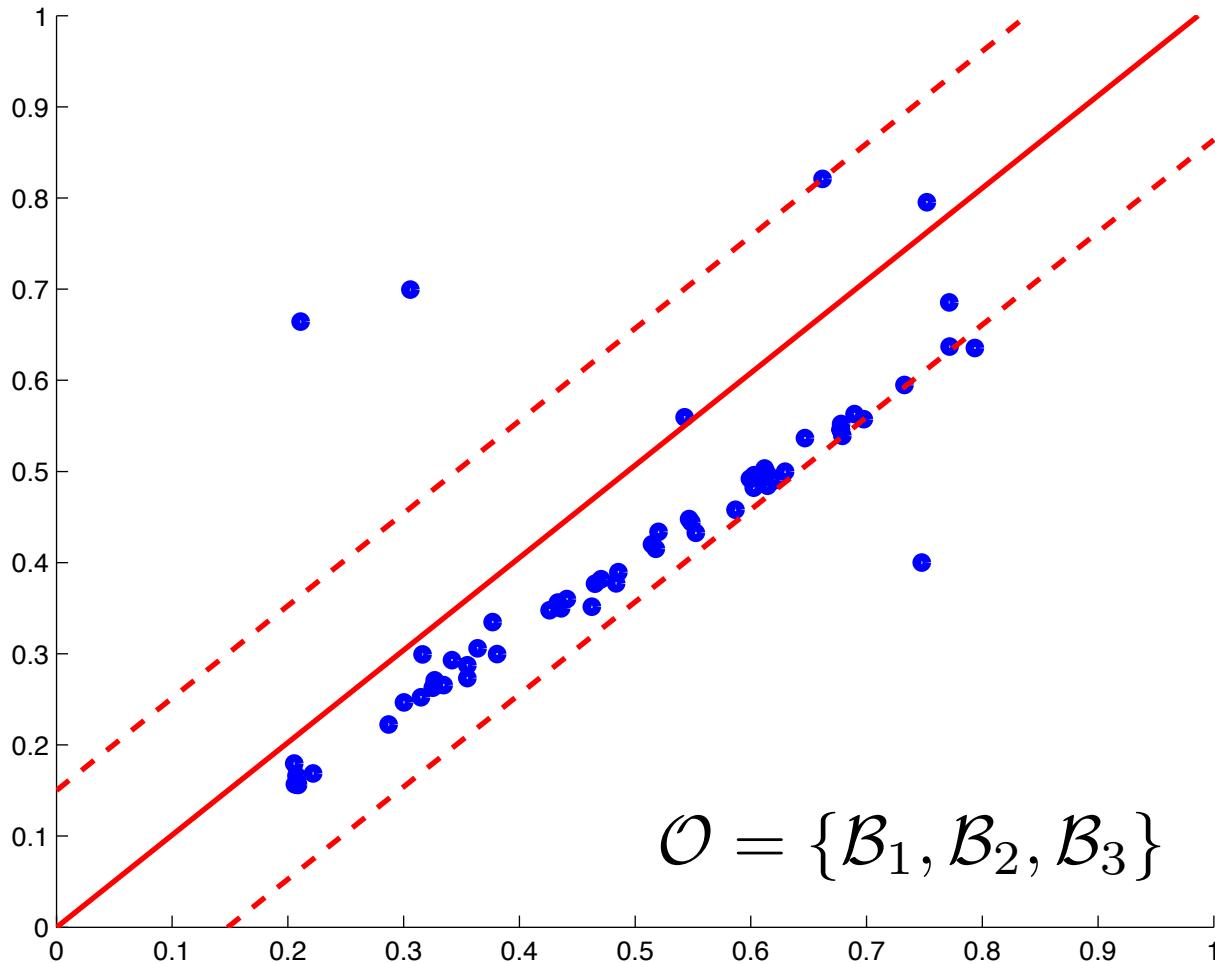
Heuristic function



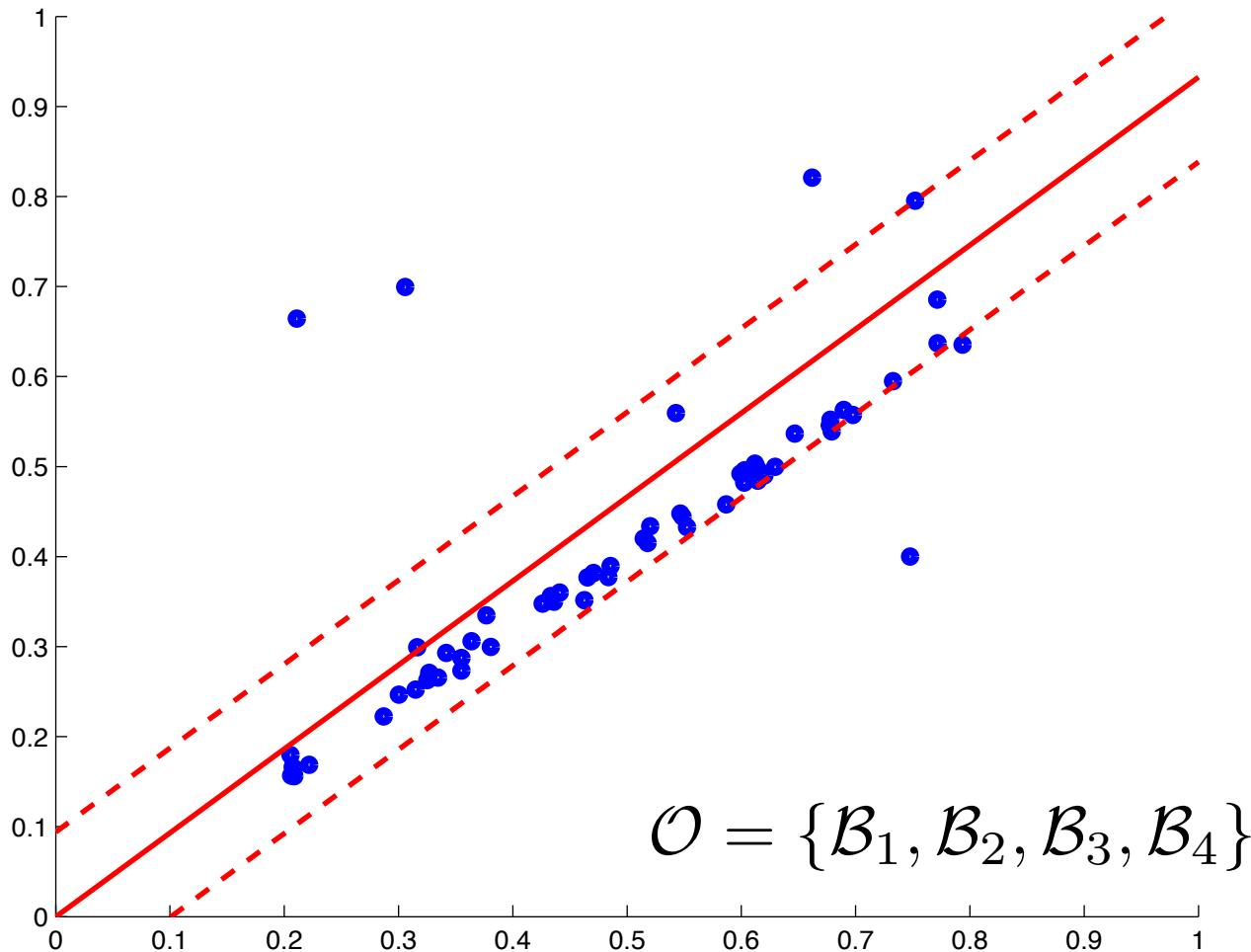
Heuristic function



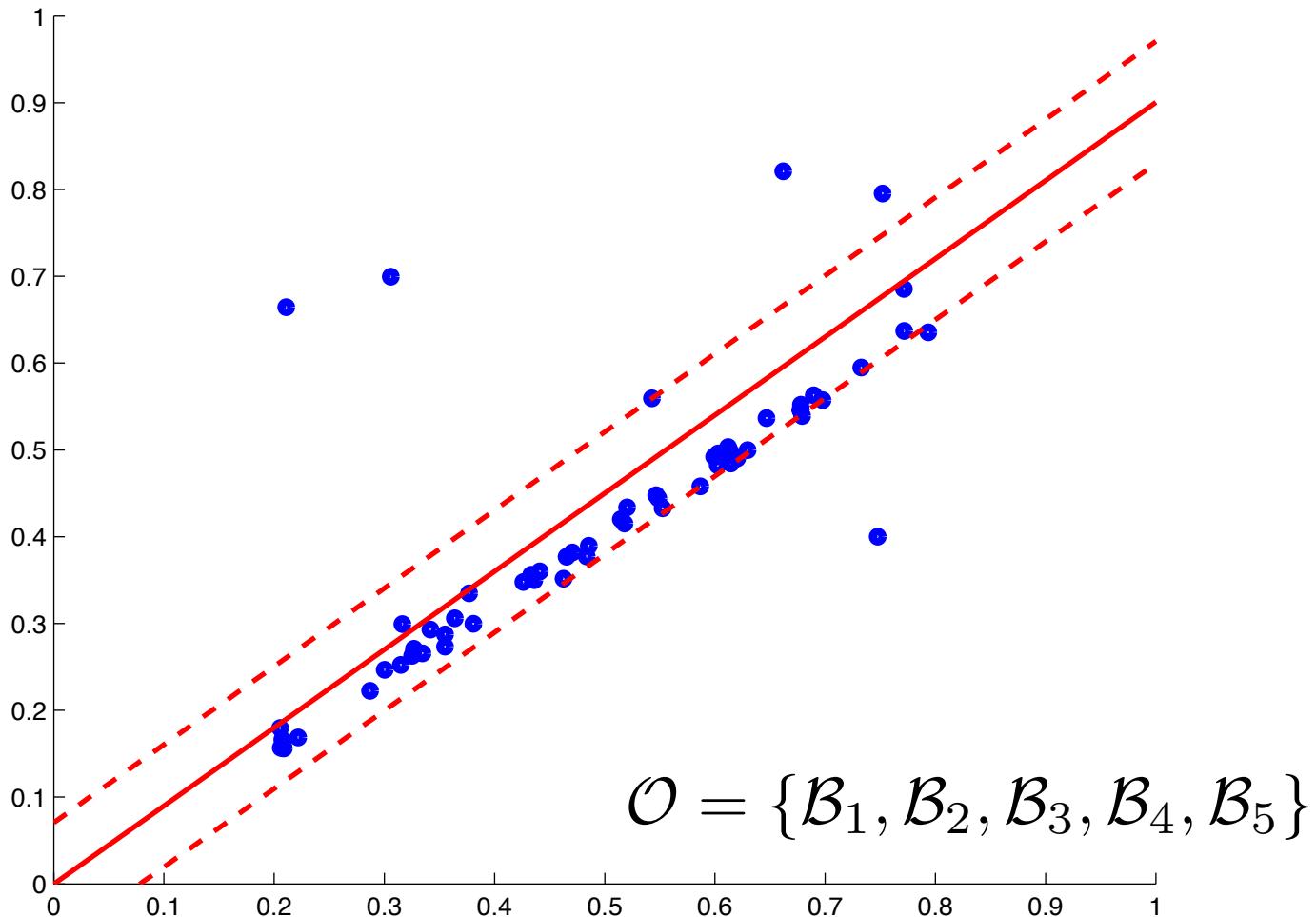
Heuristic function



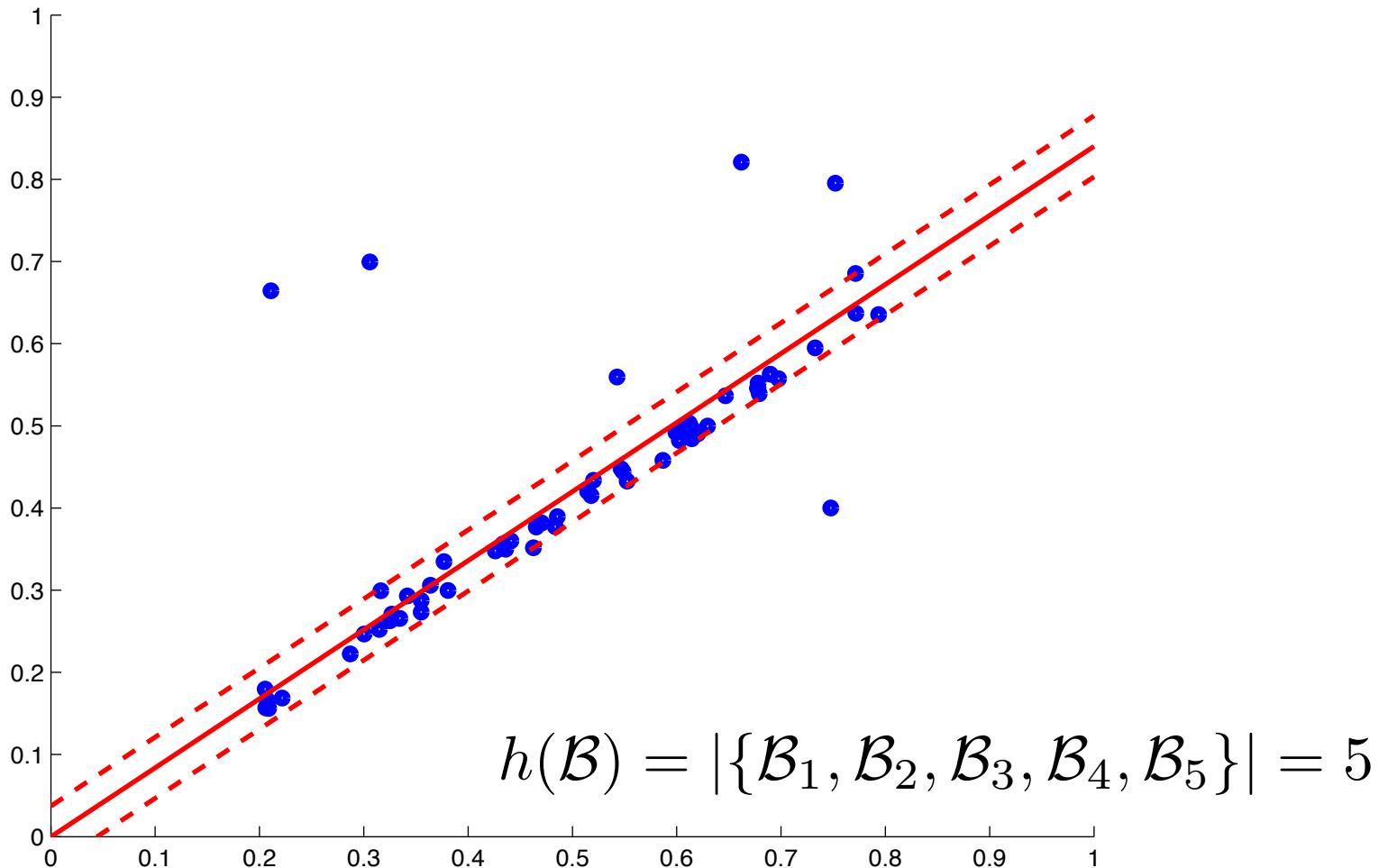
Heuristic function



Heuristic function



Heuristic function



Definition (Admissibility):

A heuristic is admissible if it satisfies

$$h(\mathcal{B}) \geq 0 \quad \text{and} \quad h(\mathcal{B}) \leq h^*(\mathcal{B})$$

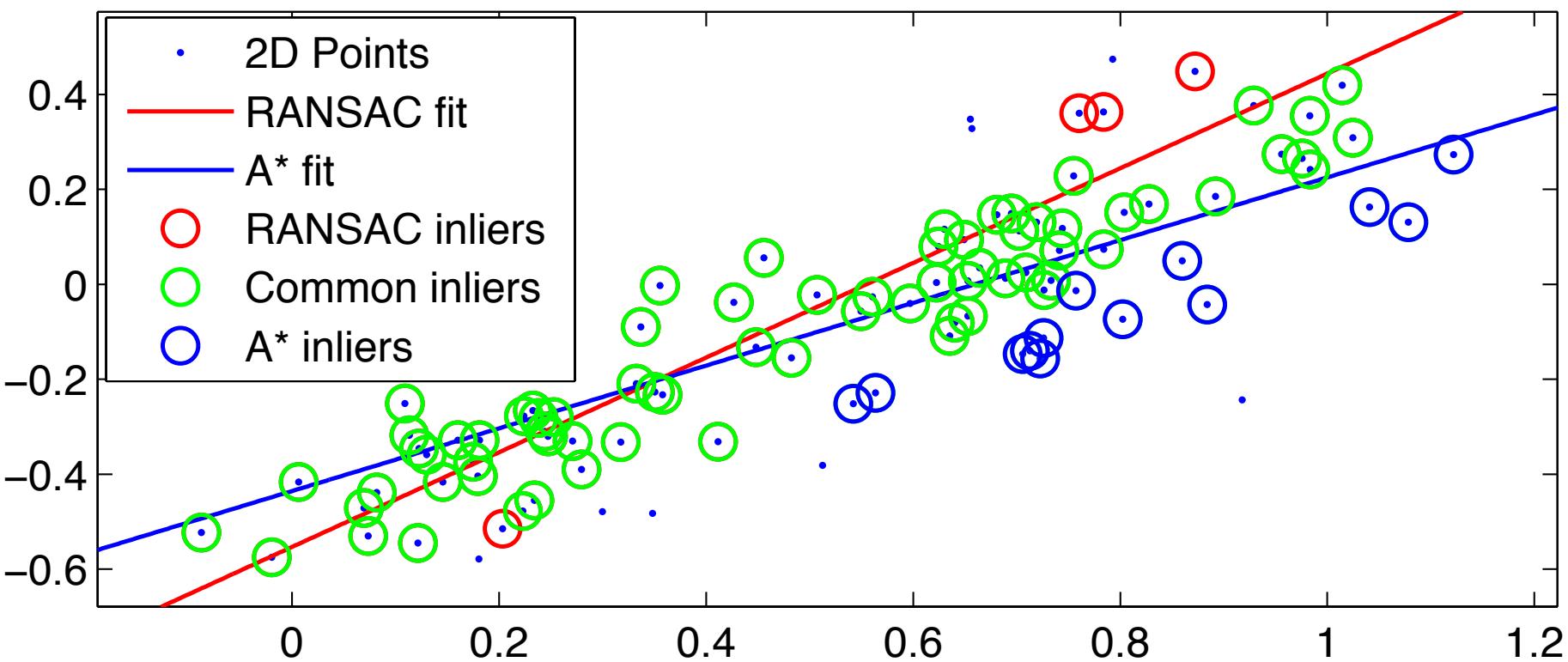
where $h^*(\mathcal{B})$ is the true remaining cost from \mathcal{B} to feasibility.

Theorem:

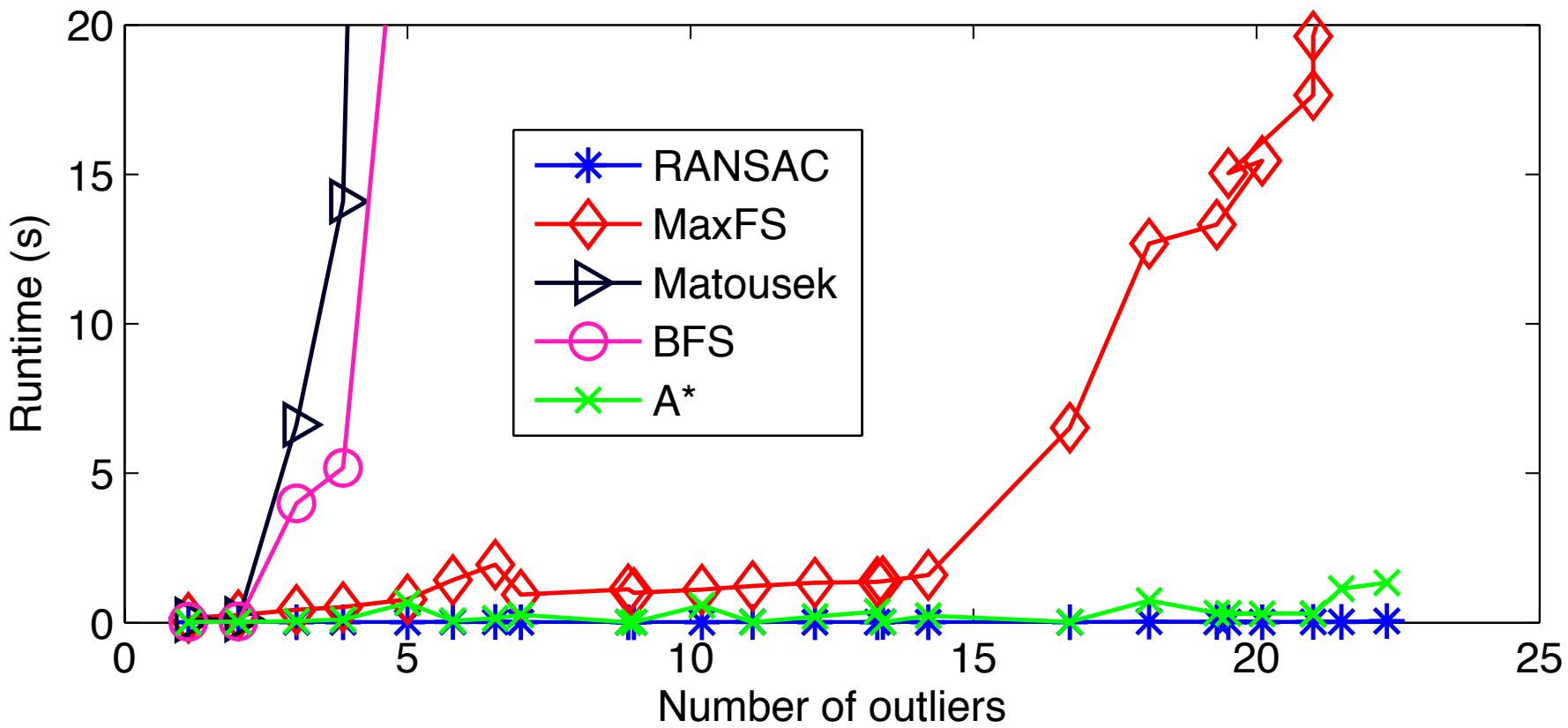
A* algorithm is optimal if $h(\mathcal{B})$ is admissible.

Results

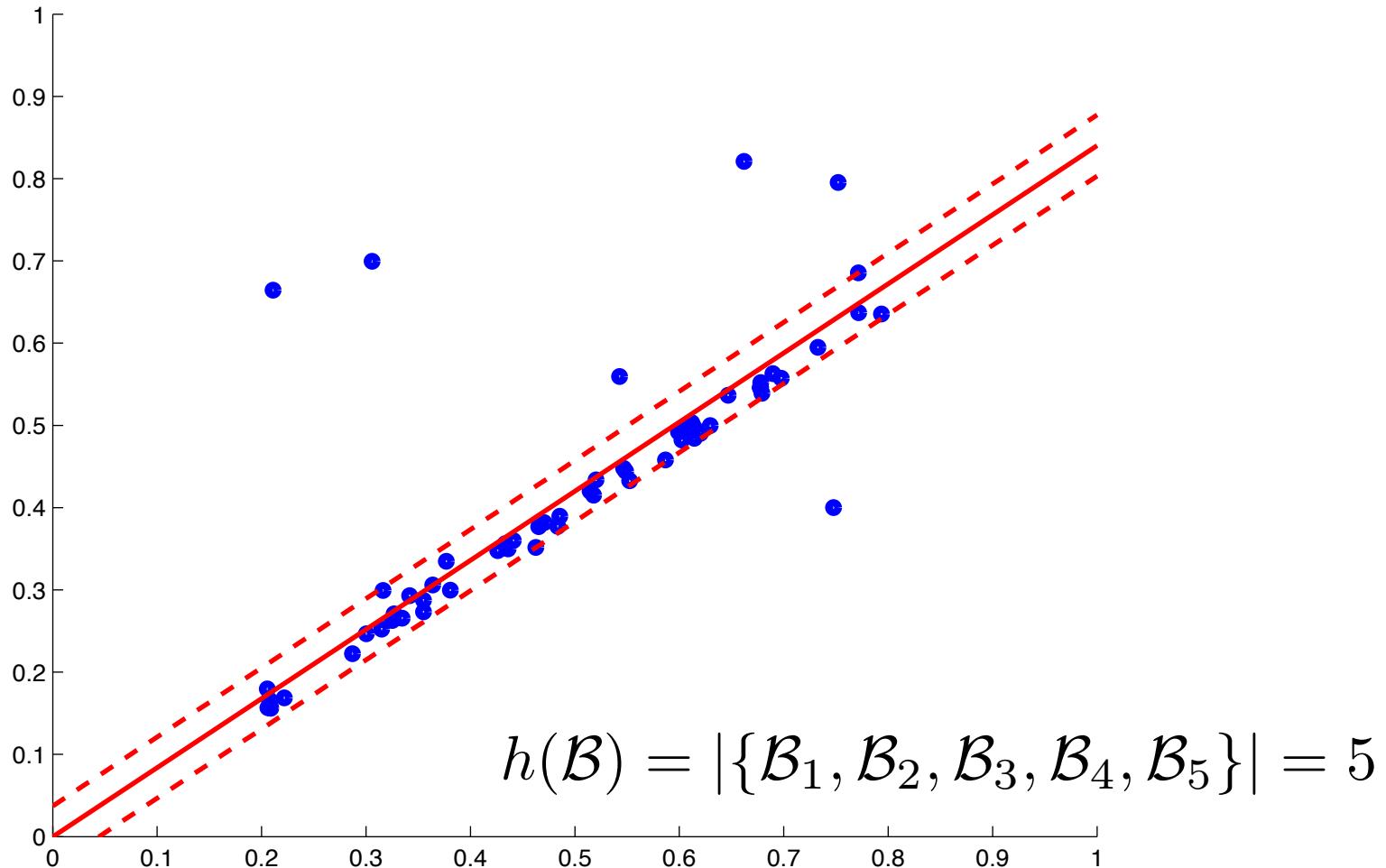
A* result (global) vs RANSAC result



Results

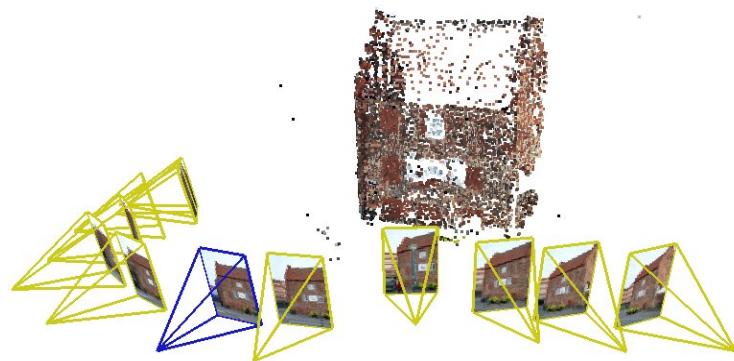


Limitation: Outlier ratio $\leq \frac{1}{p+1}$



Other residual functions?

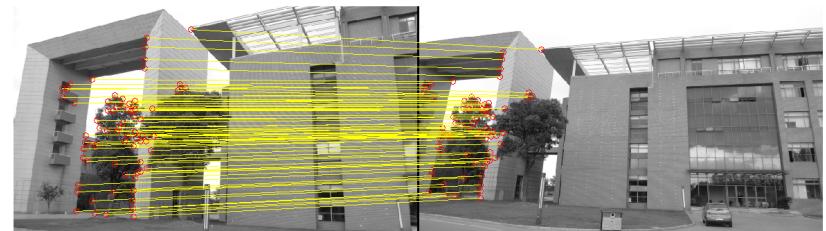
Triangulation



Reprojection error:

$$r_i(\theta) = \frac{\|(\mathbf{P}_{i,1:2} - \mathbf{x}_i \mathbf{P}_{i,3})\tilde{\theta}\|}{\mathbf{P}_{i,3}\tilde{\theta}}$$

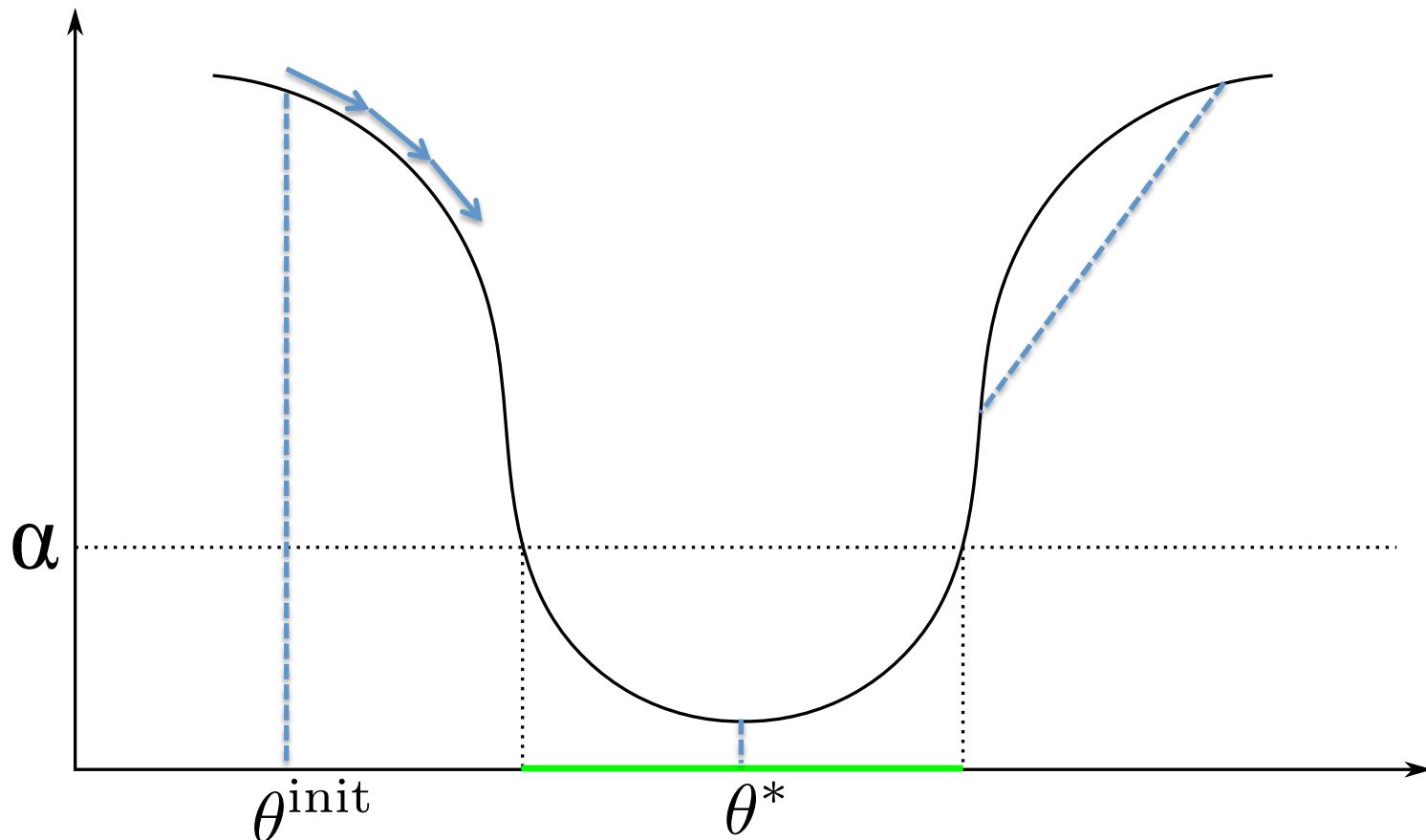
Homography fitting



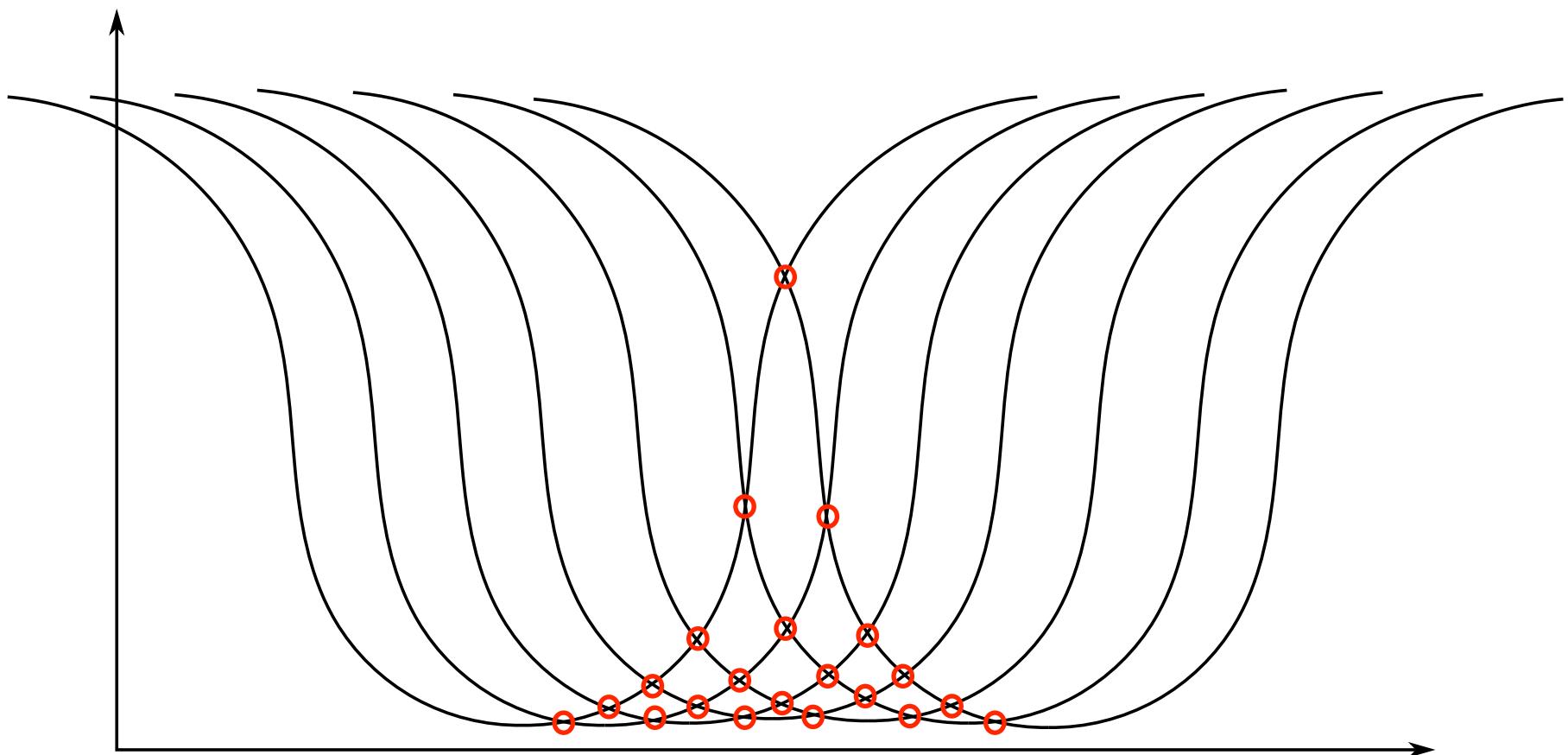
Transfer error:

$$r_i(\theta) = \frac{\|(\theta_{1:2} - \mathbf{u}'_i \theta_3)\tilde{\mathbf{u}}_i\|}{\theta_3 \tilde{\mathbf{u}}_i}$$

Pseudoconvex residual



Combinatorial dimension = p+1



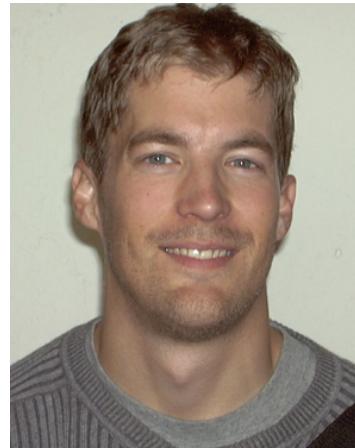
Thank you!



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