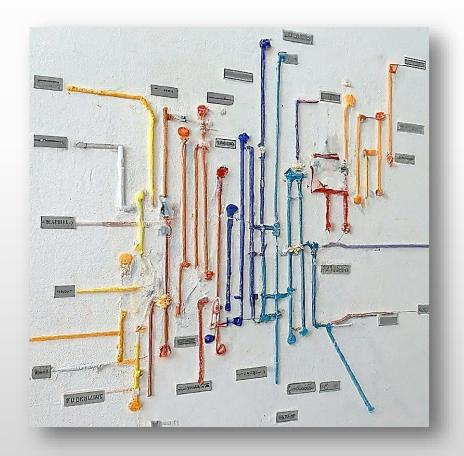


# Bölüm 4: Çizge Algoritmaları Algoritmaları



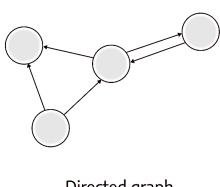


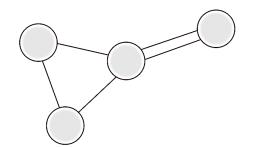
- Dünya aslında bir ağ gibidir.
  - Şehirler yollarla,
  - İnsanlar ilişkilerle,
  - Bilgisayarlar kablolarla birbirine bağlıdır.
- Çizge algoritmaları bu ağları inceler ve anlamlandırır.

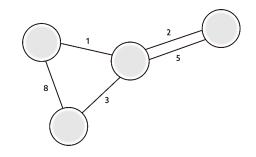


## Çizge Türleri





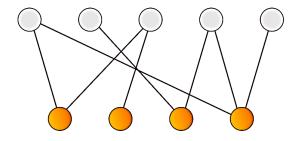


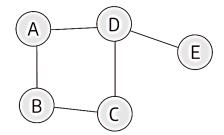


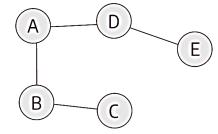
Directed graph

Undirected

Weighted







Bipartite graph

Cyclic graph

Acyclic graph





- Birbirine bağlı noktalar (düğüm) ve bu noktaları birleştiren çizgiler (kenar) ile temsil edilen ağ yapılarını inceler.
- Ağlarda en kısa yolu hesaplama, gruplama gibi işlemleri gerçekleştirir.
- Sosyal ağlar, harita uygulamaları, navigasyon gibi birçok alanda kullanılır.





- Farklı çizge algoritmaları, farklı işlemler için kullanılır.
- Derinlik Öncelikli Arama (DFS):
  - Bir düğümden başlar, dallanarak tüm ağı gezer.
- Genişlik Öncelikli Arama (BFS):
  - Bir düğümden başlar, katman katman tüm ağı gezer.
- Dijkstra Algoritması:
  - Başlangıç düğümünden diğer düğümlere en kısa yolları bulur.
- Kruskal Algoritması:
  - Bir ağı minimum maliyetle birbirine bağlayan kenarları seçer.





- DFS bir labirentten çıkış yolu ararken kullanılabilir.
- BFS bir haberin tüm şehire yayılma sürecini modelleyebilir.
- Dijkstra en kısa sürede teslimat yapmak için kullanılabilir.







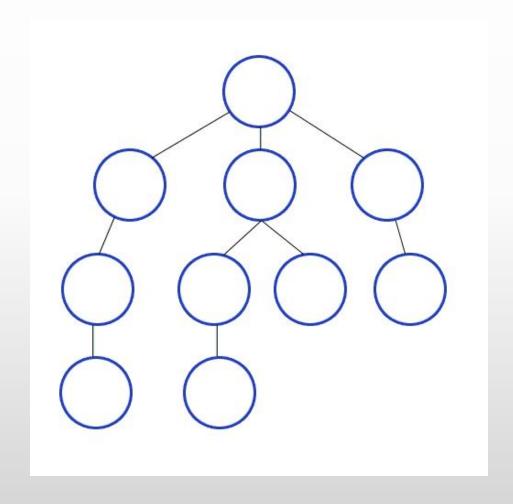
- Çizge gezinme algoritmaları (Graph traversal)
- En kısa yol algoritmaları (Shortest path)
- Minimum yayılan ağaç algoritmaları (Minimum spanning tree)
- Ağ akış algoritmaları (Network flow)



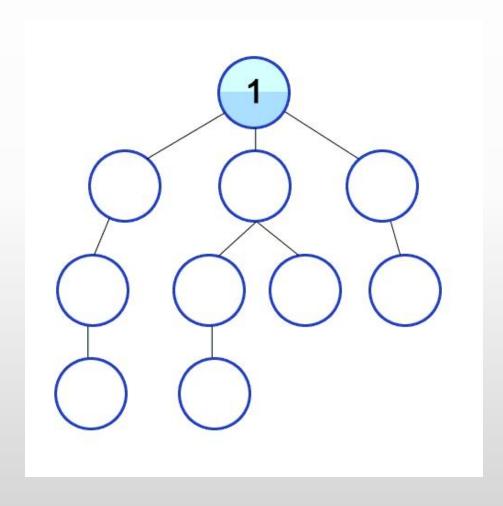


- Çizgenin yapısını sistematik bir şekilde keşfetmek için kullanılır.
- İki ana kategoriye ayrılır:
  - derinlik öncelikli arama (DFS) ve
  - genişlik öncelikli arama (BFS).
- DFS, bir düğümden başlayarak mümkün olduğunca derinlere iner ve tüm komşularını ziyaret ettikten sonra geri döner.
- BFS, bir kuyruk veri yapısı kullanarak seviye seviye tüm düğümleri ziyaret eder.

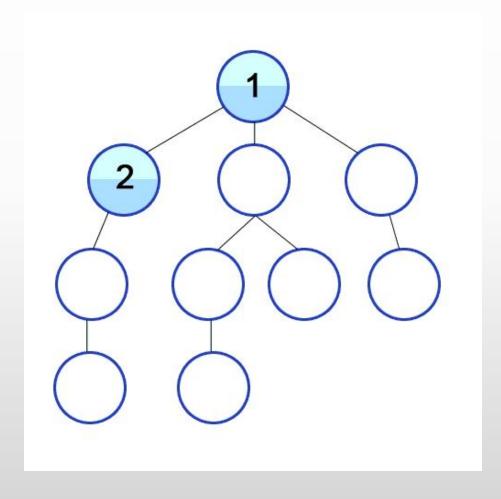




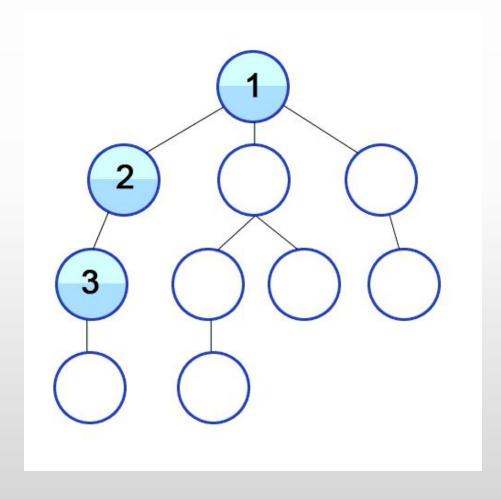




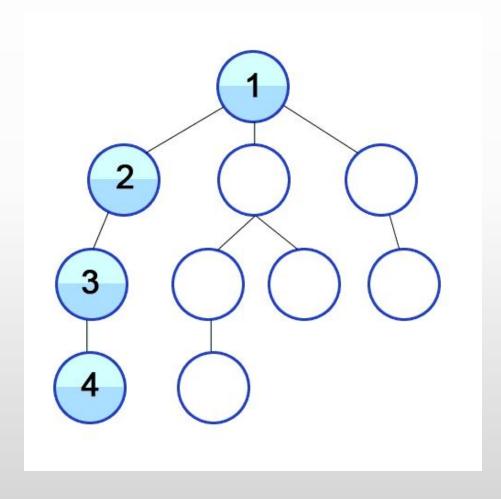




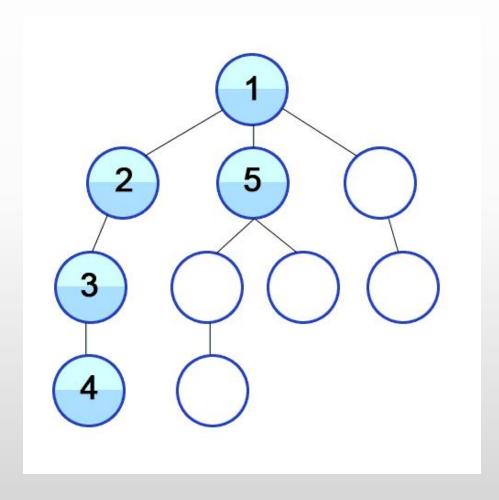




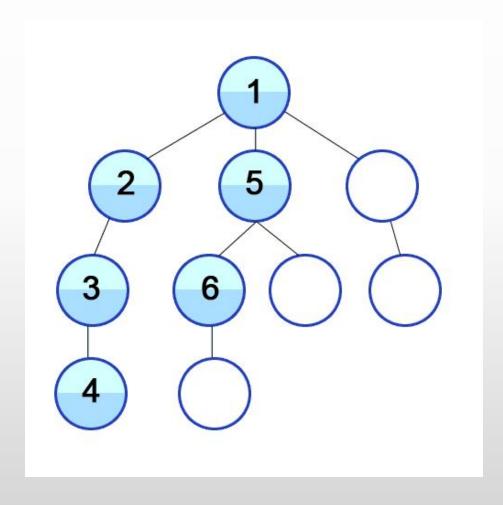




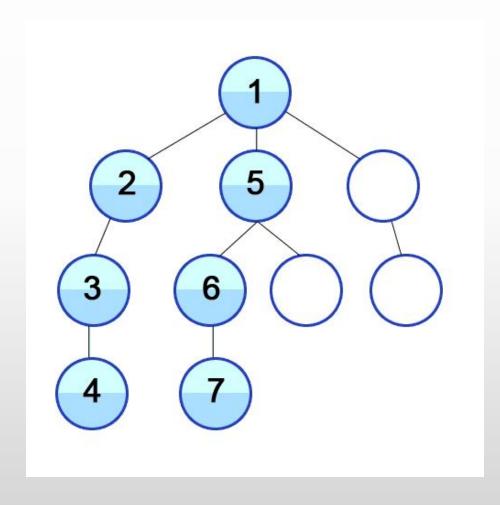




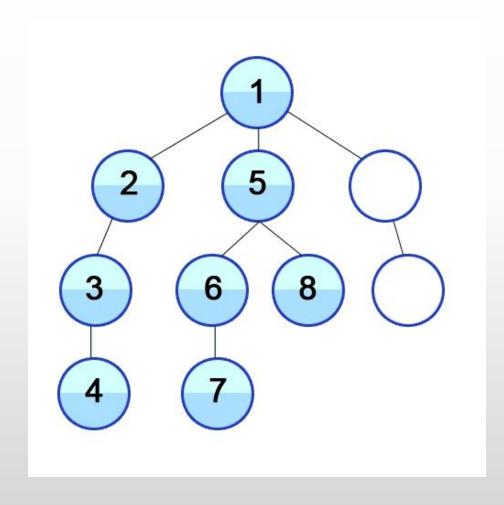




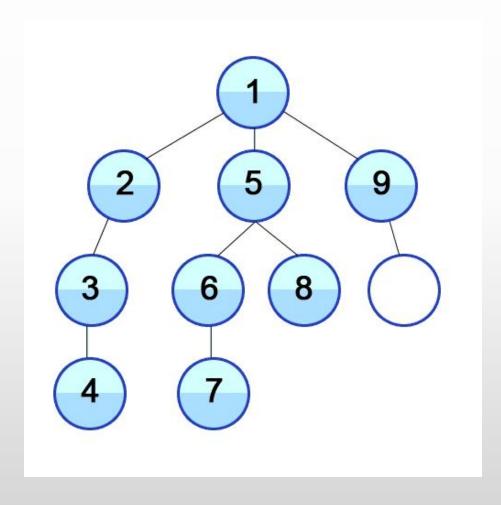




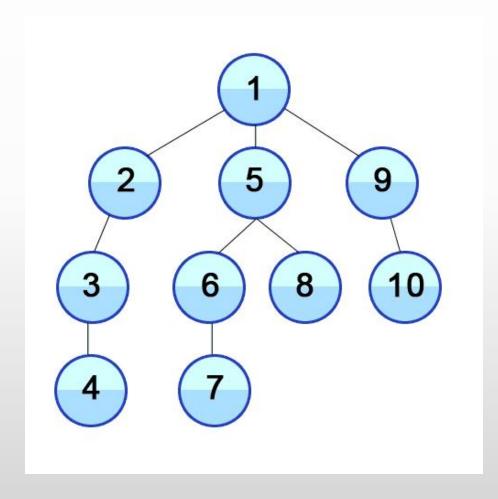
















procedure DFS(G, v) is
label v as discovered
for all directed edges from v to w that are in G.adjacentEdges(v) do
if vertex w is not labeled as discovered then
recursively call DFS(G, w)

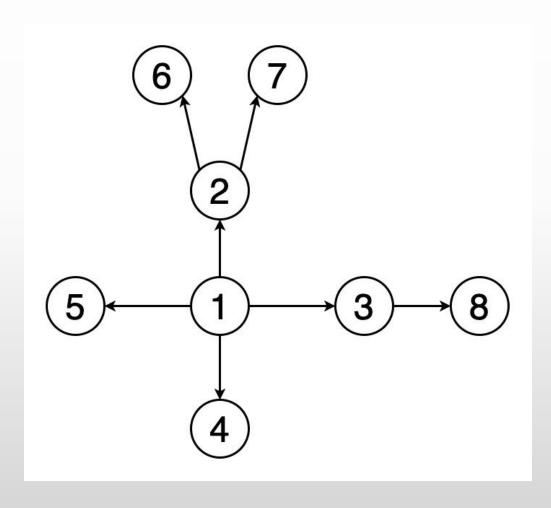




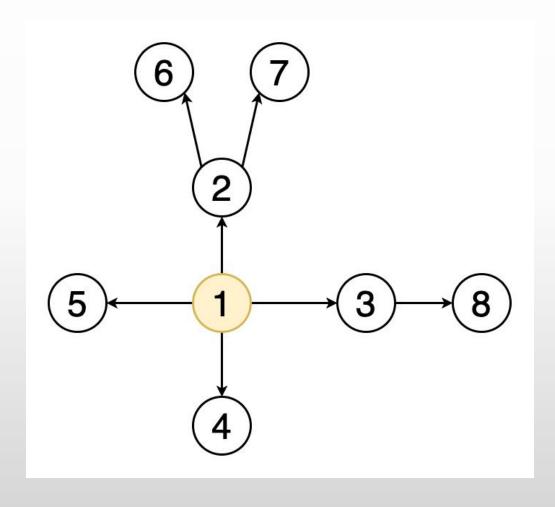
```
procedure DFS_iterative(G, v) is
  let S be a stack
  S.push(v)
  while S is not empty do
     v = S.pop()
     if v is not labeled as discovered then
       label v as discovered
       for all edges from v to w in G.adjacentEdges(v) do
          S.push(w)
```



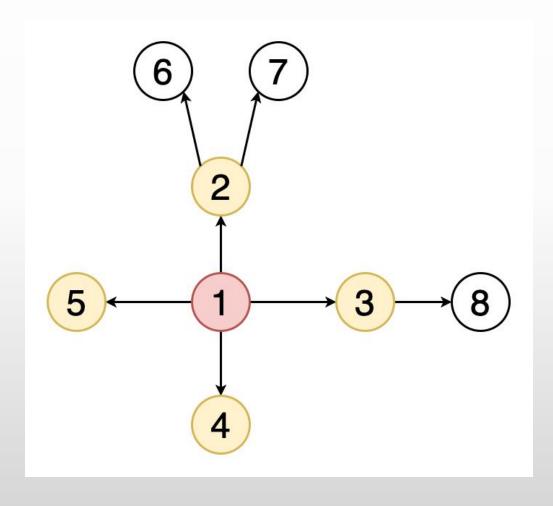




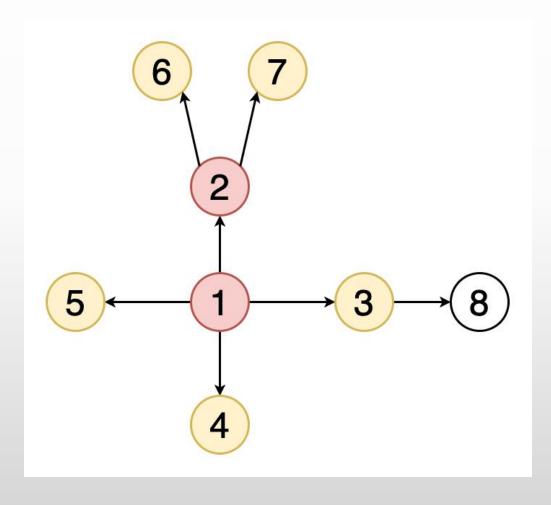




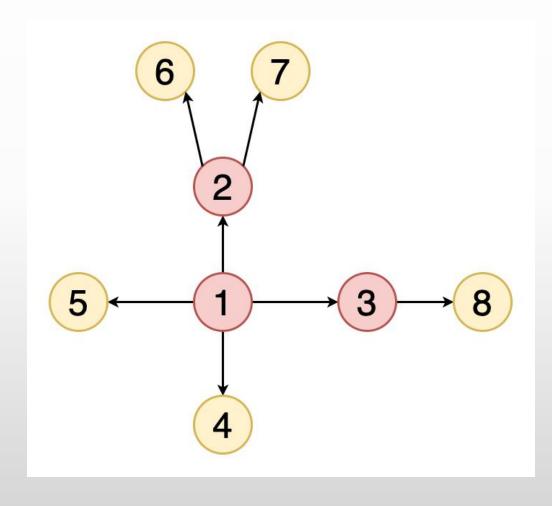






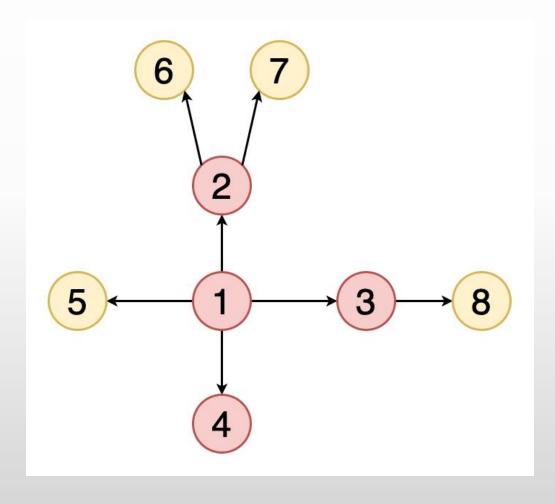








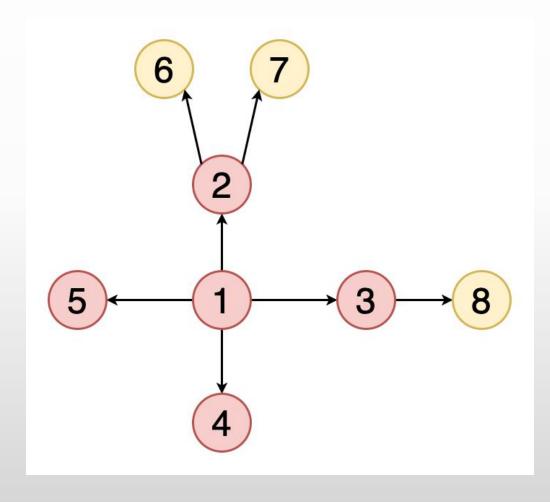




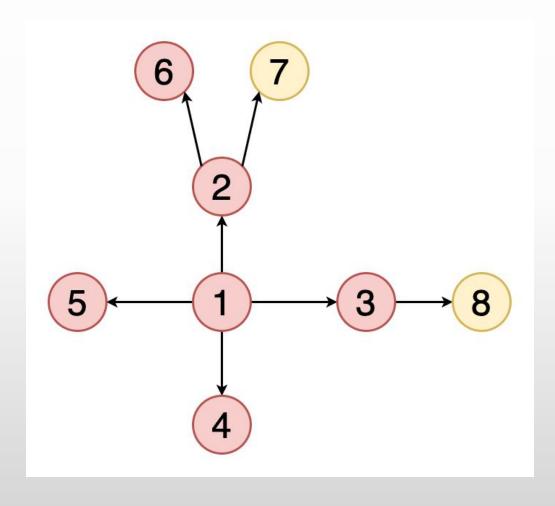




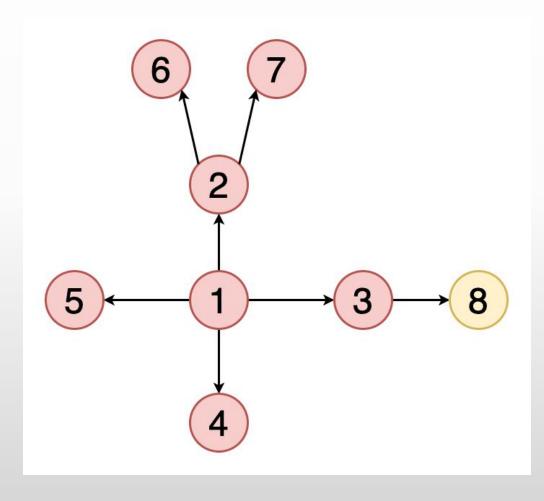
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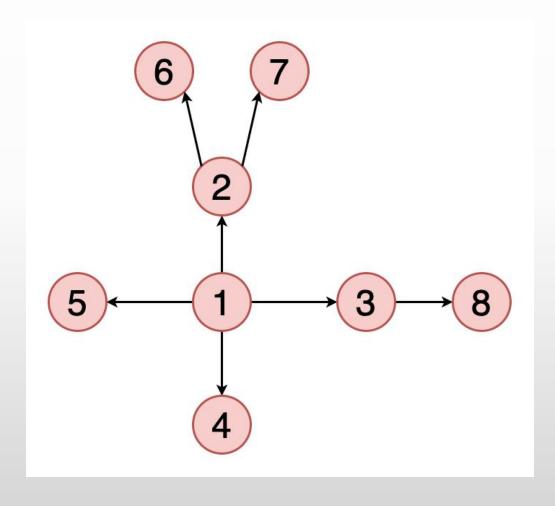












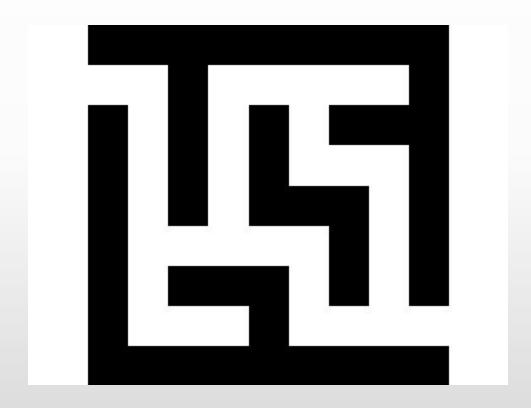




```
procedure BFS(G, root) is
   let Q be a queue
   label root as explored
   Q.enqueue(root)
   while Q is not empty do
      v := Q.dequeue()
      if v is the goal then
        return v
      for all edges from v to w in G.adjacentEdges(v) do
        if w is not labeled as explored then
           label w as explored
           w.parent := v
           Q.enqueue(w)
```

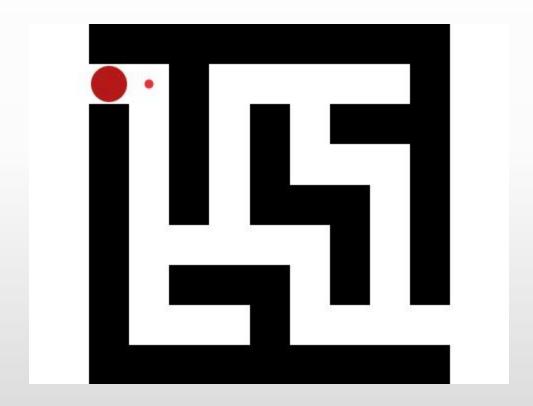




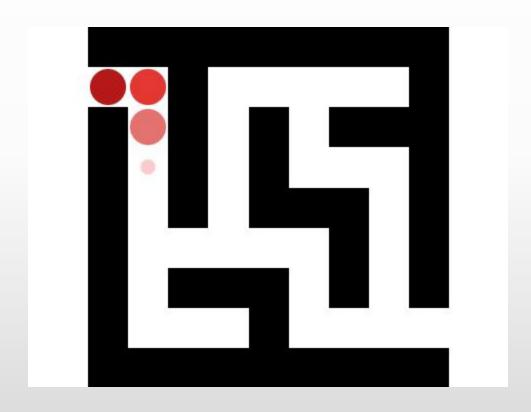




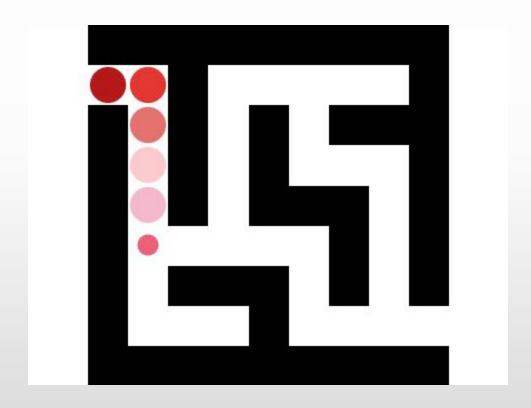




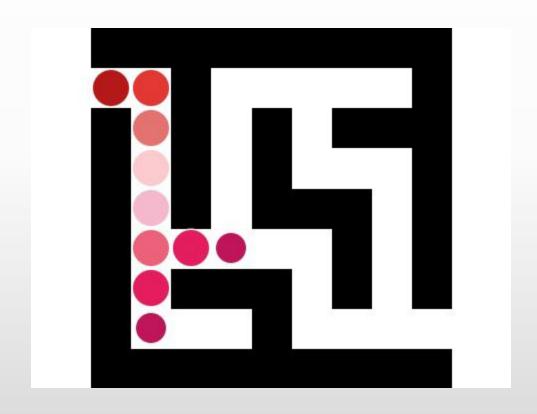




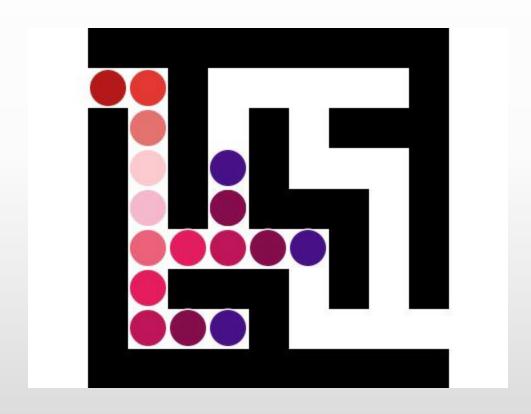






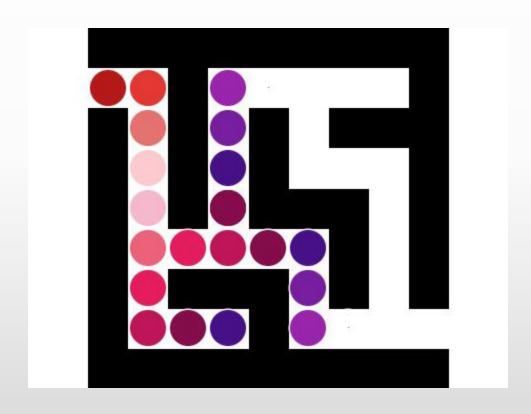




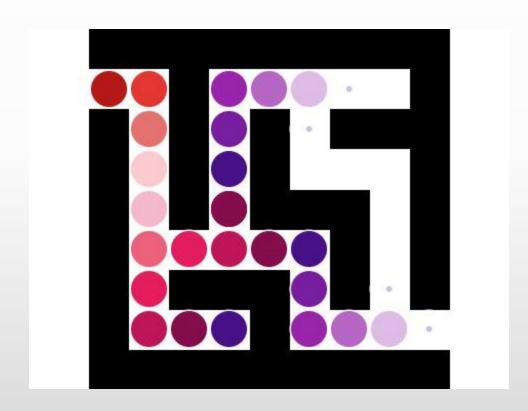














#### SON