## 算法与数据结构体系课程

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# 堆和优先队列



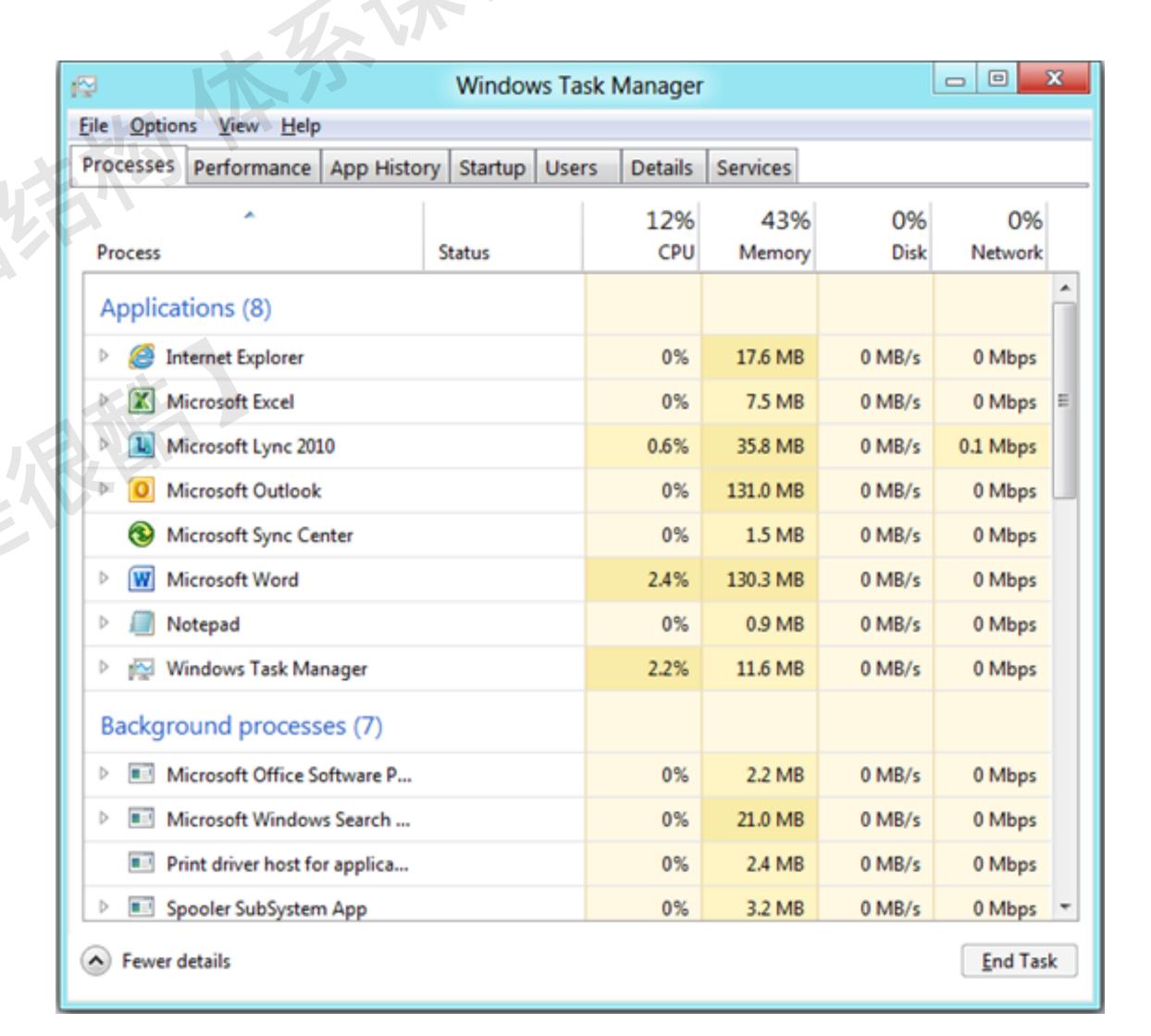
### 什么是优先队列?

普通队列: 先进先出; 后进后出

优先队列:出队顺序和入队顺序无关;和优先级相关

### 为什么使用优先队列?

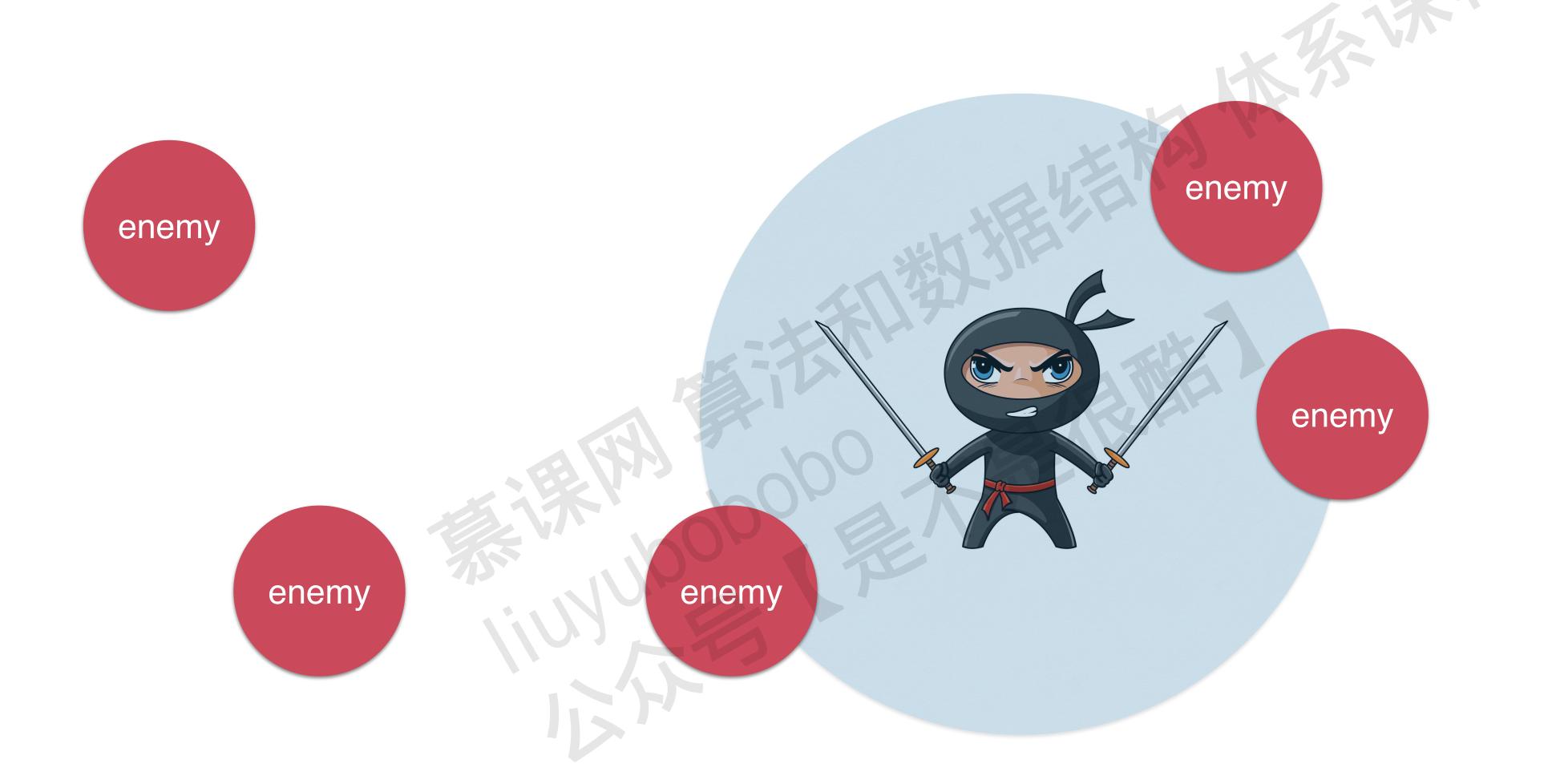
动态选择优先级最高的任务执行



### 为什么使用优先队列?



### 为什么使用优先队列?







### 优先队列

Interface Queue<E>

implement

PriorityQueue<E>

- void enqueue(E)
- E dequeue()
- E getFront()
- int getSize()
- boolean isEmpty()

可以使用不同的底层实现

### 优先队列

入队

出队(拿出最大元素)

普通线性结构

O(1)

O(n)

顺序线性结构

O(n)

O(1)

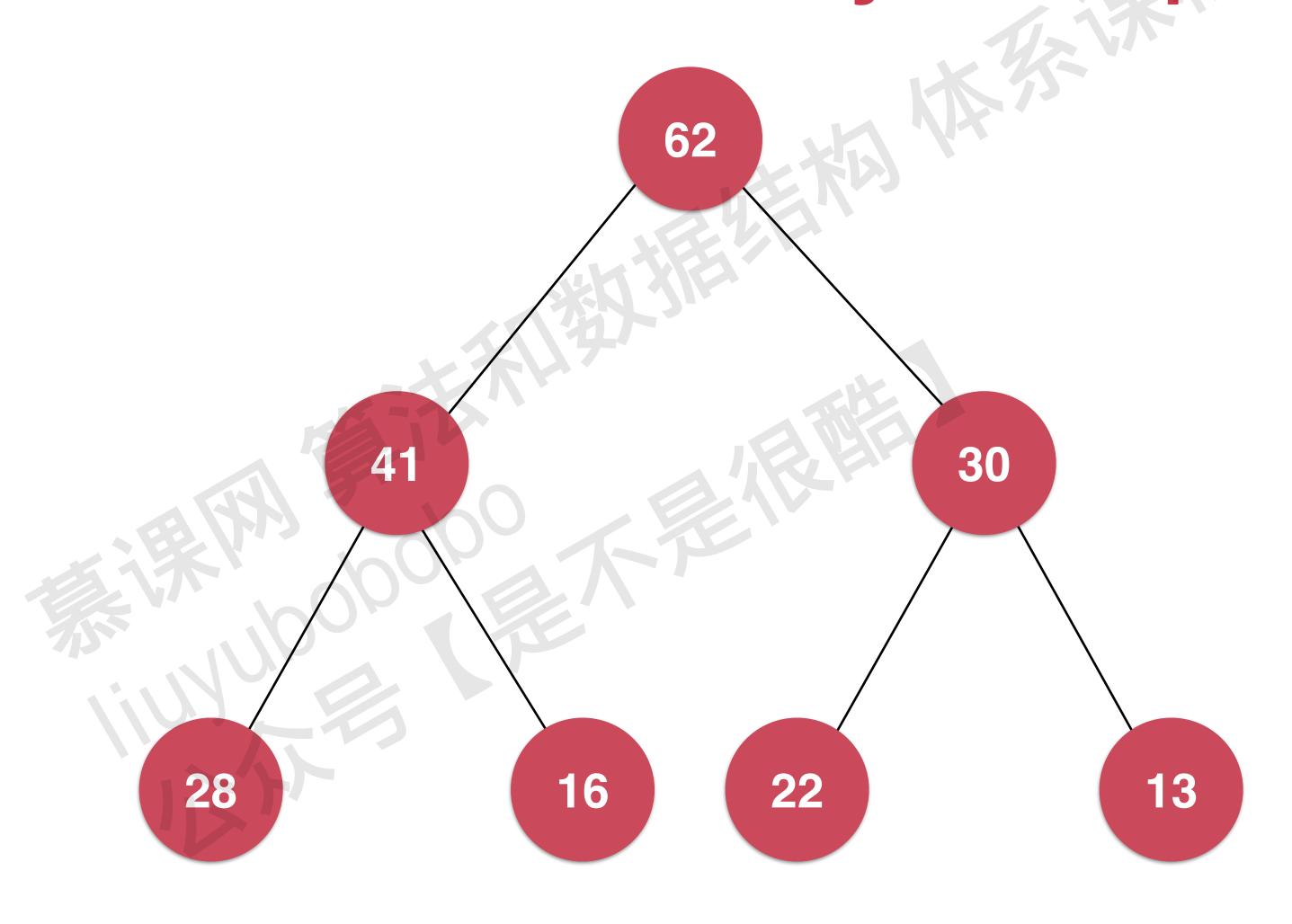
堆

O(logn)

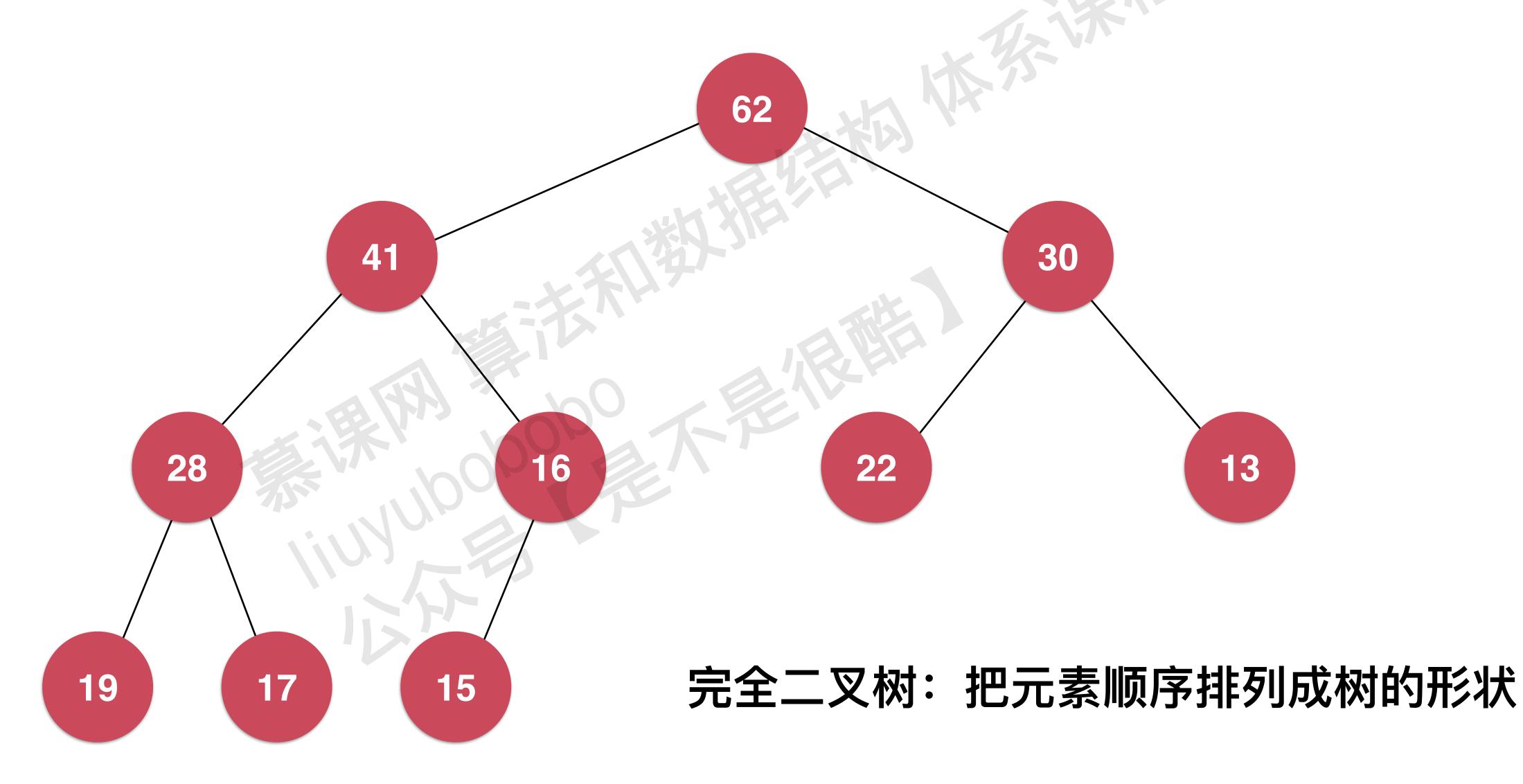
O(logn)



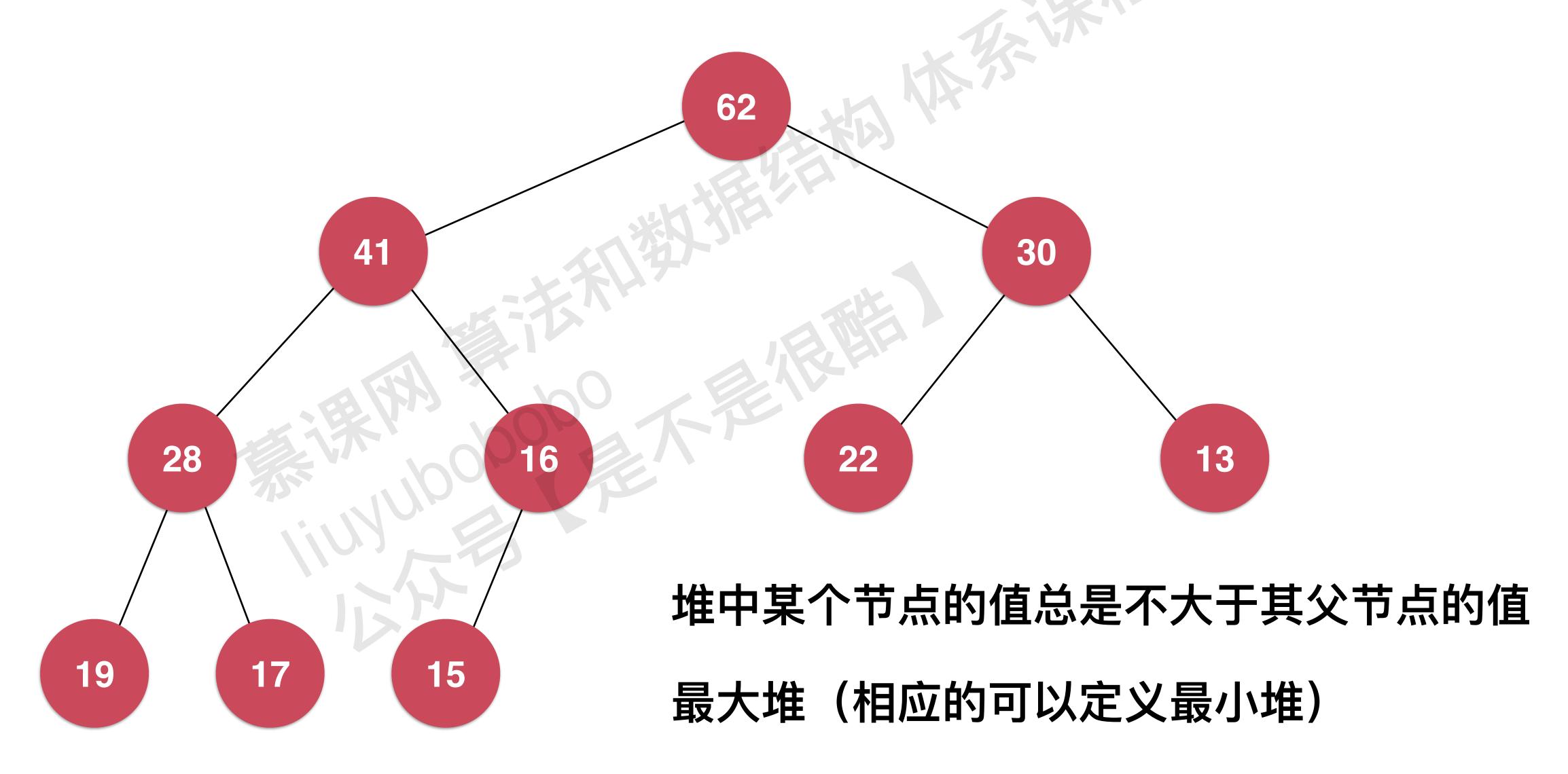
## 二叉堆 Binary Heap

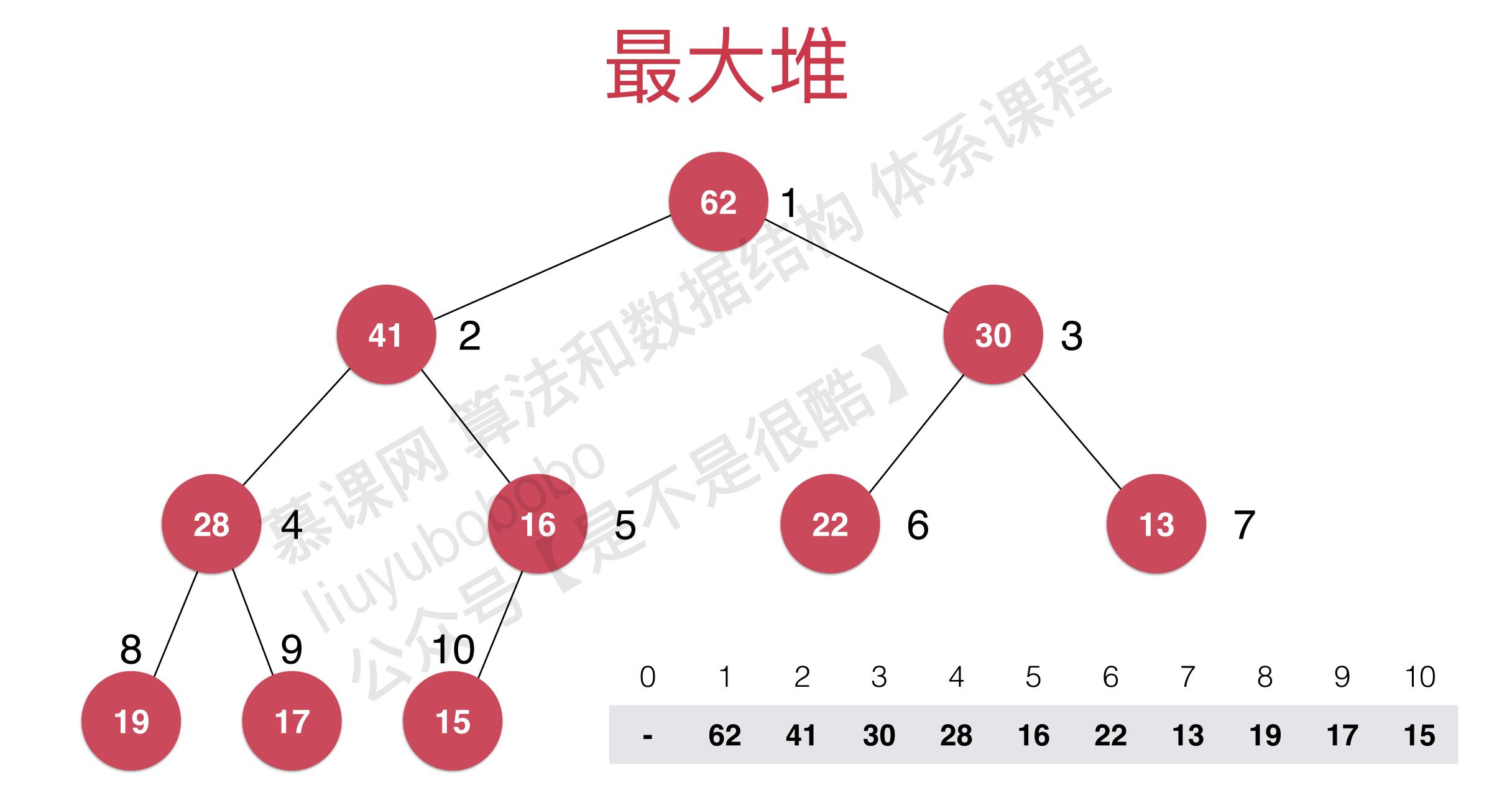


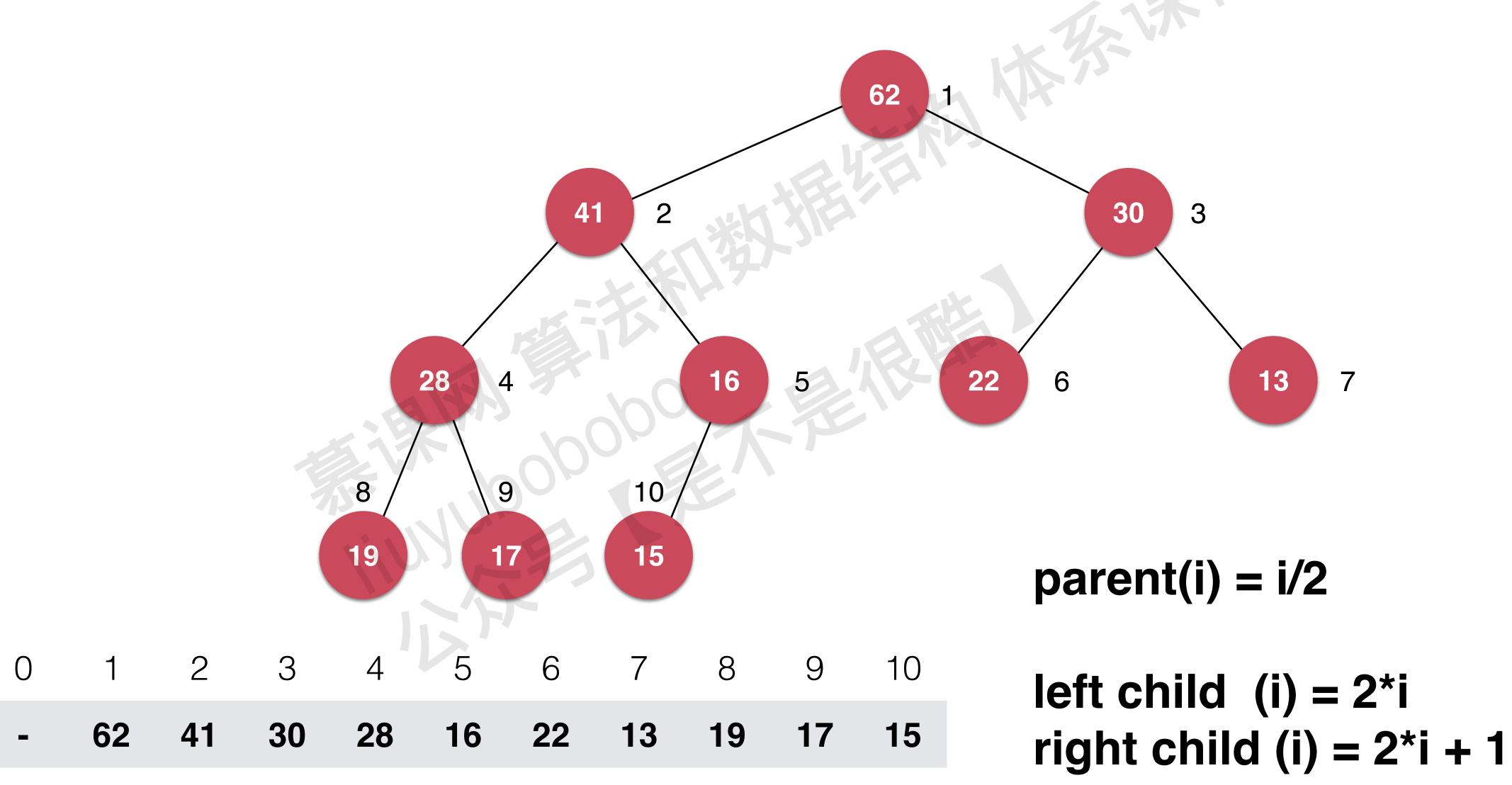
### 二叉堆是一棵完全二叉树

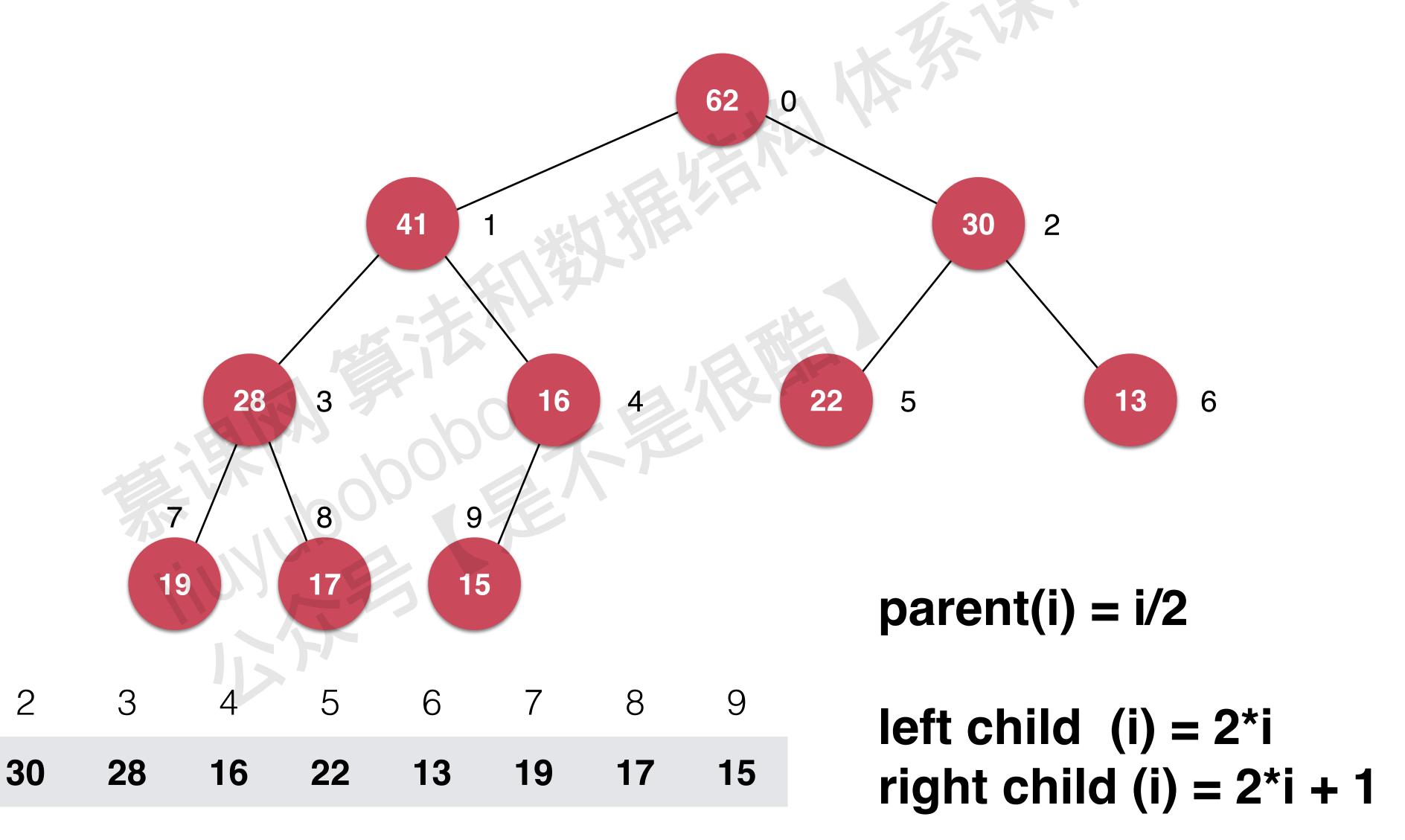


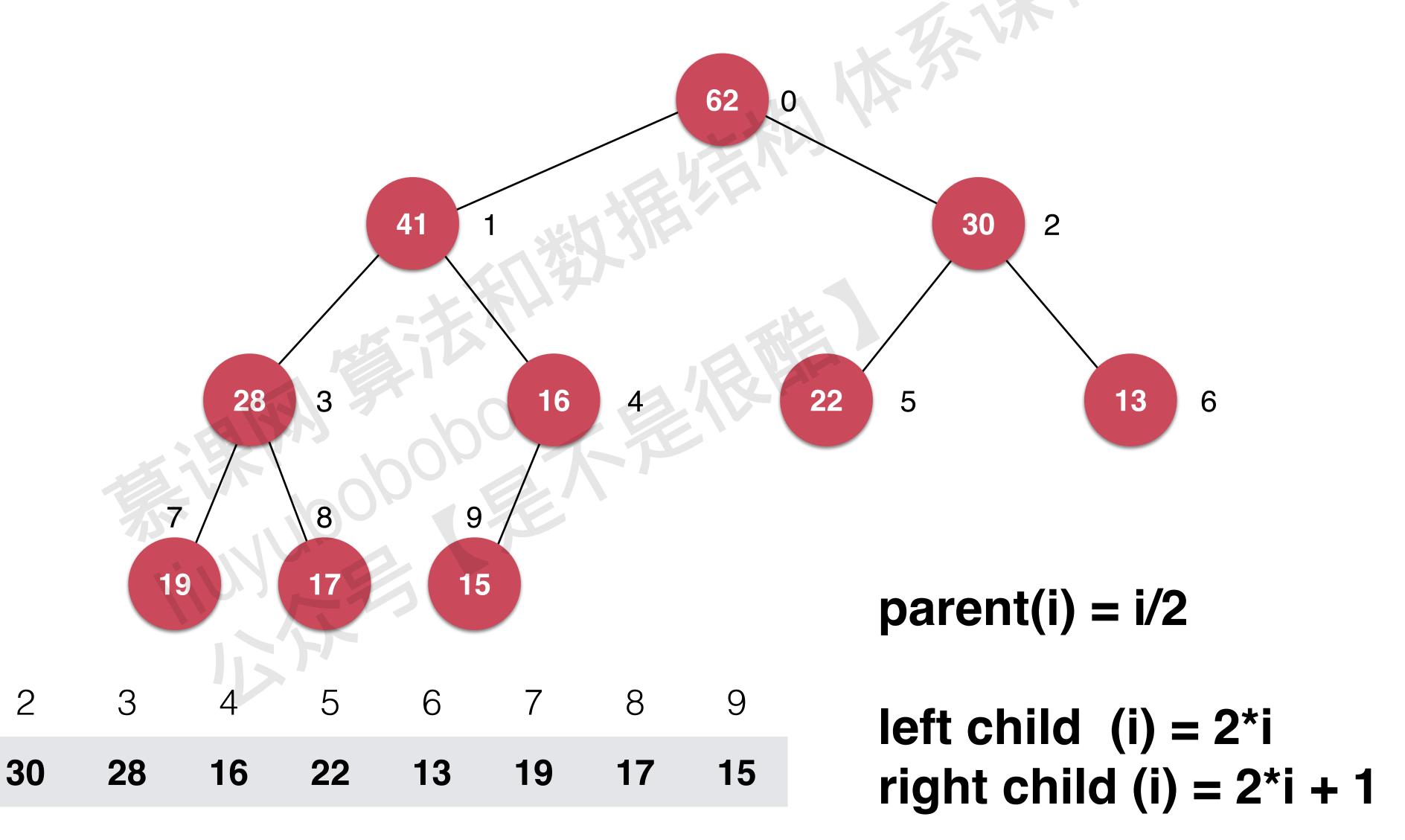
### 二叉堆的性质

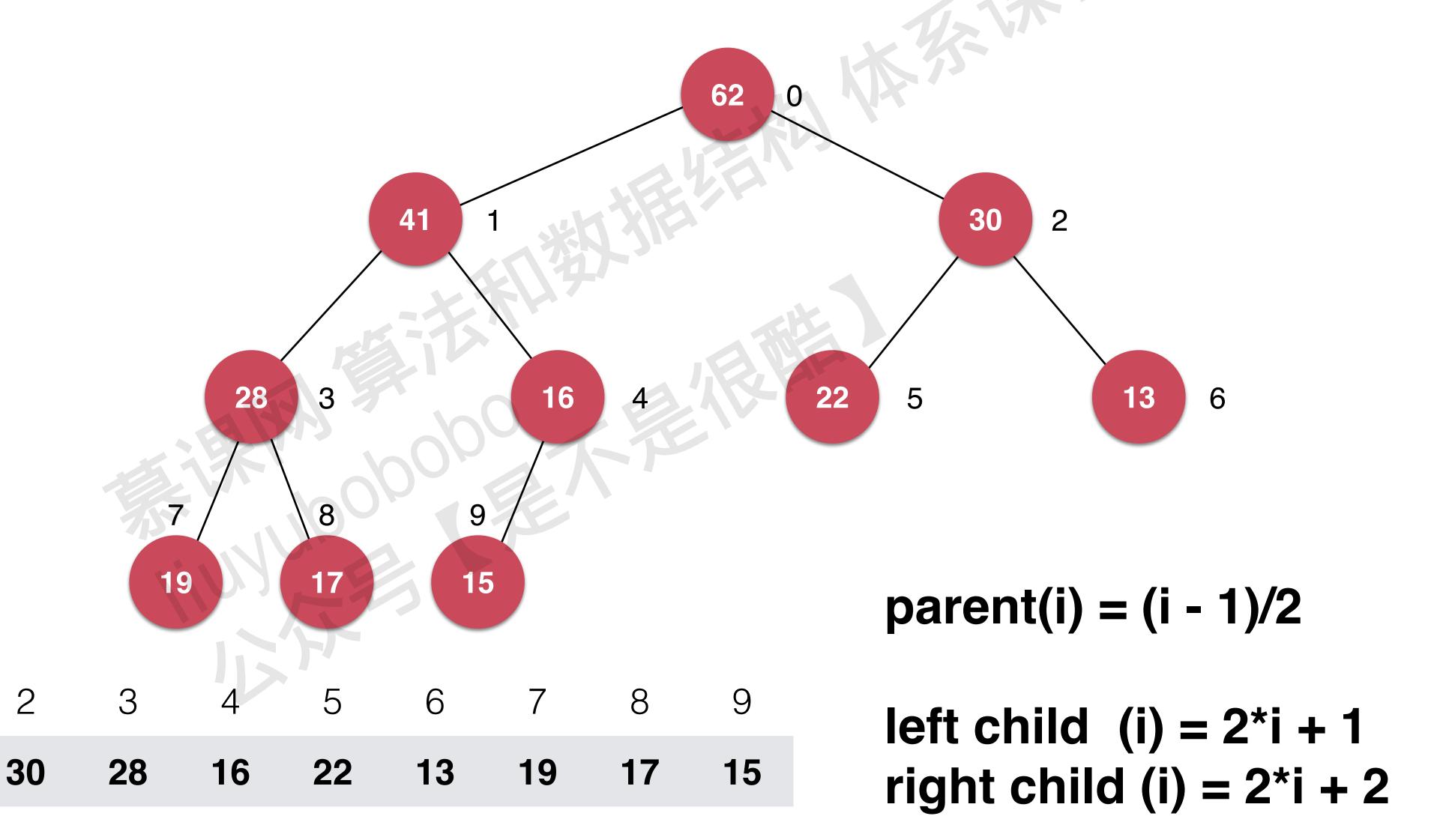






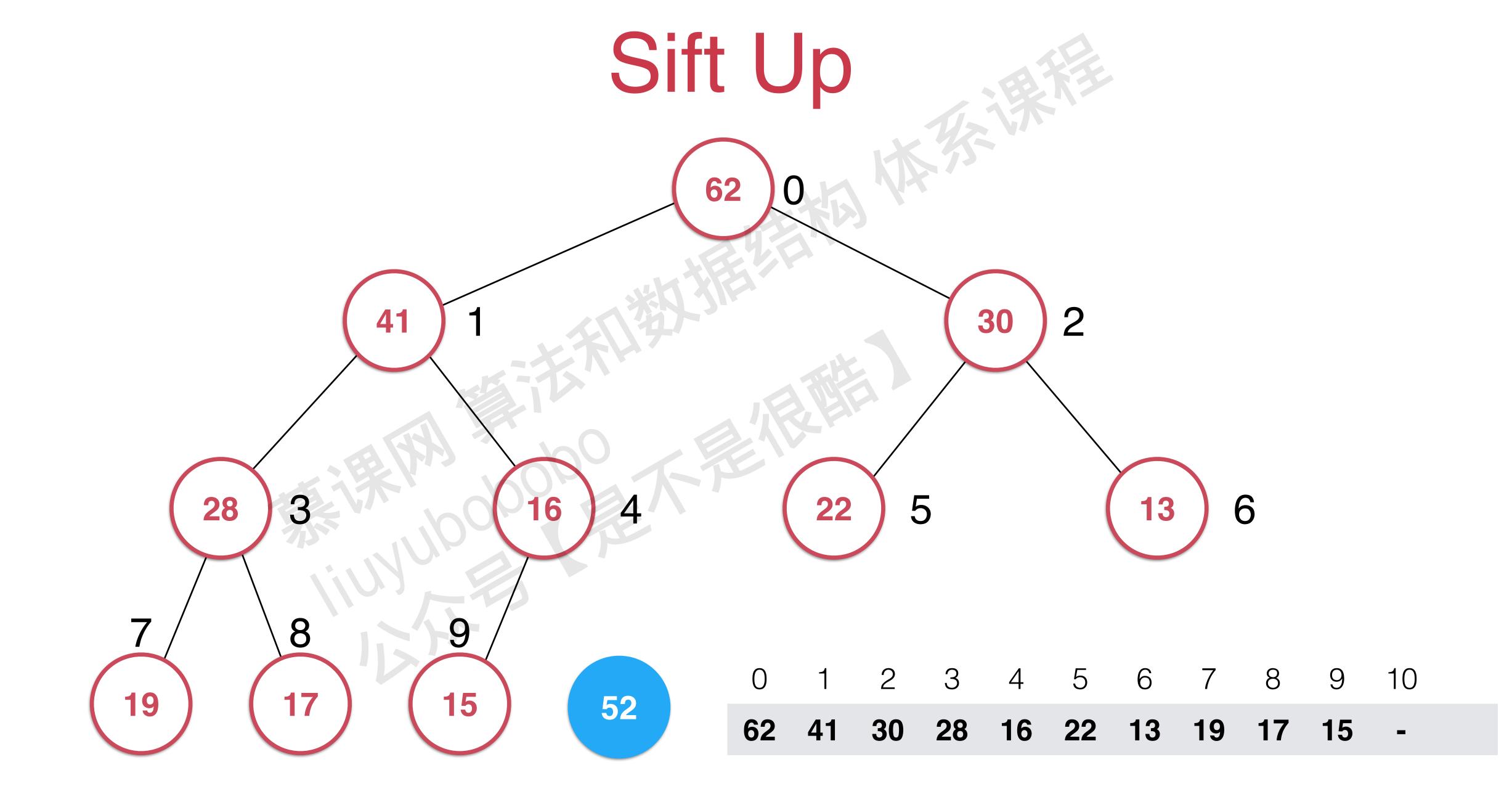


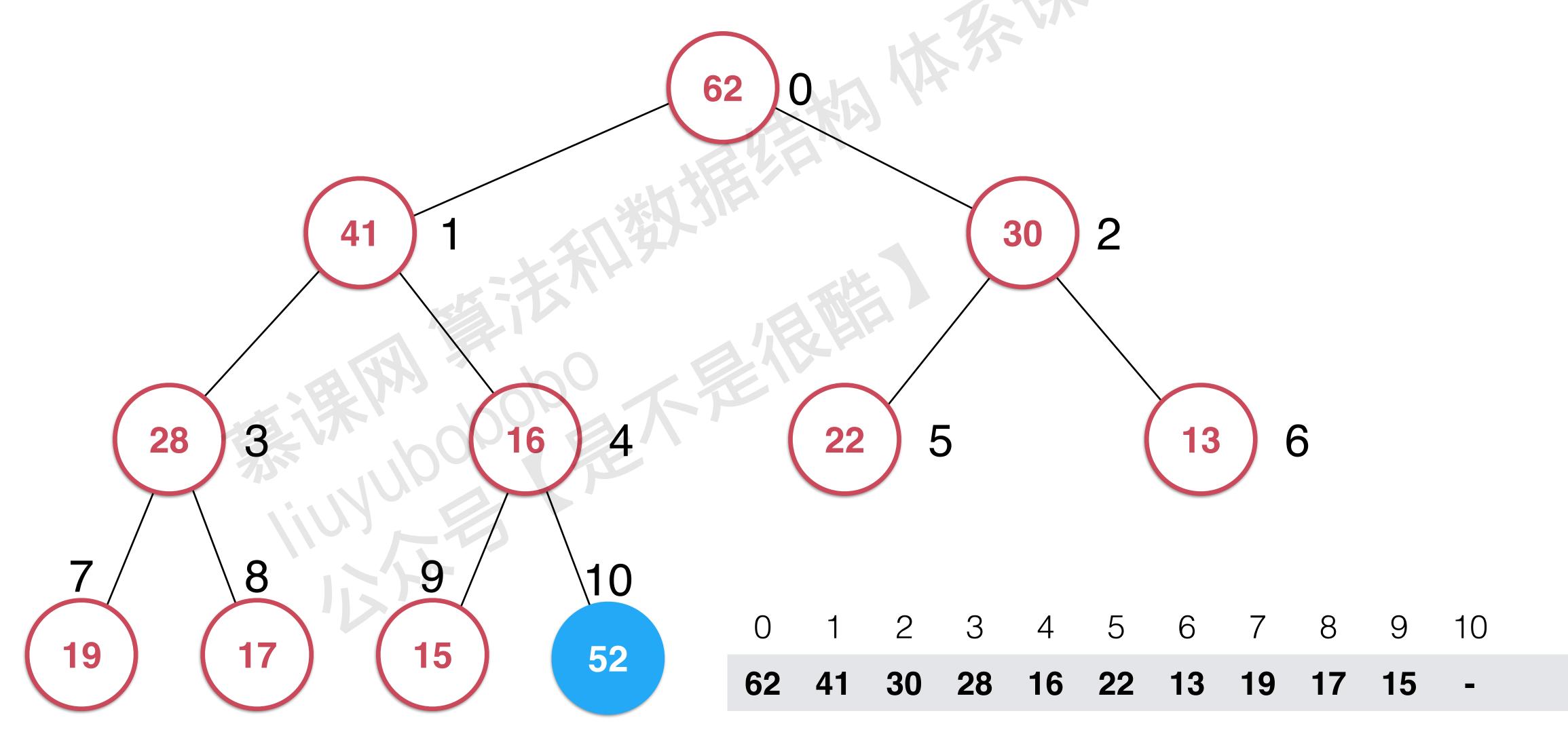


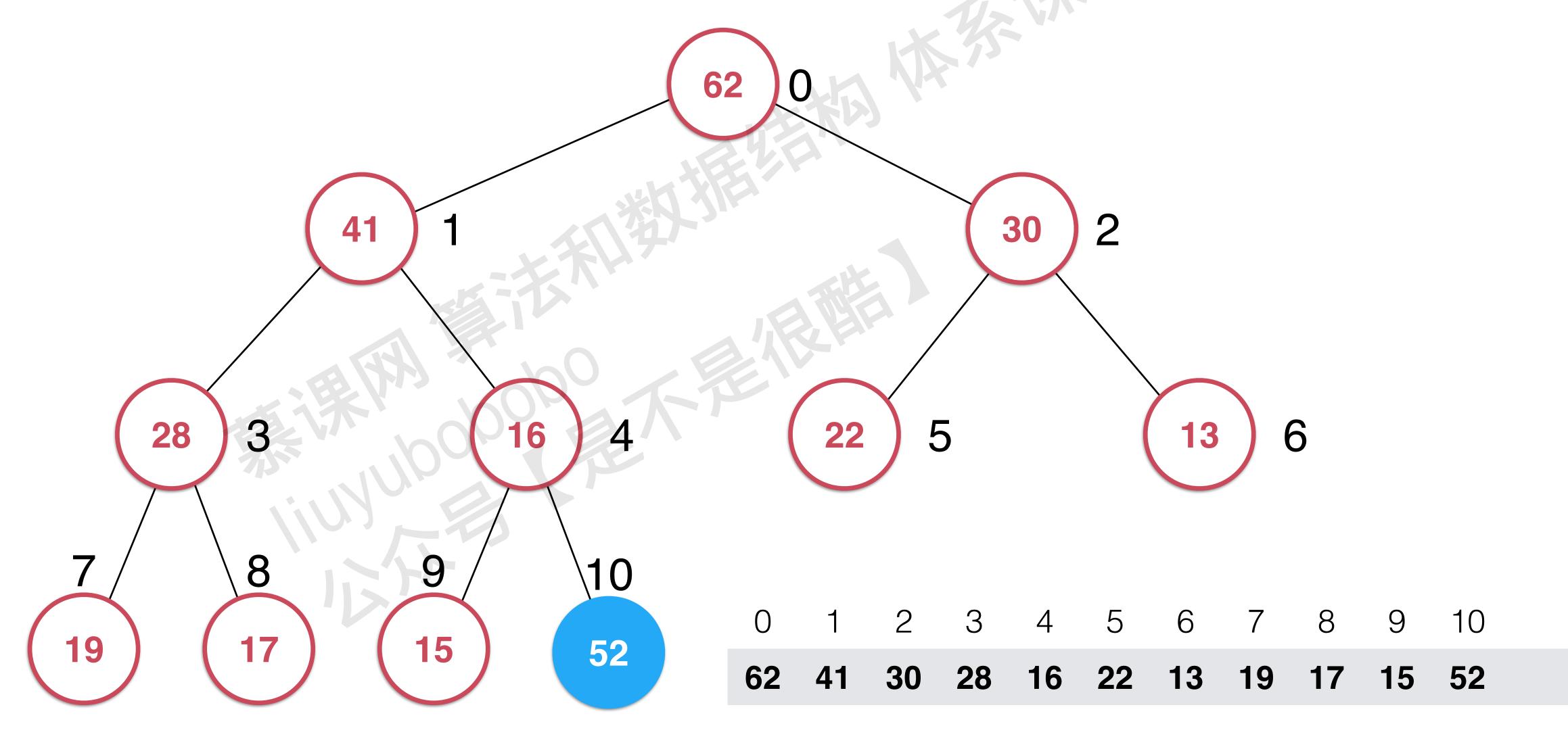


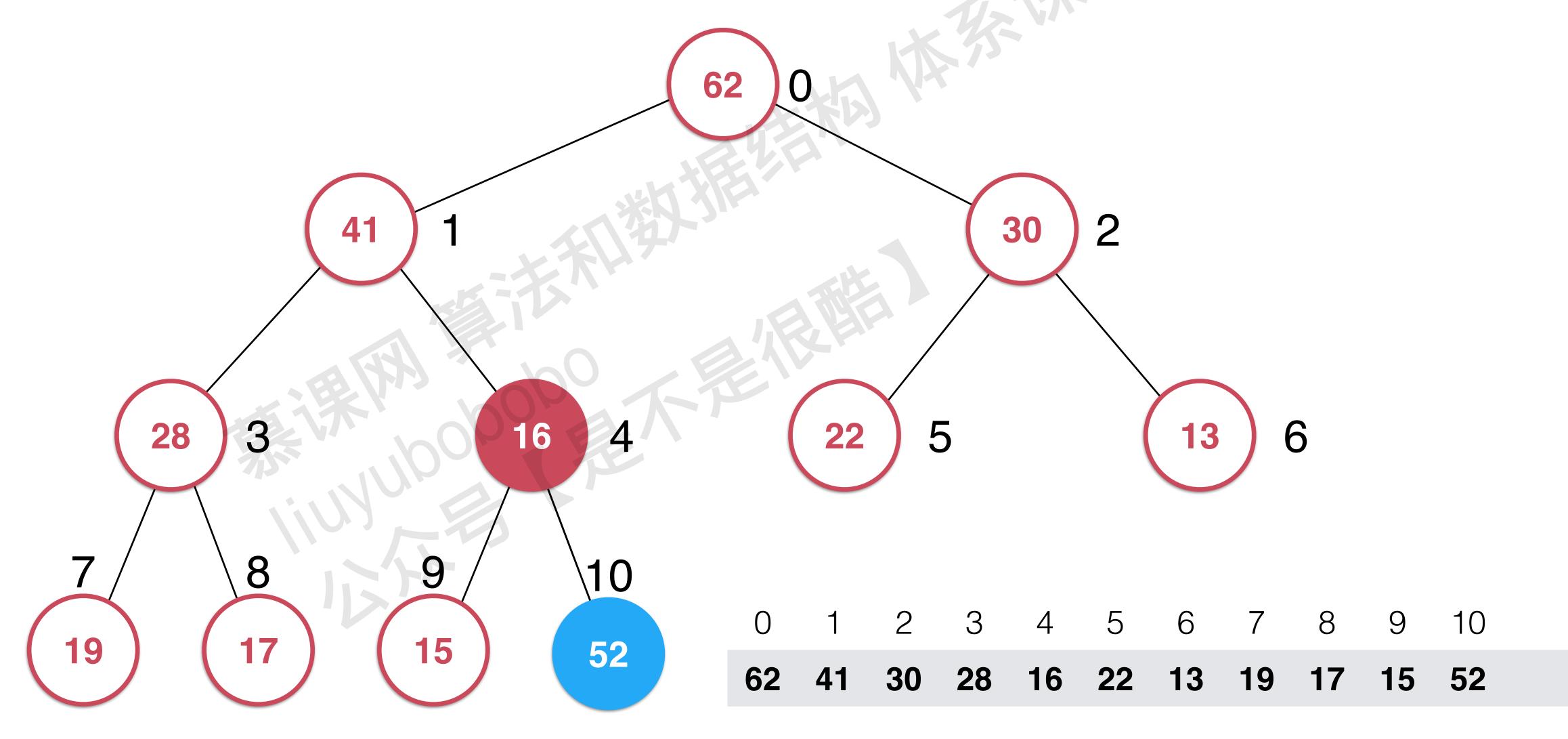
# 实践: 堆的基本框架

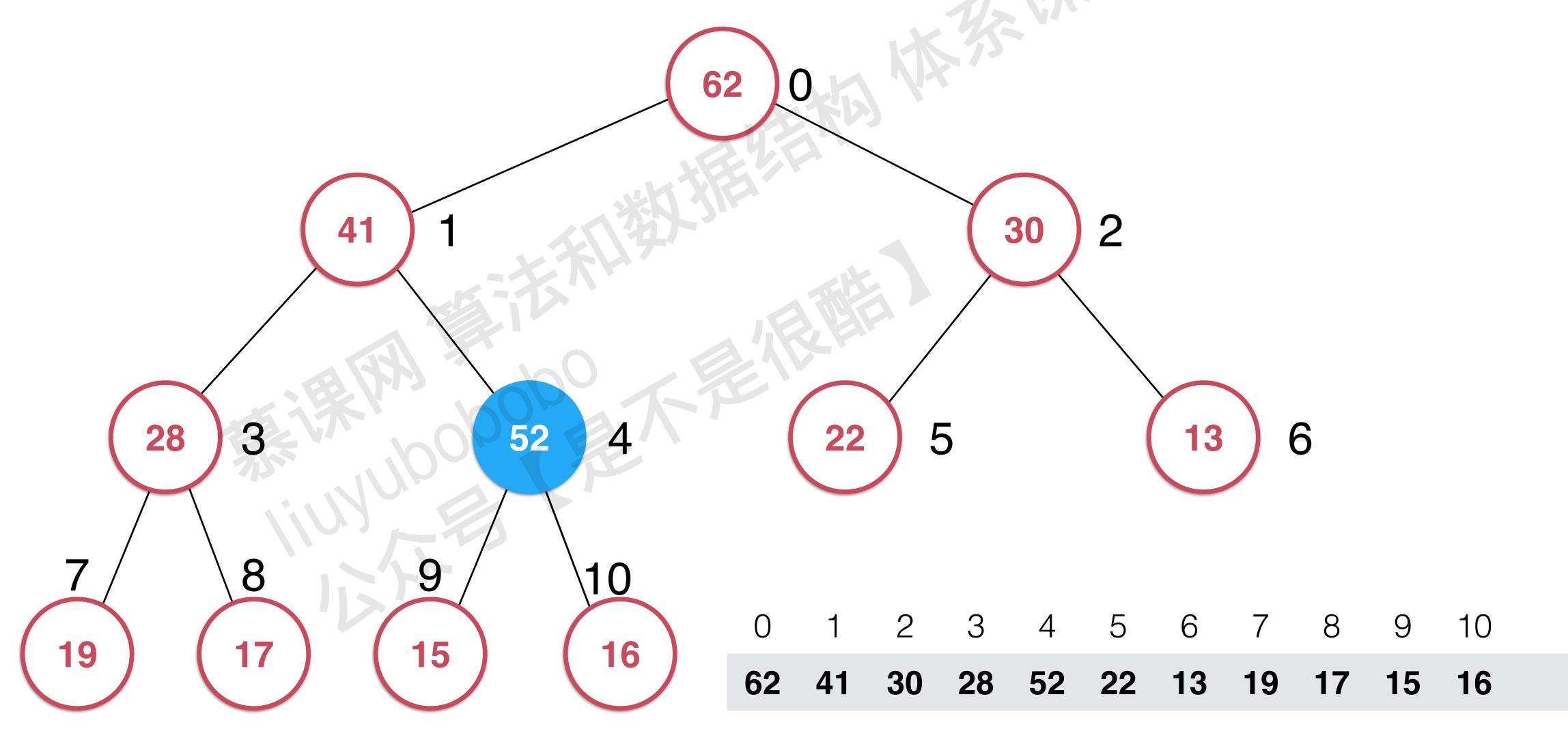
# 向堆中添加元素和Sift Up

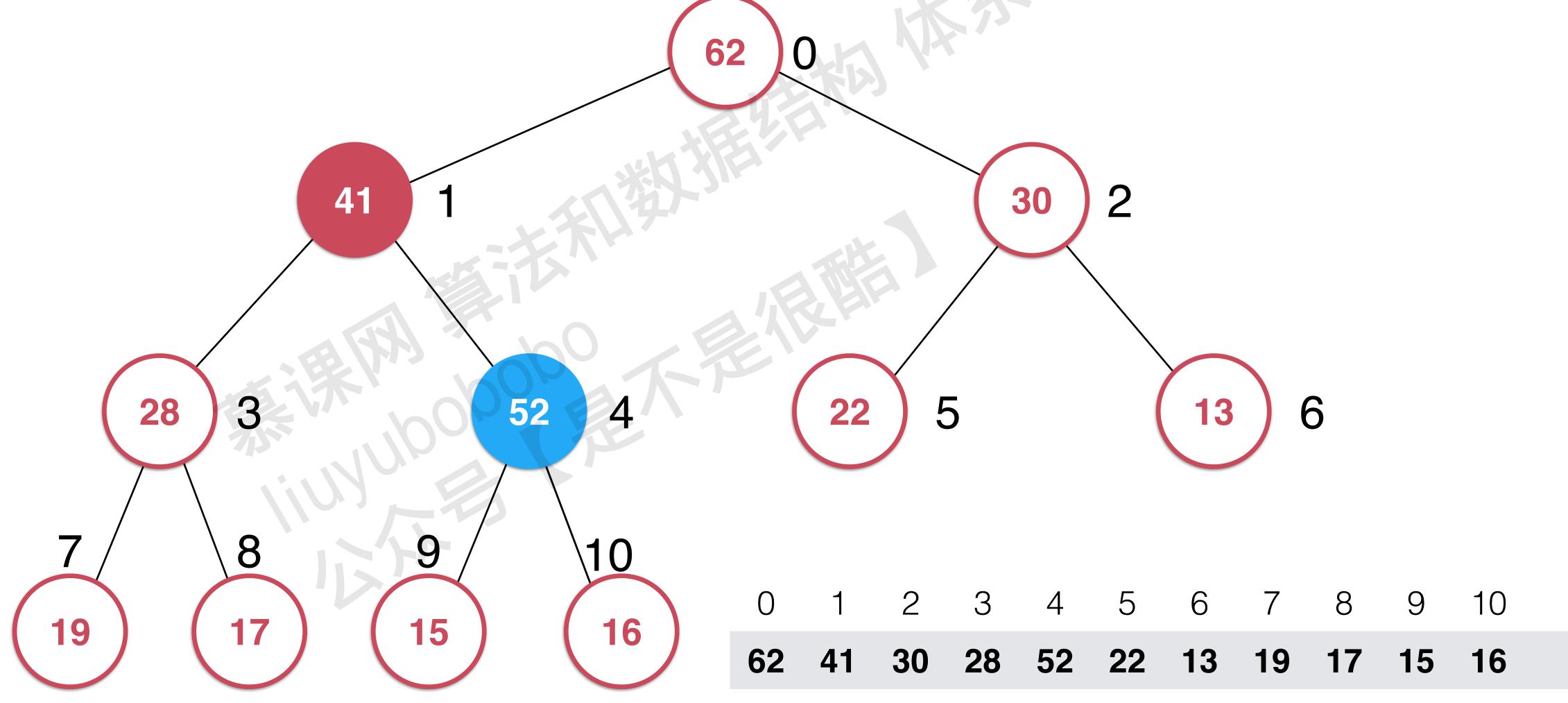


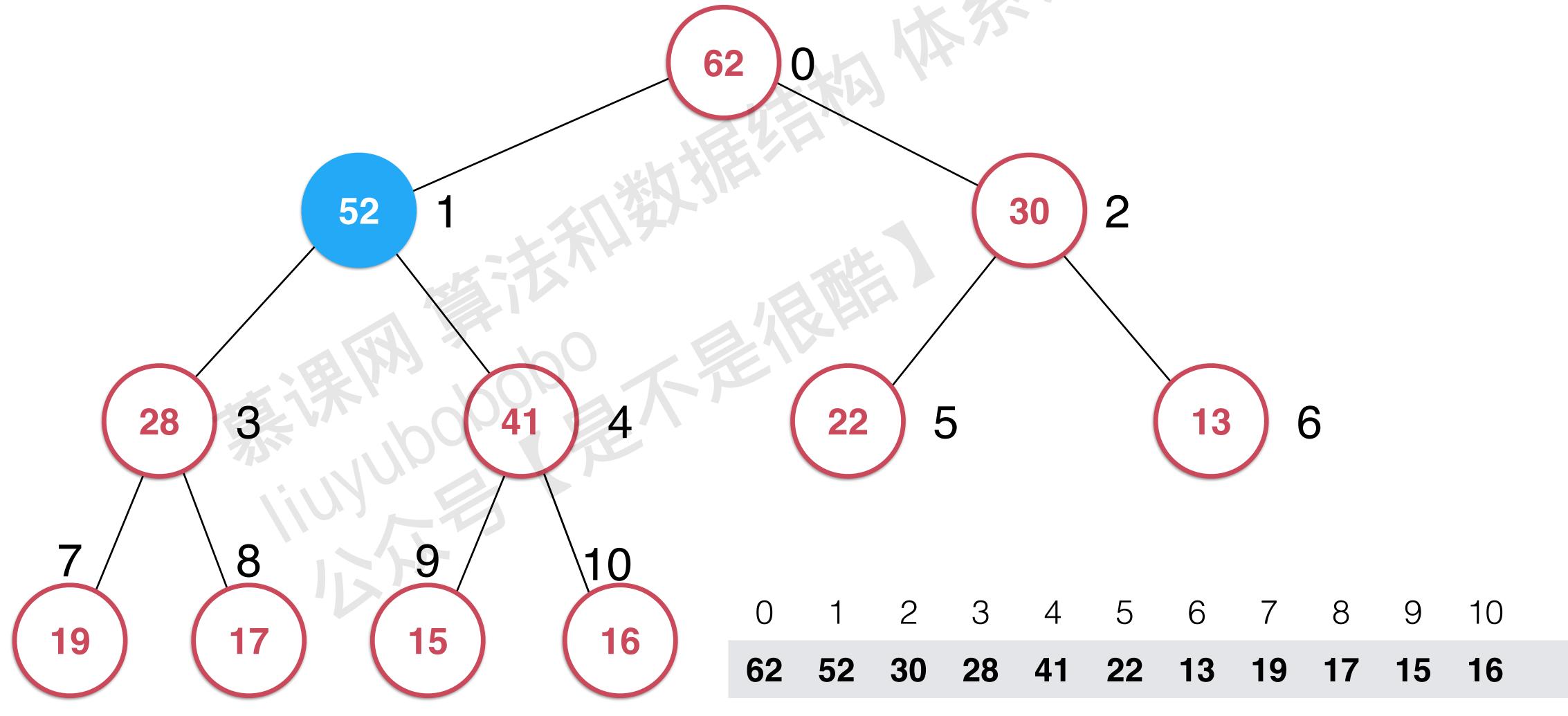


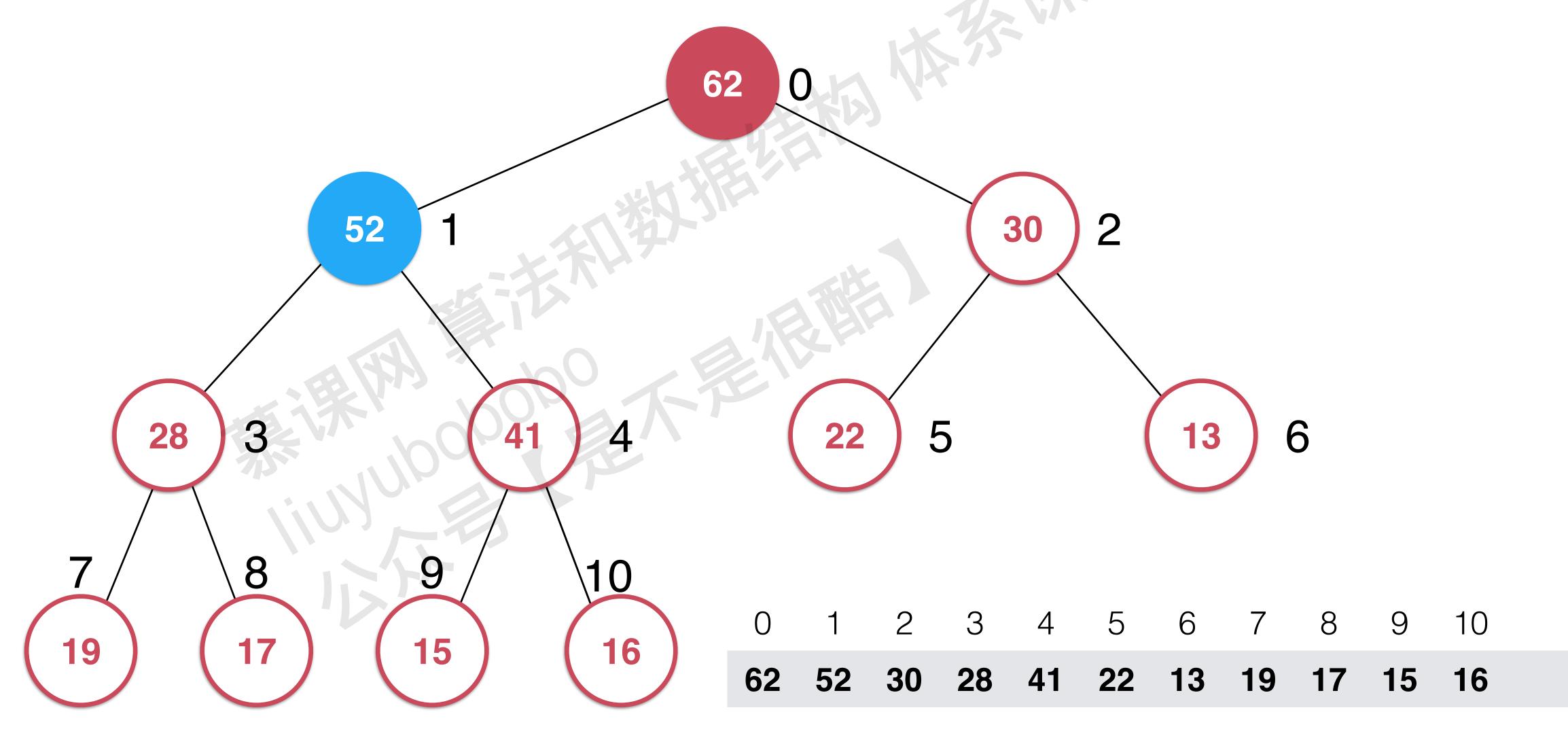


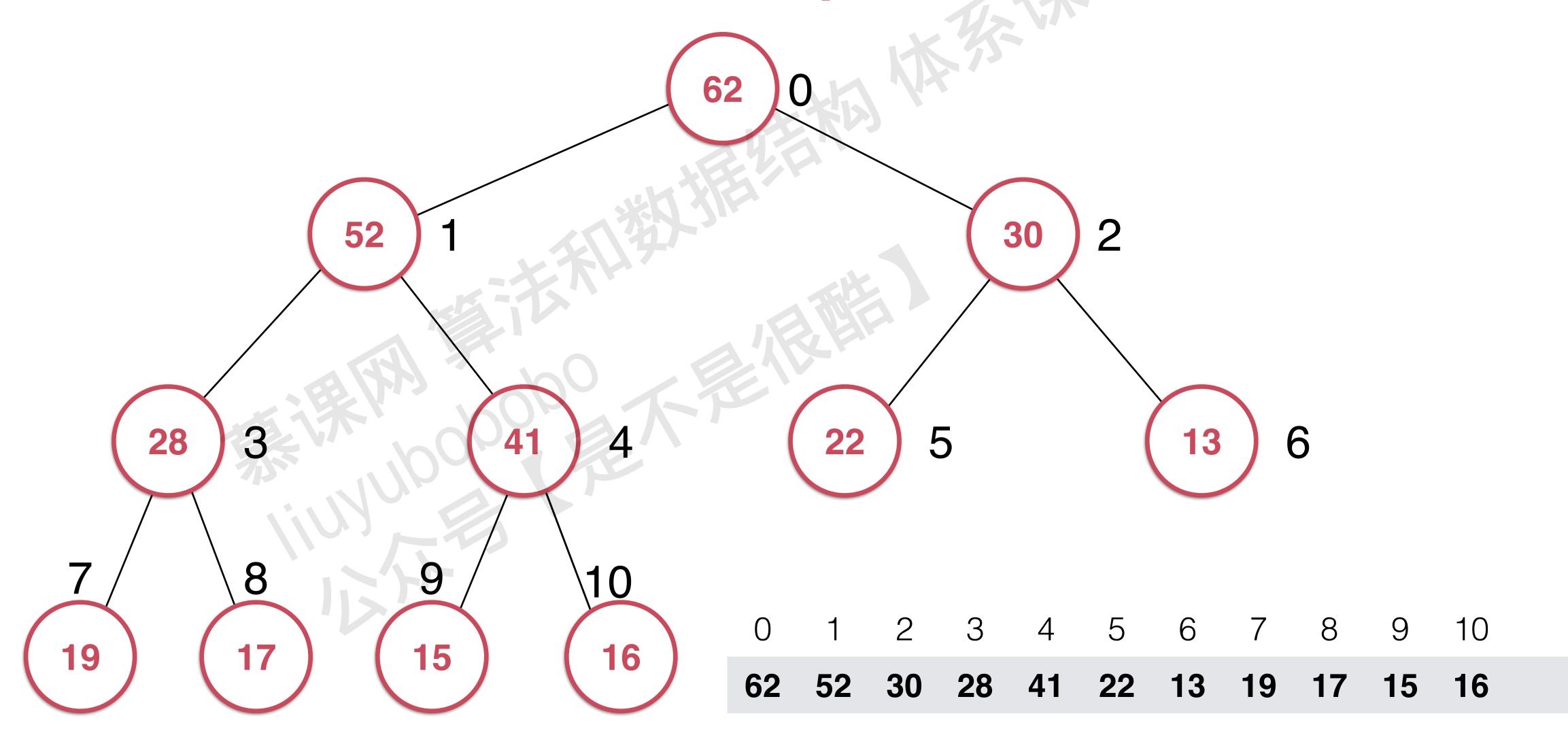






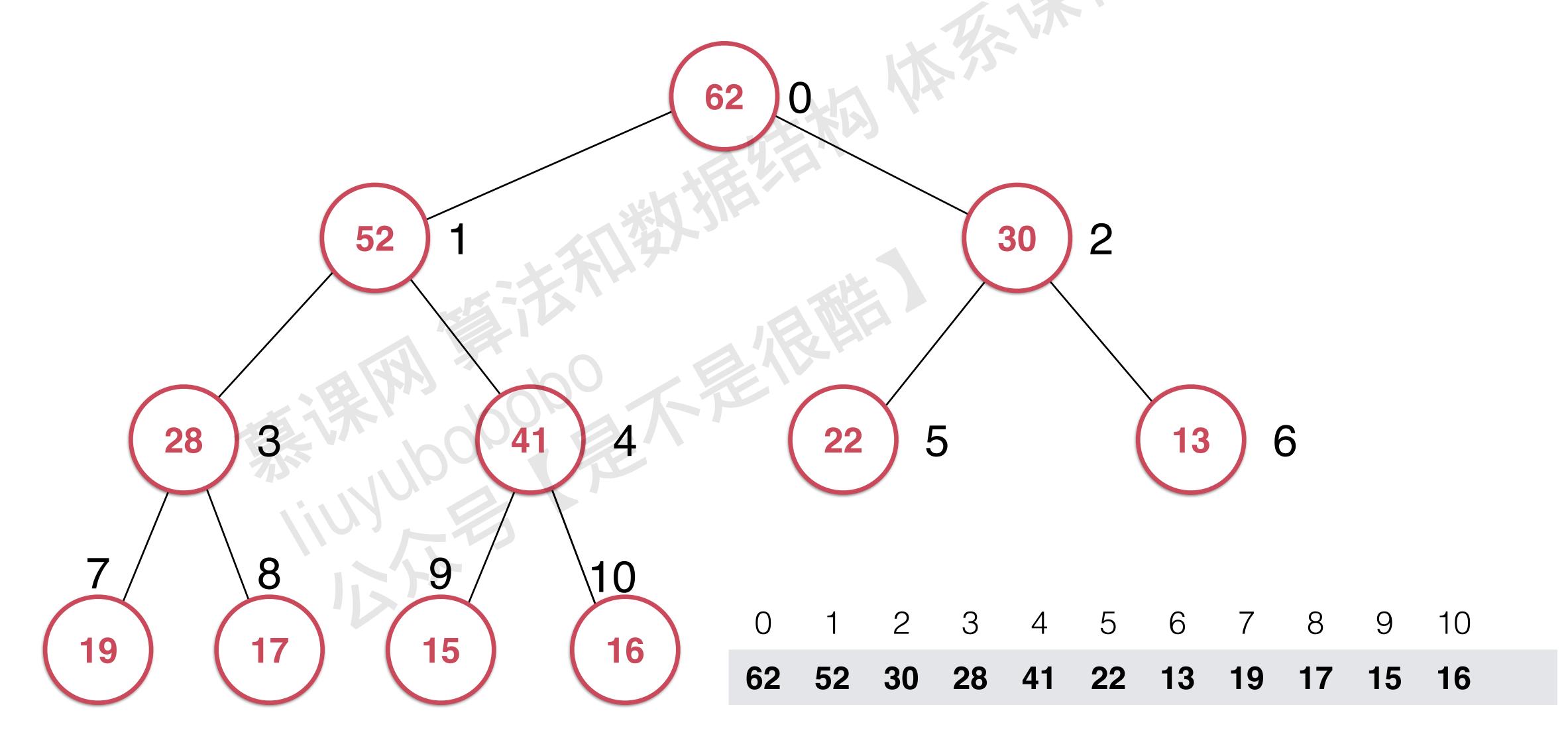


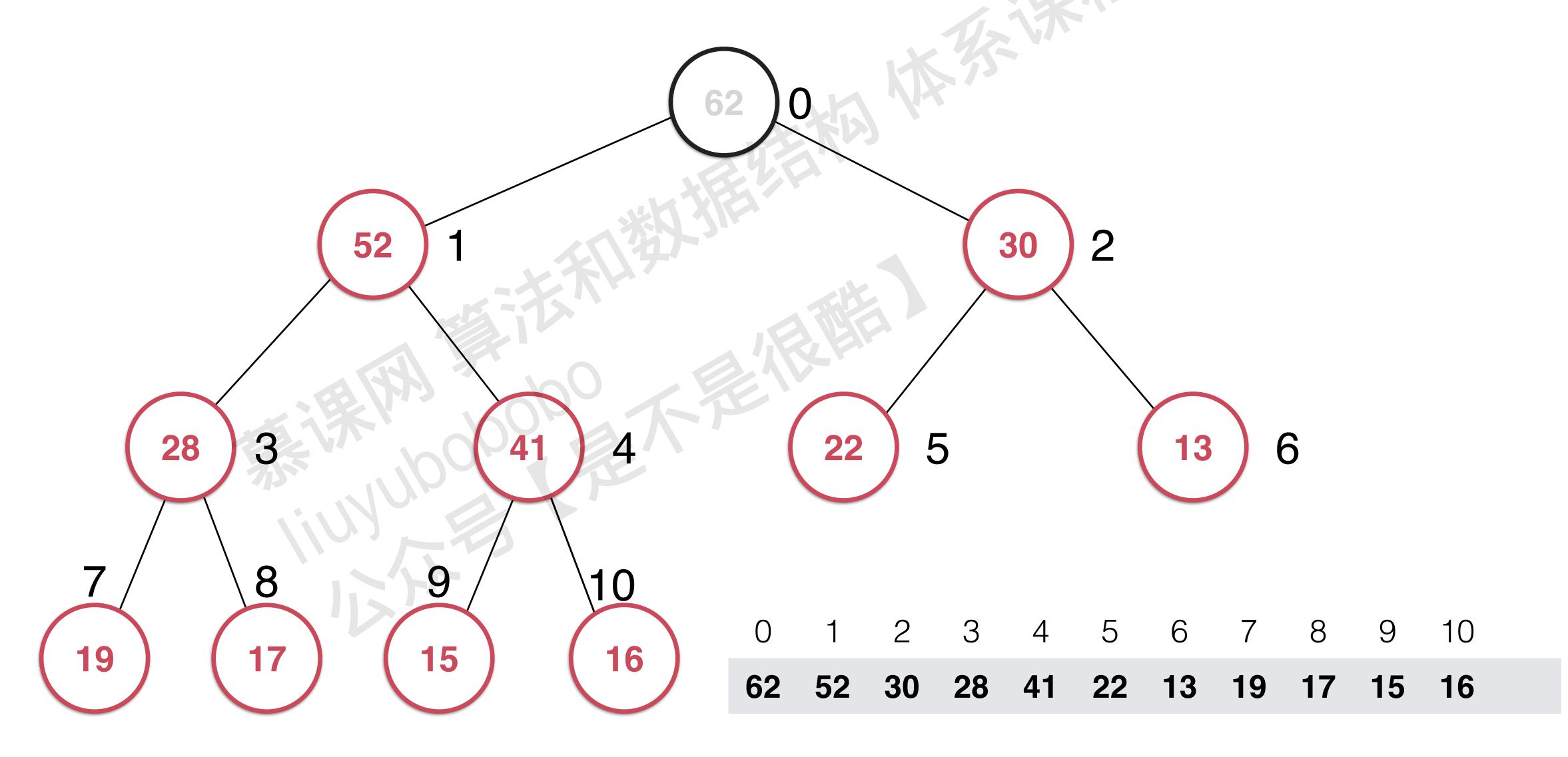


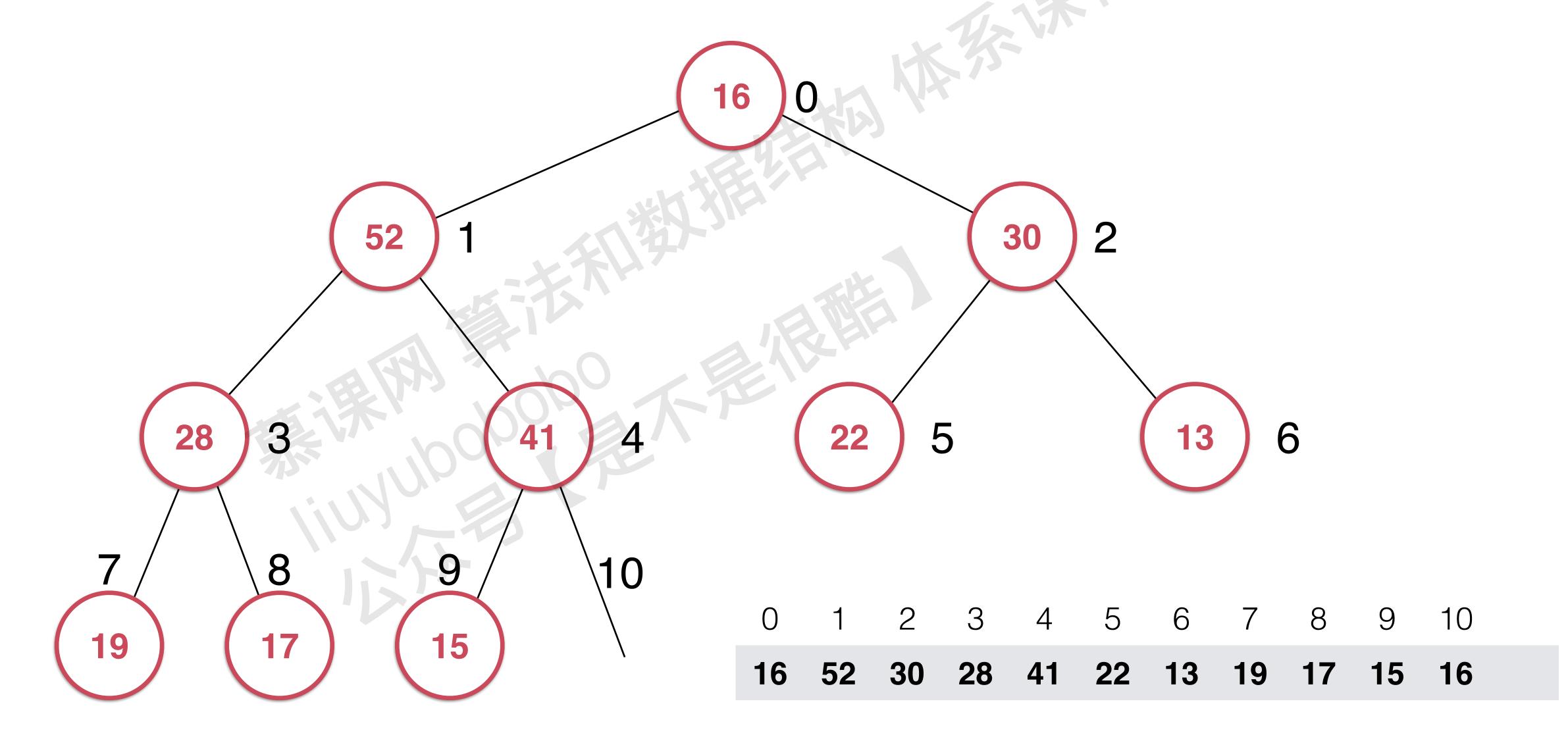


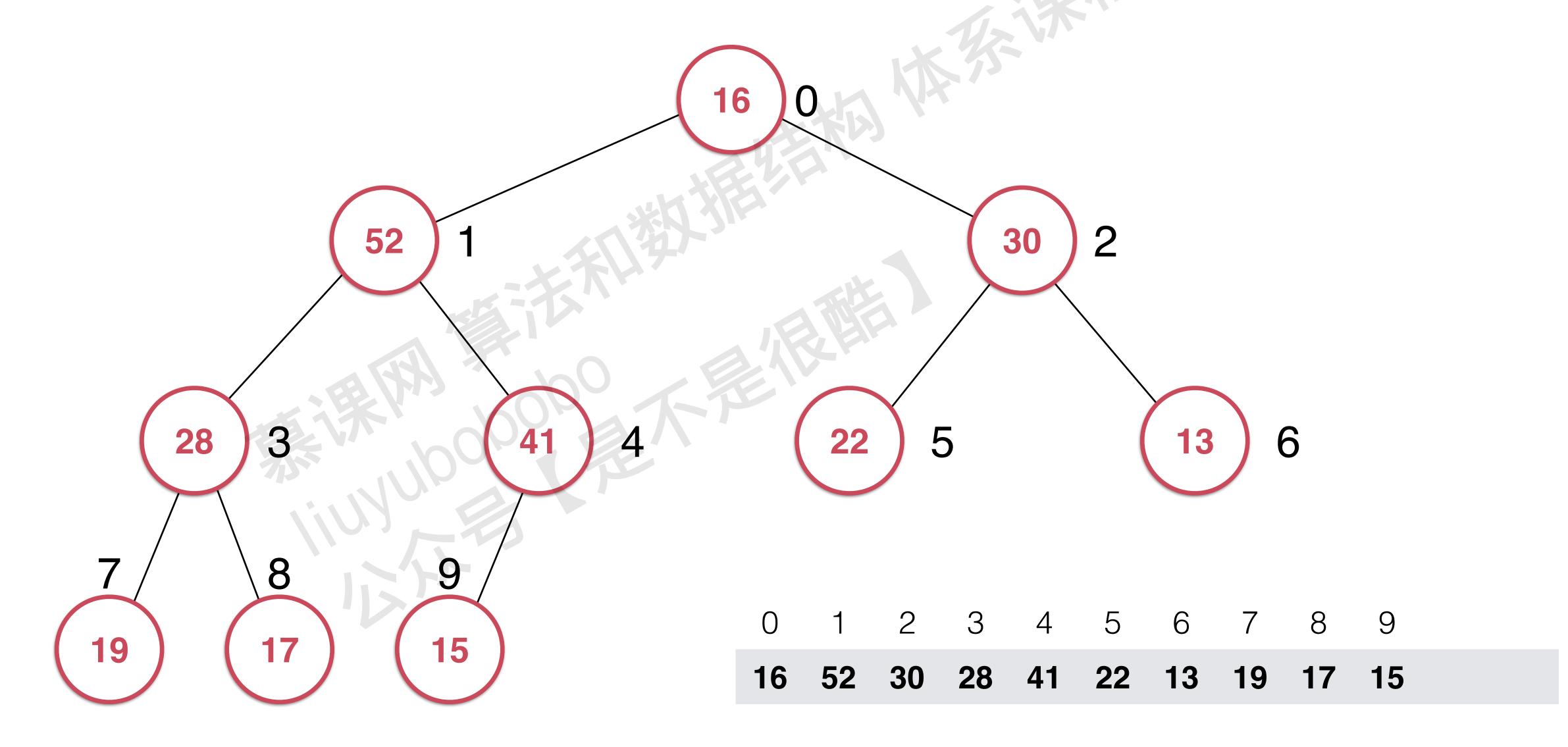
# 实践: Sift Up和 add

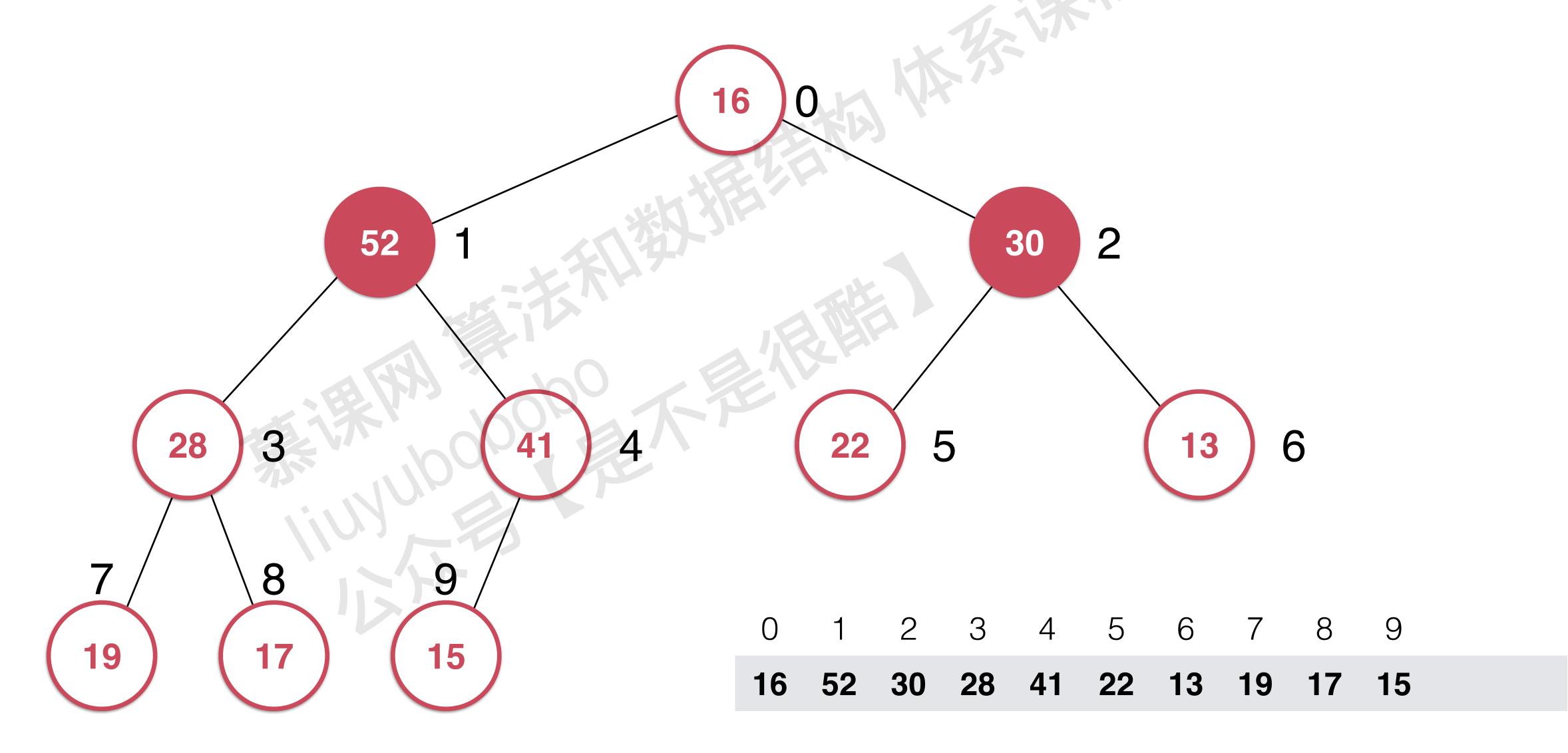
# 取出堆中的最大元素和Sift Down

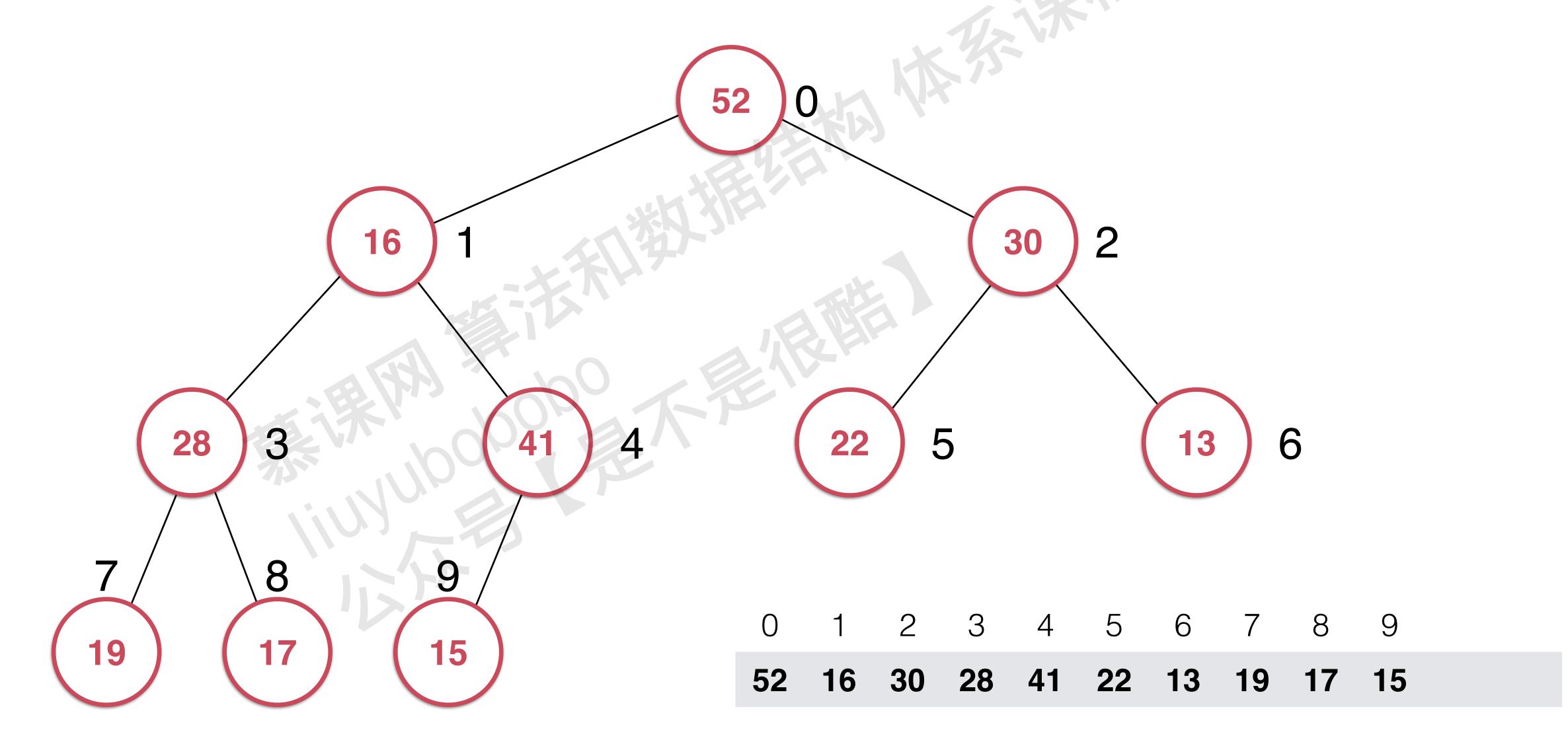


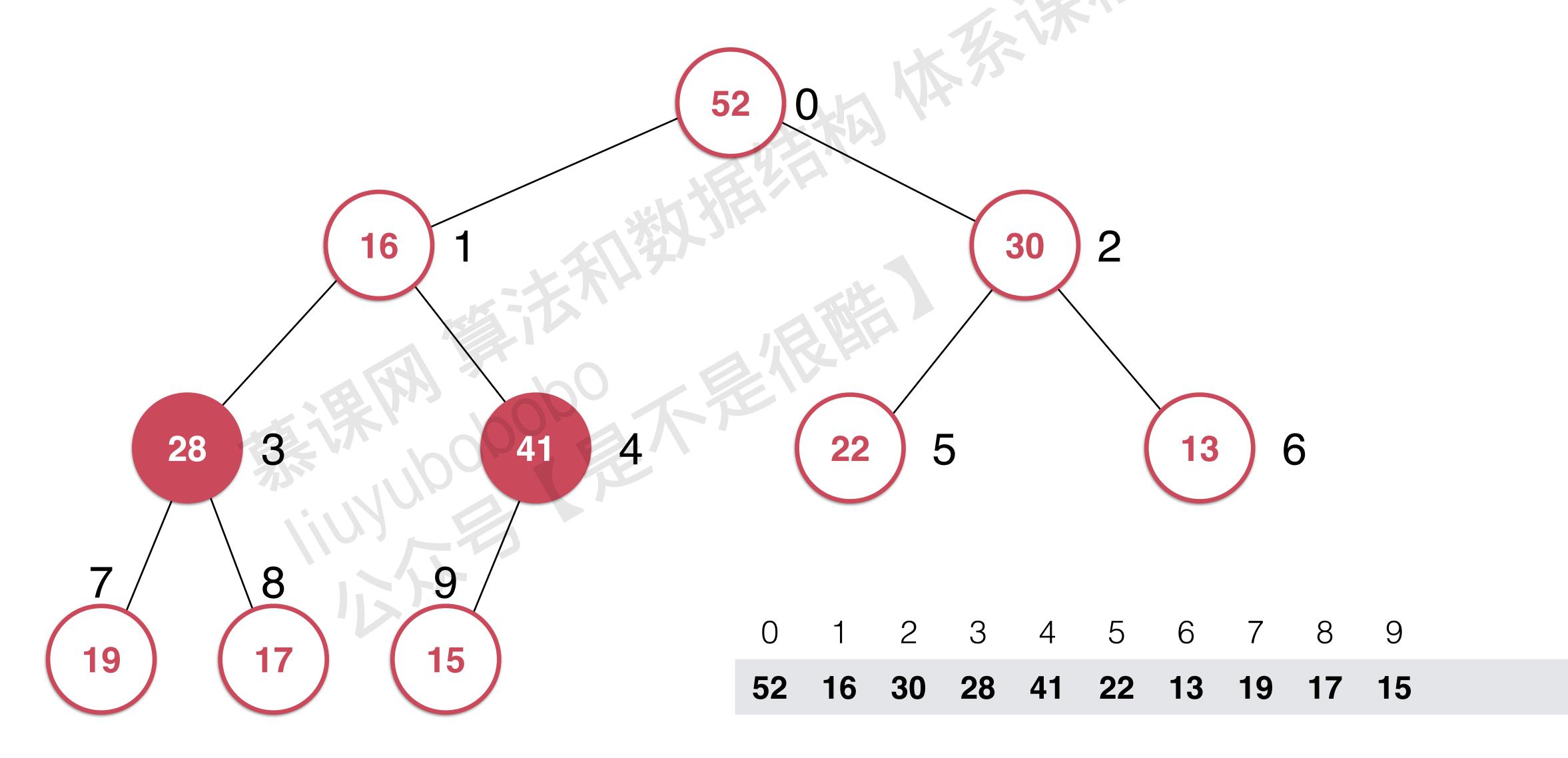


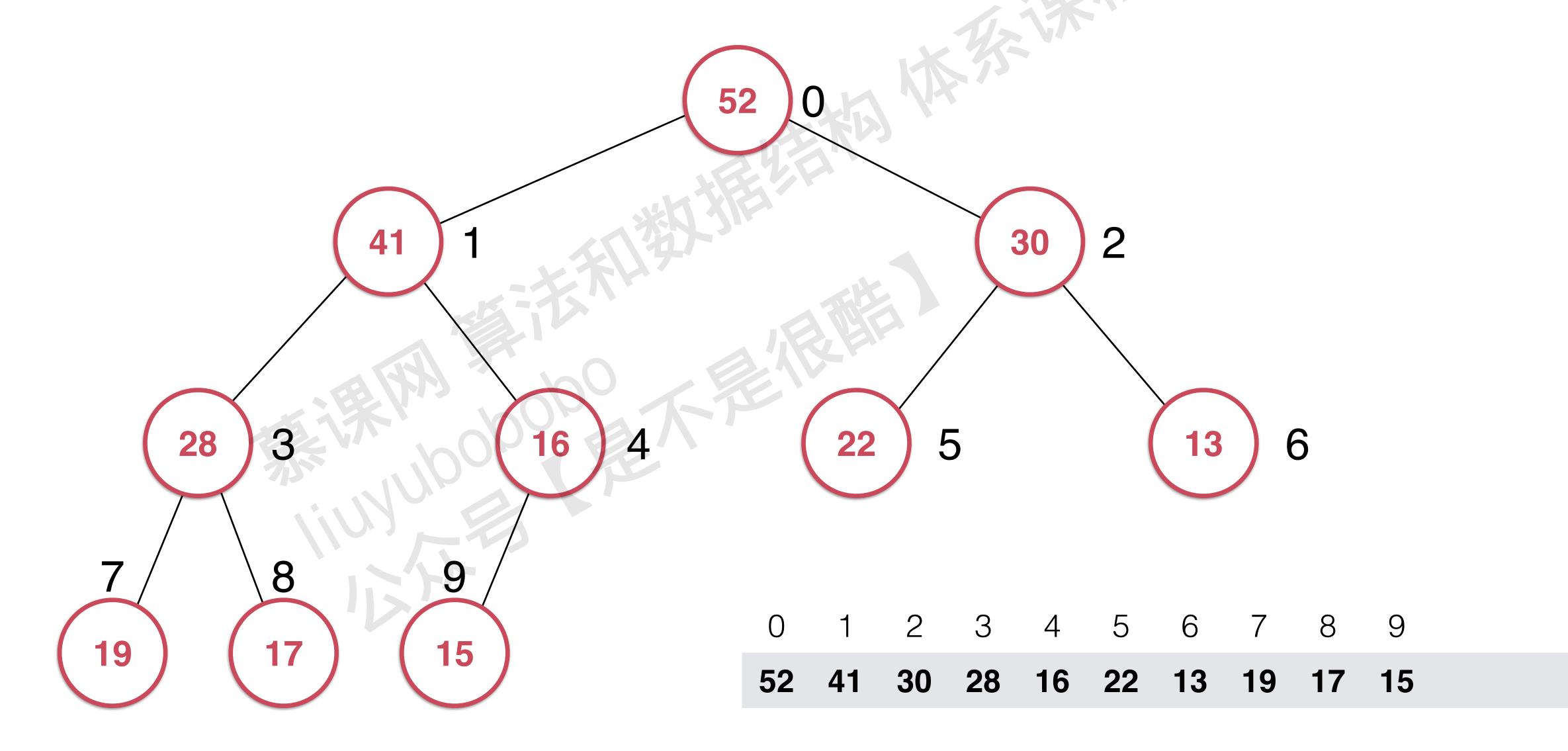


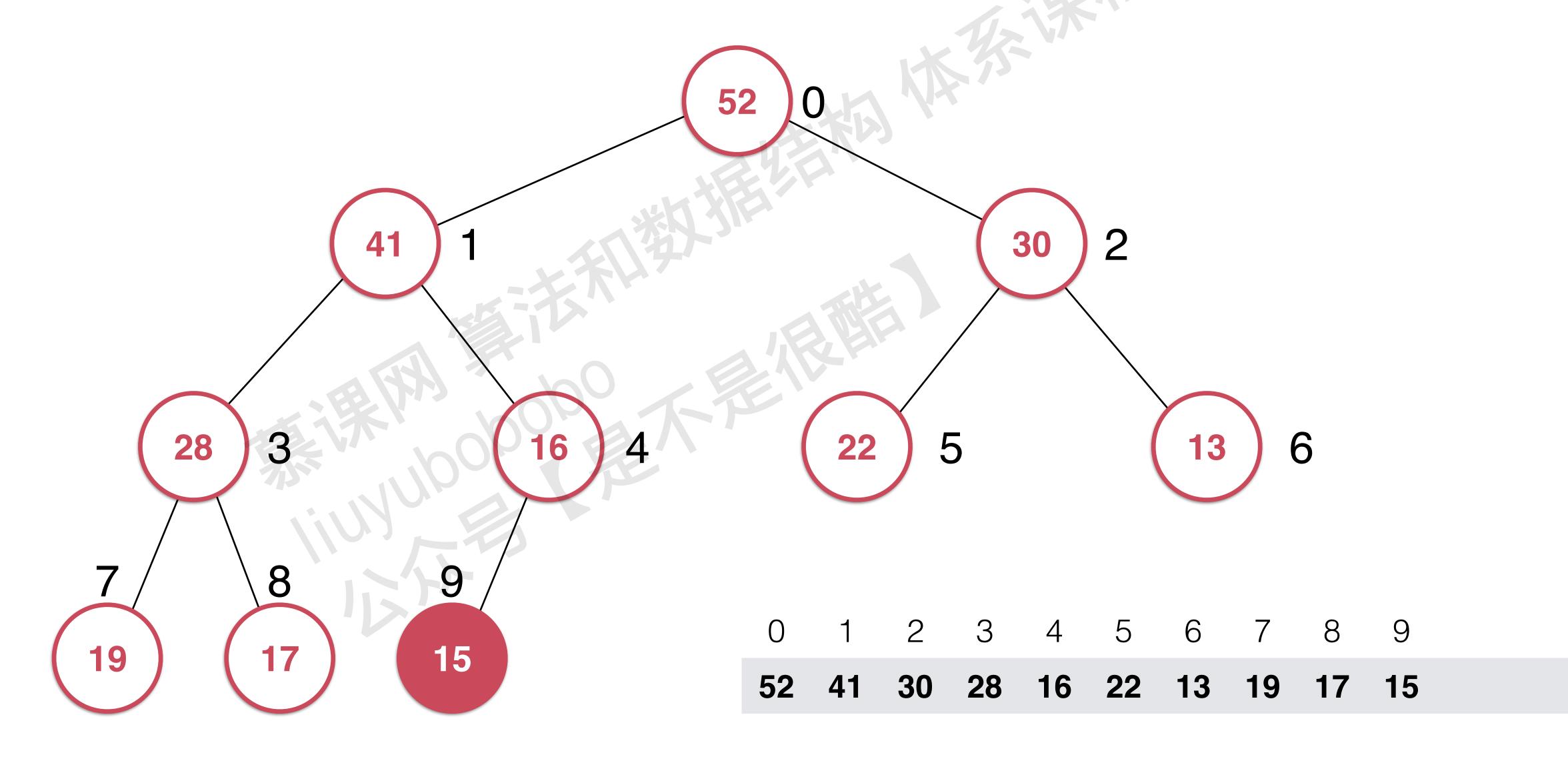


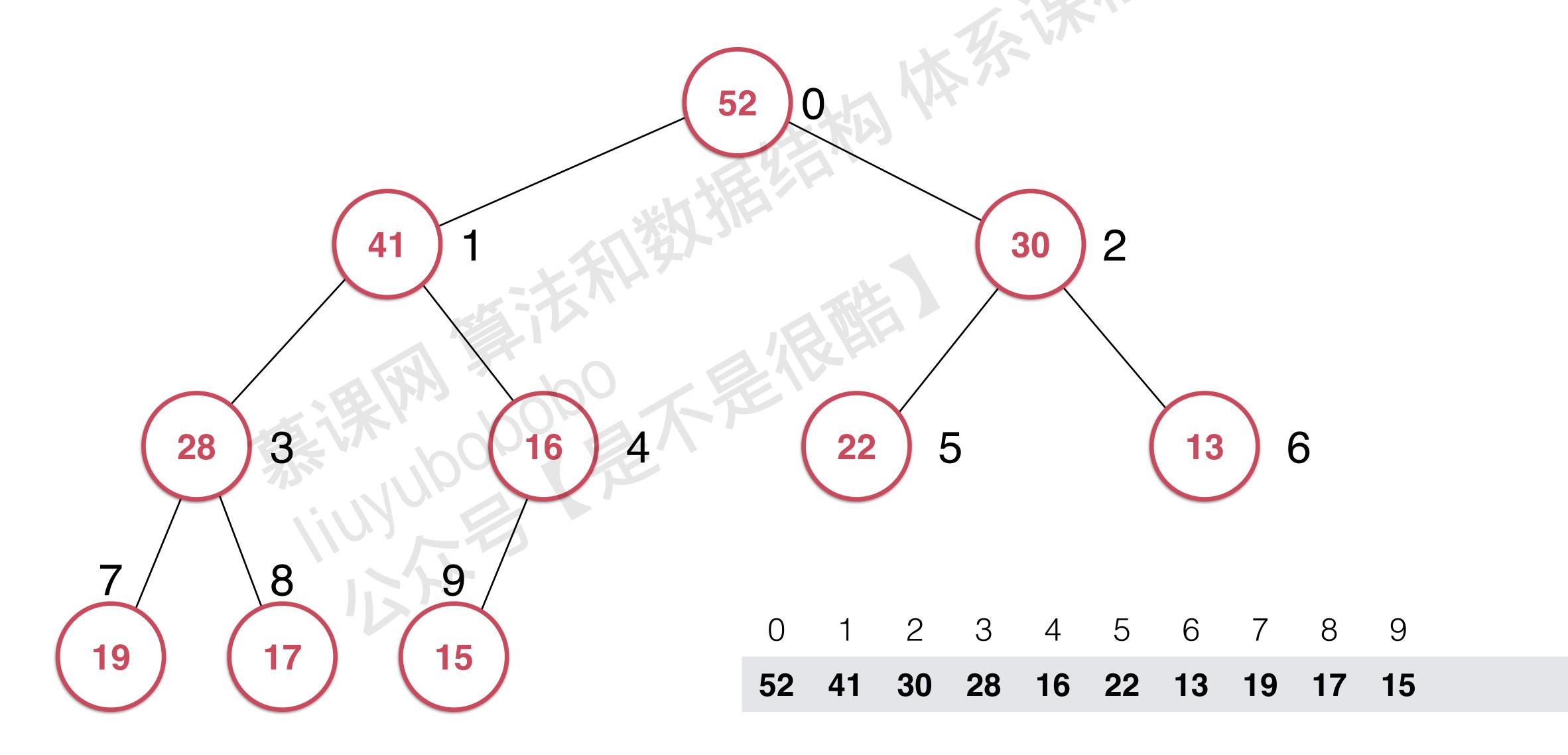








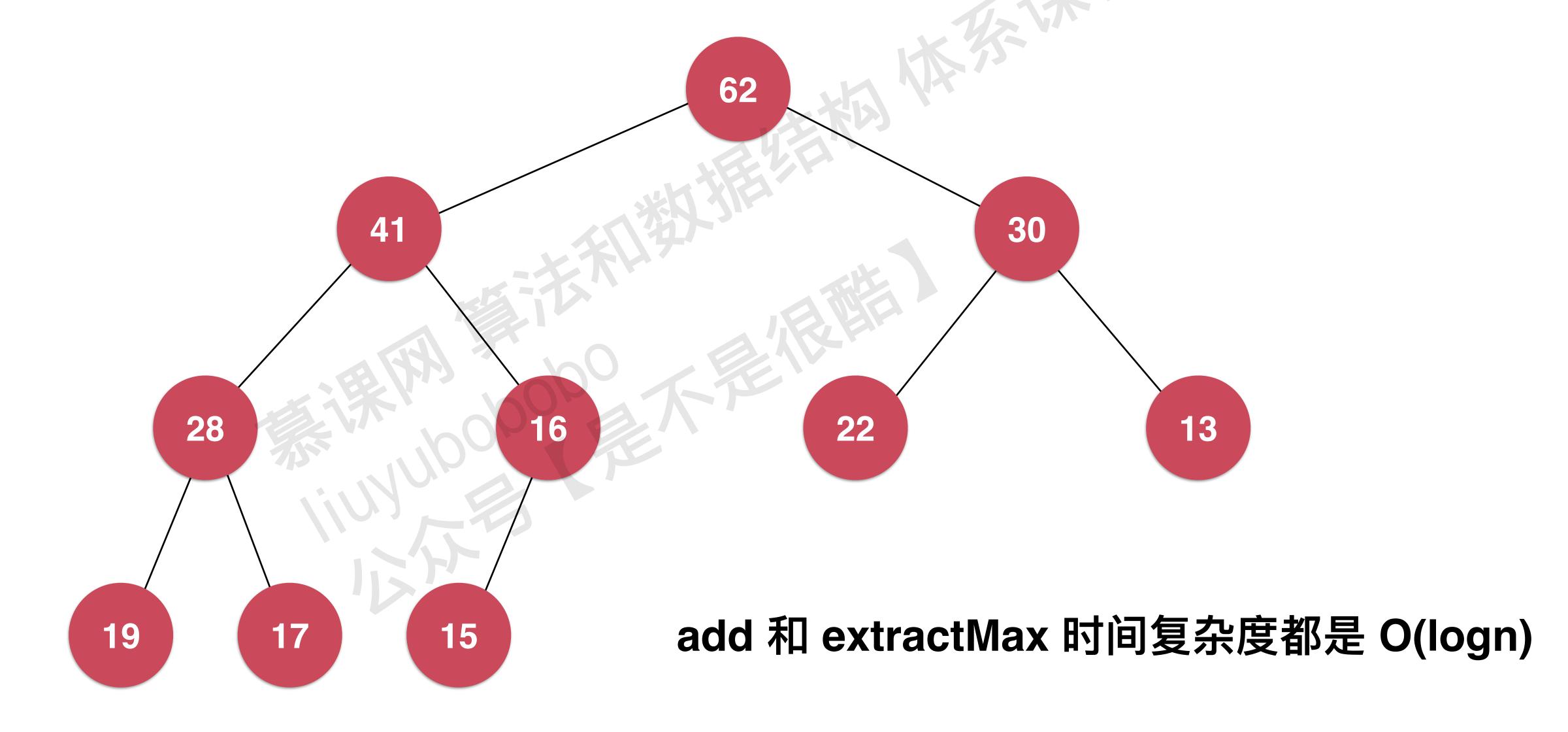




# 实践: Sift Down 和 extractMax

实践:测试堆

# 堆的时间复杂度分析



# 最简单的堆排序

# 实践:最简单的堆排序



# replace

replace:取出最大元素后,放入一个新元素

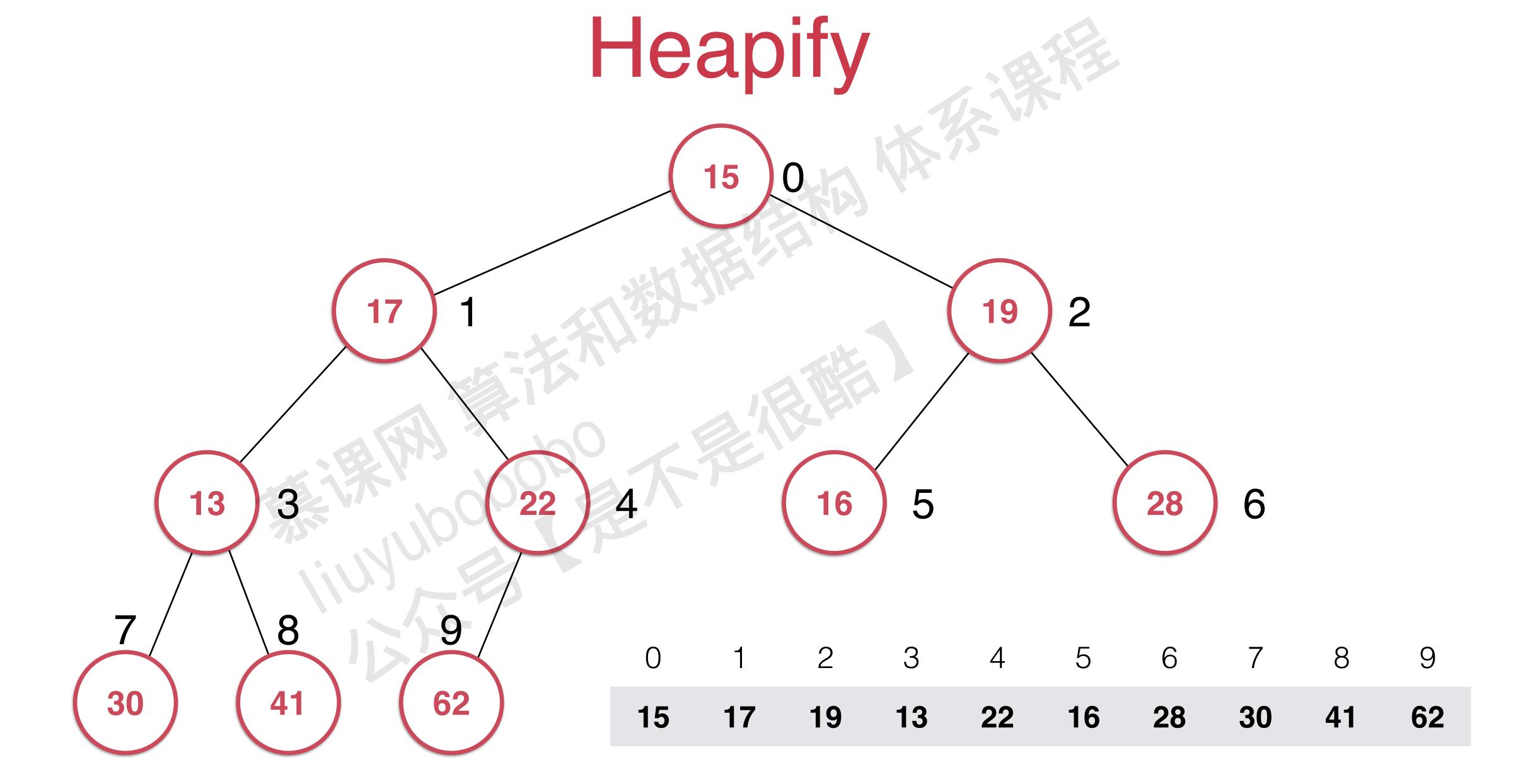
实现:可以先extractMax,再add,两次O(logn)的操作

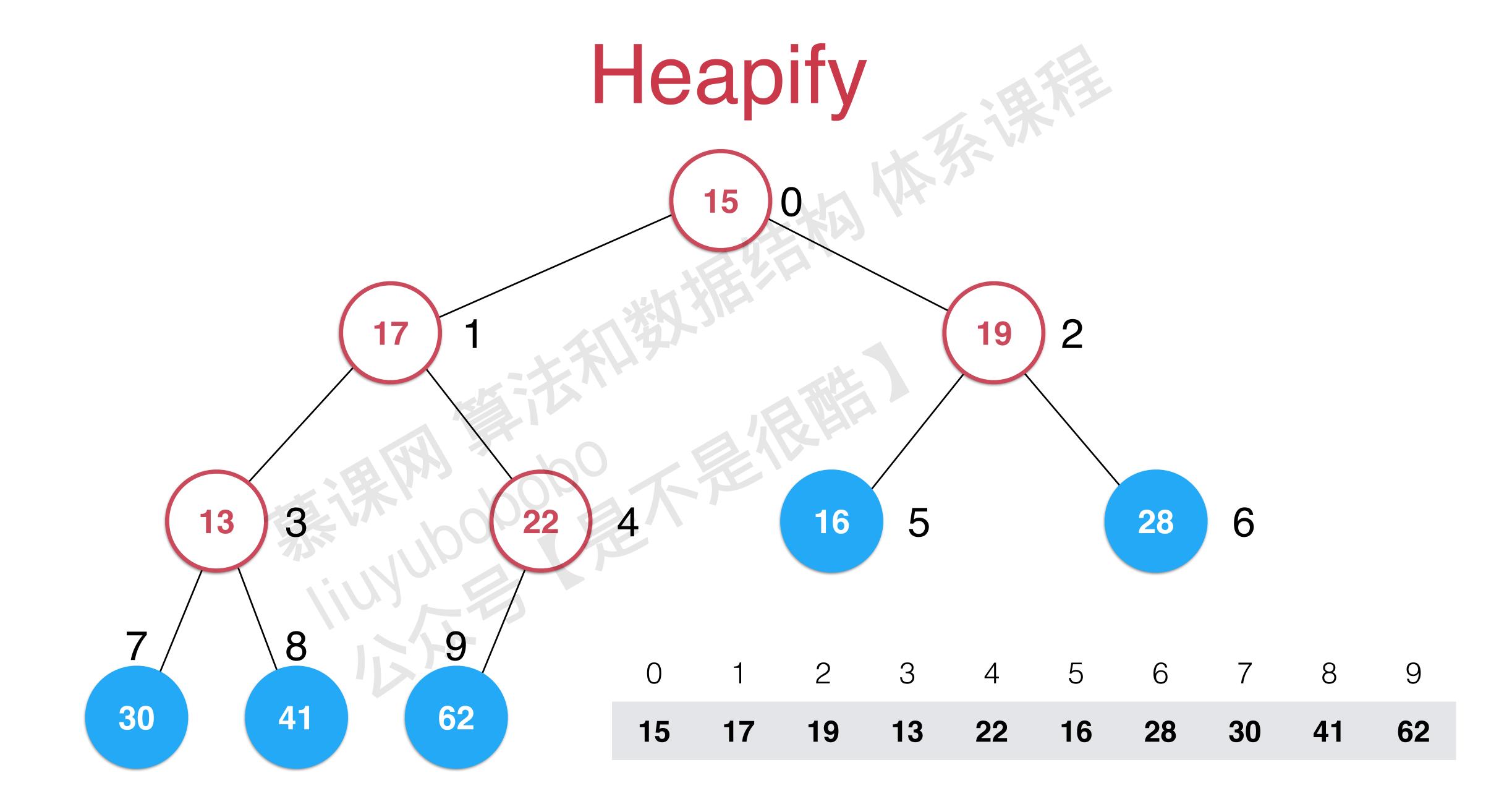
实现:可以直接将堆顶元素替换以后Sift Down,一次O(logn)的操作

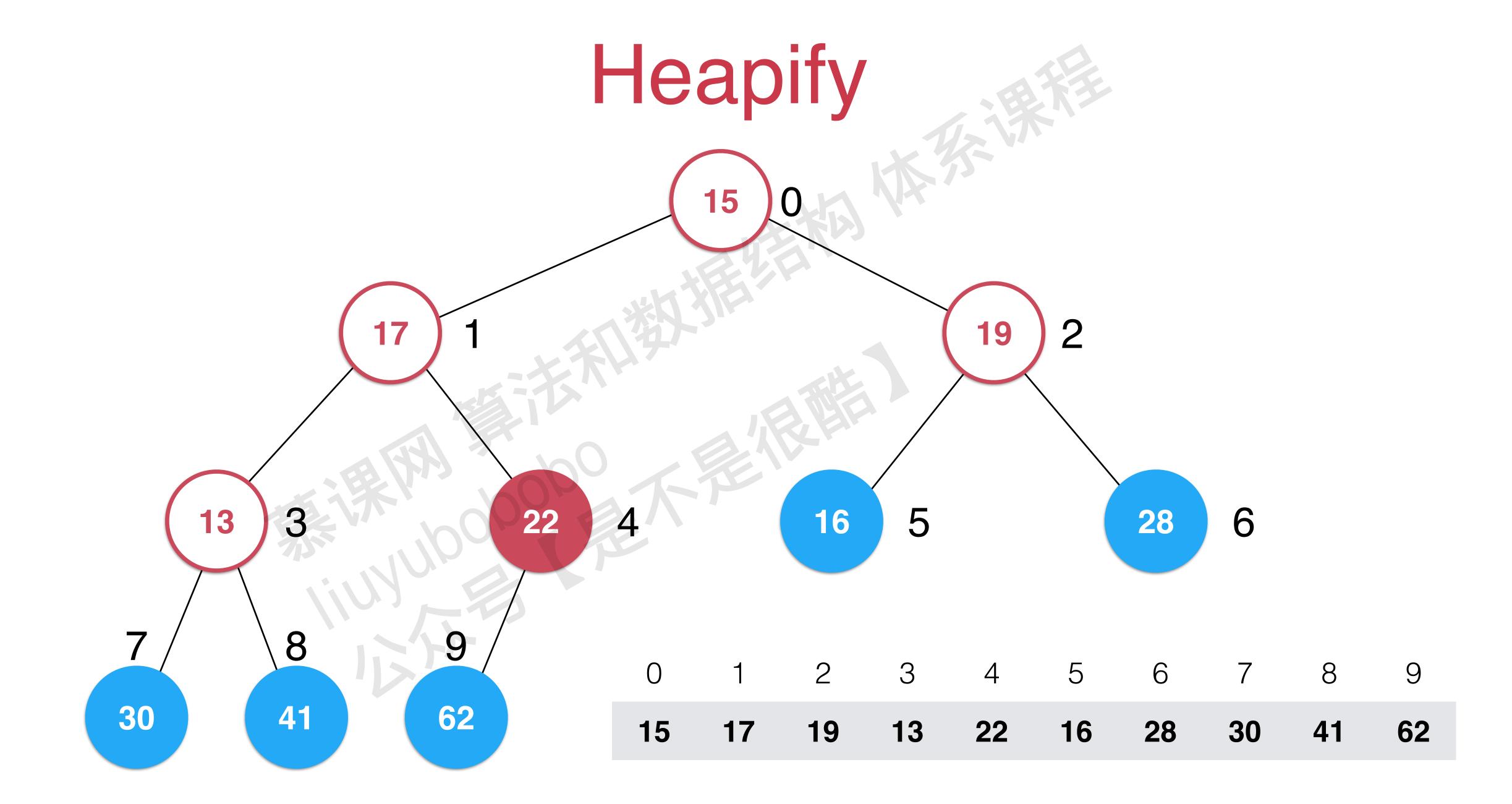
实践: replace

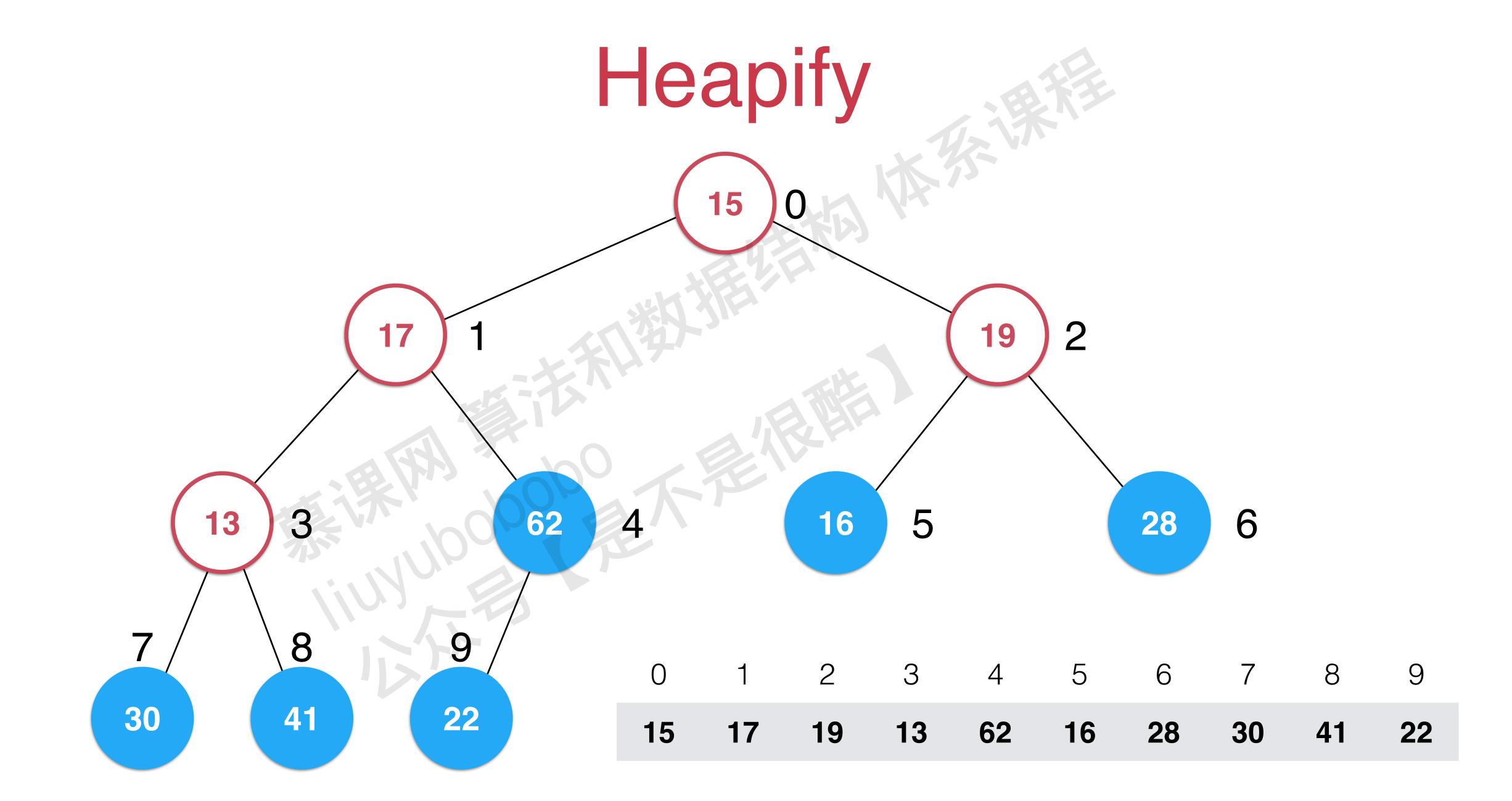
# heapify

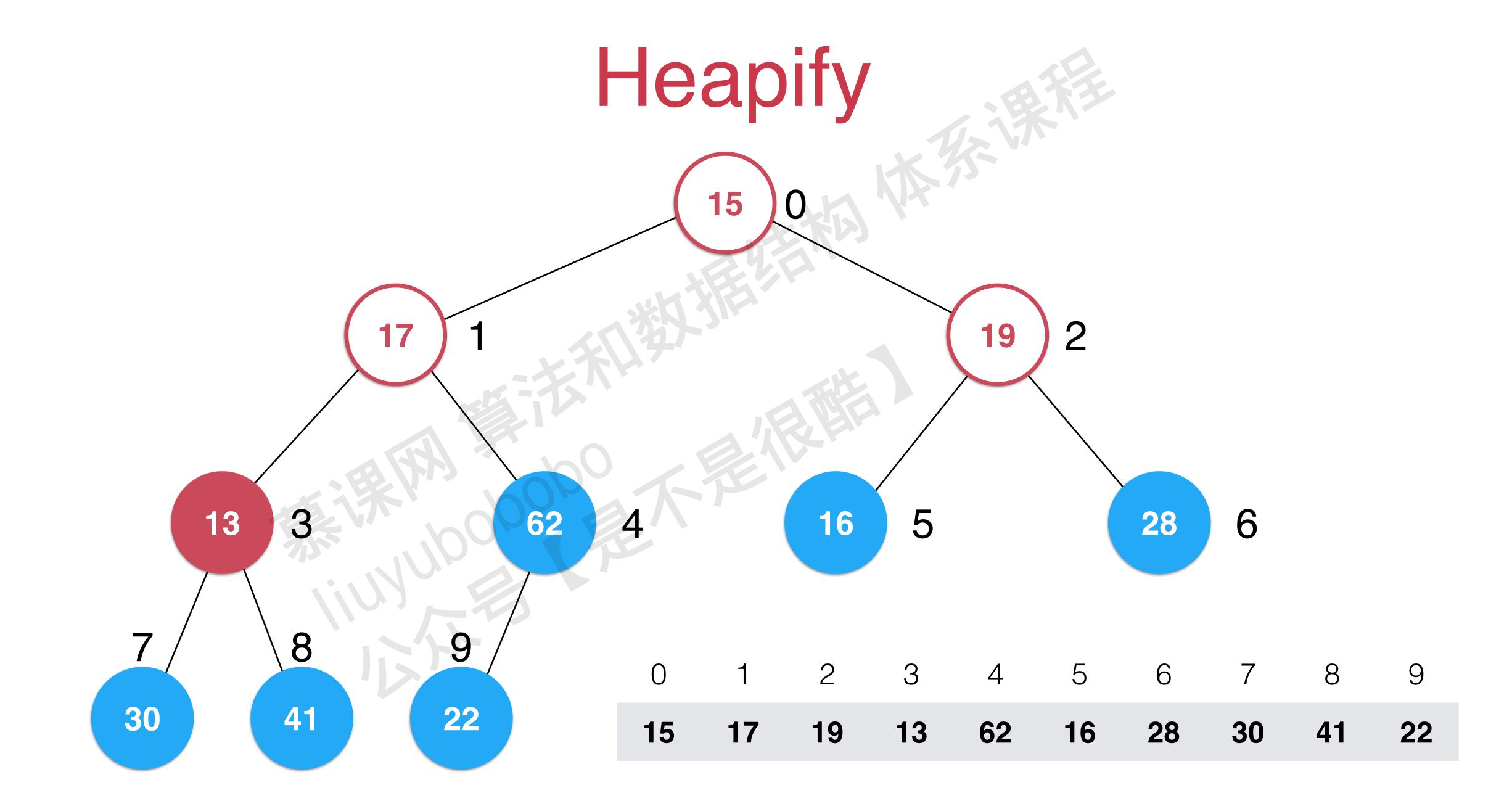
heapify: 将任意数组整理成堆的形状

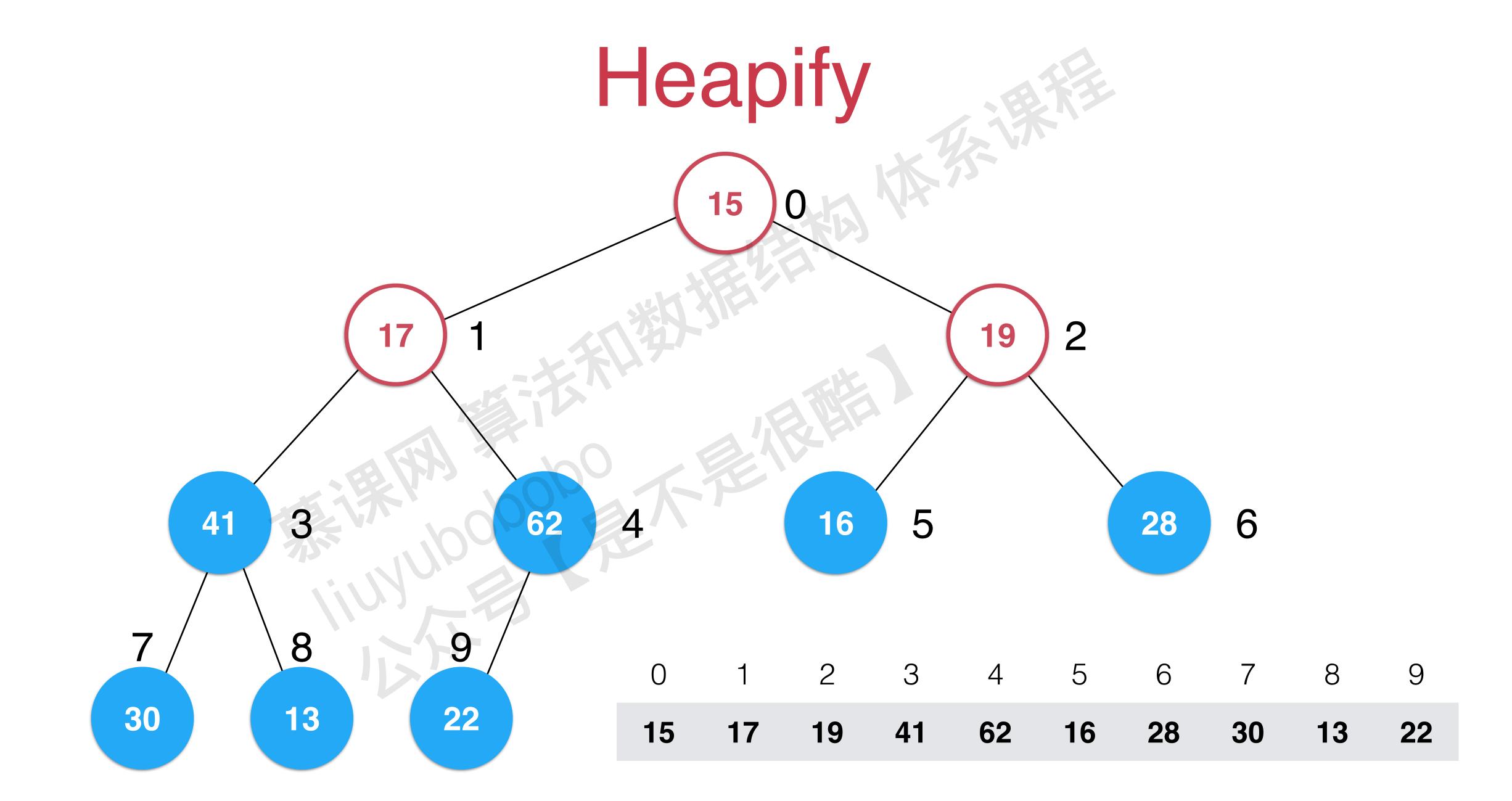


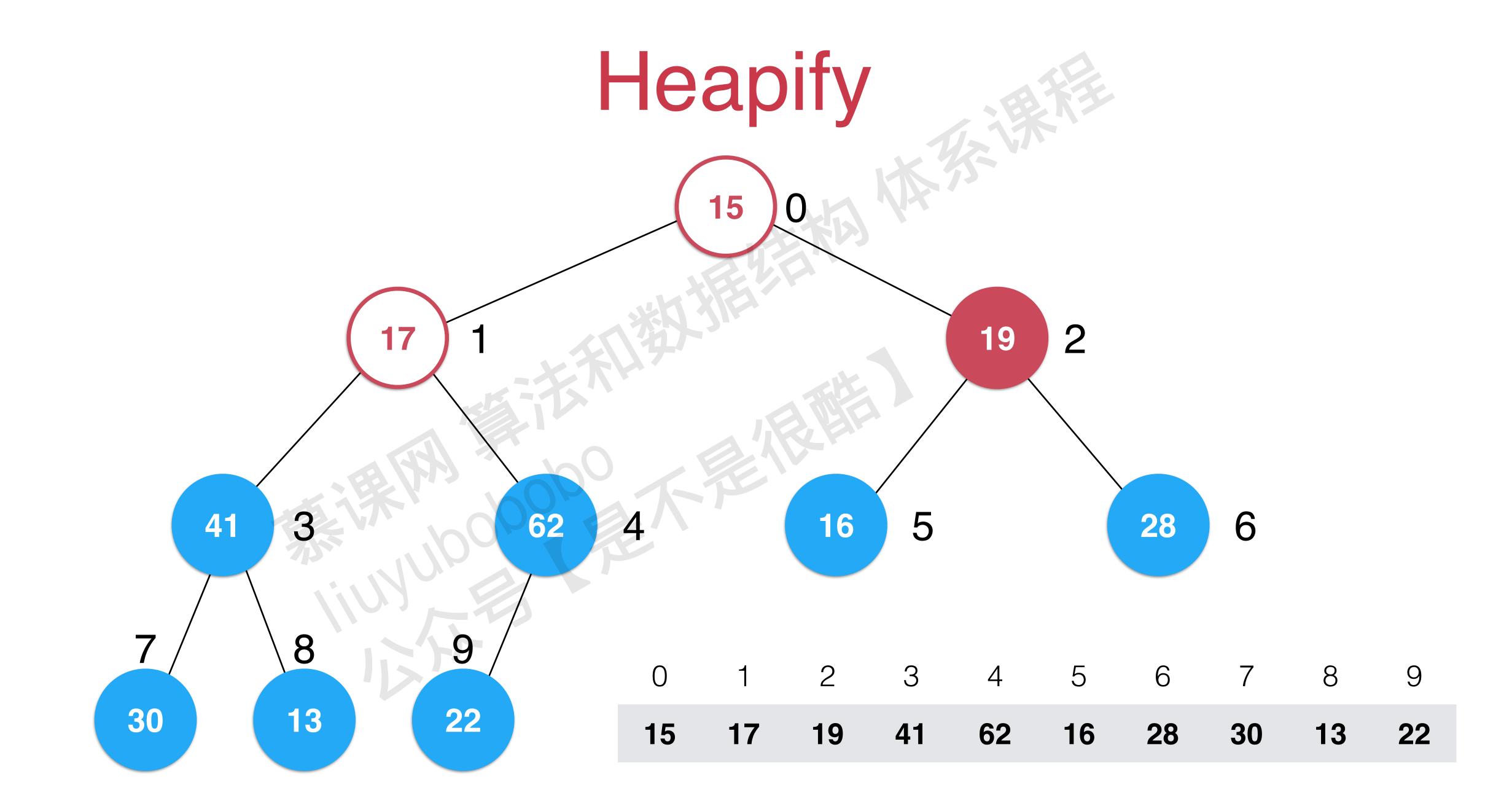


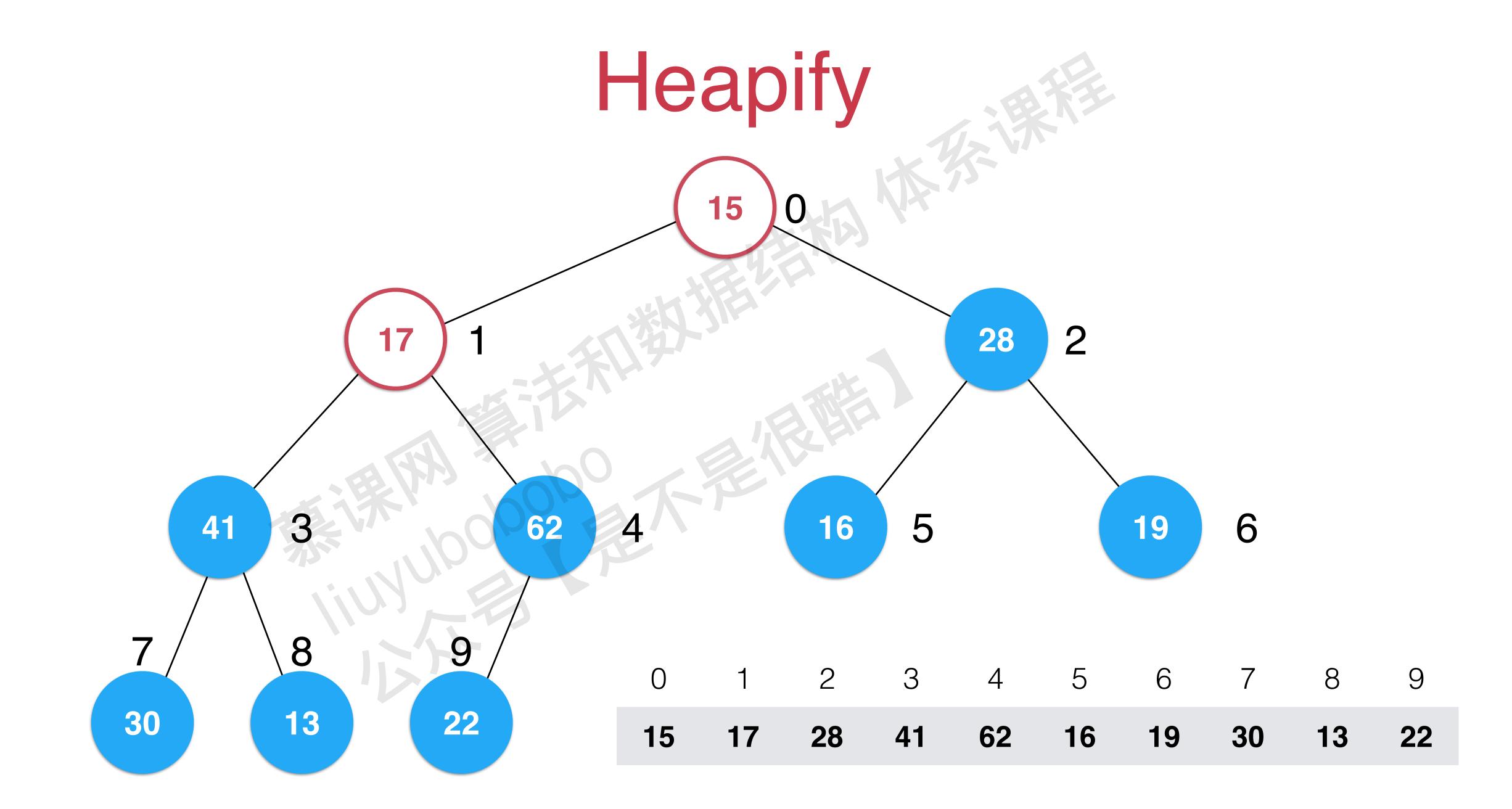


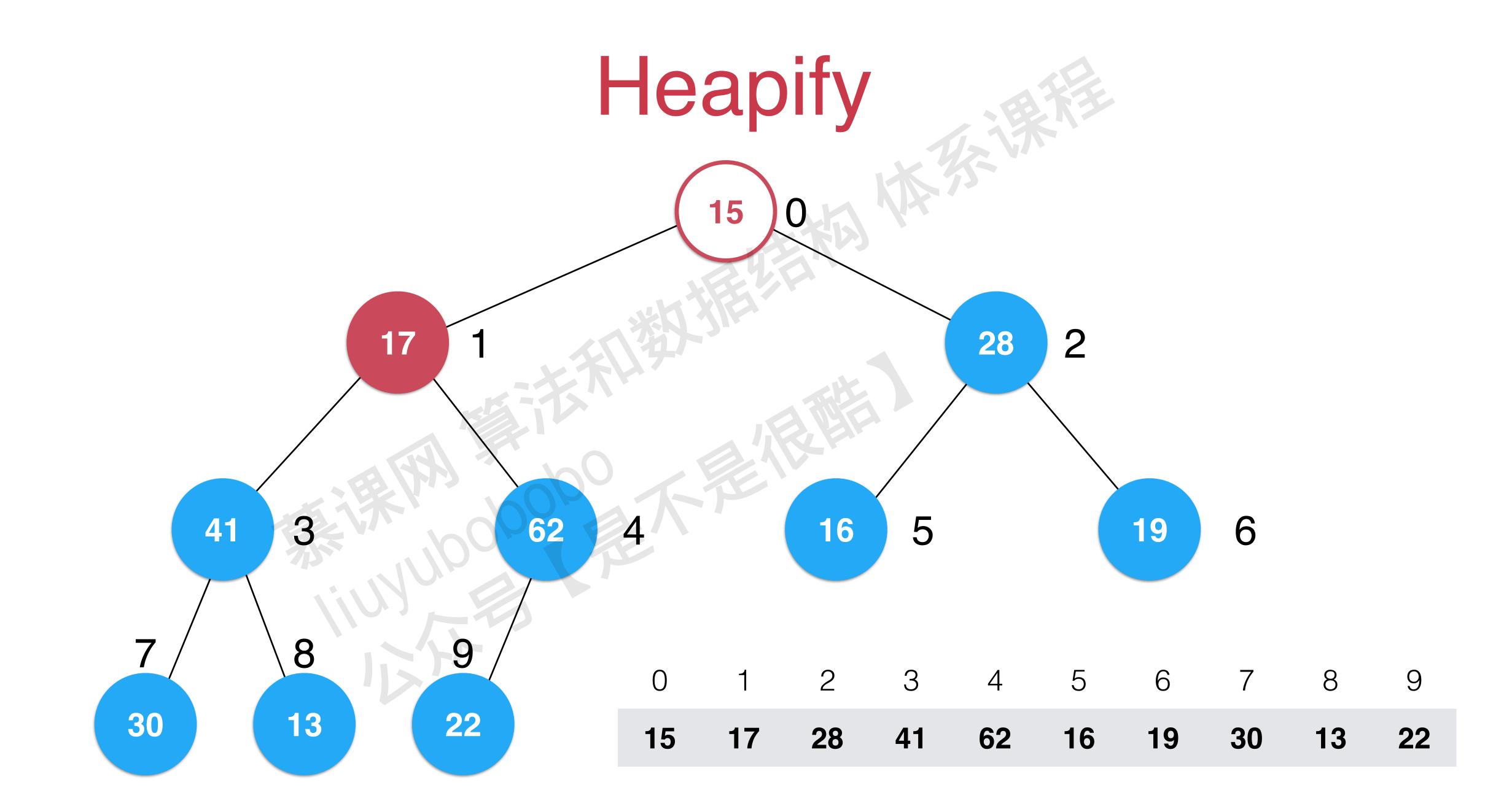


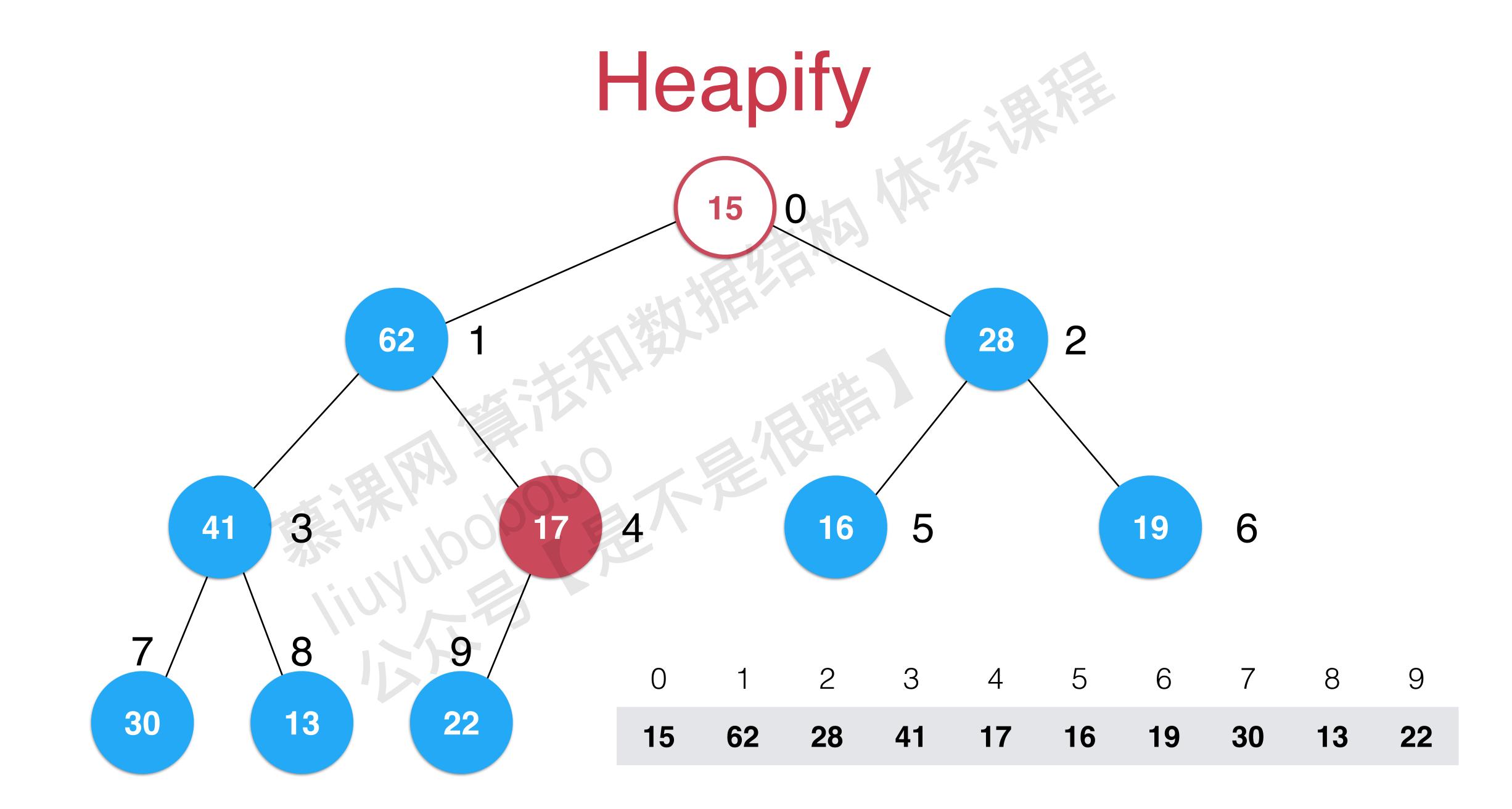


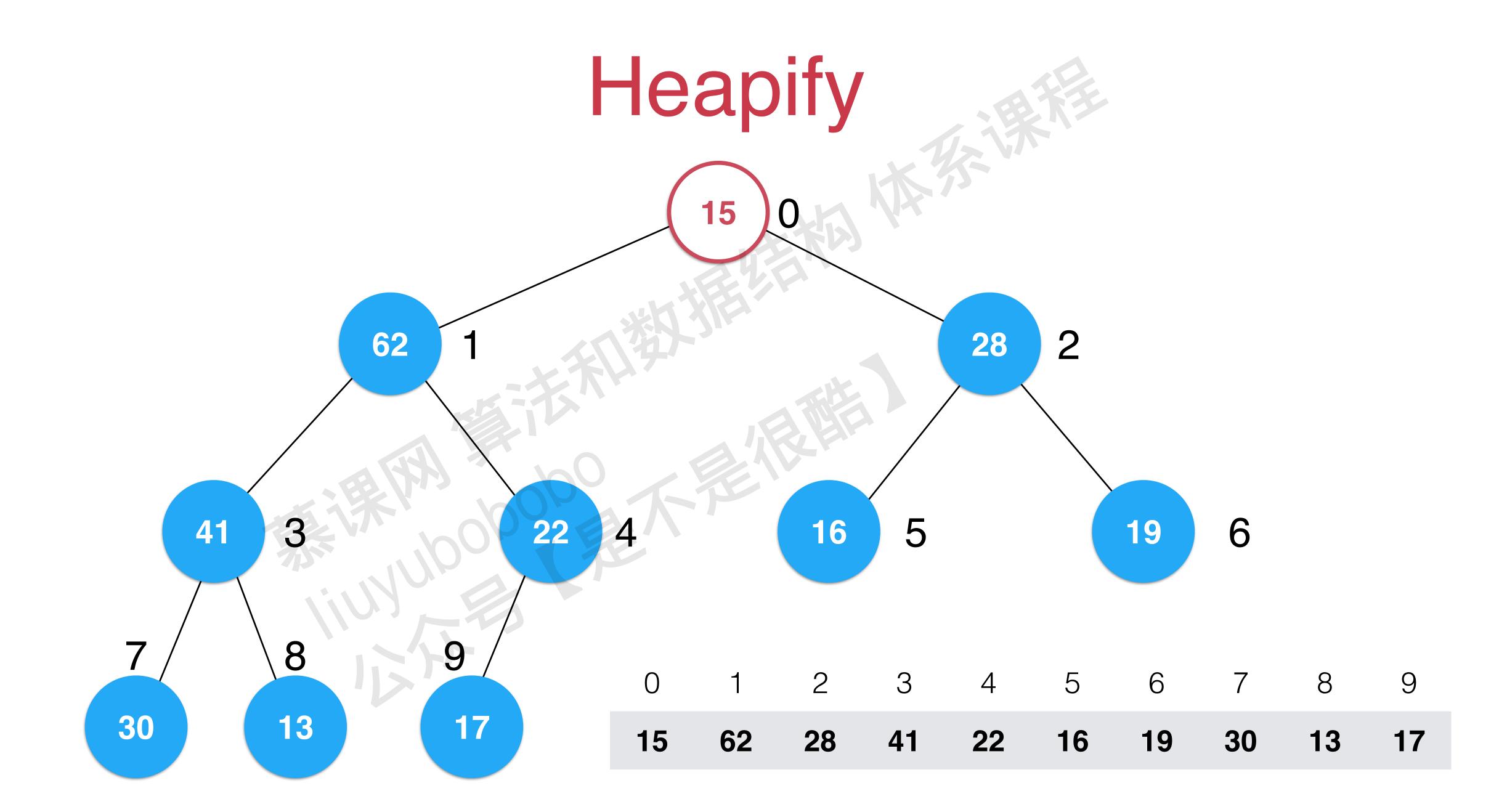


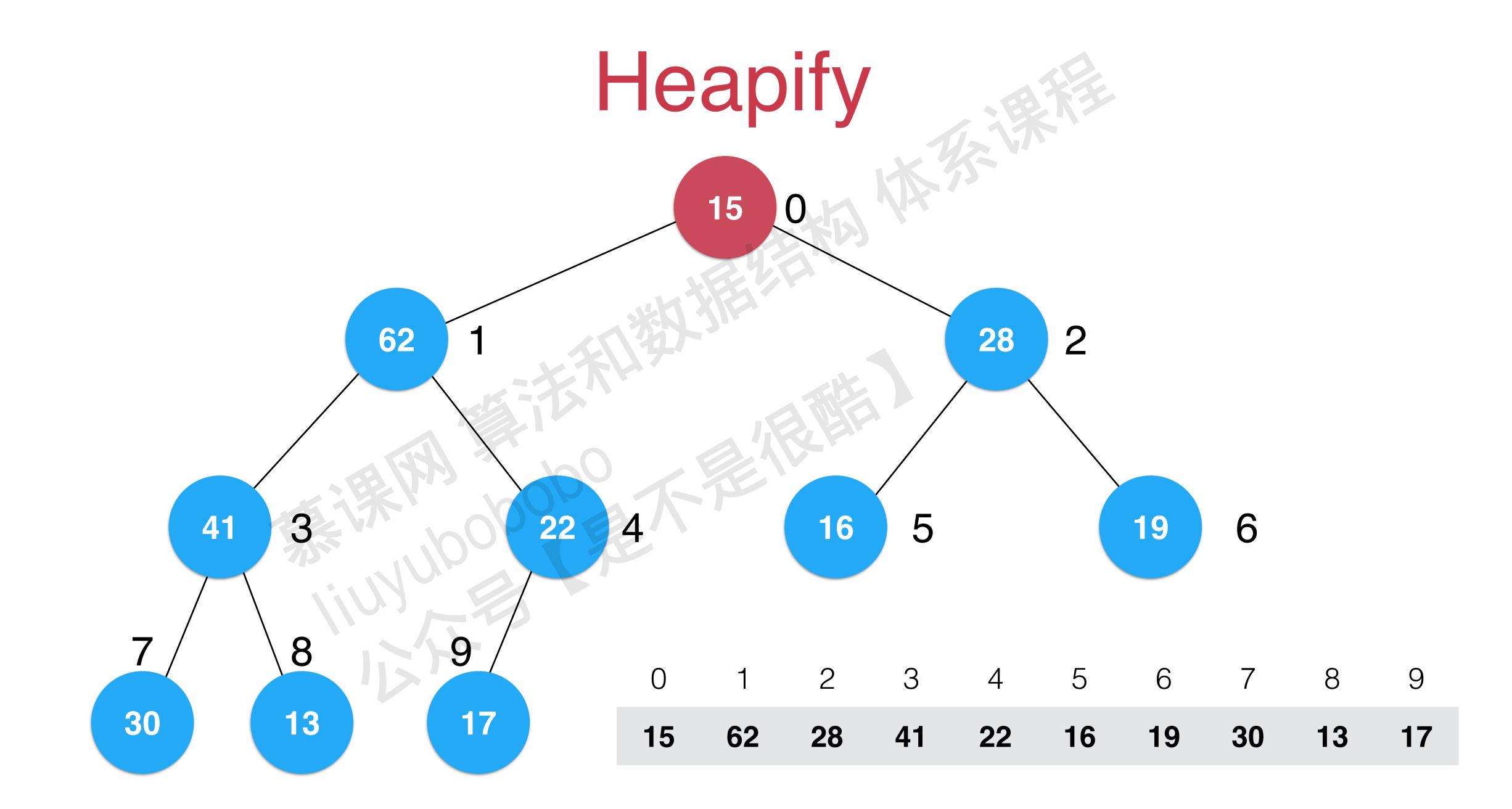


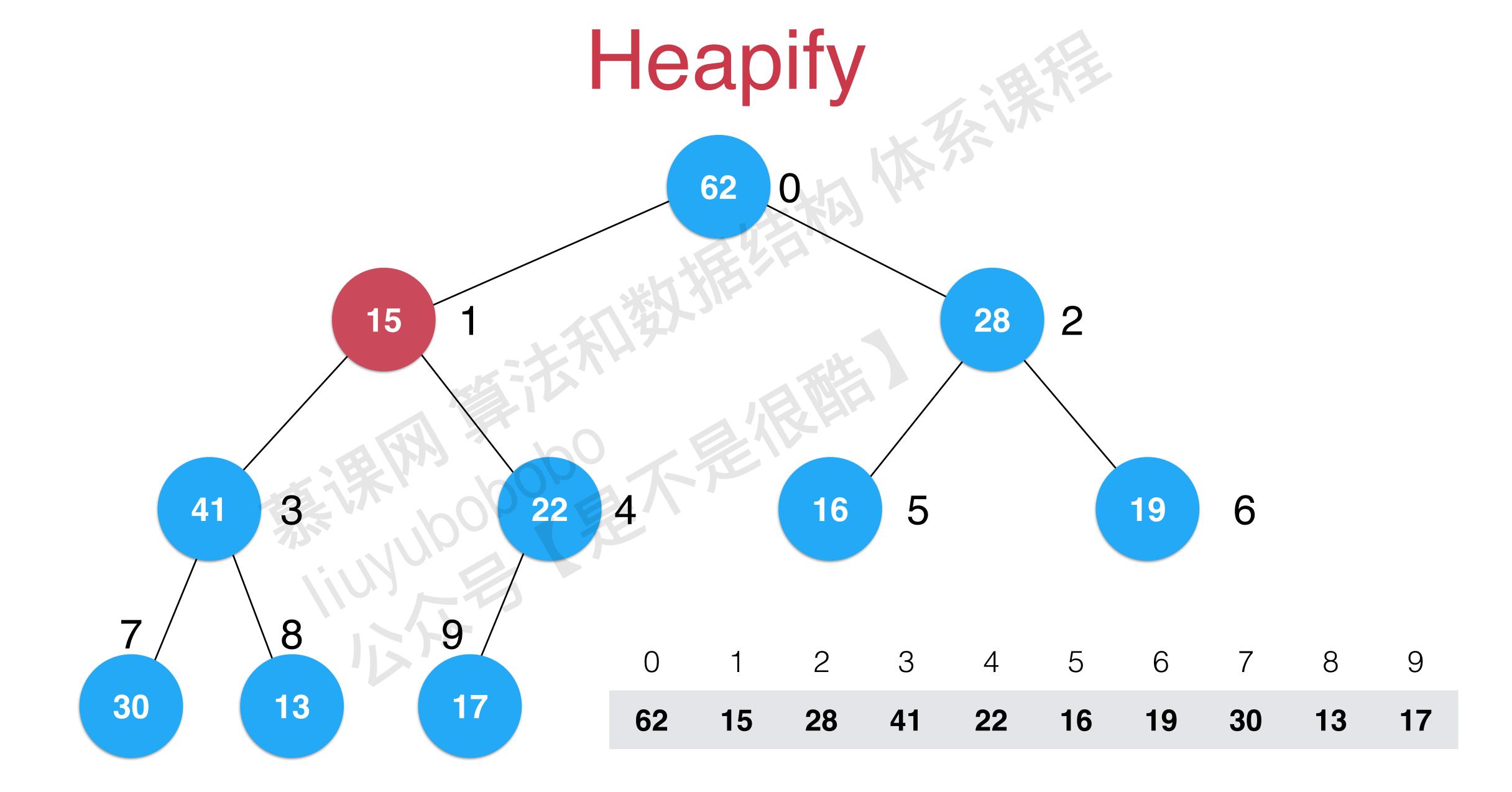


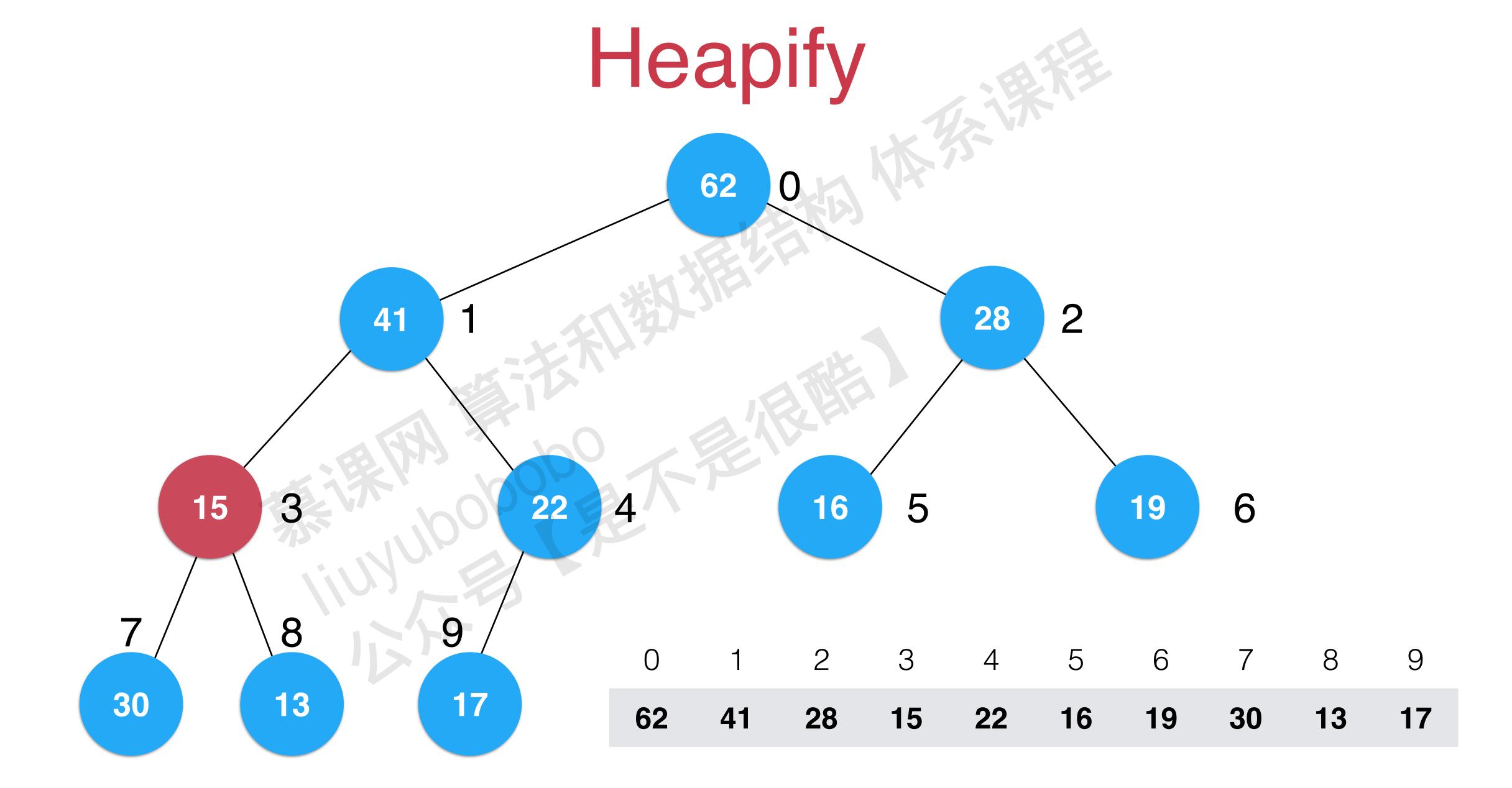


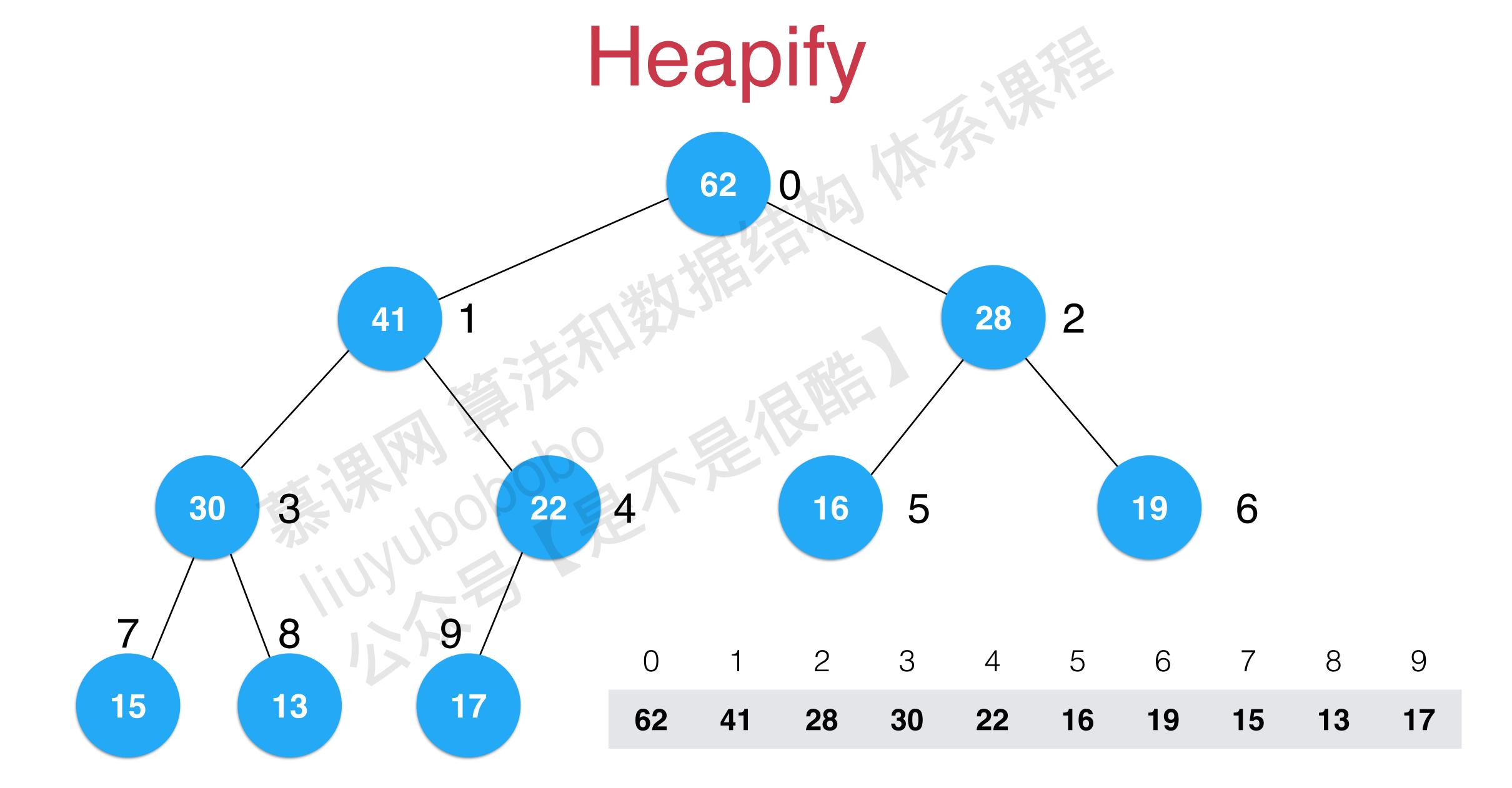










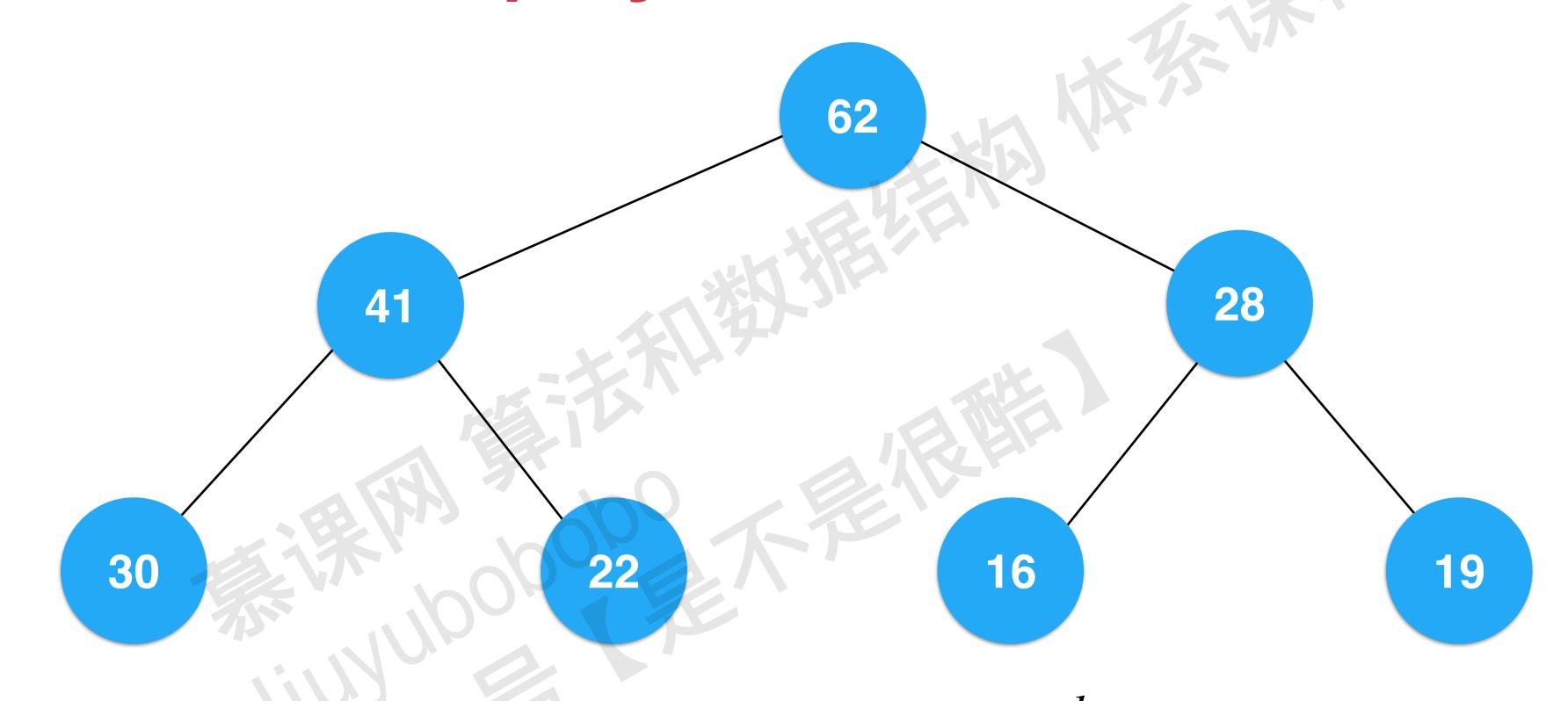


将n个元素逐个插入到一个空堆中,算法复杂度是O(nlogn)

heapify的过程,算法复杂度为O(n)

heapify的过程,算法复杂度为O(n)

最后一层最多有多少个节点?



1 + 2 + 4 + 8 + ... + 
$$2^{h-1} = \frac{1(1-2^h)}{1-2} = 2^h - 1$$

非叶子节点:  $2^{h-1}-1$  叶子节点:  $2^{h-1}$  大约是 n/2 个

heapify的过程,算法复杂度为O(n)

最后一层最多有多少个节点? n/2 n/2 \* 0

倒数第2层最多有多少个节点? n/4 n/4 \* 1

下面有 h 层节点  $\frac{h}{2h+1} = \frac{h}{2h+1} * h$ 

$$\sum_{h=0}^{\lfloor \lg n \rfloor} \left\lceil \frac{n}{2^{h+1}} \right\rceil O(h) = O\left(n \sum_{h=0}^{\lfloor \lg n \rfloor} \frac{h}{2^h}\right)$$

heapify的过程,算法复杂度为O(n)

$$\sum_{h=0}^{\lfloor \lg n \rfloor} \left\lceil \frac{n}{2^{h+1}} \right\rceil O(h) = O\left(n \sum_{h=0}^{\lfloor \lg n \rfloor} \frac{h}{2^h}\right)$$

$$= O\left(n\sum_{h=0}^{\infty} \frac{h}{2^h}\right)$$
$$= O(n).$$

$$\sum_{k=0}^{\infty} kx^k = \frac{x}{(1-x)^2}$$
for  $|x| < 1$ .

$$\sum_{h=0}^{\infty} \frac{h}{2^h} = \frac{1/2}{(1-1/2)^2} = 2.$$

heapify的过程,算法复杂度为O(n)



### 实践: 实现 Heapify

## 优化的堆排序

max



Max Heap

#### 原地维持序

max



Max Heap

W

max



W

Max Heap

V



max



Max Heap



max

Max Heap

W



W

Max Heap

# 作业:实现最小堆

### 作业解析: 实现最小堆

#### 其他

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