

ALU Implementation

Computer Architecture



CS3501 – 2025 I

PROF.: JGONZALEZ@UTEC.EDU.PE

CWILLIAMS@UTE.EDU.PE

SRC: HARRIS, HARRIS - DIGITAL DESIGN AND COMPUTER ARCHITECTURE

Executive Summary

2

- **Motivation:** We need a working ALU for the microprocessor design.
- **Problem:** We need to put together and test the building blocks for the microprocessor design.
- **Overview:**
 - ALU implementation and testing.
 - Introduce Project Logistics.
- **Conclusion:** We can implement an ALU using logic building blocks, and we can evaluate its design. We are ready to work in our microprocessor!

Outline

3

ALU Implementation

Course Logistics

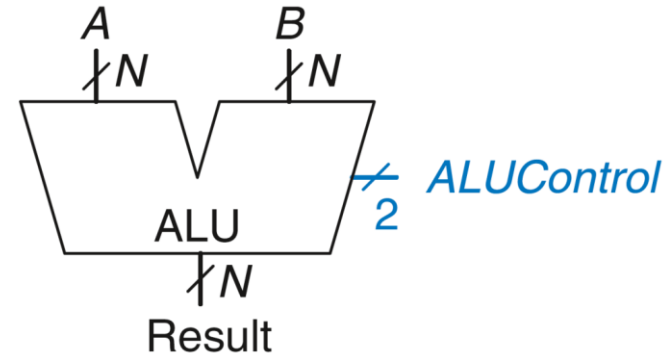
Conclusions

Recall: Arithmetic Logic Unit

4

- **ALU is the heart** of the processor.
- **ALU should perform at least:**
 - Addition
 - Subtraction
 - AND
 - OR

ALUControl _{1:0}	Function
00	Add
01	Subtract
10	AND
11	OR



Example: Perform $A + B$

$ALUControl = 00$

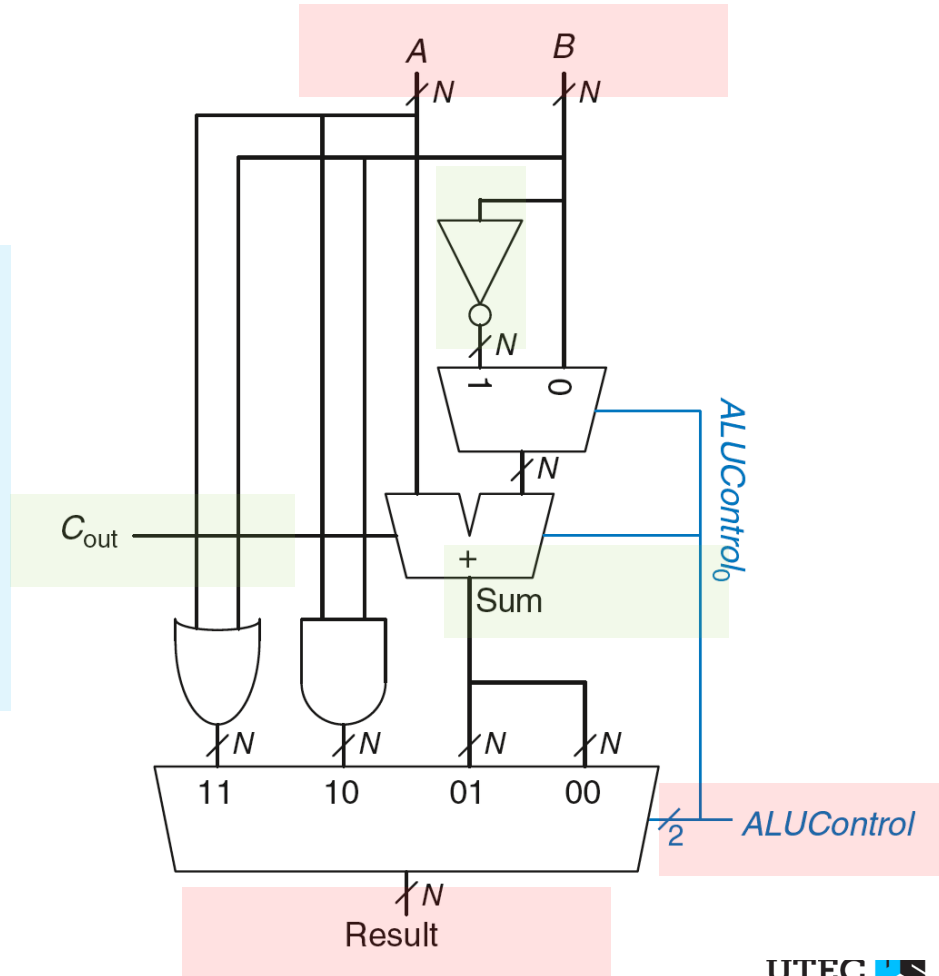
$Result = A + B$

Inputs and Outputs

5

```
//32-bit ALU for ARM processor  
module alu(input  [31:0] a, b,  
          input  [1:0] ALUControl,  
          output reg [31:0] Result,  
          output wire [3:0] ALUFlags);
```

```
    wire      neg, zero, carry, overflow;  
    wire [31:0] condinvb;  
    wire [32:0] sum;
```



6

UTEC
UNIVERSIDAD DE INGENIERÍA
Y TECNOLOGÍA

ALU ADD

7

ALUControl _{1:0}	Function
00	Add
01	Subtract
10	AND
11	OR

Example: Perform $A + B$

$ALUControl_{1:0} = 00$

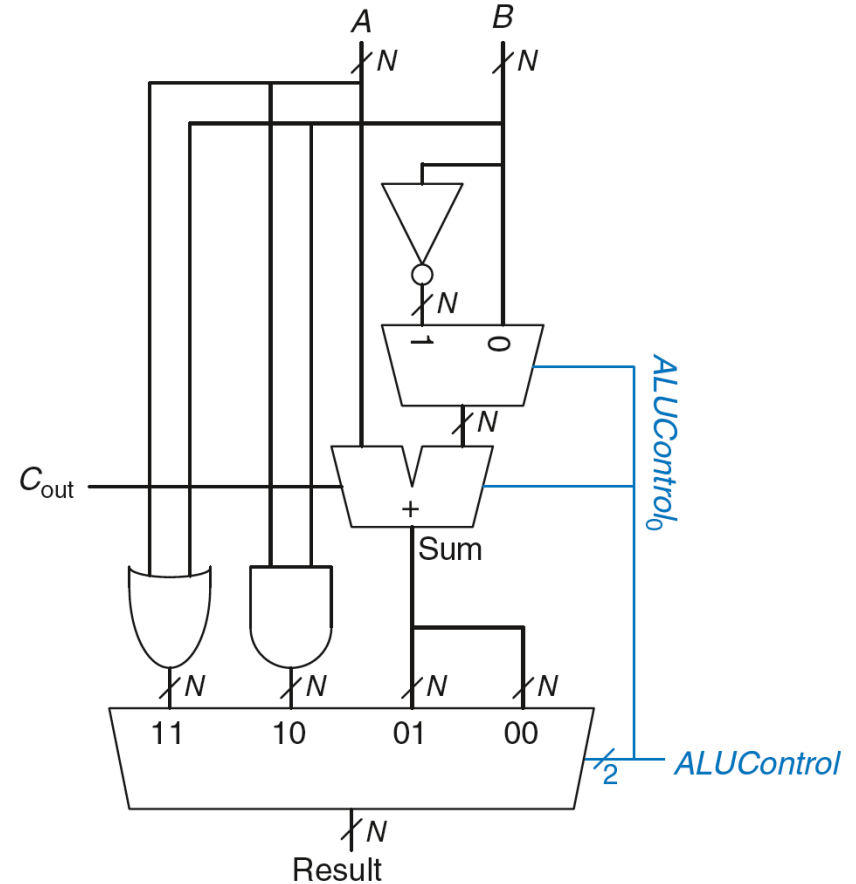
$ALUControl_0 = 0$, so:

Cin to adder = 0

2nd input to adder is B

Mux selects *Sum* as *Result*, so

Result = $A + B$

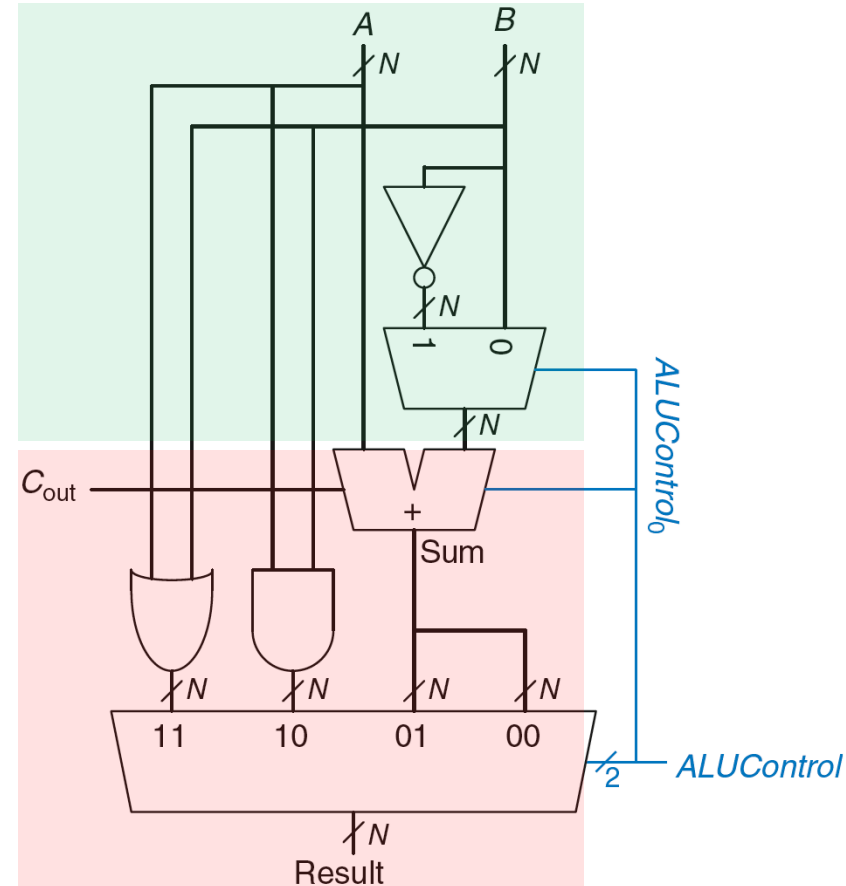


Operation Logic

8

```
assign condinvb = ALUControl[0] ? ~b : b;  
assign sum = a + condinvb + ALUControl[0];
```

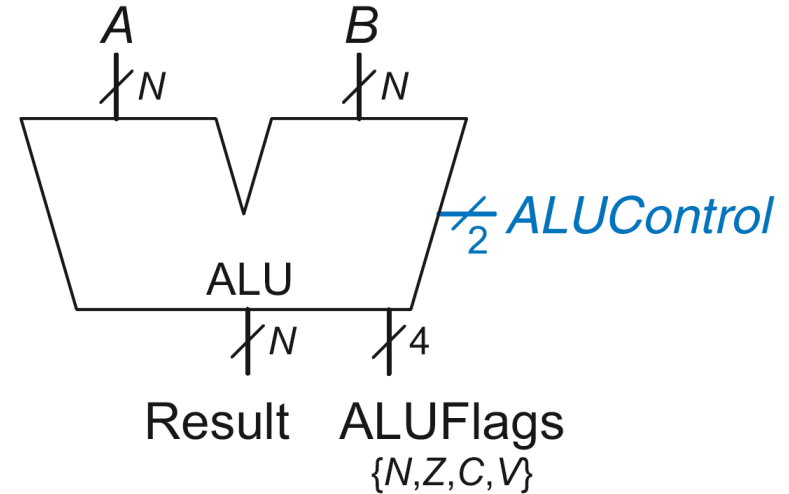
```
always @(*)  
begin  
    casex (ALUControl[1:0])  
        2'b0?: Result = sum;  
        2'b10: Result = a & b;  
        2'b11: Result = a | b;  
    endcase  
end
```



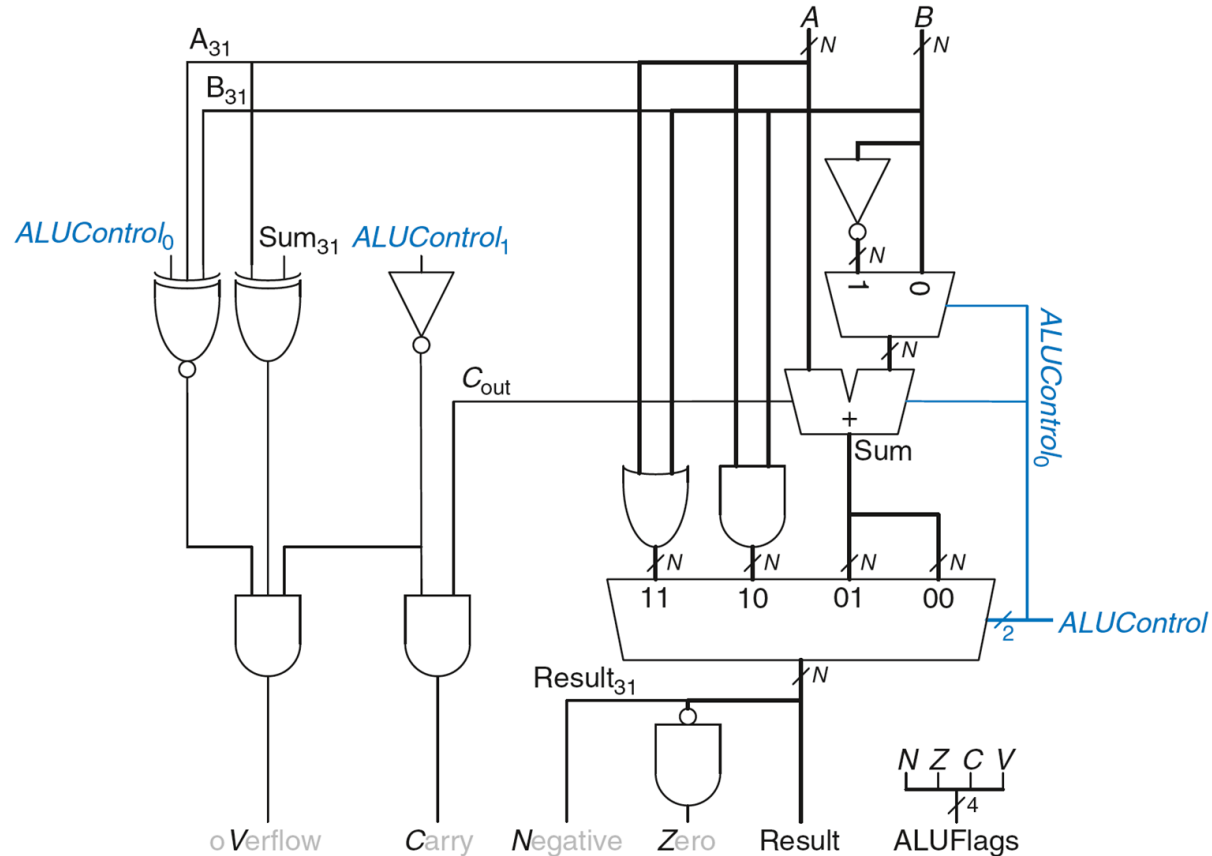
ALU with Status Flags

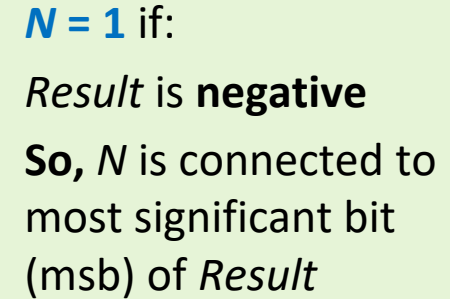
9

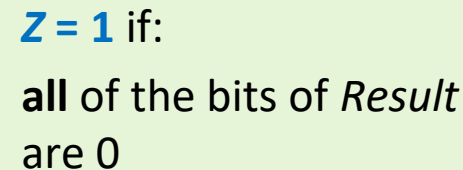
Flag	Description
<i>N</i>	Result is N egative
<i>Z</i>	Result is Z ero
<i>C</i>	Adder produces C arry out
<i>V</i>	Adder o V erflowed



ALU with Status Flags

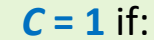






E.g.: Result = 0101
~Result = 1010
Z=AND(Result bits) = 0

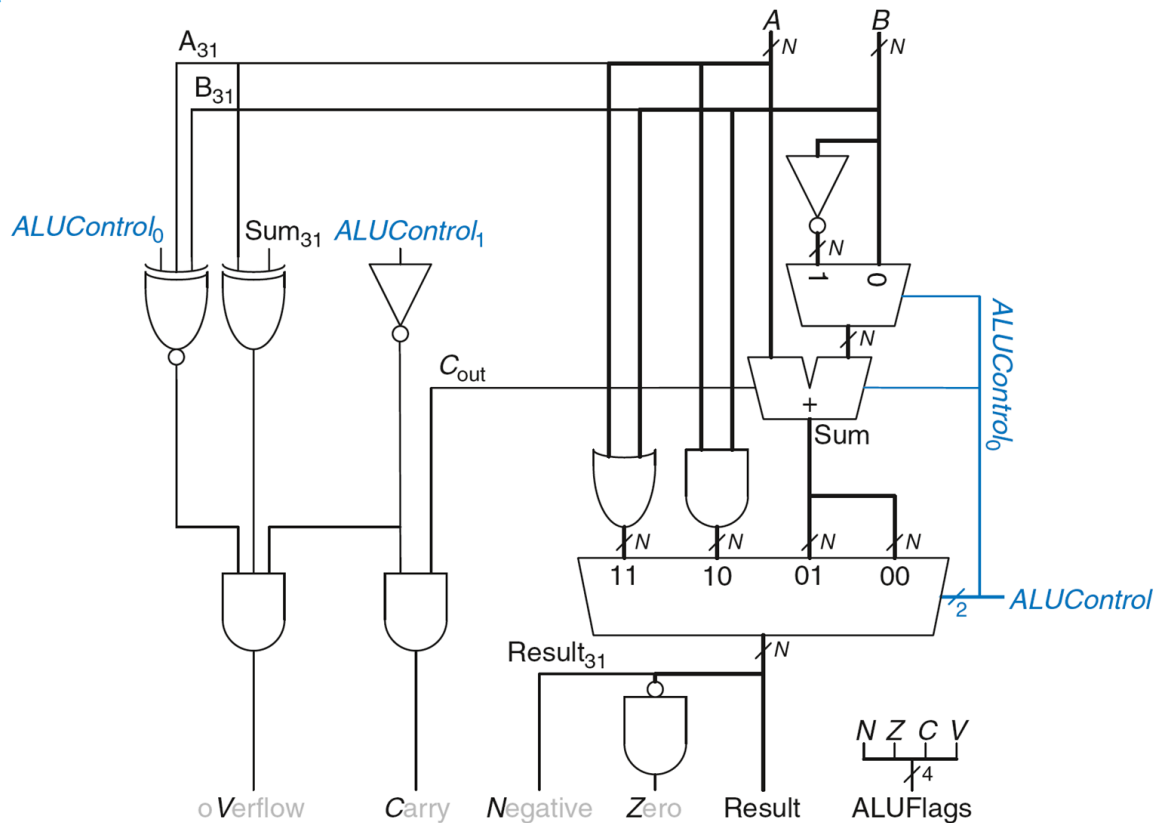
E.g.: Result = 0000
 ~Result = 1111
 Z=AND(Result bits) = 1



AND

ALU is adding or subtracting
(ALUControl is 00 or 01)

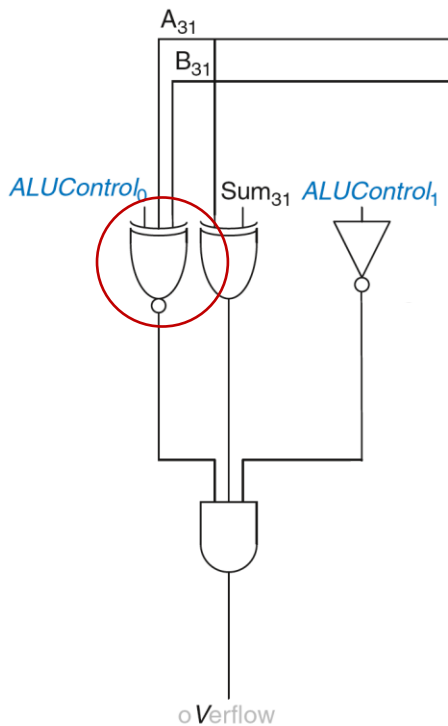
Flag Overflow (V)



$V = 1$ if:

The addition of 2 **same-signed numbers** produces a result with the **opposite sign**

Flag Overflow (V)



V = 1 if:

ALU is performing addition or subtraction
($ALUControl_1 = 0$)

AND

A and Sum have opposite signs

AND

{ A and B have same signs upon addition ($ALUControl_0 = 0$)

OR

A and B have different signs upon subtraction
($ALUControl_0 = 1$) }

Outline

17

ALU Implementation

Course Logistics

Conclusions

Outline

19

ALU Implementation

Course Logistics

Conclusions

Conclusions

20

- **We implemented an ALU** that can be used as building blocks for our microprocessor.
- **We evaluated the ALU using the testvector testbench.**
- **We conclude that the ALU enables our microprocessor design.**

ALU Implementation

Computer Architecture



CS3501 – 2025 I

PROF.: JGONZALEZ@UTEC.EDU.PE
CWILLIAMS@UTEC.EDU.PE

