算法与数据结构体系课程

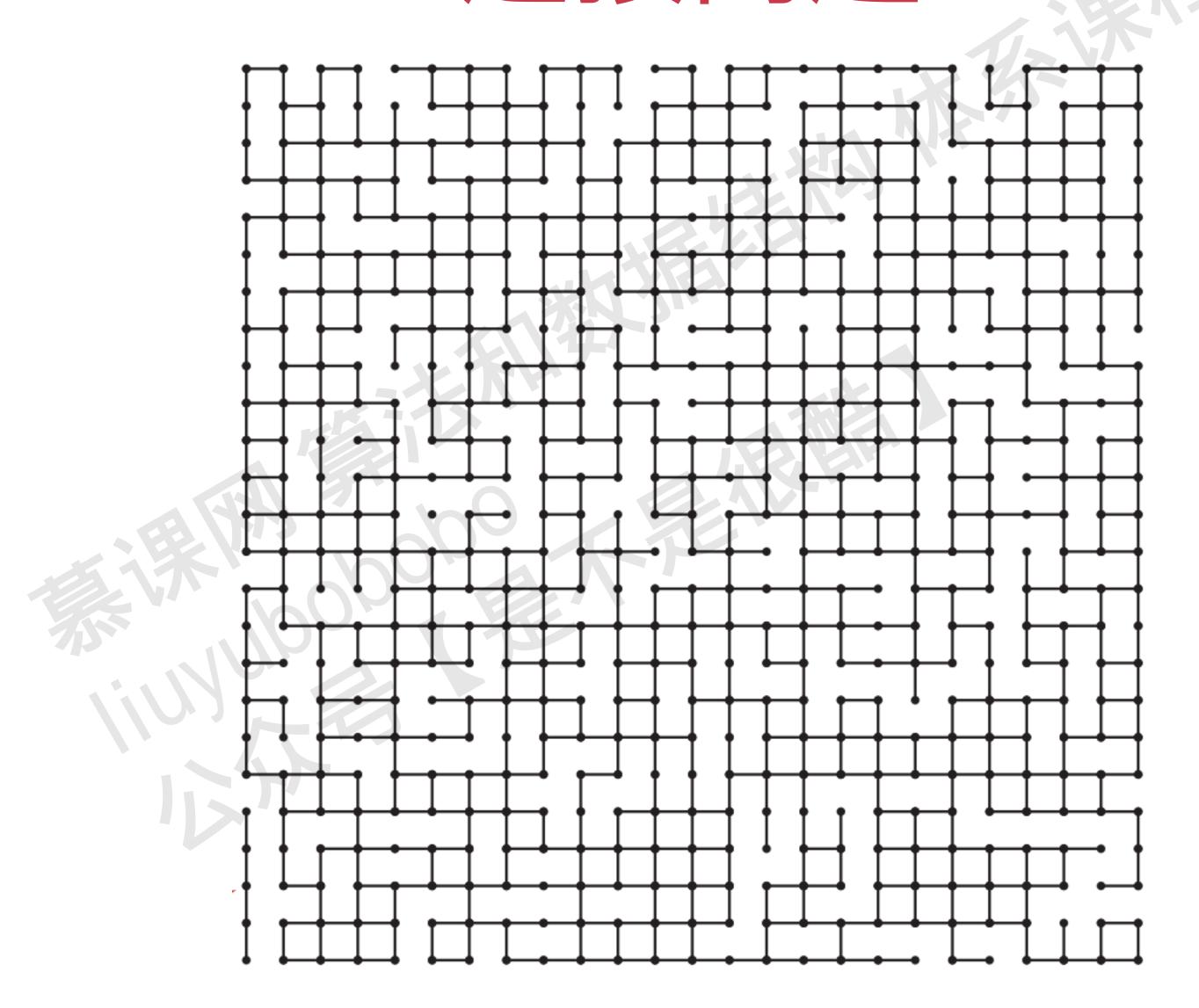
liuyubobobo

并查集 Union Find

一种很不一样的树形结构

连接问题 Connectivity Problem

连接问题



连接问题

网络中节点间的连接状态

• 网络是个抽象的概念: 用户之间形成的网络

数学中的集合类实现

连接问题和路径问题

比路径问题要回答的问题少

• 和堆作比较

并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q)

实践: 并查集的接口设计



并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q)

并查集的基本数据表示

0 1 2 3 4 5 6 7 8 9

0 0 0 0 1 1 1 1 1

并查集的基本数据表示

id 0 1 2 3 4 5 6 7 8 9 0 1 0 1 0 1 0 1

并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)
- isConnected(p,q)

并查集 Union Find

对于一组数据,主要支持两个动作:

- union(p,q)isConnected(p,q) find(p) == find(q)

Quick Find

id 0 1 2 3 4 5 6 7 8 9 0 1 0 1 0 1 0 1

Quick Find 时间复杂度 O(1)

Quick Find 下的 Union

union(1, 4)

id 0 1 2 3 4 5 6 7 8 9

0 1 0 1 0 1 0 1

Quick Find 下的 Union

union(1, 4)

id 0 1 2 3 4 5 6 7 8 9

1 1 1 1 1 1 1

Quick Find 下的 Union

id 0 1 2 3 4 5 6 7 8 9 1 1 1 1 1 1 1 1 1

Quick Find 下的 Union 时间复杂度 O(n)

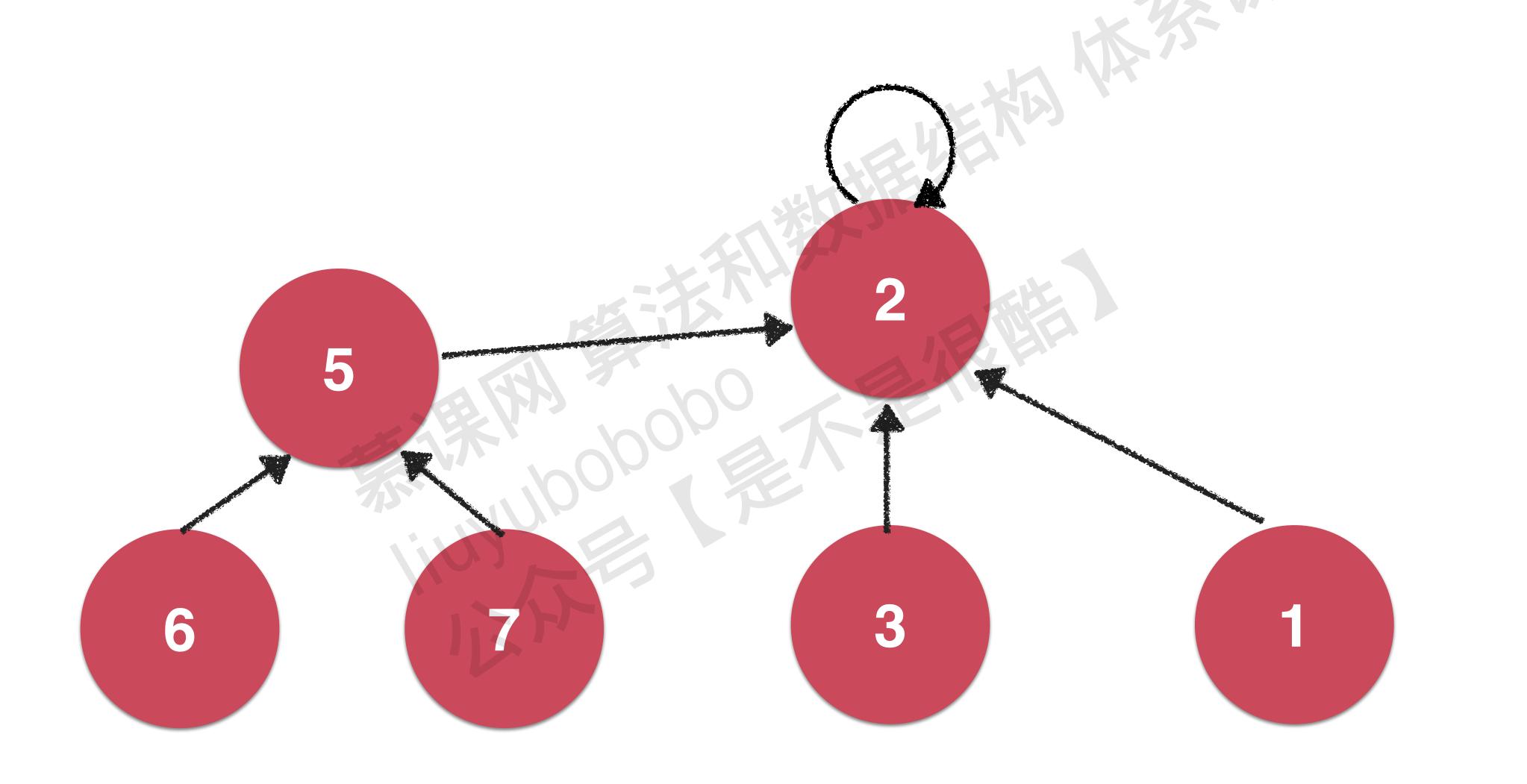
实践: Quick Find

Quick Find

- unionElements(p,q)O(n
- isConnected(p,q) O(1



将每一个元素,看做是一个节点



Quick Union 下的数据表示

parent 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9

Quick Union



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 3 4 5 6 7 8 9

union 4, 3



parent 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9

union 4, 3



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 3 4 5 6 7 8 9

union 4,3



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 3 3 5 6 7 8 9

union 3,8



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 3 3 5 6 7 8 9

union 3,8

parent 0 1 2 3 4 5 6 7 8 9

0 1 2 3 3 5 6 7 8 9

union 3,8

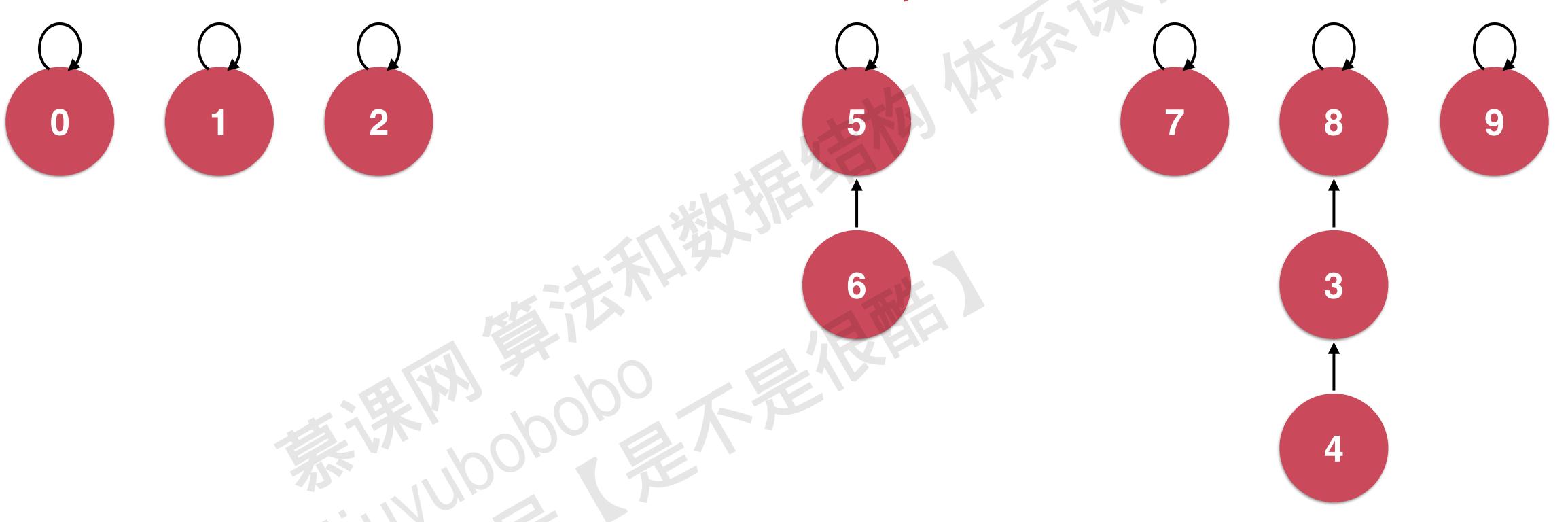
parent 0 1 2 3 4 5 6 7 8 9

0 1 2 8 3 5 6 7 8 9

union 6,5

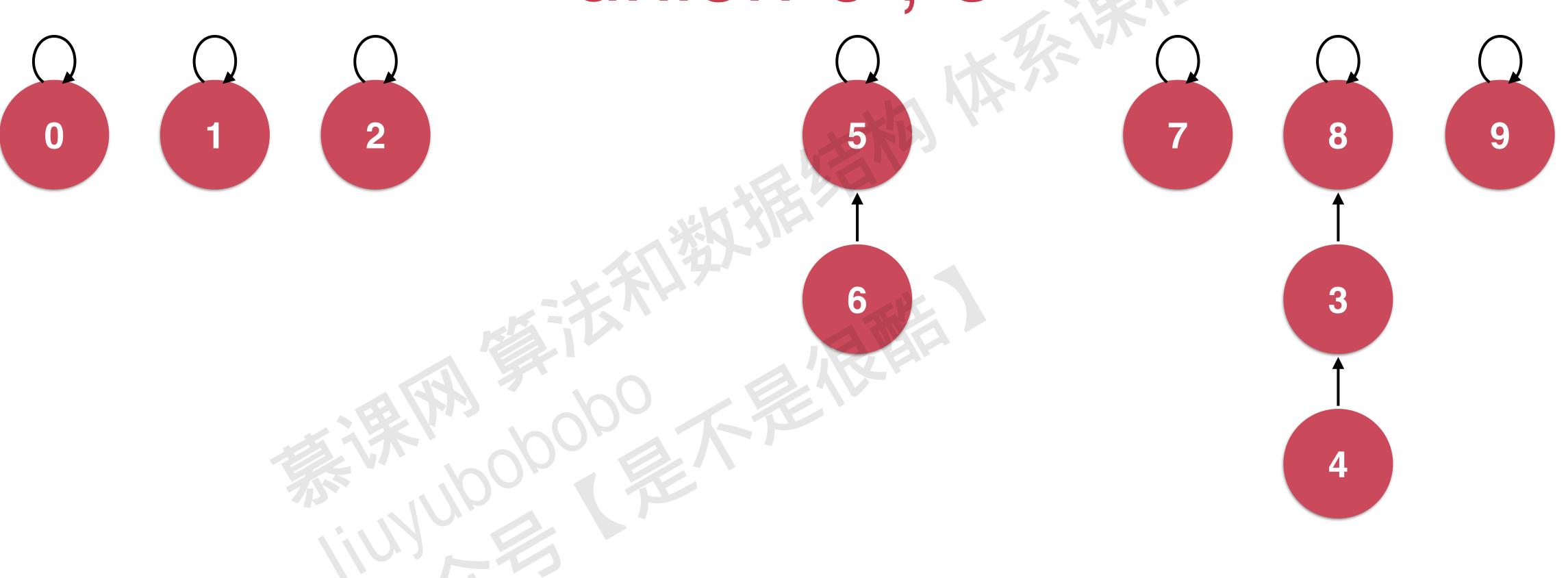
parent 0 1 2 3 4 5 6 7 8 9
0 1 2 8 3 5 6 7 8 9

union 6,5



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 8 3 5 6 7 8 9

union 6,5



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 8 3 5 5 7 8 9

union 9

parent 0 1 2 3 4 5 6 7 8 9
0 1 2 8 3 5 5 7 8 9



parent 0 1 2 3 4 5 6 7 8 9

0 1 2 8 3 5 5 7 8 9



0 1 2 8 3 5 5 7 8 8

union 2,

parent 0 1 2 3 4 5 6 7 8 9

0 1 2 8 3 5 5 7 8 8

union 2,

parent 0 1 2 3 4 5 6 7 8 9

0 1 2 8 3 5 5 7 8 8

union 2,

parent 0 1 2 3 4 5 6 7 8 9

0 1 1 8 3 5 5 7 8 8

union 5, 0

parent 0 1 2 3 4 5 6 7 8 9
0 1 1 8 3 5 5 7 8 8

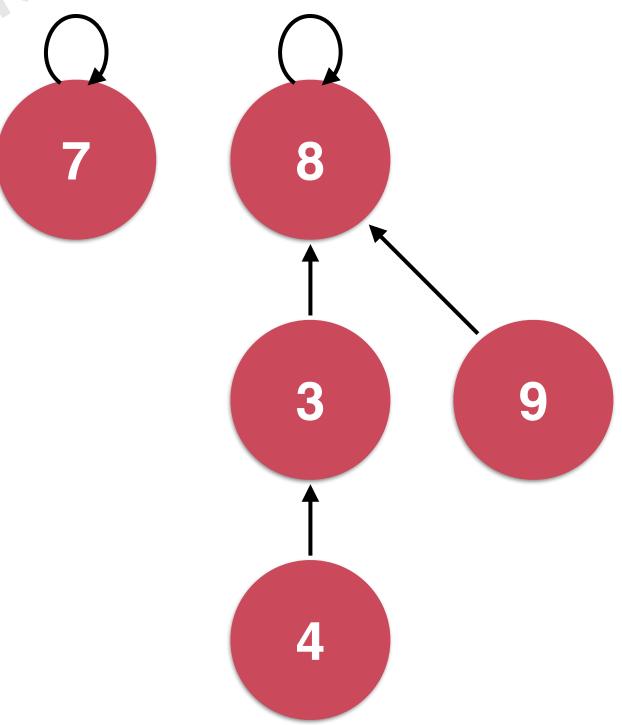
union 5,0



parent 0 1 2 3 4 5 6 7 8 9
0 1 1 8 3 5 5 7 8 8

union 5, 0





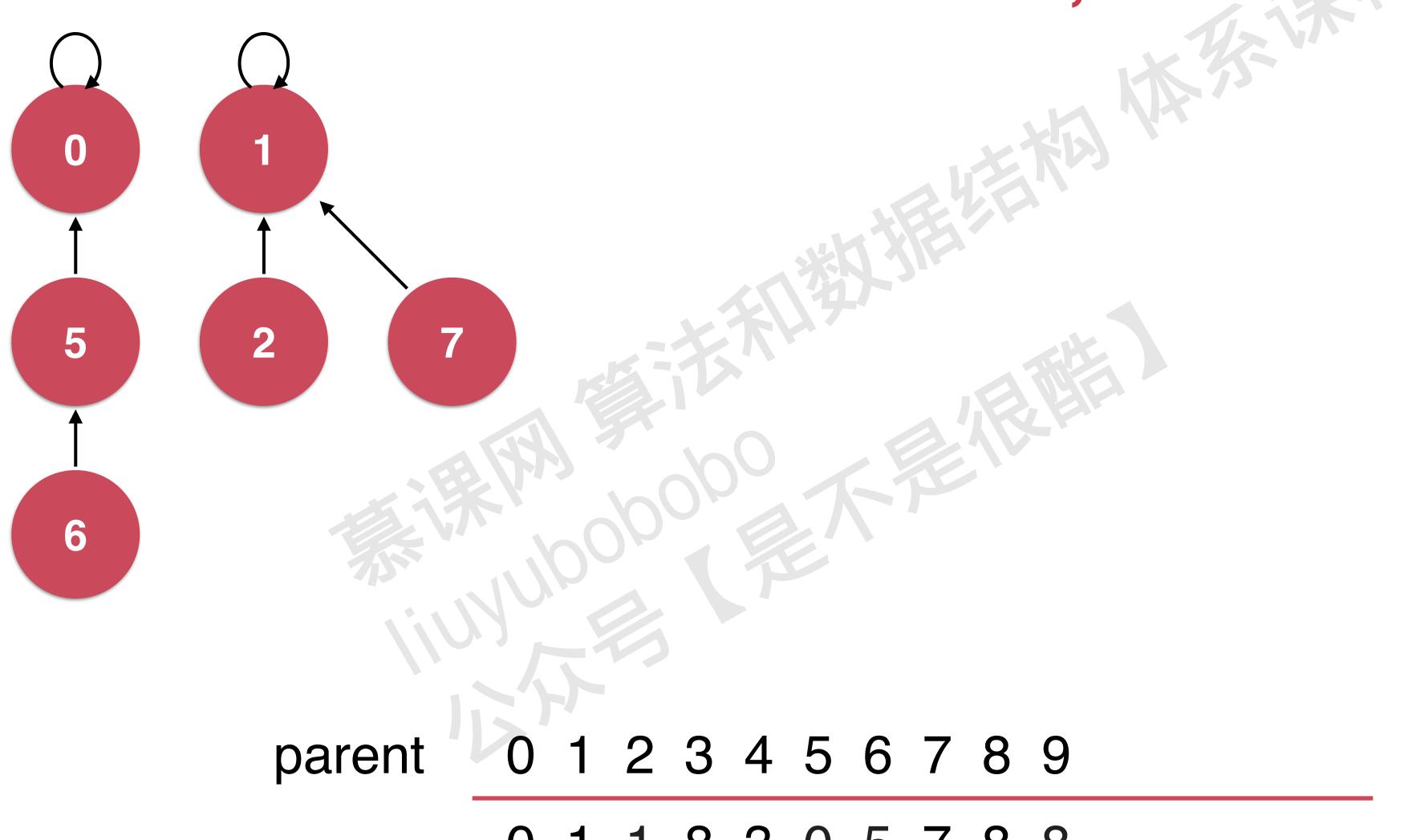
0 1 1 8 3 0 5 7 8 8

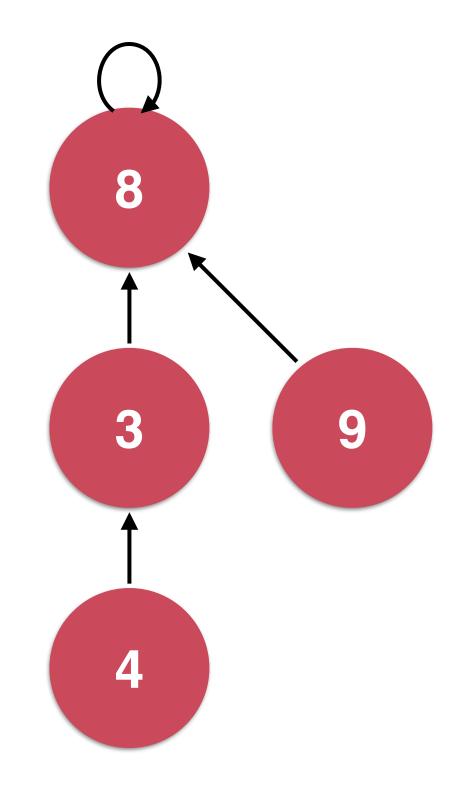
union 7, 2



parent 0 1 2 3 4 5 6 7 8 9
0 1 1 8 3 0 5 7 8 8

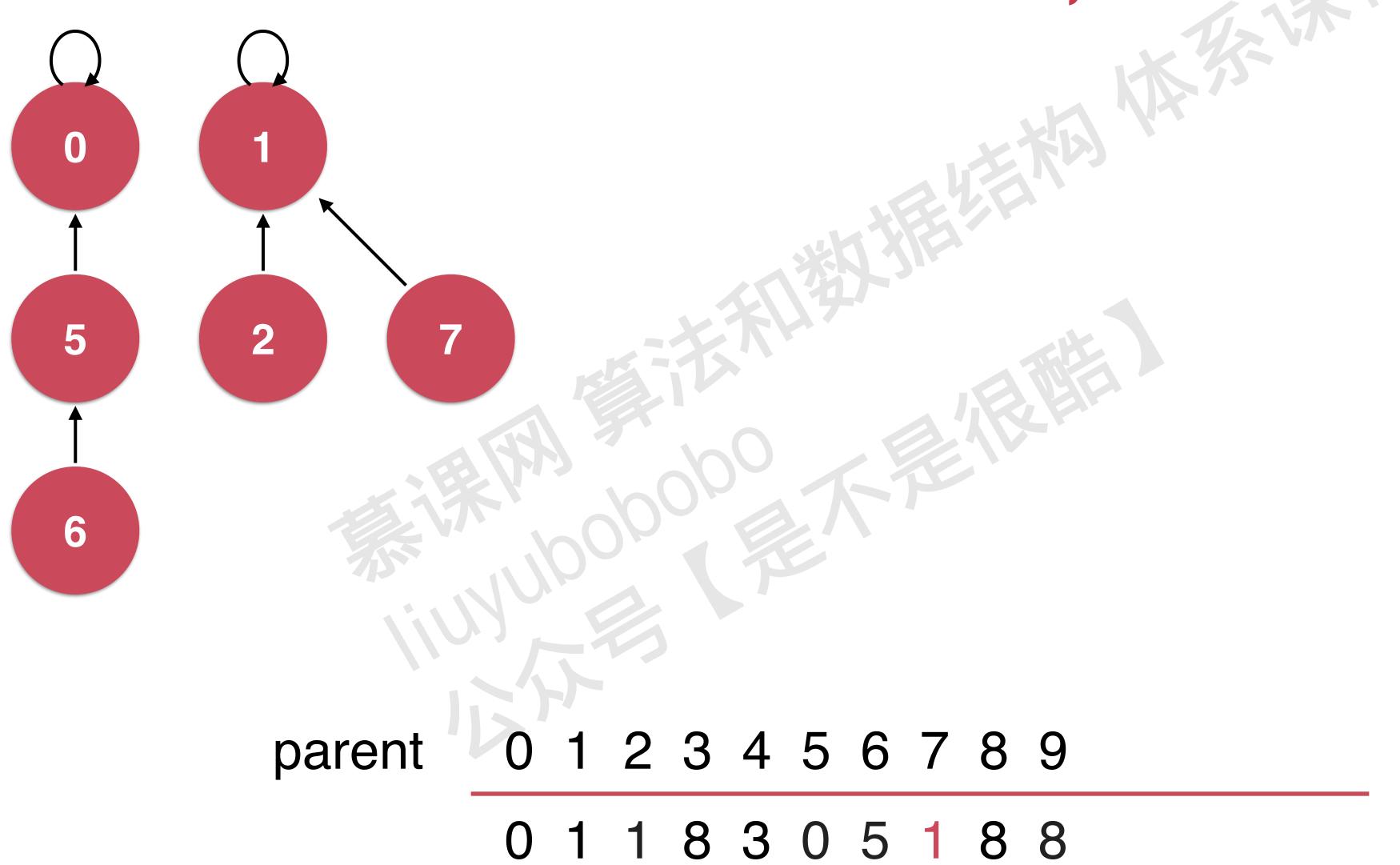
union 7, 2

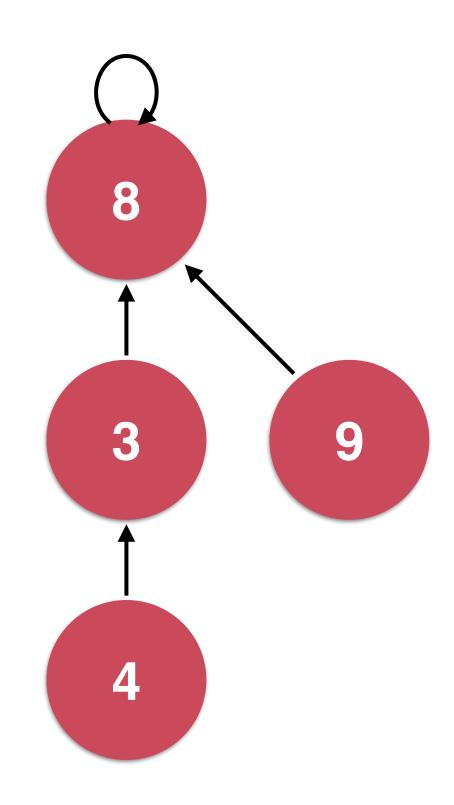




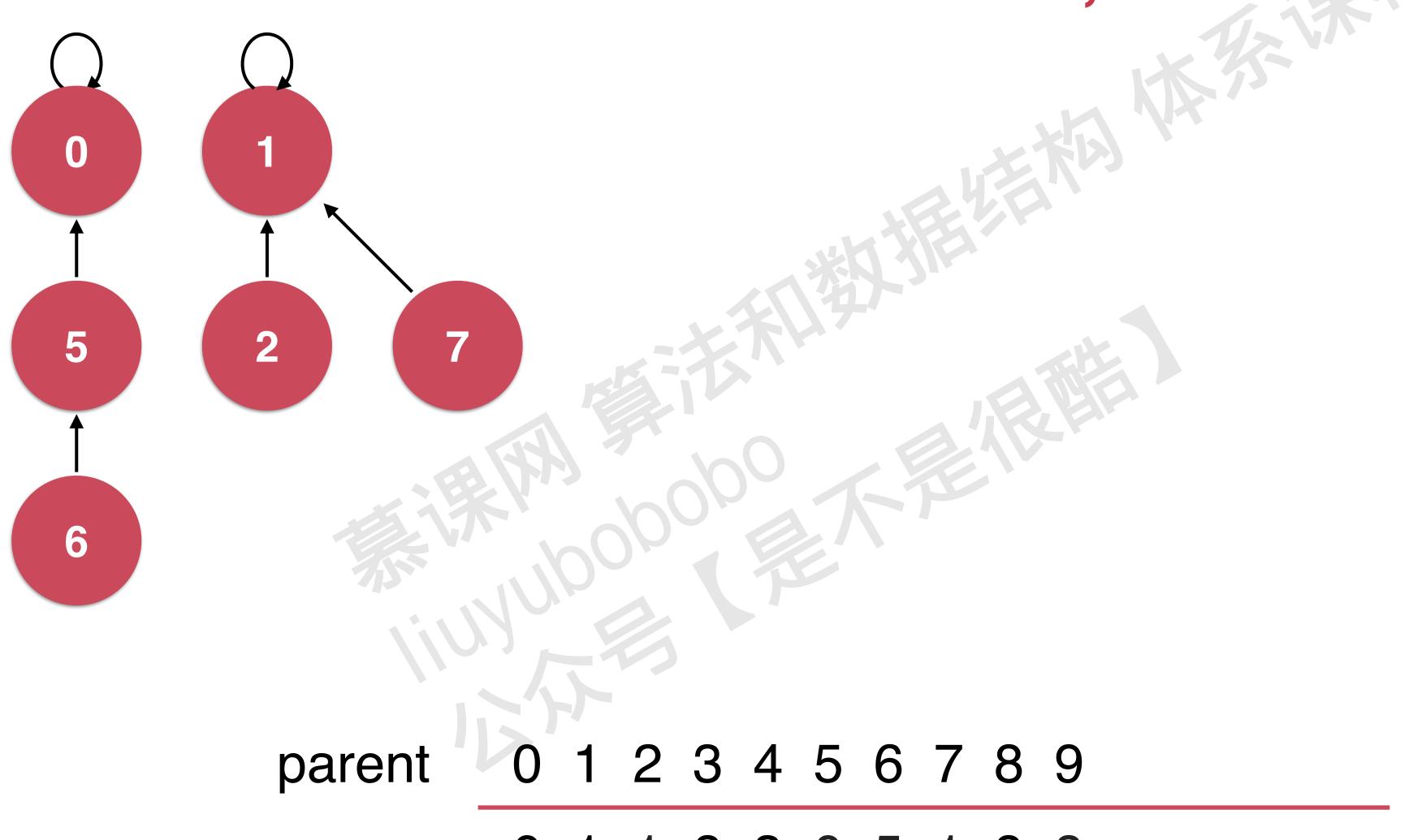
0 1 1 8 3 0 5 7 8 8

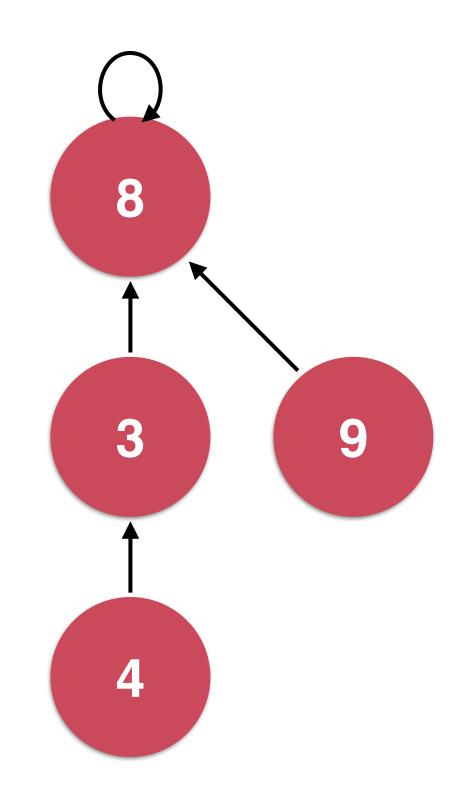
union 7, 2





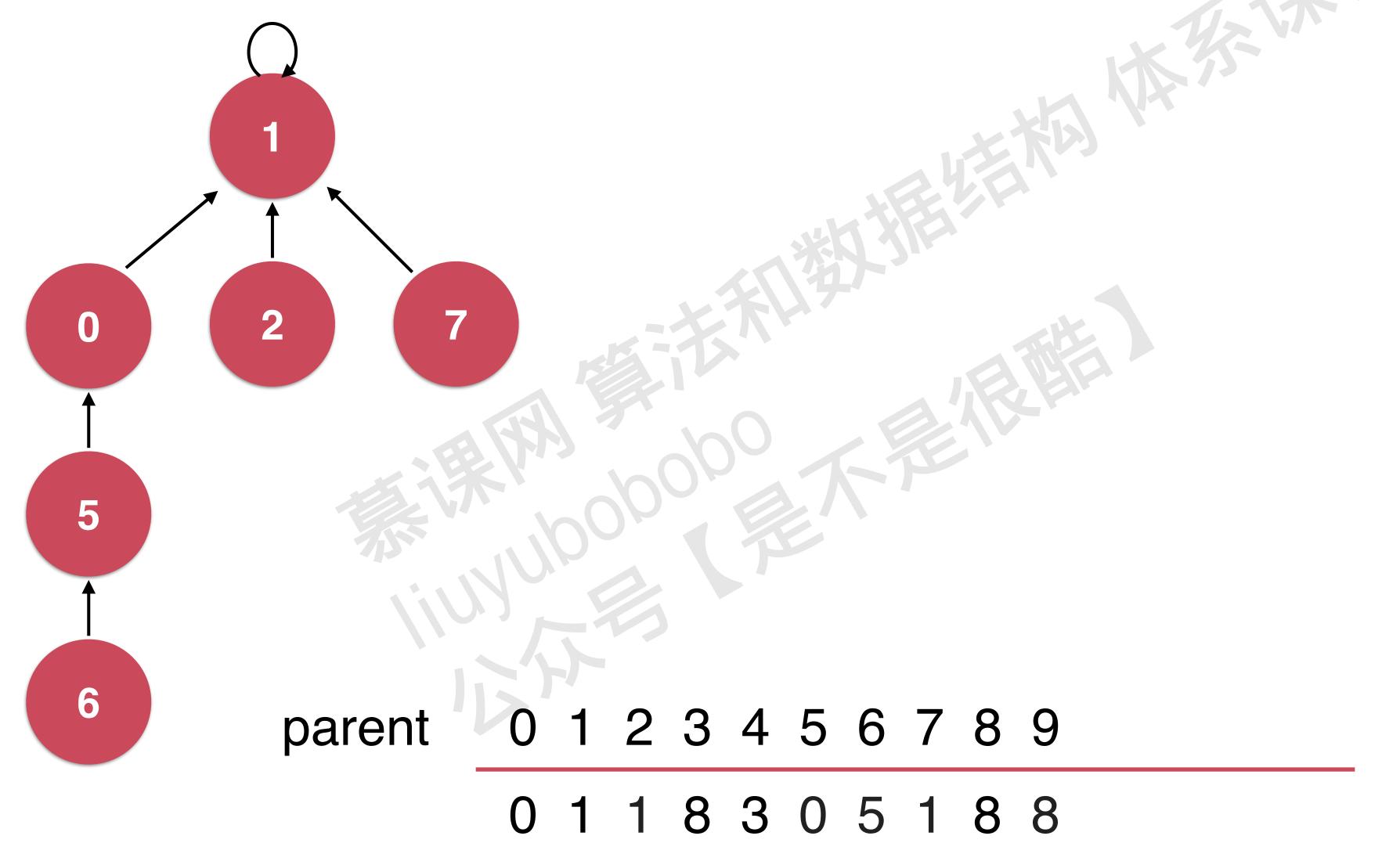
union 6, 2

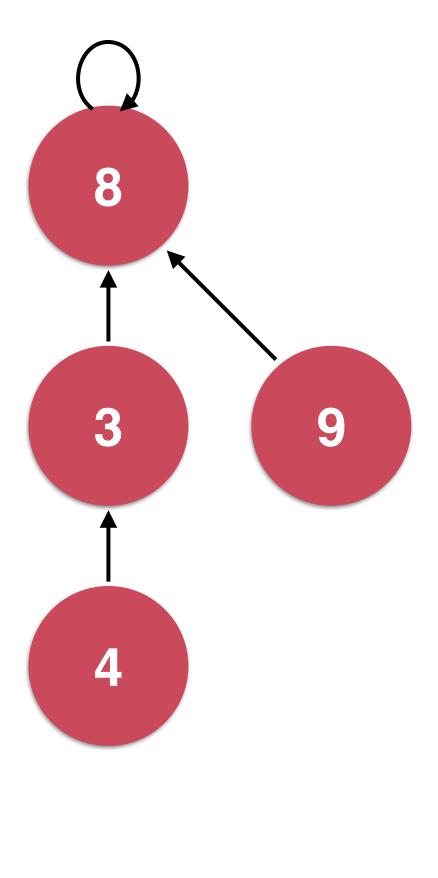




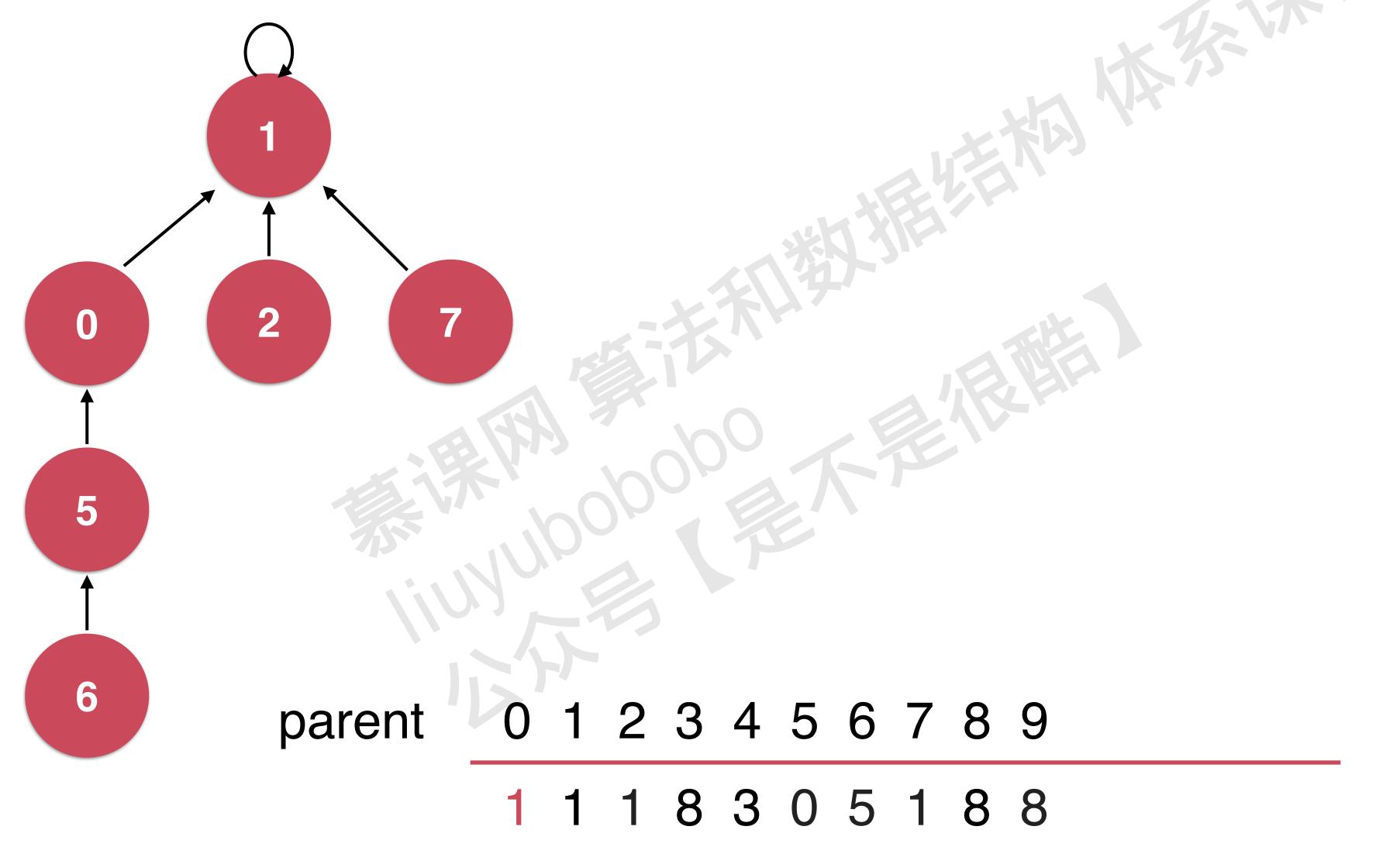
0 1 1 8 3 0 5 1 8 8

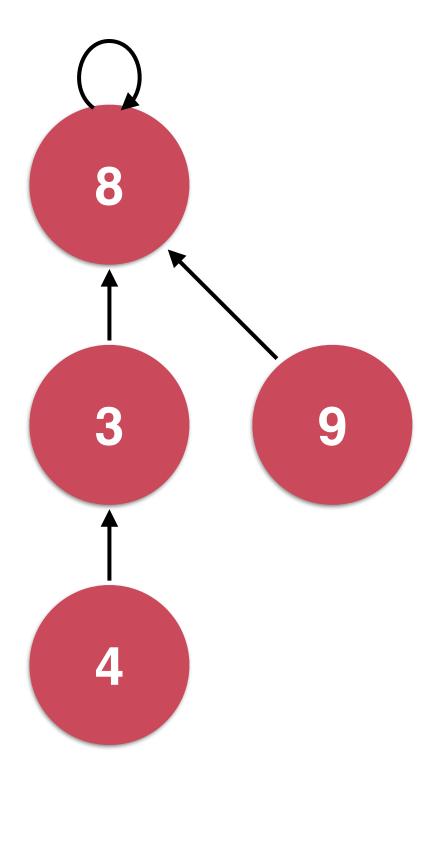
union 6, 2

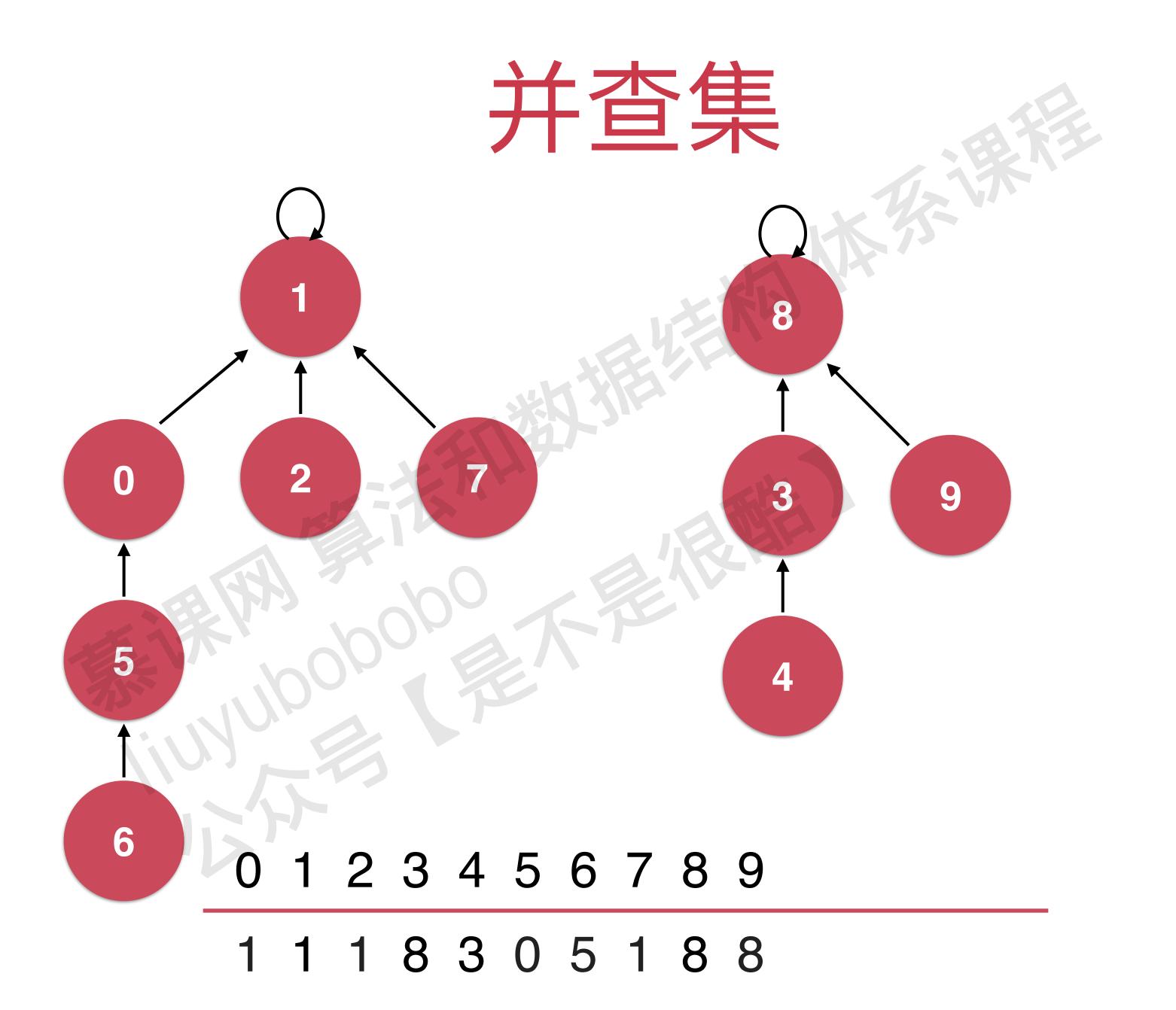




union 6, 2







实践: Quick Union 的并查集实现

并查集的优化



之前Quick Union的问题

Quick Union



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 3 4 5 6 7 8 9

union 0, 1



parent 0 1 2 3 4 5 6 7 8 9
0 1 2 3 4 5 6 7 8 9

parent 0 1 2 3 4 5 6 7 8 9

0 1 2 3 4 5 6 7 8 9

parent 0 1 2 3 4 5 6 7 8 9

1 1 2 3 4 5 6 7 8 9

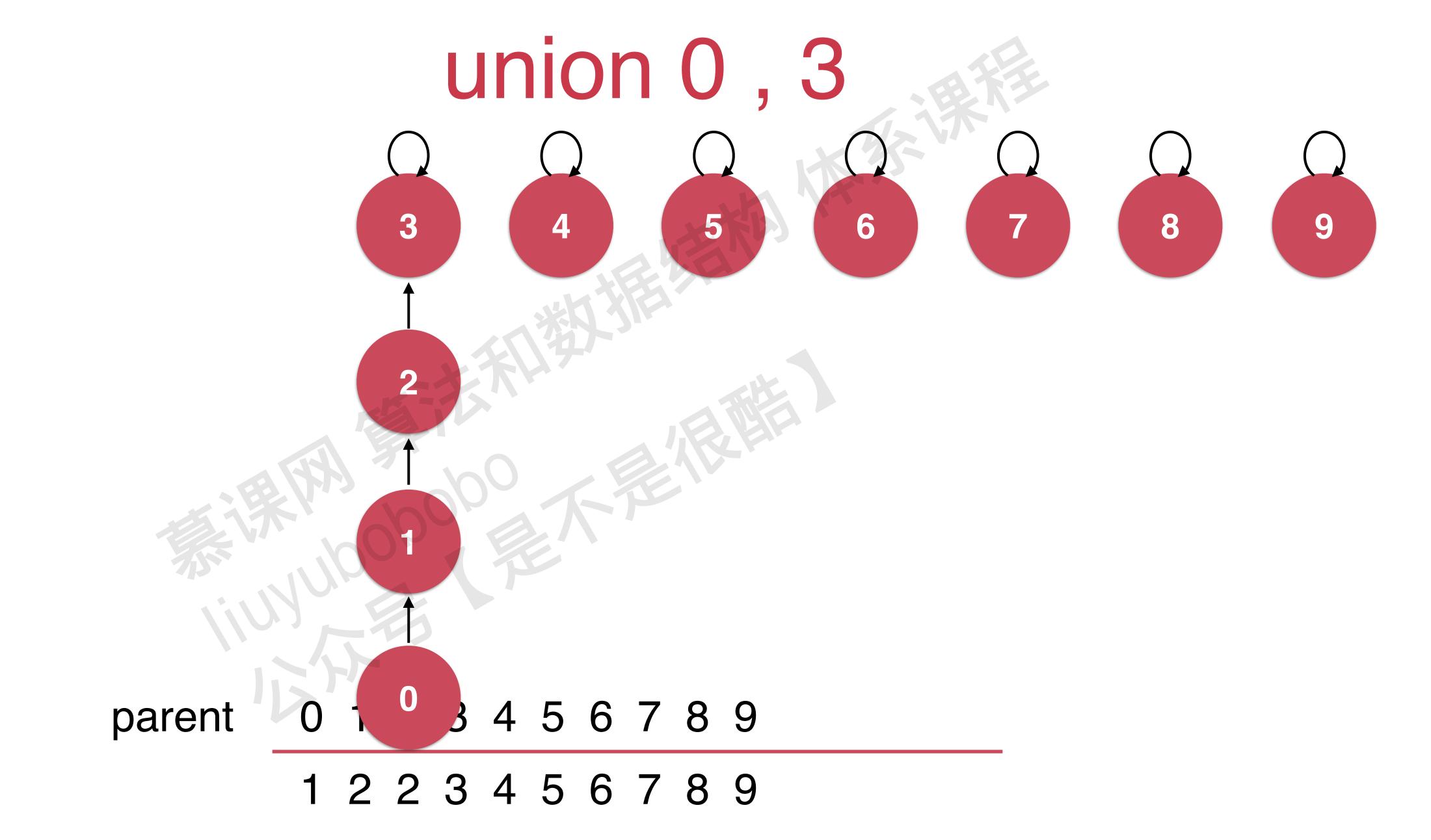
parent 0 1 2 3 4 5 6 7 8 9

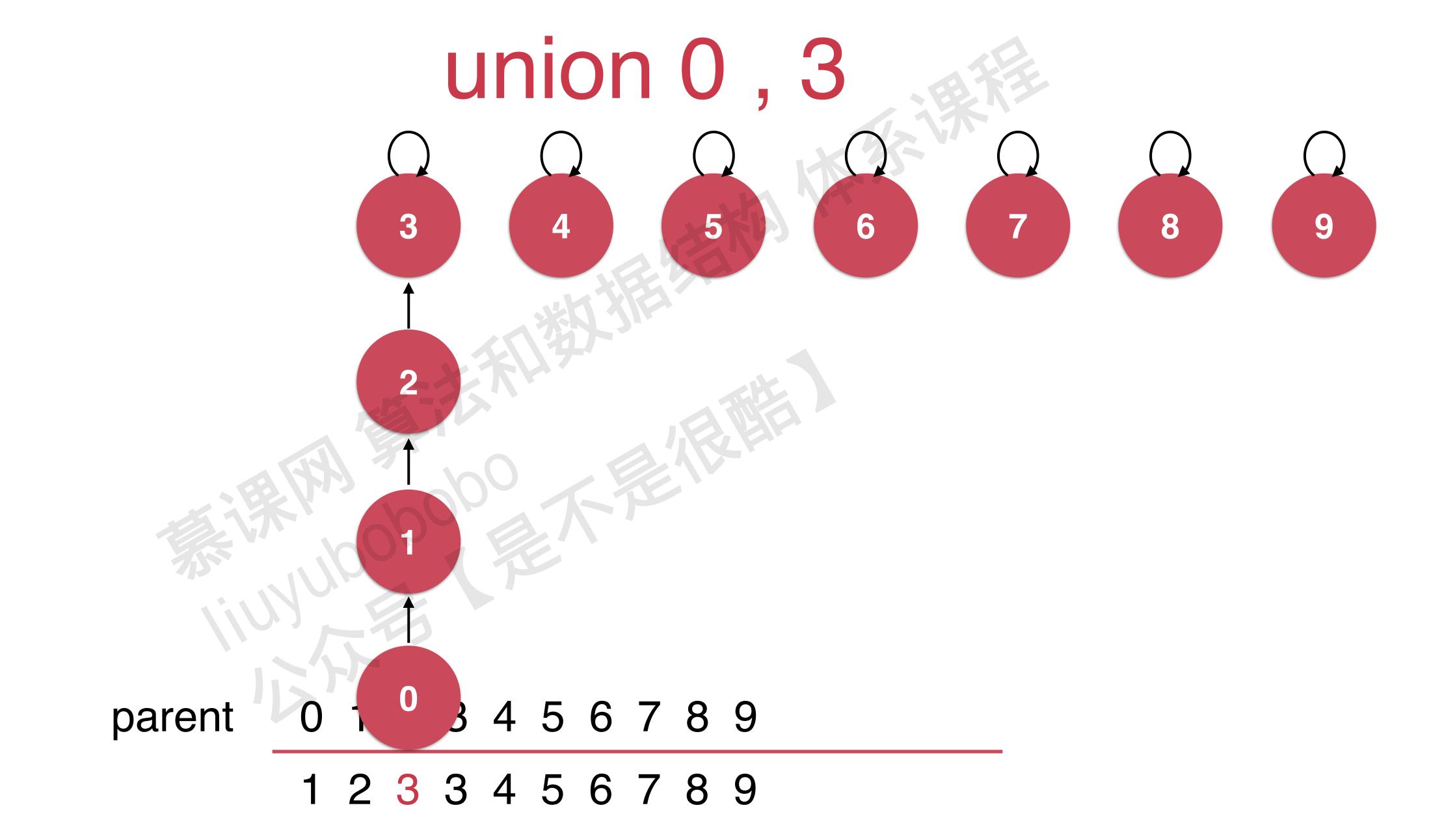
1 1 2 3 4 5 6 7 8 9

parent 0 1 2 3 4 5 6 7 8 9 1 1 2 3 4 5 6 7 8 9

parent 0 1 2 3 4 5 6 7 8 9 1 2 2 3 4 5 6 7 8 9

union 0,3 parent 0 1 2 3 4 5 6 7 8 9 1 2 2 3 4 5 6 7 8 9





解决方案:考虑size

union 4, 9 0 1 2 3 4 5 6 7 8 9 parent 0 1 2 8 3 5 5 7 8 9

union 4,9 0 1 2 3 4 5 6 7 8 9 parent 0 1 2 8 3 5 5 7 8 9

union 4,9 0 1 2 3 4 5 6 7 8 9 parent 0 1 2 8 3 5 5 7 9 9

union 4, 9 0 1 2 3 4 5 6 7 8 9 parent 0 1 2 8 3 5 5 7 8 9

union 4,9 0 1 2 3 4 5 6 7 8 9 parent 0 1 2 8 3 5 5 7 8 9

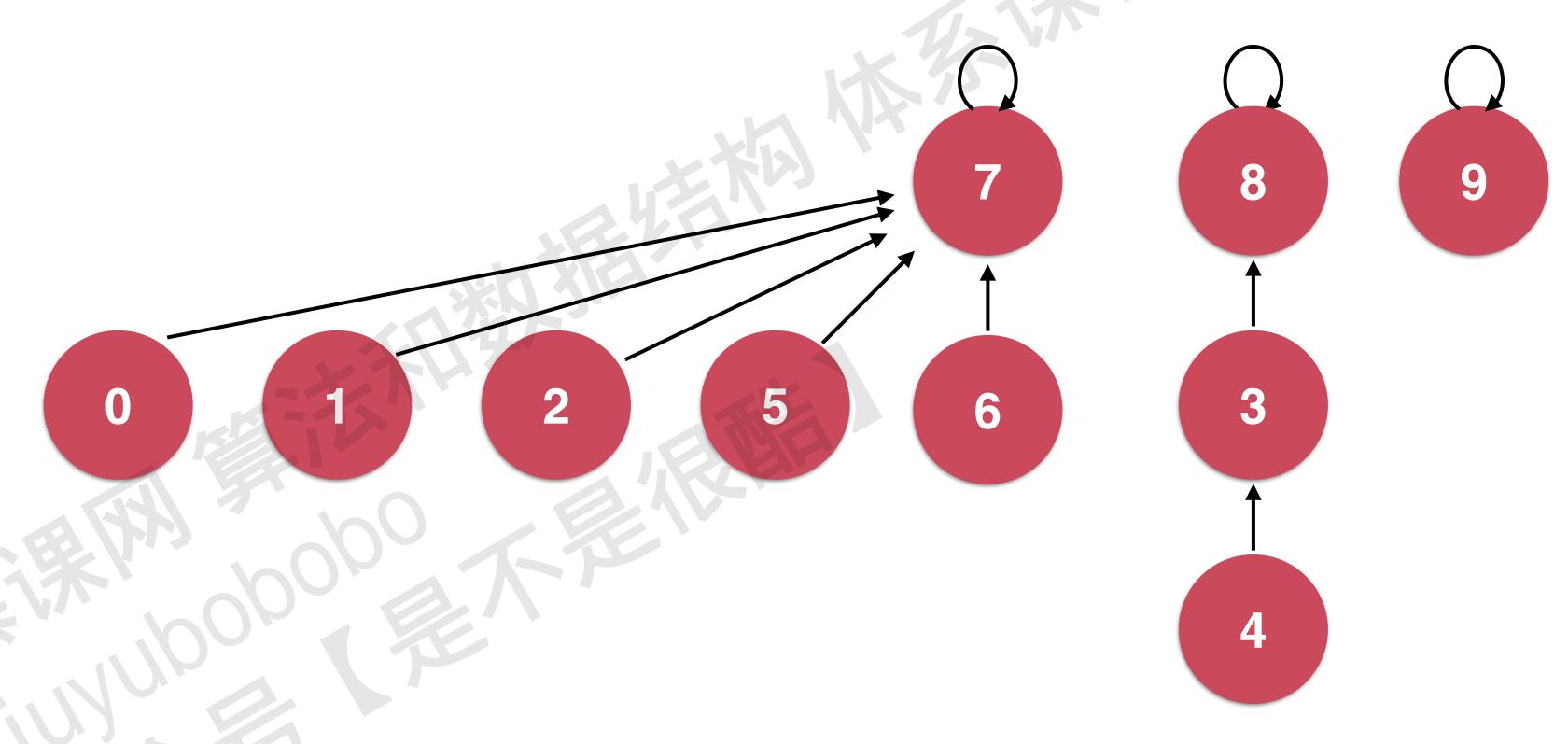
union 4,9 0 1 2 3 4 5 6 7 8 9 parent 0 1 2 8 3 5 5 7 8 8

实践:基于size的优化

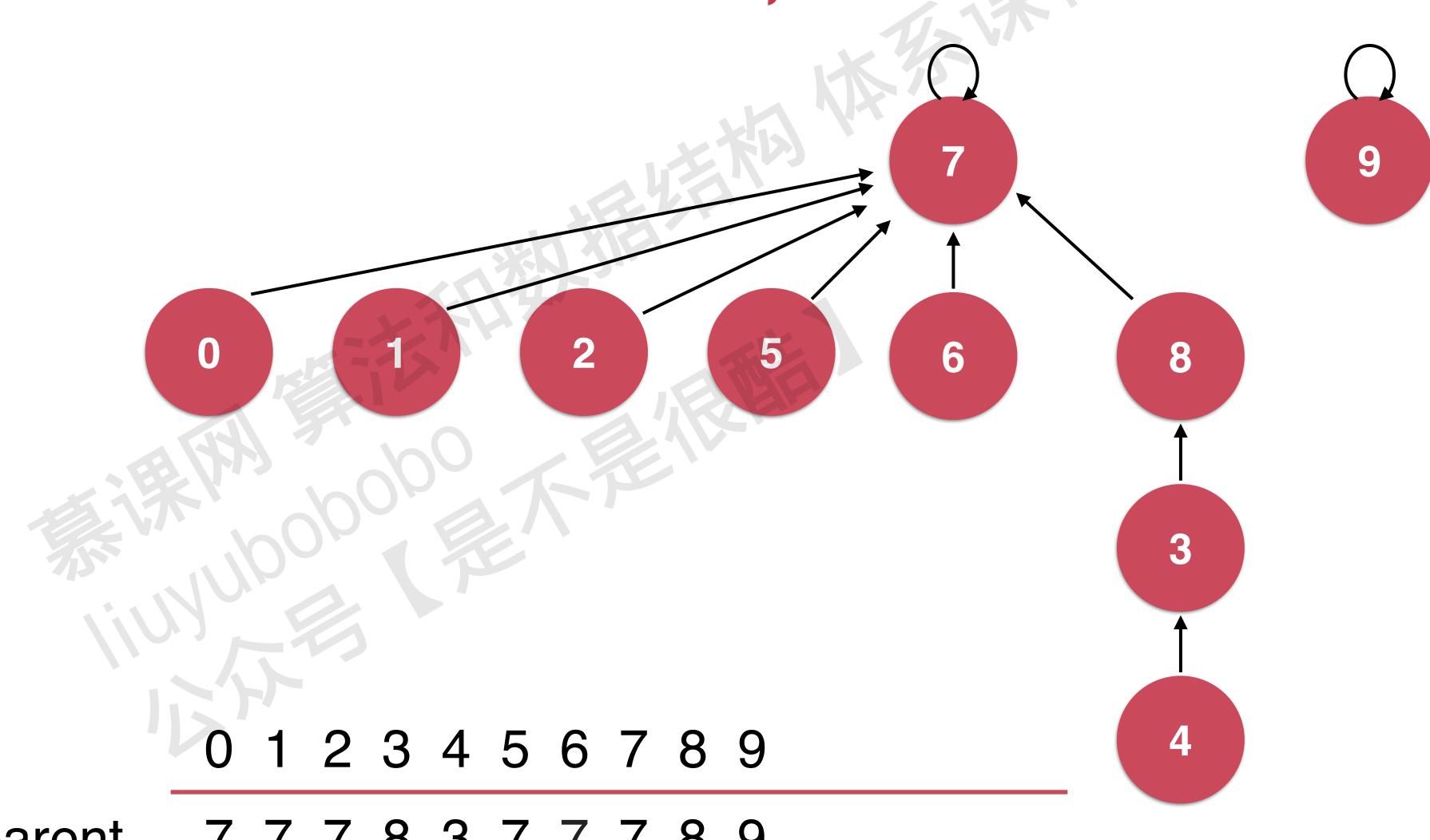


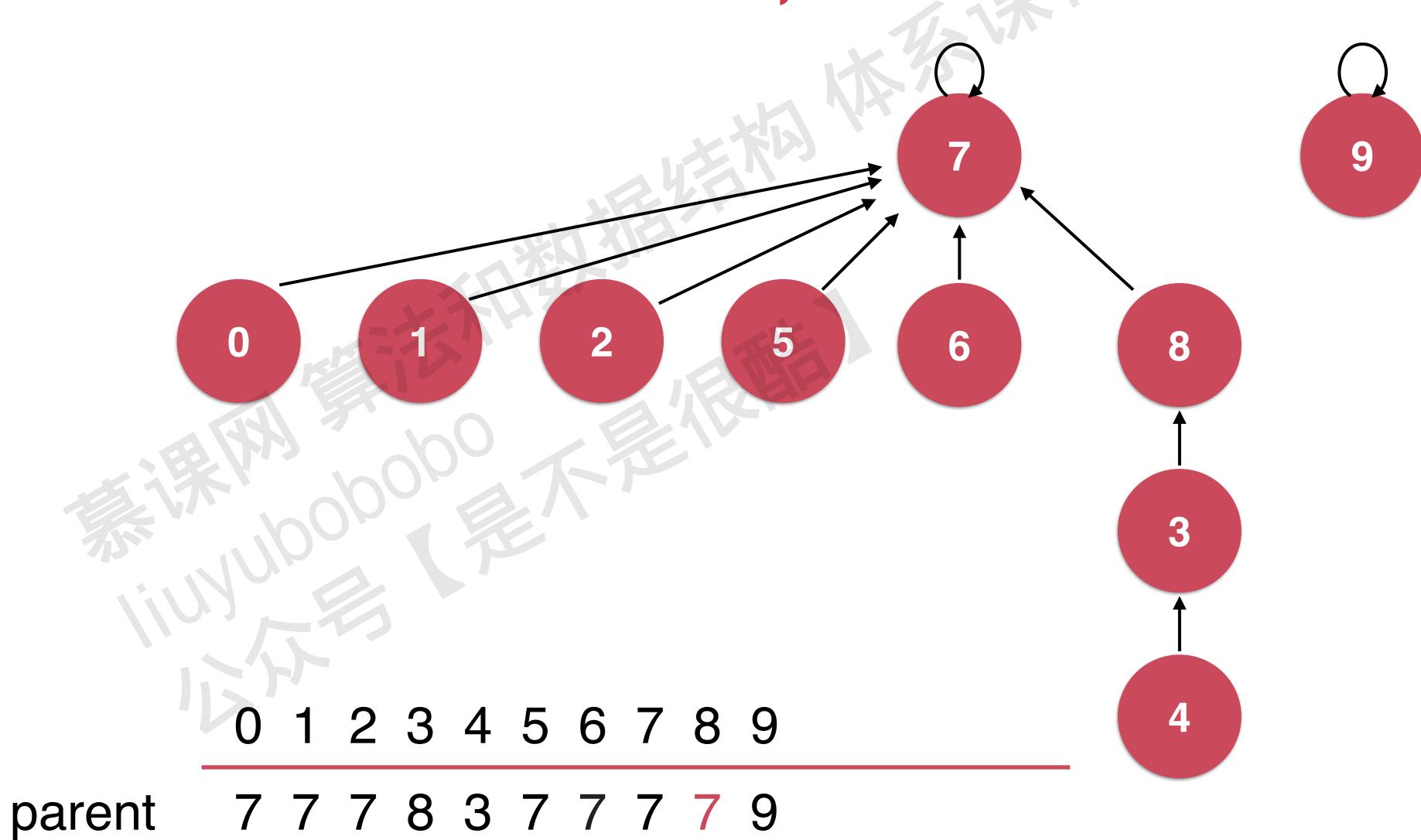
实践: Quick Union 和优化后的时间效率比较

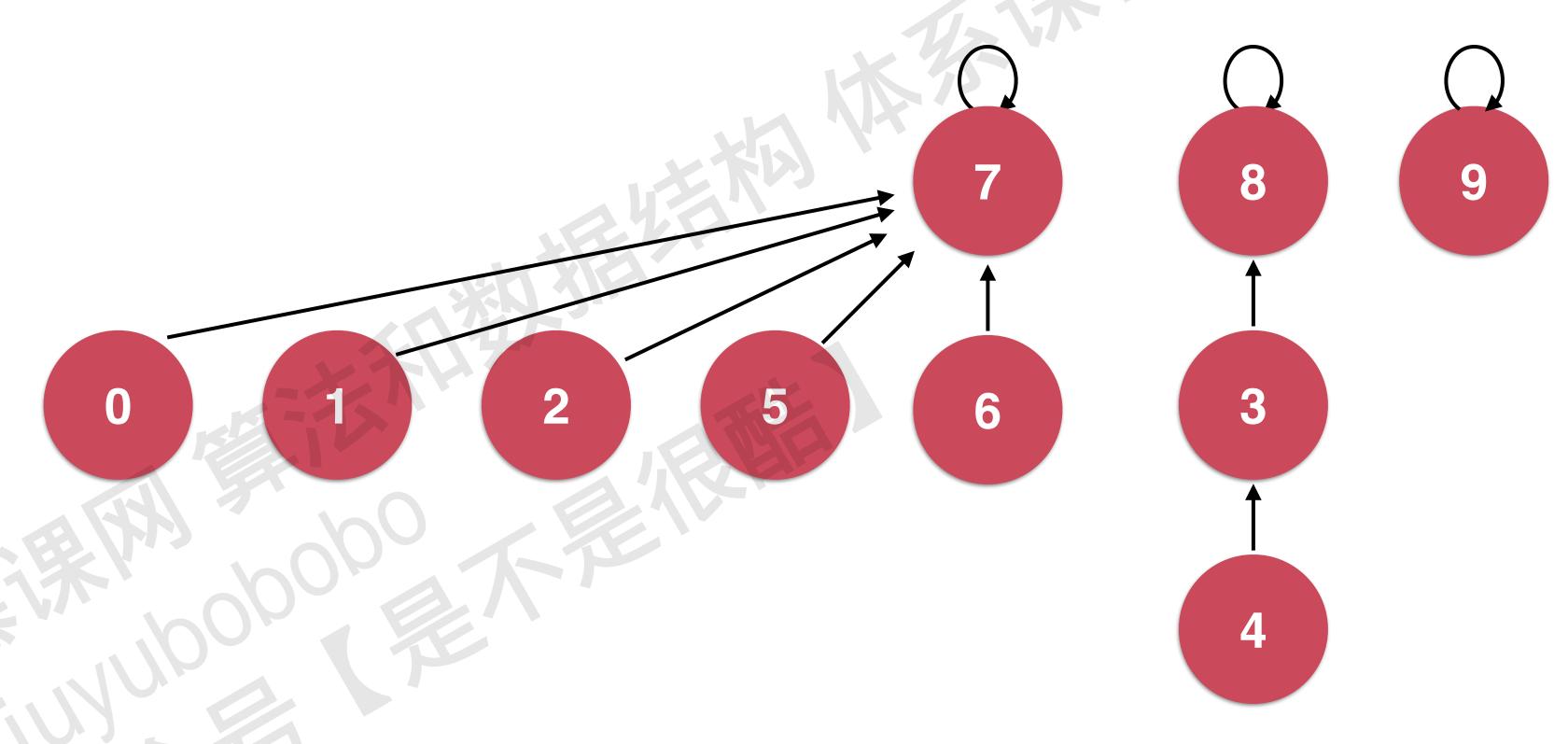
基于rank的优化



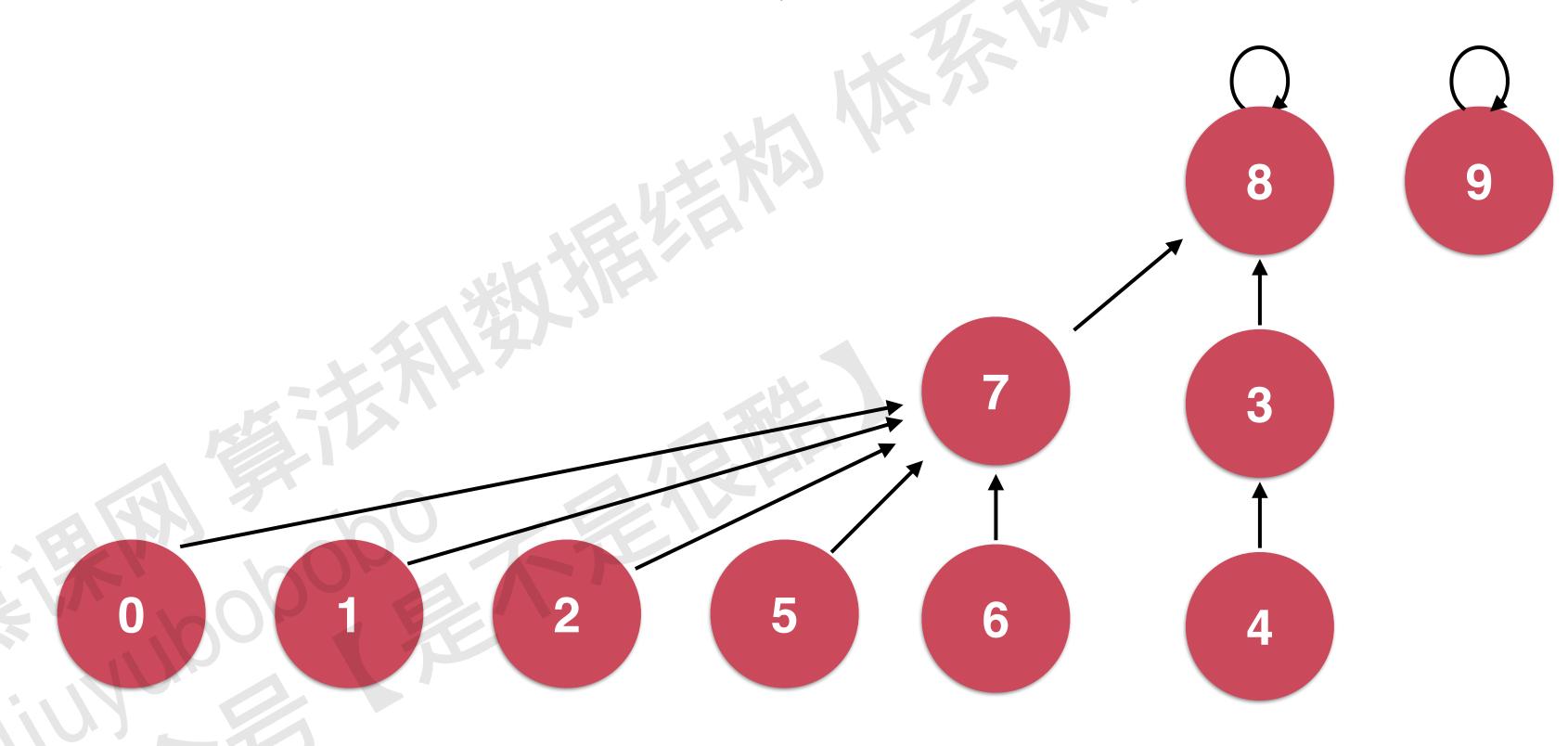
0 1 2 3 4 5 6 7 8 9



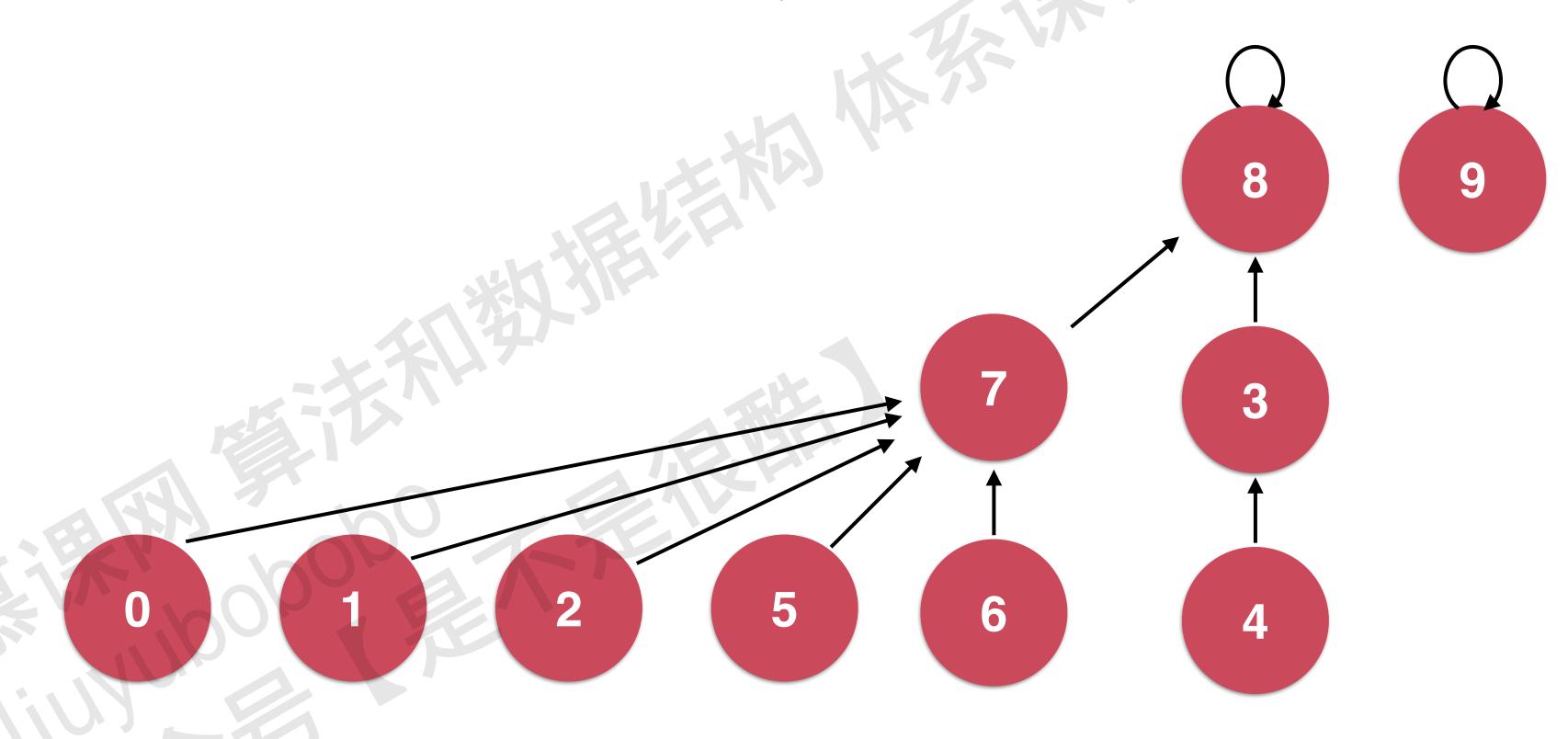




0 1 2 3 4 5 6 7 8 9



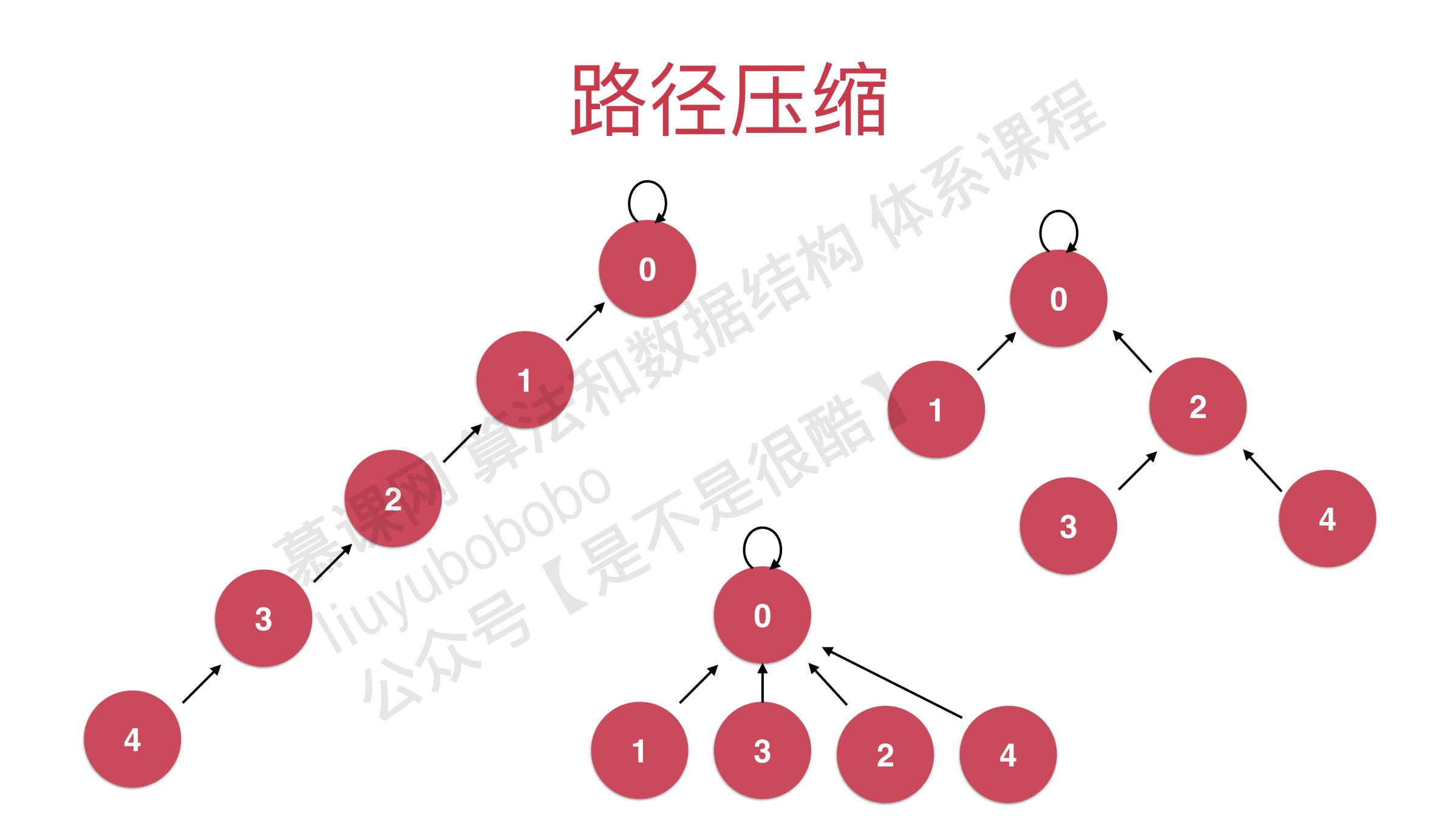
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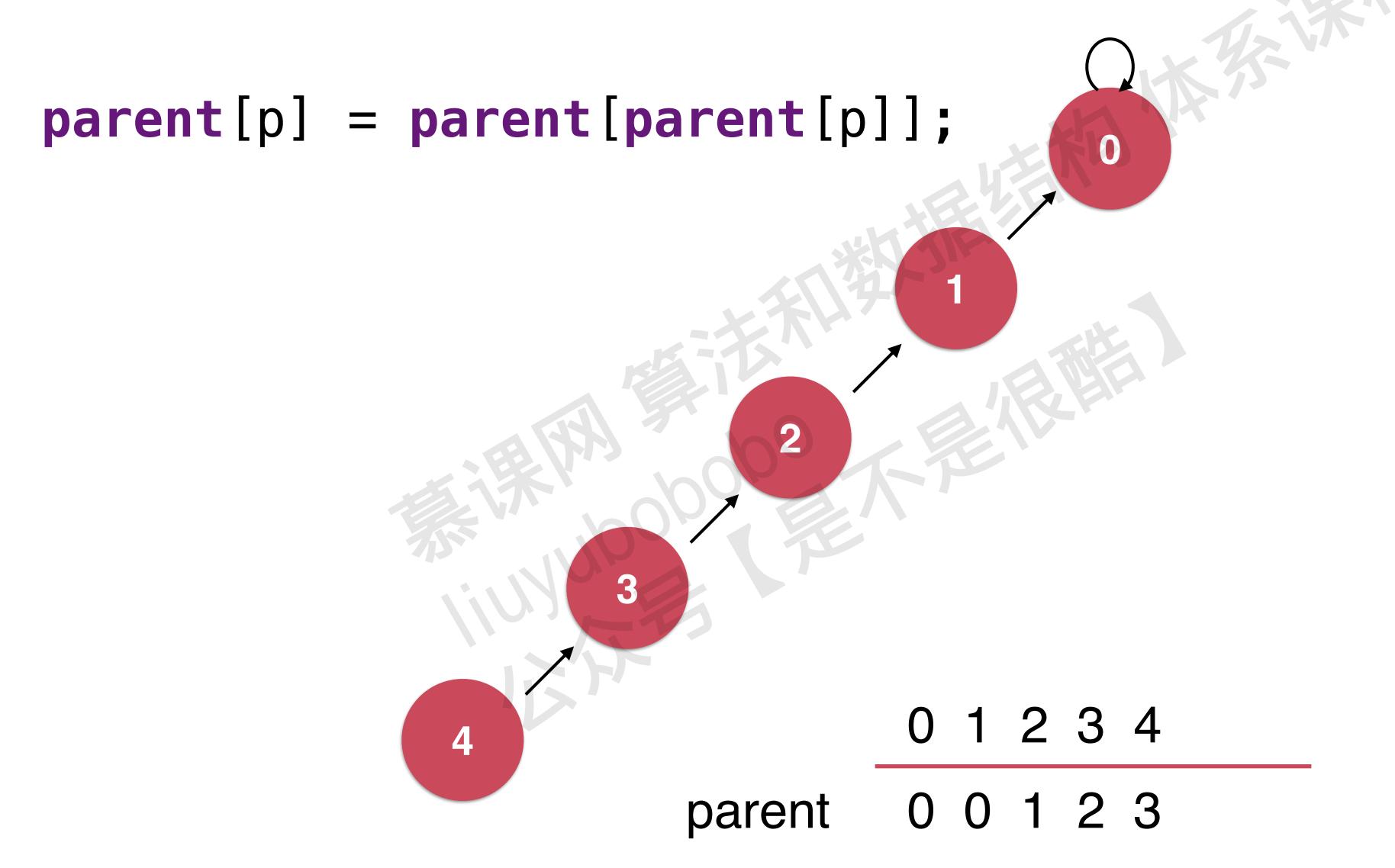


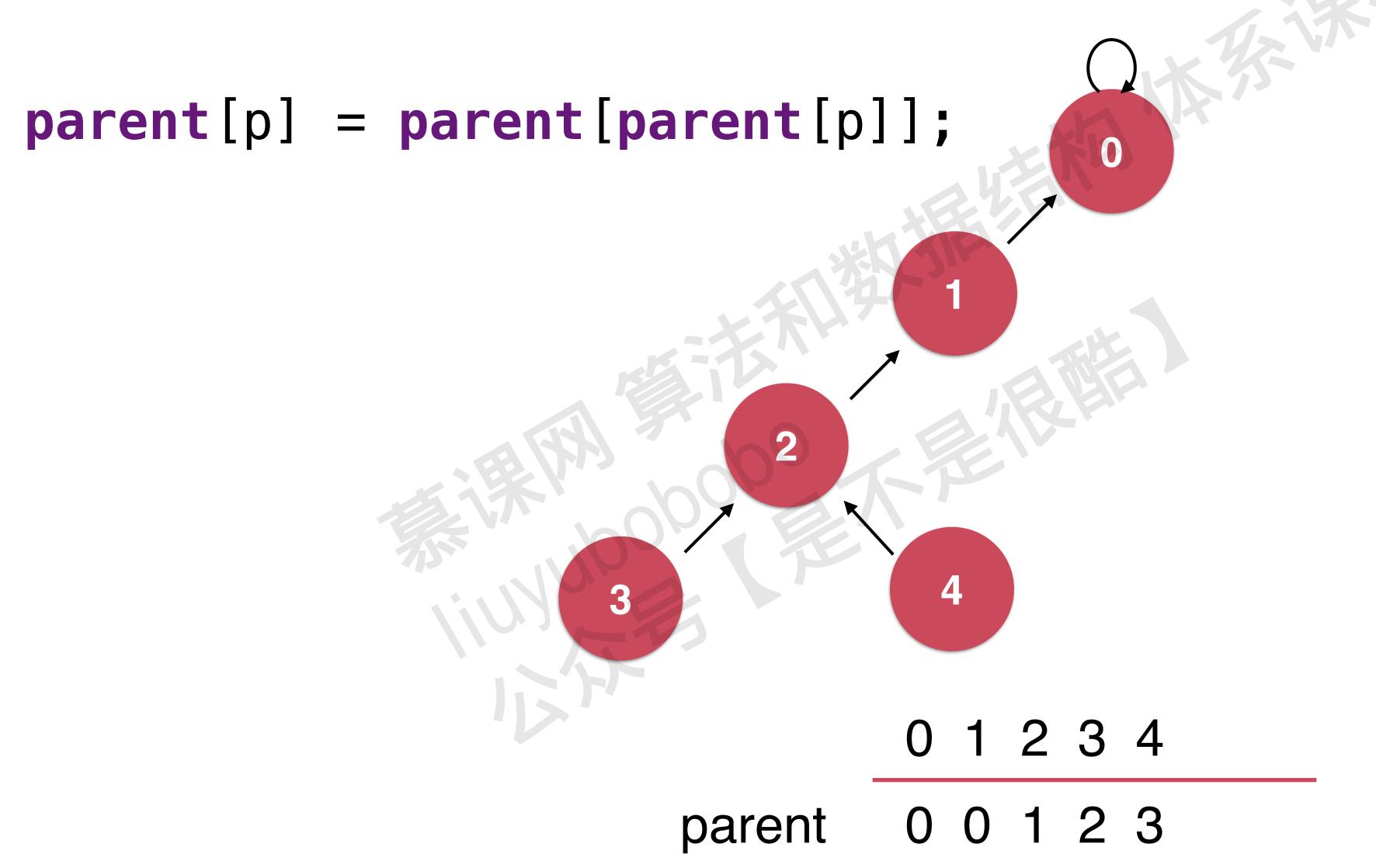
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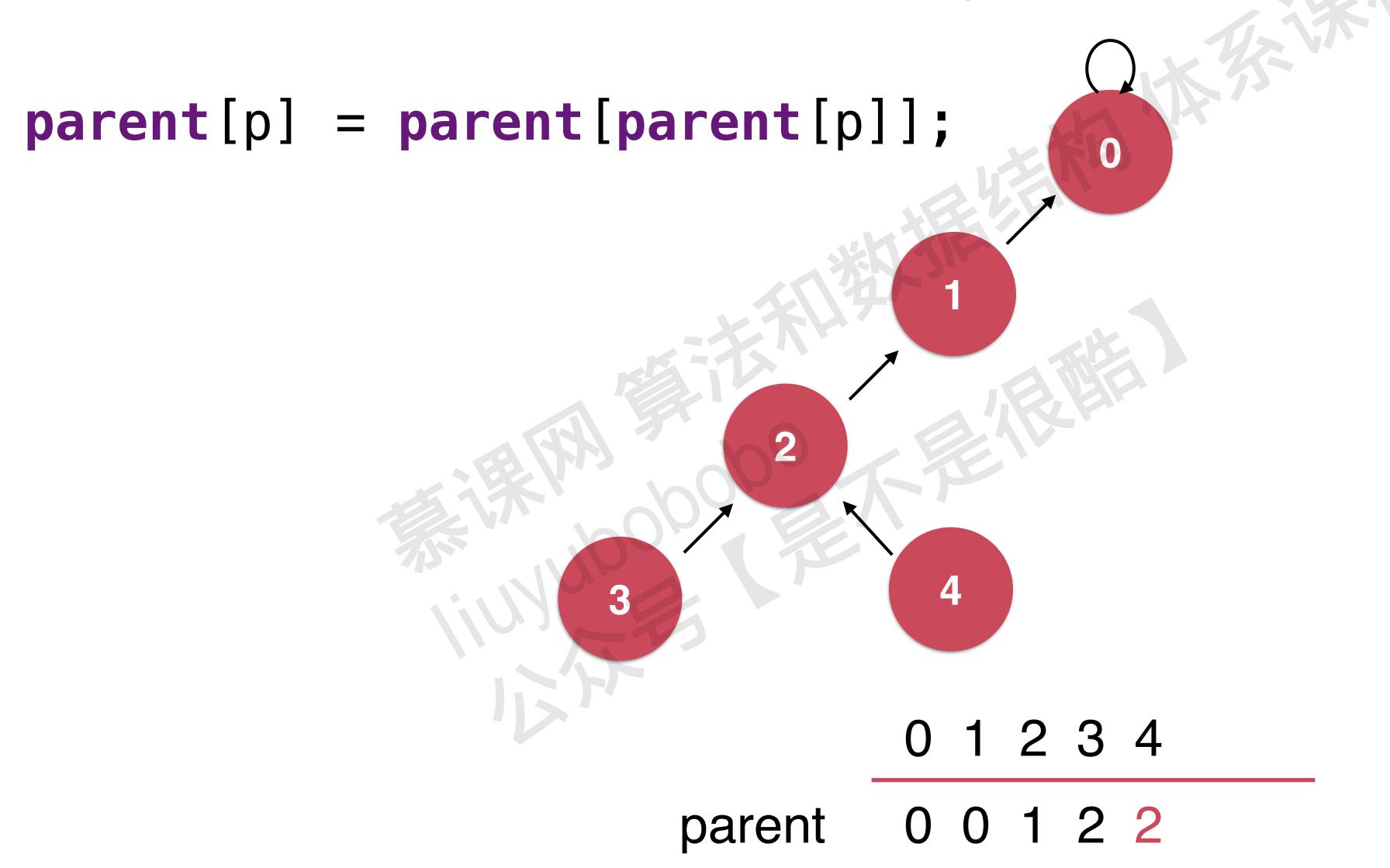
基于rank的优化 rank[i] 表示根节点为i的树的高度

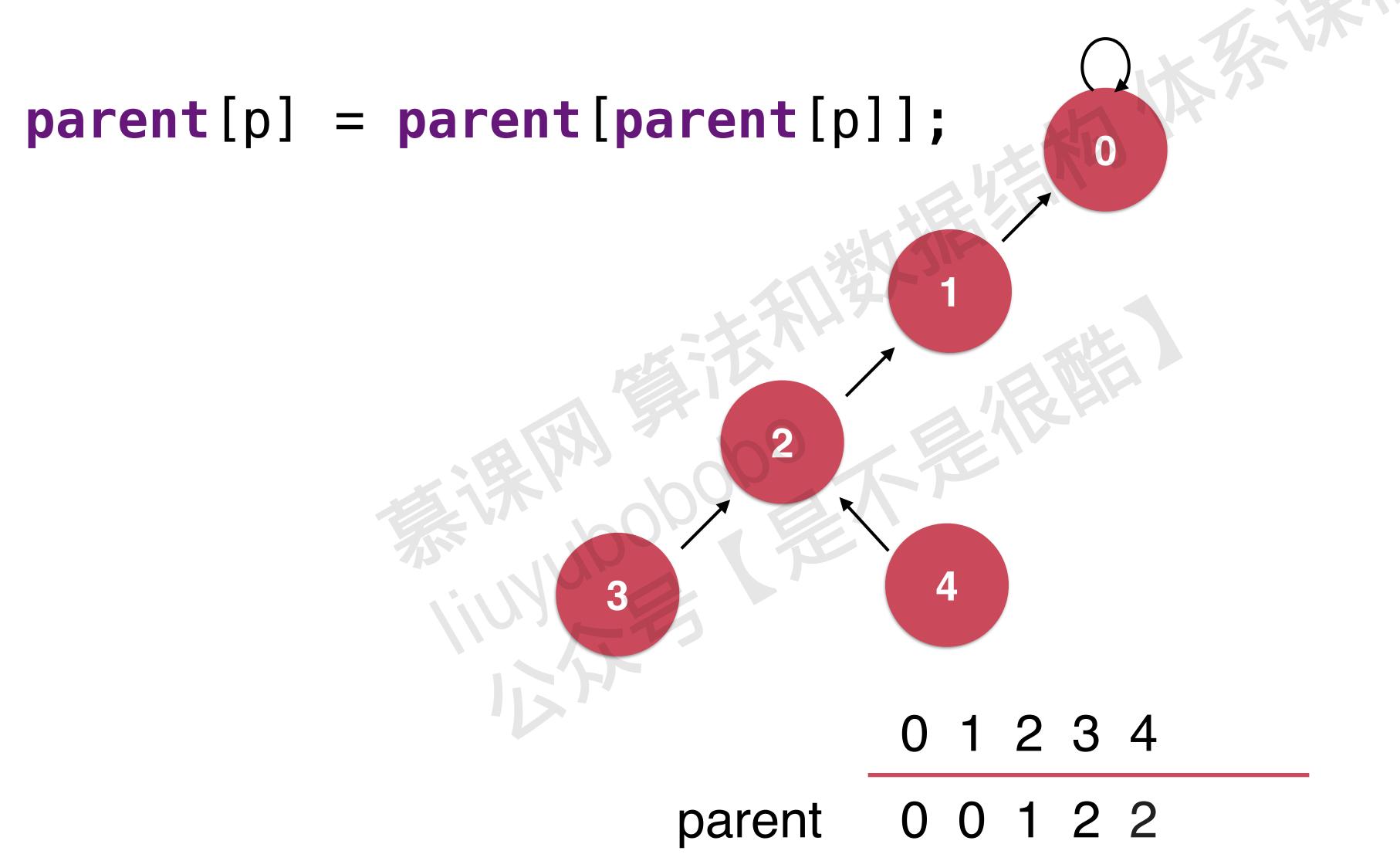
实践:基于rank的优化

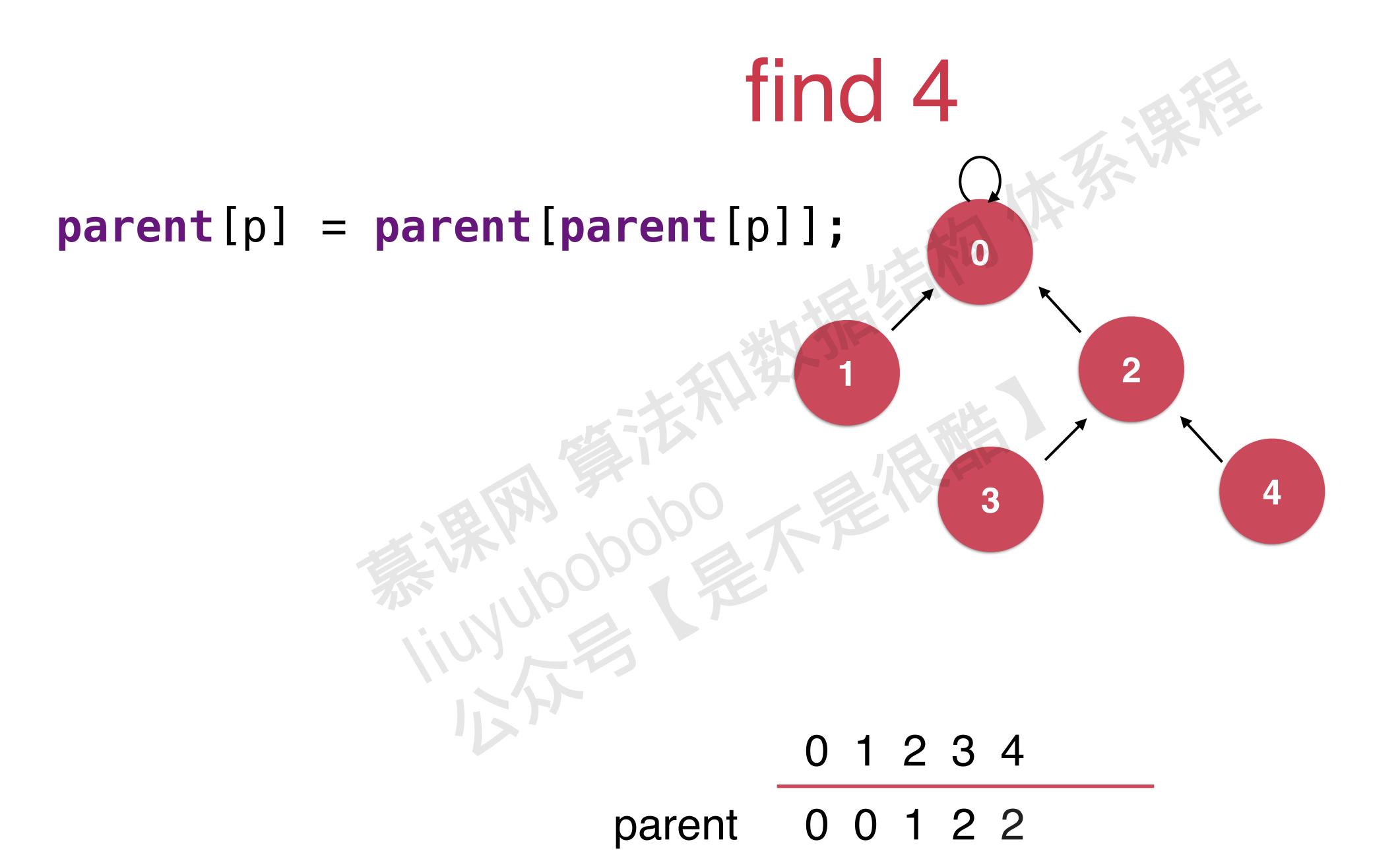


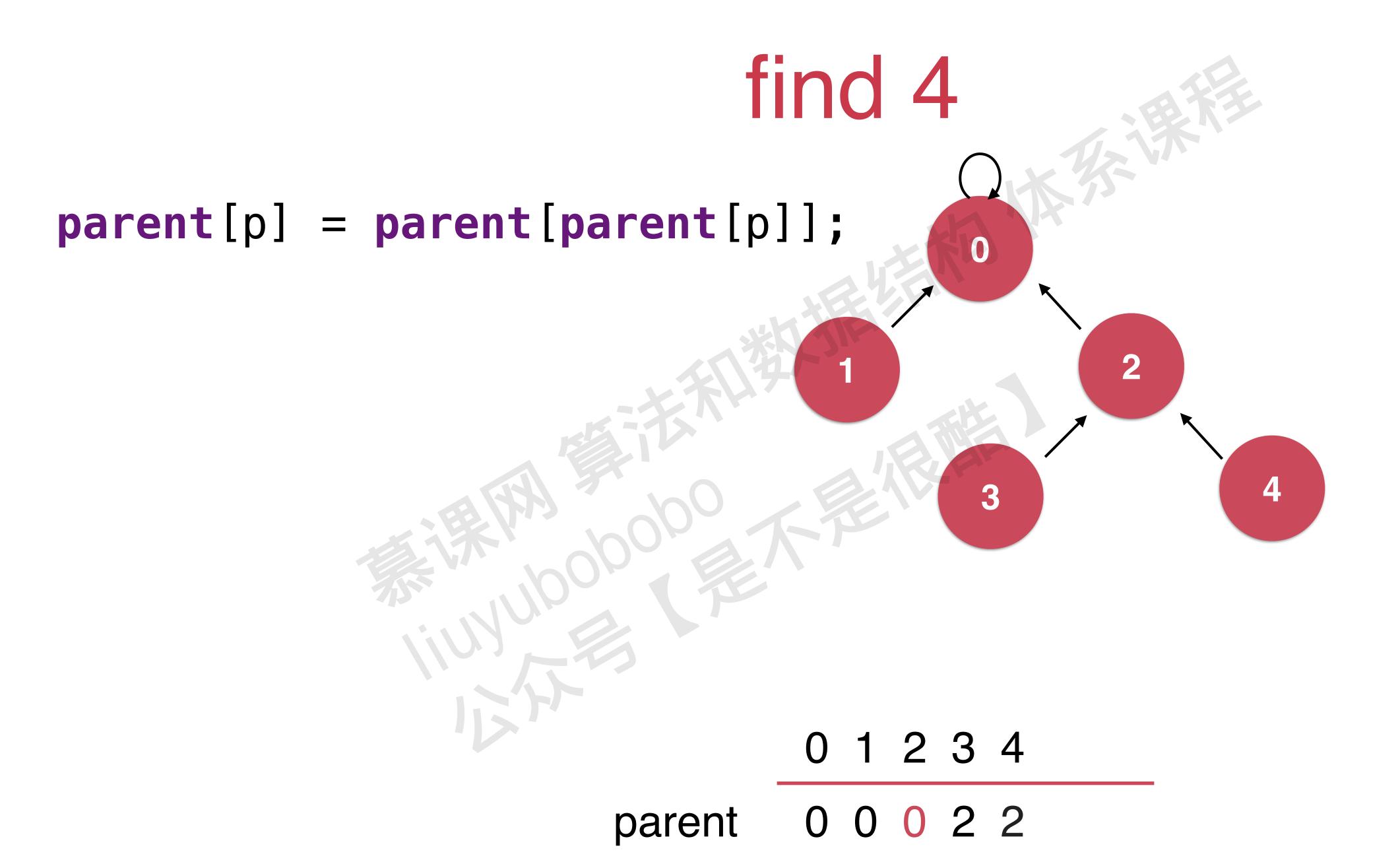


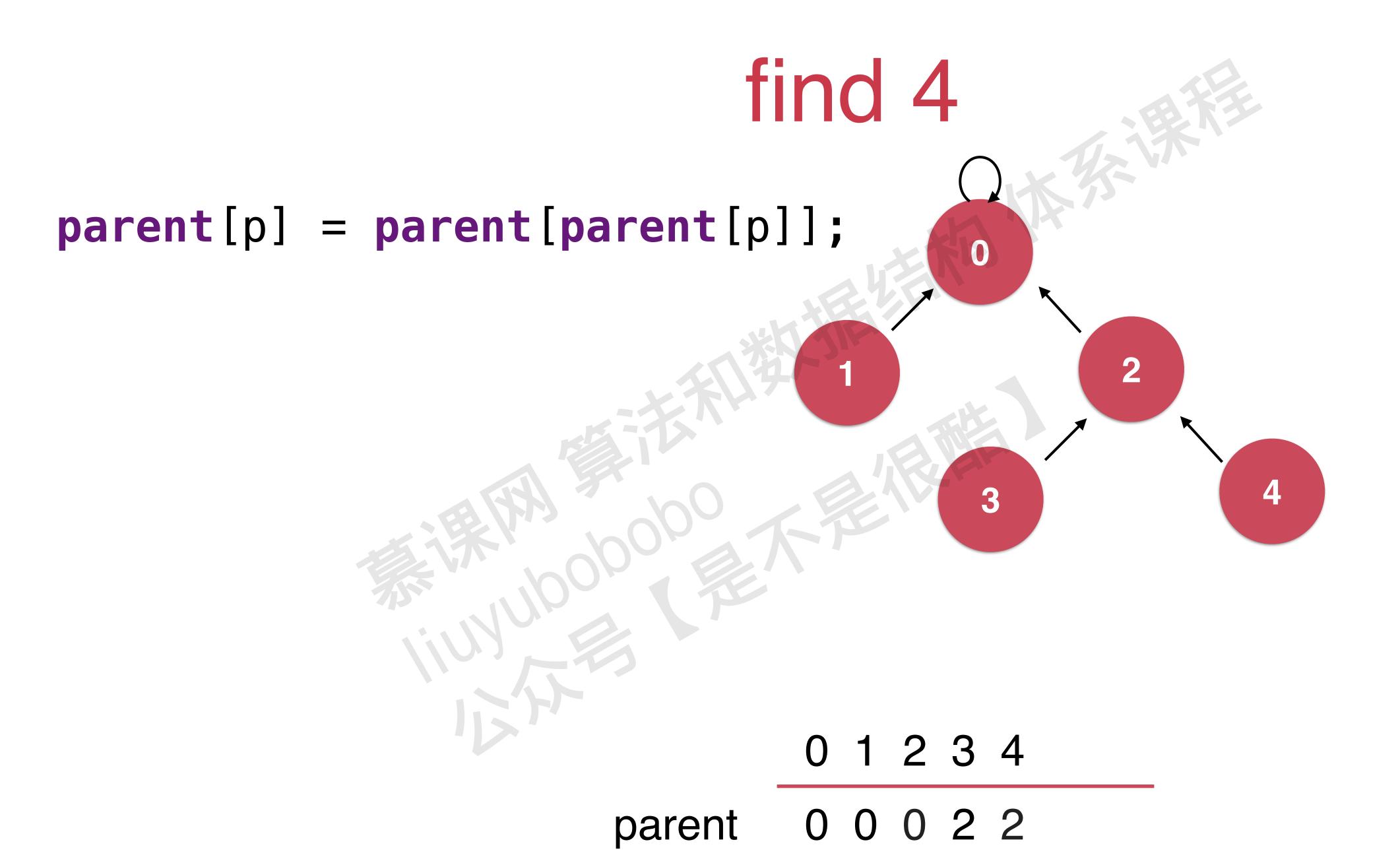


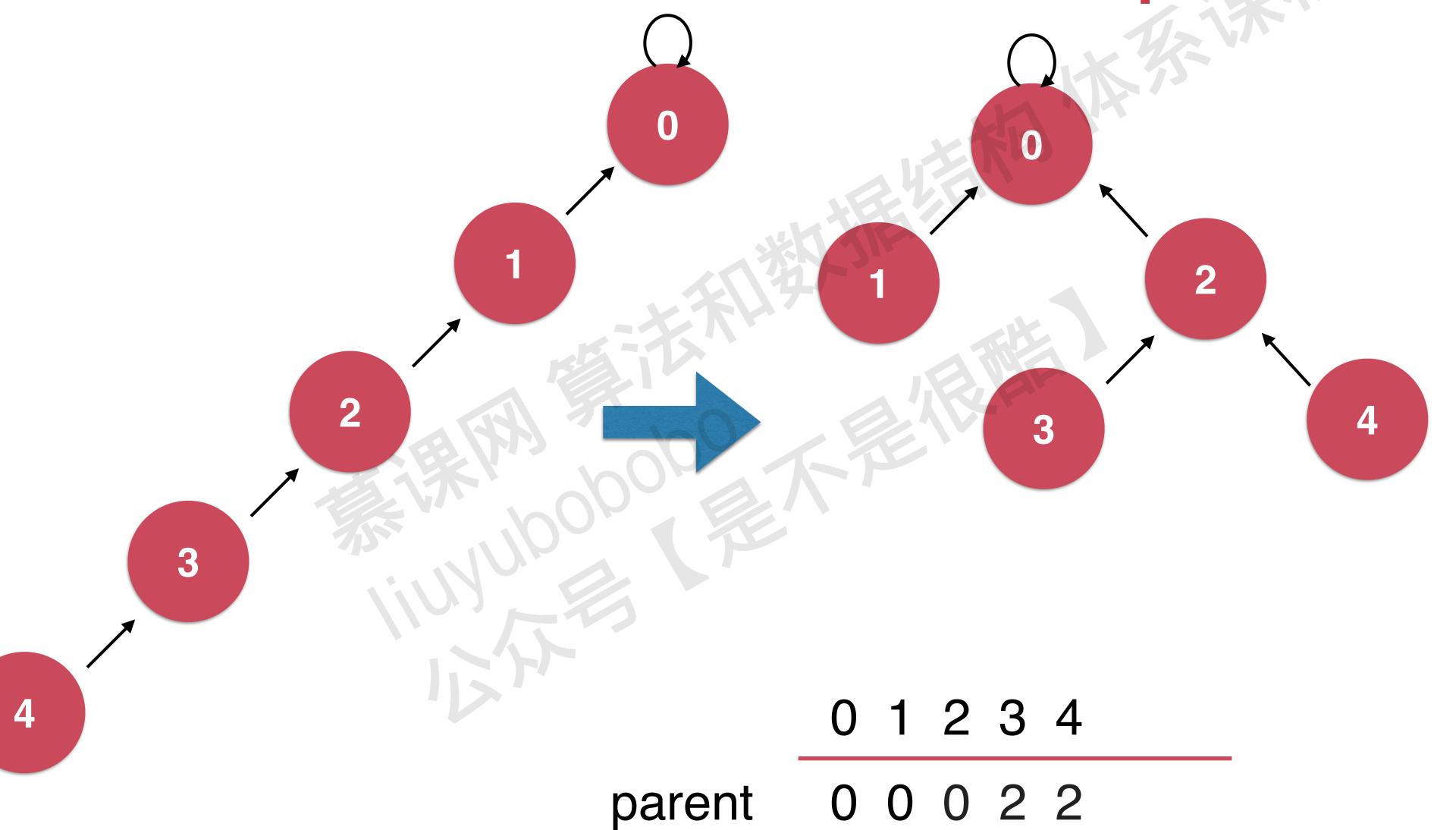










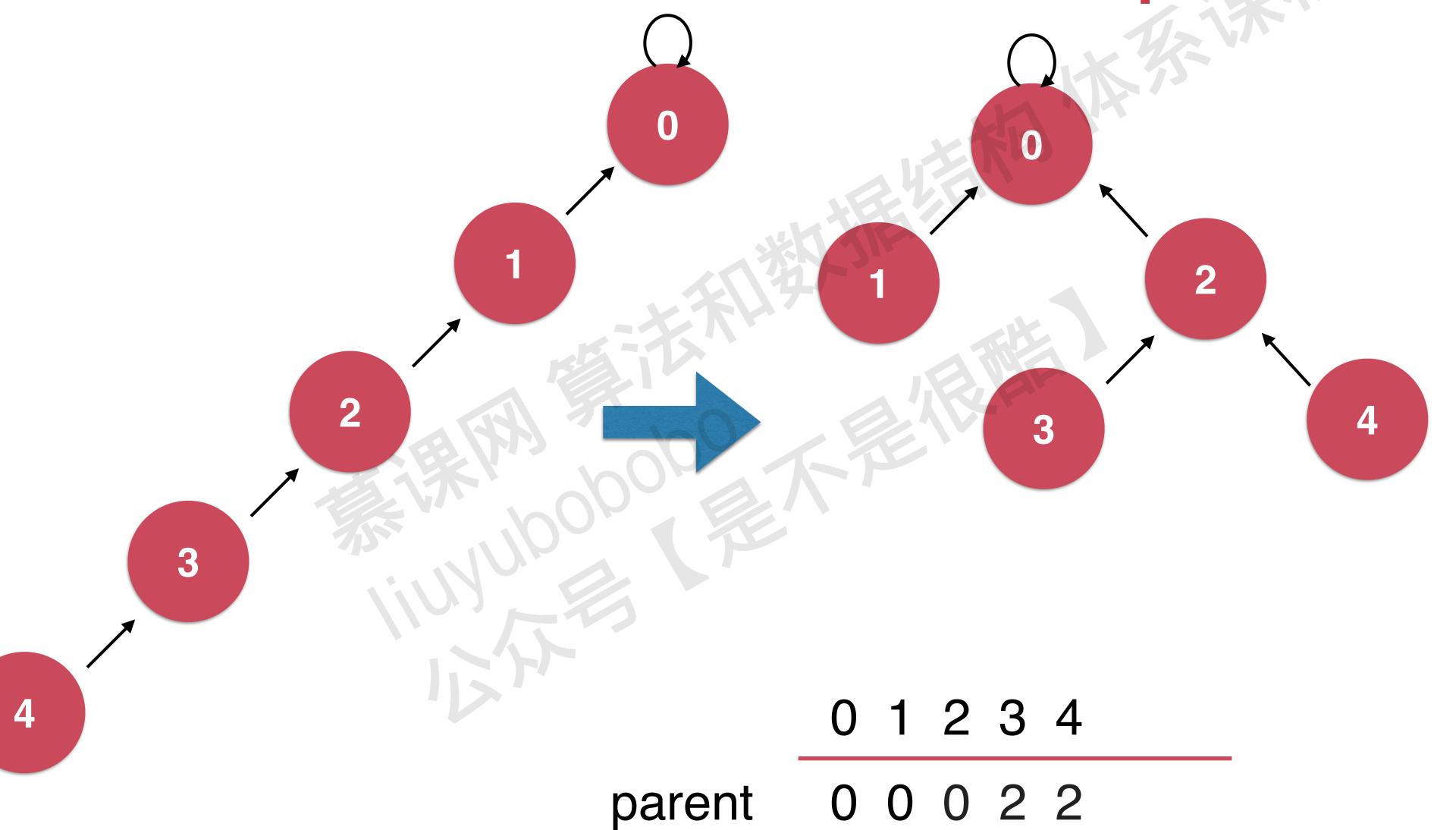


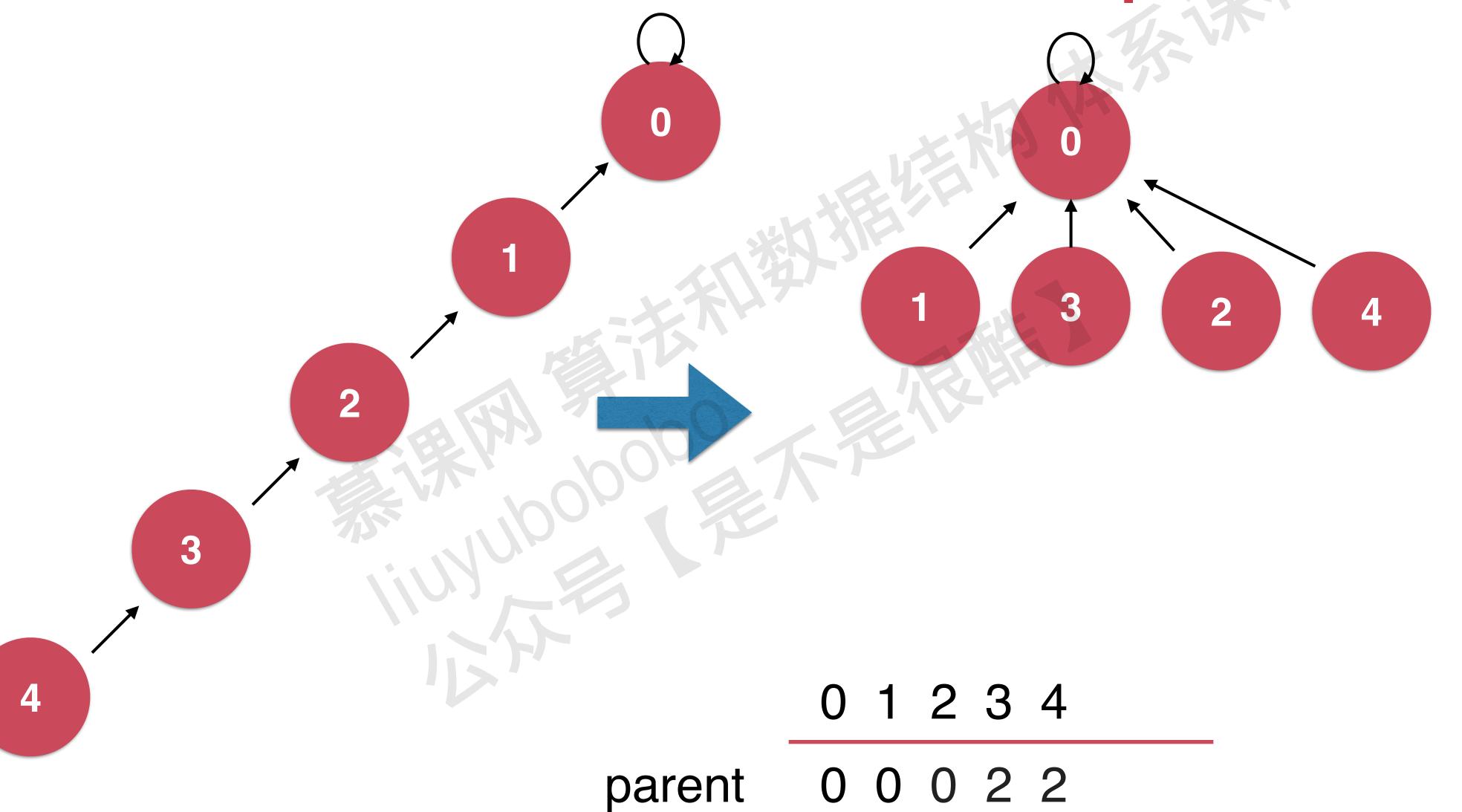
实践:路径压缩

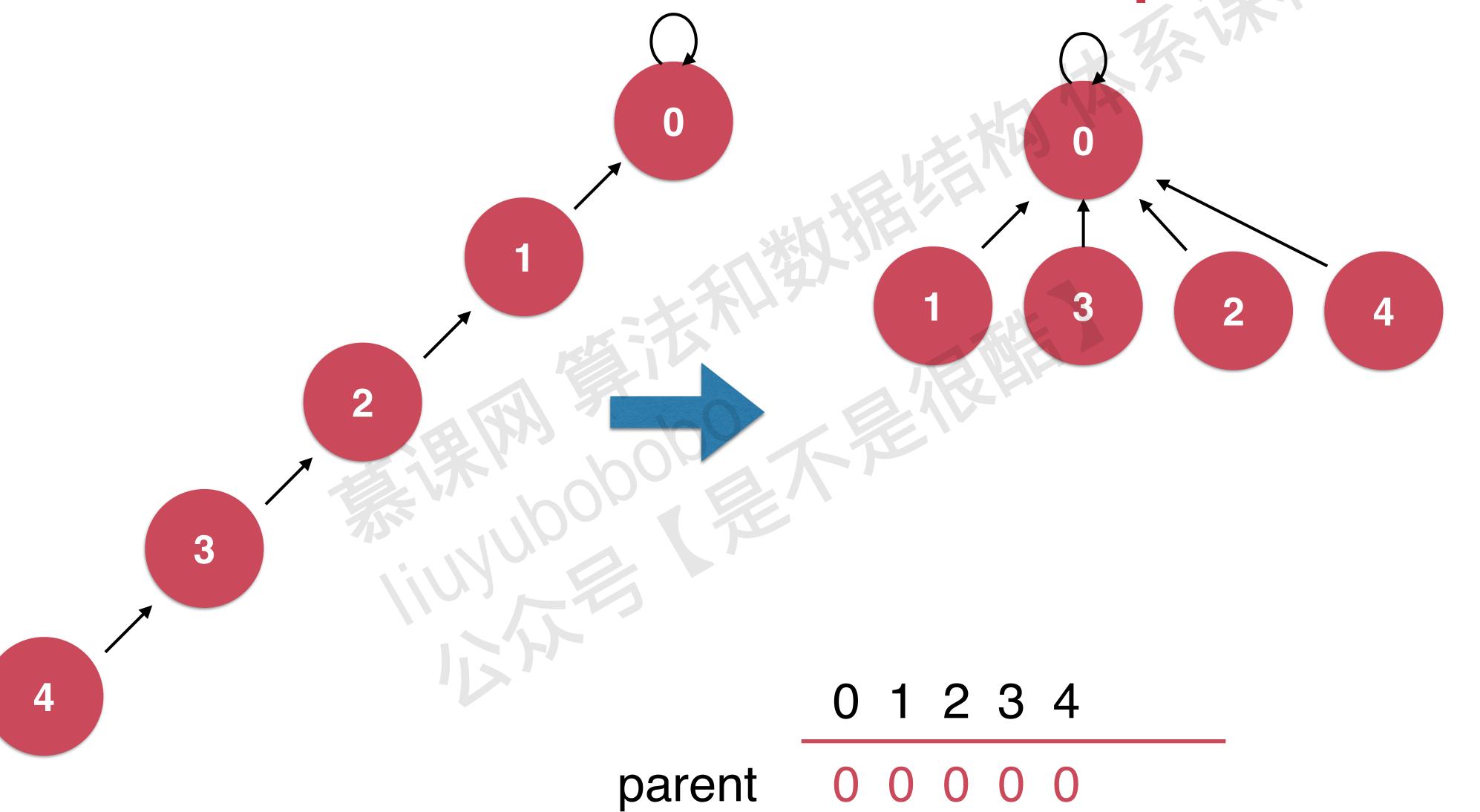


实践:比较路径压缩和非路径压缩的效率

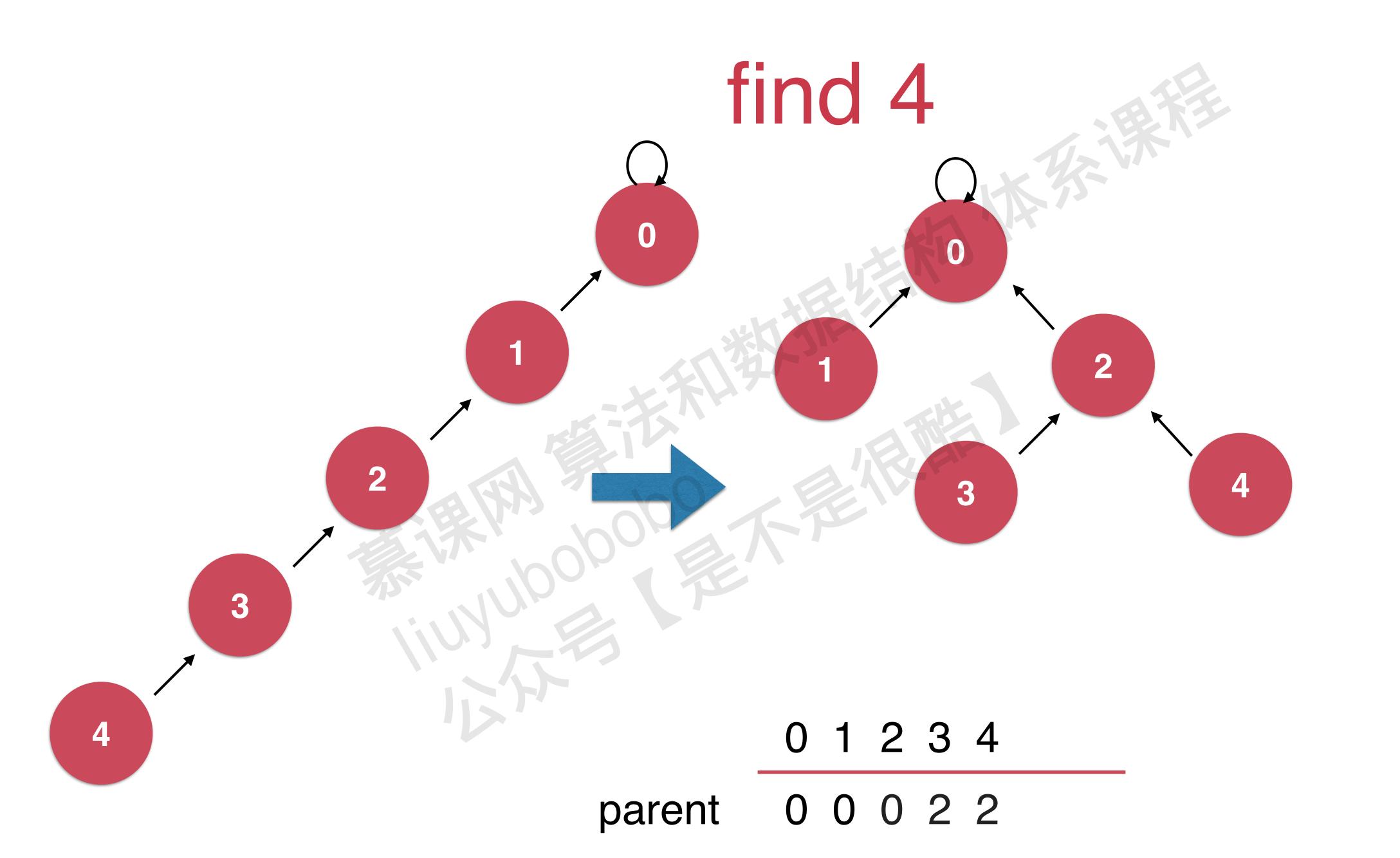
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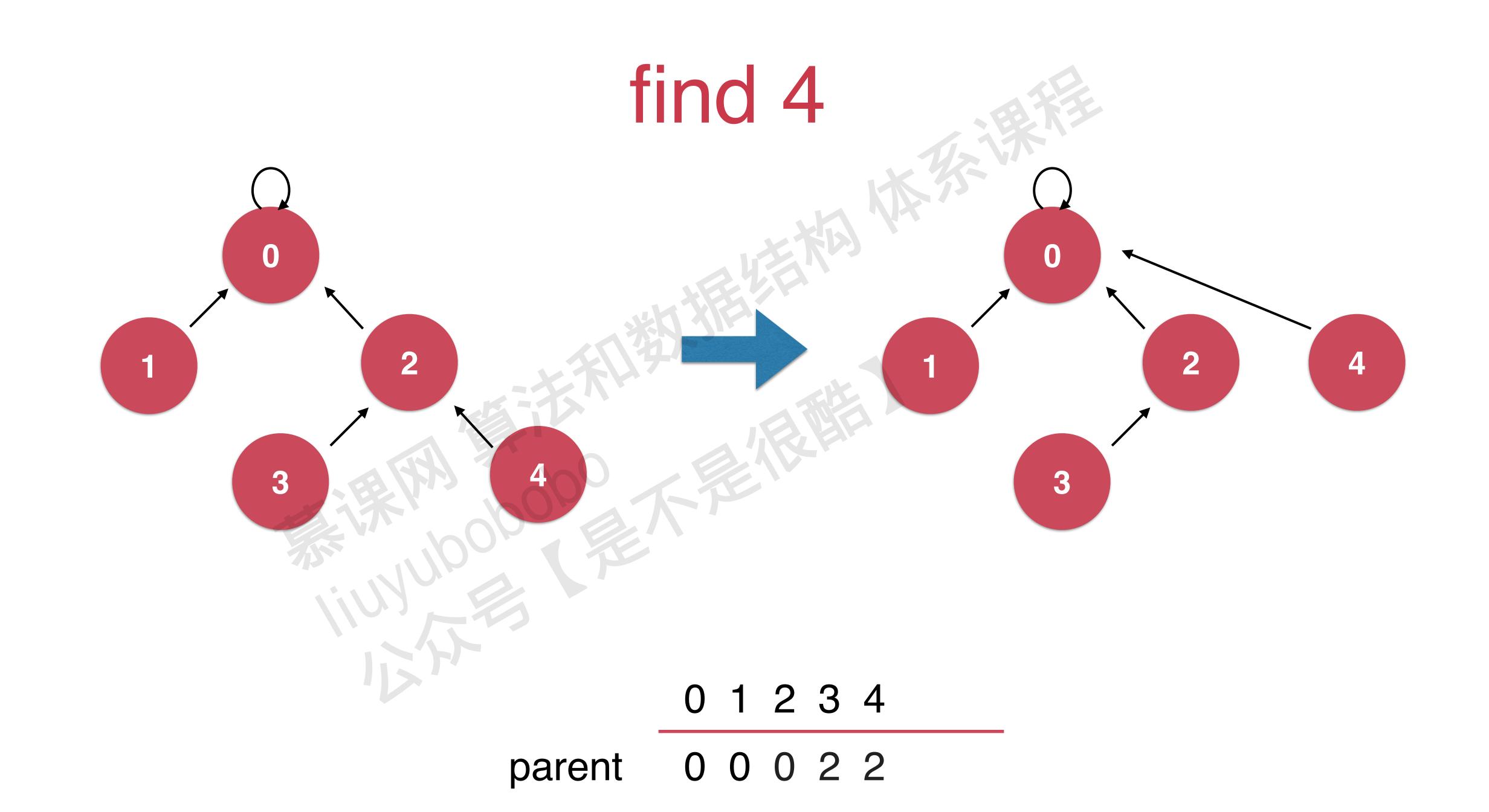


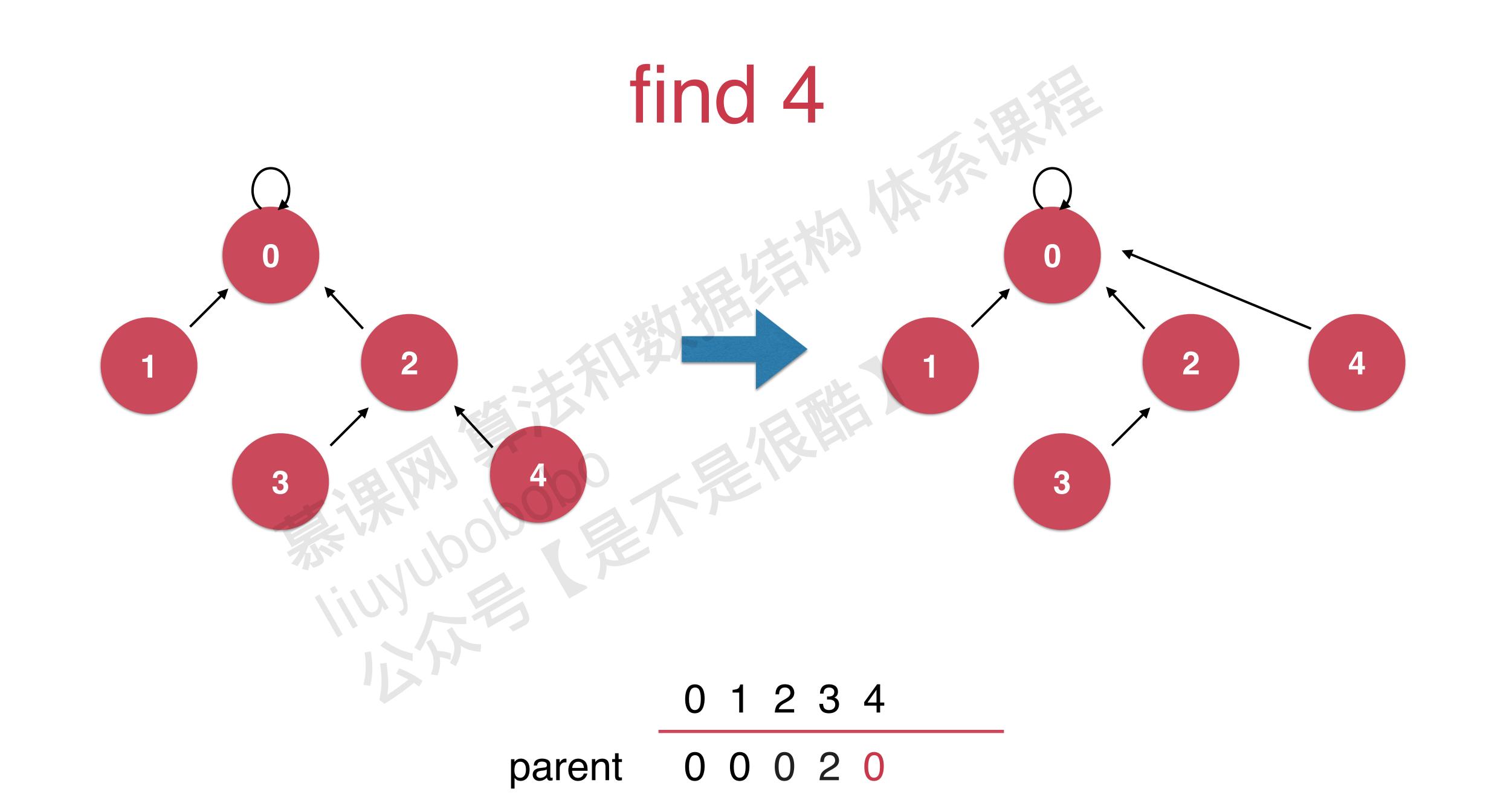


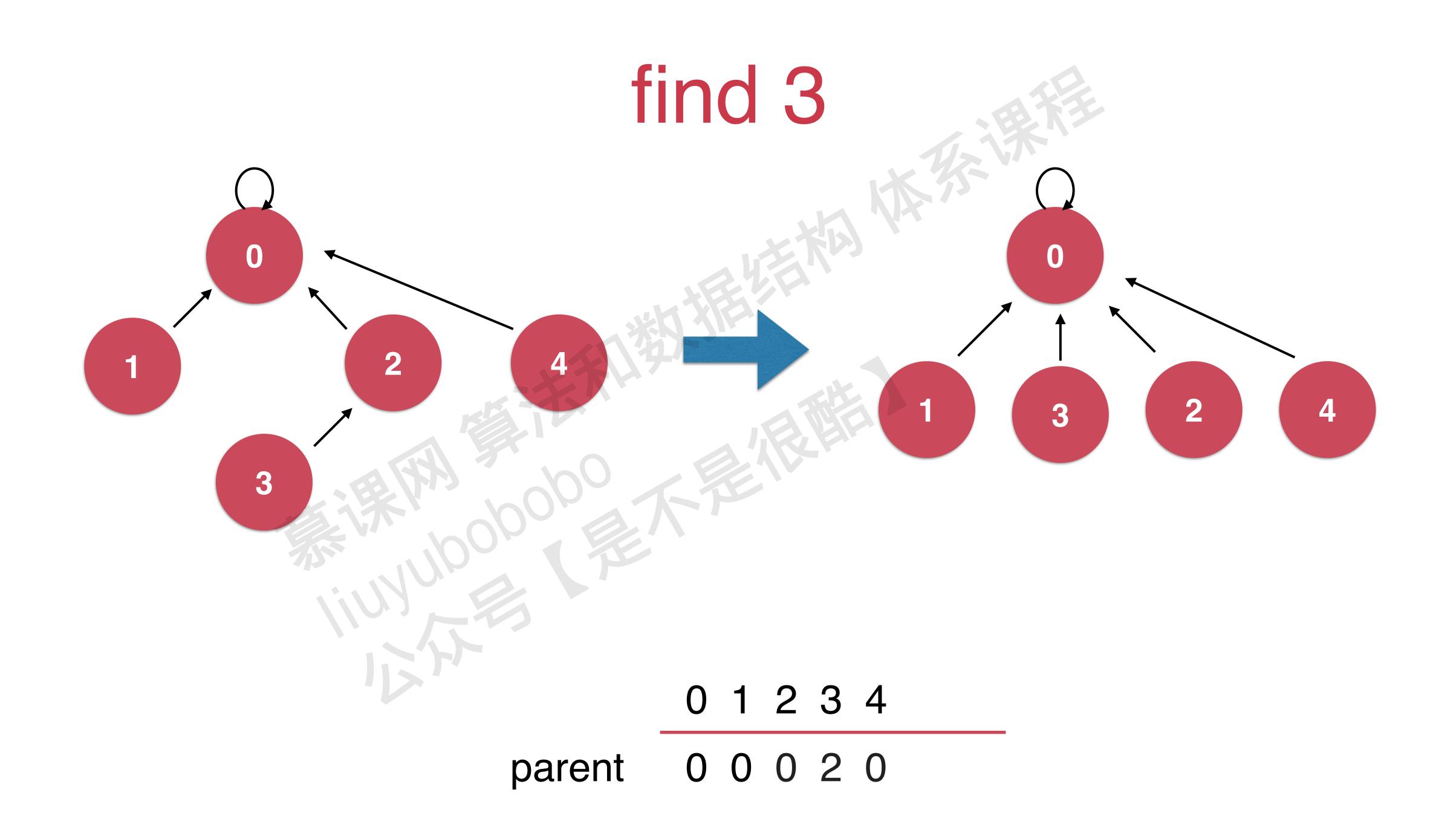


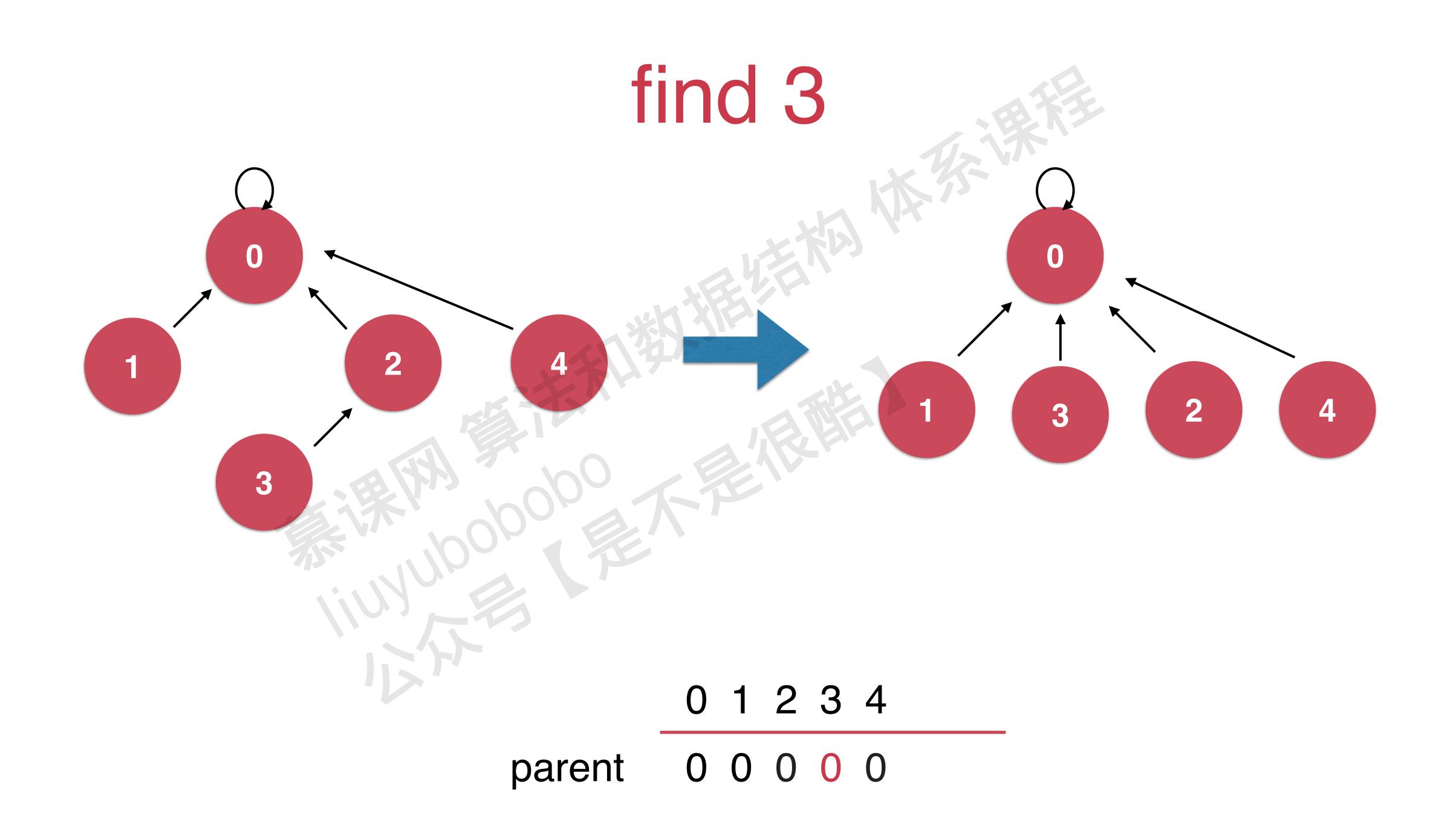
实践:递归的路径压缩











并查集的时间复杂度分析

O(h)

严格意义上: O(log*n) iterated logarithm

并查集的时间复杂度分析

严格意义上: O(log*n) iterated logarithm

$$\log^* n = \begin{cases} 0 & if (n \le 1) \\ 1 + \log^* (\log n) & if (n > 1) \end{cases}$$

近乎是O(1)级别的

更多Leetcode上Union Find相关的问题

并查集 Union Find

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