

**Supplementary document of paper:**  
**Joint source, channel and space-time coding of progressive bitstream in MIMO channels**

PROOF OF LEMMA 2

From (13), it follows that

$$\begin{aligned}
 & D_{1,2,\dots,L-i+1}^P(r_i^*, r_{i+1}^*, \dots, r_L^*; c_i^*, c_{i+1}^*, \dots, c_L^*; \alpha) \\
 & \leq D_{1,2,\dots,L-i+1}^P(r_i^*, r_{i+1}, \dots, r_L; c_i^*, c_{i+1}, \dots, c_L; \alpha) \\
 & \text{for any } r_{i+1}, r_{i+2}, \dots, r_L \in \mathcal{R} \text{ and } c_{i+1}, c_{i+2}, \dots, c_L \in \mathcal{C}.
 \end{aligned} \tag{49}$$

From (31), the inequality given by (49) can be rewritten as

$$\begin{aligned}
 & \sigma^2 p(r_i^*, c_i^*) + g(b(r_i^*, c_i^*)) (1 - p(r_i^*, c_i^*)) D_{1,2,\dots,L-i}^P(r_{i+1}^*, r_{i+2}^*, \dots, r_L^*; c_{i+1}^*, c_{i+2}^*, \dots, c_L^*; \alpha) \\
 & \leq \sigma^2 p(r_i^*, c_i^*) + g(b(r_i^*, c_i^*)) (1 - p(r_i^*, c_i^*)) D_{1,2,\dots,L-i}^P(r_{i+1}, r_{i+2}, \dots, r_L; c_{i+1}, c_{i+2}, \dots, c_L; \alpha) \\
 & \text{for any } r_{i+1}, r_{i+2}, \dots, r_L \in \mathcal{R} \text{ and } c_{i+1}, c_{i+2}, \dots, c_L \in \mathcal{C}.
 \end{aligned} \tag{50}$$

Since  $p(r_i^*, c_i^*) < 1$ ,  $g(b(r_i^*, c_i^*)) = 2^{-\alpha r_i^* c_i^* T_{\text{pkt}} W_{\text{pkt}}} > 0$ , and from (50), we have

$$\begin{aligned}
 & D_{1,2,\dots,L-i}^P(r_{i+1}^*, r_{i+2}^*, \dots, r_L^*; c_{i+1}^*, c_{i+2}^*, \dots, c_L^*; \alpha) \\
 & \leq D_{1,2,\dots,L-i}^P(r_{i+1}, r_{i+2}, \dots, r_L; c_{i+1}, c_{i+2}, \dots, c_L; \alpha) \\
 & \text{for any } r_{i+1}, r_{i+2}, \dots, r_L \in \mathcal{R} \text{ and } c_{i+1}, c_{i+2}, \dots, c_L \in \mathcal{C}.
 \end{aligned} \tag{51}$$

Based on (51), we will prove (14) by induction on the number of packets: We first consider  $L - i$  packets. Eq. (51) is identical to (14) when we let  $j = i + 1$  in (14) (i.e.,  $L - i$  packets). We next suppose that (14) holds for  $j = n$  ( $\geq i + 1$ ). In other words, for  $L - n + 1$  ( $\leq L - i$ ) packets, we have

$$\begin{aligned}
 & D_{1,2,\dots,L-n+1}^P(r_n^*, r_{n+1}^*, \dots, r_L^*; c_n^*, c_{n+1}^*, \dots, c_L^*; \alpha) \\
 & \leq D_{1,2,\dots,L-n+1}^P(r_n, r_{n+1}, \dots, r_L; c_n, c_{n+1}, \dots, c_L; \alpha) \\
 & \text{for any } r_n, r_{n+1}, \dots, r_L \in \mathcal{R} \text{ and } c_n, c_{n+1}, \dots, c_L \in \mathcal{C}.
 \end{aligned} \tag{52}$$

Eq. (52), which is the induction hypothesis, is identical to (13) when we let  $i = n$  in (13). Since

(13) implies (51), (51) holds for  $i = n$ , i.e.,

$$\begin{aligned}
& D_{1,2,\dots,L-n}^p(r_{n+1}^*, r_{n+2}^*, \dots, r_L^*; c_{n+1}^*, c_{n+2}^*, \dots, c_L^*; \alpha) \\
& \leq D_{1,2,\dots,L-n}^p(r_{n+1}, r_{n+2}, \dots, r_L; c_{n+1}, c_{n+2}, \dots, c_L; \alpha) \\
& \text{for any } r_{n+1}, r_{n+2}, \dots, r_L \in \mathcal{R} \text{ and } c_{n+1}, c_{n+2}, \dots, c_L \in \mathcal{C}.
\end{aligned} \tag{53}$$

Letting  $j = n + 1$  in (14) (i.e.,  $L - n$  packets), we obtain a result that is identical to (53). Hence, (14) holds for  $j = n + 1$ . We have thus shown that (14) is valid for  $j \geq i + 1$ .  $\square$

### PROOF OF LEMMA 3

From (9), it can be shown that  $D_{1,2,\dots,L}^p(r_1, r_2, \dots, r_L; c_1, c_2, \dots, c_L; \alpha)$  is rewritten as

$$\begin{aligned}
& D_{1,2,\dots,L}^p(r_1, r_2, \dots, r_L; c_1, c_2, \dots, c_L; \alpha) \\
& = \sum_{n=0}^{L-2} \sigma^2 \left( \prod_{i=1}^n g(b(r_i, c_i)) \right) p(r_{n+1}, c_{n+1}) \prod_{i=1}^n (1 - p(r_i, c_i)) \\
& \quad + \sigma^2 \prod_{i=1}^{L-1} g(b(r_i, c_i)) \prod_{i=1}^{L-1} (1 - p(r_i, c_i)) \left\{ p(r_L, c_L) + g(b(r_L, c_L)) (1 - p(r_L, c_L)) \right\}.
\end{aligned} \tag{54}$$

In addition, from (9),  $D_{1,2,\dots,L-1}^p(r_1, r_2, \dots, r_{L-1}; c_1, c_2, \dots, c_{L-1}; \alpha)$  is given by

$$\begin{aligned}
& D_{1,2,\dots,L-1}^p(r_1, r_2, \dots, r_{L-1}; c_1, c_2, \dots, c_{L-1}; \alpha) \\
& = \sum_{n=0}^{L-2} \sigma^2 \left( \prod_{i=1}^n g(b(r_i, c_i)) \right) p(r_{n+1}, c_{n+1}) \prod_{i=1}^n (1 - p(r_i, c_i)) \\
& \quad + \sigma^2 \prod_{i=1}^{L-1} g(b(r_i, c_i)) \prod_{i=1}^{L-1} (1 - p(r_i, c_i)).
\end{aligned} \tag{55}$$

From (54) and (55), we obtain

$$\begin{aligned}
& D_{1,2,\dots,L}^p(r_1, r_2, \dots, r_L; c_1, c_2, \dots, c_L; \alpha) \\
& = D_{1,2,\dots,L-1}^p(r_1, r_2, \dots, r_{L-1}; c_1, c_2, \dots, c_{L-1}; \alpha) \\
& \quad + \sigma^2 \prod_{i=1}^{L-1} g(b(r_i, c_i)) \prod_{i=1}^{L-1} (1 - p(r_i, c_i)) \left\{ p(r_L, c_L) + g(b(r_L, c_L)) (1 - p(r_L, c_L)) - 1 \right\} \\
& < D_{1,2,\dots,L-1}^p(r_1, r_2, \dots, r_{L-1}; c_1, c_2, \dots, c_{L-1}; \alpha),
\end{aligned} \tag{56}$$

where the inequality follows from  $p(r_L, c_L) < 1$ , and  $0 < g(b(r_L, c_L)) = 2^{-\alpha r_L c_L T_{\text{pkt}} W_{\text{pkt}}} < 1$ .  $\square$

### PROOF OF LEMMA 4

From (15) of Lemma 3, it is clear that for  $1 \leq i \leq L - 1$ , we have

$$\begin{aligned}
 & D_{1,2,\dots,L-i+1}^p(r_{i+1}^*, r_{i+2}^*, \dots, r_L^*, r_k; c_{i+1}^*, c_{i+2}^*, \dots, c_L^*, c_k; \alpha) \\
 & < D_{1,2,\dots,L-i}^p(r_{i+1}^*, r_{i+2}^*, \dots, r_L^*; c_{i+1}^*, c_{i+2}^*, \dots, c_L^*; \alpha) \\
 & \text{for any } r_k \in \mathcal{R} \text{ and } c_k \in \mathcal{C}.
 \end{aligned} \tag{57}$$

From the condition of this lemma given by (13) and (57), we can derive

$$\begin{aligned}
 & D_{1,2,\dots,L-i+1}^p(r_i^*, r_{i+1}^*, \dots, r_L^*; c_i^*, c_{i+1}^*, \dots, c_L^*; \alpha) \\
 & \leq D_{1,2,\dots,L-i+1}^p(r_{i+1}^*, r_{i+2}^*, \dots, r_L^*, r_k; c_{i+1}^*, c_{i+2}^*, \dots, c_L^*, c_k; \alpha) \\
 & < D_{1,2,\dots,L-i}^p(r_{i+1}^*, r_{i+2}^*, \dots, r_L^*; c_{i+1}^*, c_{i+2}^*, \dots, c_L^*; \alpha) \\
 & \text{for any } r_k \in \mathcal{R} \text{ and } c_k \in \mathcal{C},
 \end{aligned} \tag{58}$$

where the first inequality follows from (13), and the second inequality follows from (57).

Based on (58), we will prove (16) by induction on the number of packets: We first consider  $L - i$  packets. Eq. (58) is identical to (16) when we let  $j = i + 1$  in (16) (i.e.,  $L - i$  packets). We next suppose that (16) holds for  $j = n$  ( $\geq i + 1$ ). In other words, for  $L - n + 1$  ( $\leq L - i$ ) packets, we have the following induction hypothesis.

$$\begin{aligned}
 & D_{1,2,\dots,L-i+1}^p(r_i^*, r_{i+1}^*, \dots, r_L^*; c_i^*, c_{i+1}^*, \dots, c_L^*; \alpha) \\
 & < D_{1,2,\dots,L-n+1}^p(r_n^*, r_{n+1}^*, \dots, r_L^*; c_n^*, c_{n+1}^*, \dots, c_L^*; \alpha).
 \end{aligned} \tag{59}$$

Note that the right hand side of (59) is also a parametric distortion-based optimum for  $L - n + 1$  progressive packets, because Lemma 2 indicates that the condition of this lemma, which is given by (13), implies (14) for some integer  $j \geq i + 1$ . From the fact that a parametric distortion-based optimal solution satisfies (58), and that the right hand side of (59) equals the first line of (58) when setting  $i = n$  in (58), it follows that

$$\begin{aligned}
 & D_{1,2,\dots,L-n+1}^p(r_n^*, r_{n+1}^*, \dots, r_L^*; c_n^*, c_{n+1}^*, \dots, c_L^*; \alpha) \\
 & < D_{1,2,\dots,L-n}^p(r_{n+1}^*, r_{n+2}^*, \dots, r_L^*; c_{n+1}^*, c_{n+2}^*, \dots, c_L^*; \alpha).
 \end{aligned} \tag{60}$$

From the induction hypothesis given by (59) and (60), we have

$$\begin{aligned} & D_{1,2,\dots,L-i+1}^p(r_i^*, r_{i+1}^*, \dots, r_L^*; c_i^*, c_{i+1}^*, \dots, c_L^*; \alpha) \\ & < D_{1,2,\dots,L-n}^p(r_{n+1}^*, r_{n+2}^*, \dots, r_L^*; c_{n+1}^*, c_{n+2}^*, \dots, c_L^*; \alpha). \end{aligned} \quad (61)$$

Letting  $j = n + 1$  in (16) (i.e.,  $L - n$  packets), we obtain a result identical to (61). Hence, (16) holds for  $j = n + 1$ . We have thus shown that (16) holds for  $j \geq i + 1$ .  $\square$

### PROOF OF COROLLARY 6

The condition of this corollary, given by (18), is identical to (13) of Lemma 2 when we let  $i = 1$  in (13). Thus, (14) of Lemma 2 holds for some integer  $k$  in the range of  $2 \leq k \leq L$  as follows:

$$\begin{aligned} & D_{1,2,\dots,L-k+1}^p(r_k^*, r_{k+1}^*, \dots, r_L^*; c_k^*, c_{k+1}^*, \dots, c_L^*; \alpha) \\ & \leq D_{1,2,\dots,L-k+1}^p(r_k, r_{k+1}, \dots, r_L; c_k, c_{k+1}, \dots, c_L; \alpha) \\ & \text{for any } r_k, r_{k+1}, \dots, r_L \in \mathcal{R} \text{ and } c_k, c_{k+1}, \dots, c_L \in \mathcal{C}. \end{aligned} \quad (62)$$

If we let  $i = k$  in the condition of Theorem 5, given by (13), then it equals (62) in the range of  $2 \leq k \leq L - 1$ . Note that this range of  $k$  is a subset of  $1 \leq k \leq L - 1$  and  $2 \leq k \leq L$  given by (13) and (62), respectively. As a result, it follows from Theorem 5 that, for  $2 \leq k \leq L - 1$  and  $k + 1 \leq j \leq L$ , (19) holds with  $k$  being substituted into  $i$ . In addition, from Theorem 5 and (18), it follows immediately that for  $2 \leq j \leq L$ , (19) holds with  $i = 1$ . We have thus shown that (19) holds for  $1 \leq i \leq L - 1$  and  $i + 1 \leq j \leq L$ . Letting  $i = 1, 2, \dots, L - 1$  and  $j = i + 1, i + 2, \dots, L$  in (19), we obtain at least  $(L^2 - L)/2$  ( $= \sum_{i=1}^{L-1} L - i$ ) constraints on  $r_1^*, r_2^*, \dots, r_{L-1}^*$  or  $c_1^*, c_2^*, \dots, c_{L-1}^*$  of all the  $L$  packets except the last one.  $\square$