



# Data Structures and Algorithms Course

## Searching Algorithms

---



# Today's Content

- Searching Algorithms
- Algorithm Questions

# Searching Algorithms - Agenda

- Linear (Sequential) Search
- Binary Search
- Ternary Search
- Jump Search
- Exponential Search
- Algorithm questions with search

# Linear Search (Sequential)

- Unordered or order is not known
- Scan complete array and see if element is there.

For example: I'm searching value "3"



	Best	Worst
Time Complexity	$O(1)$	$O(n)$

# Linear Search (Sequential)

```
public int linearSearch (int[] Array, int data){  
    for (int i=0;i<Array.length;i++) if(Array[i]==data) return i;  
    return -1;  
}
```

**Lets switch to IntelliJ for implementation of  
Linear Search**



# Binary Search

**Binary Search** is a searching algorithm used in a sorted array by **repeatedly dividing the search interval in half**.

The idea of binary search is to use the information that the array is sorted and reduce the time complexity to  $O(\log n)$ .

## **Algorithm:**

- Calculate middle index.
- Compare  $x$  with the middle element.
- If  $x$  matches with the middle element, we return the mid index. ( $==$ )
- Else If  $x$  is greater than the mid element, then we recur for the right half. ( $>$ )
- Else ( $x$  is smaller) recur for the left half. ( $<$ )

- **Caution:** Works on sorted lists!



# Binary Search

3	5	6	9	11	18	20	25	30	45
0	1	2	3	4	5	6	7	8	9

1 Calculate middle index  $\text{middle} = (\text{left} + \text{right}) / 2$   
 $= (0 + 9) / 2 = 4$

2 Compare value searched with the middle value  
if ( $\text{data} < \text{array}[\text{middle}]$ ) branch into left partition

3 Calculate new middle index  
 $\text{middle} = (\text{left} + \text{right}) / 2$   
 $= (0 + 3) / 2 = 1$

4 Compare value searched with the middle value  
if ( $\text{data} > \text{array}[\text{middle}]$ ) branch into right partition

5 Calculate new middle index  
 $\text{middle} = (\text{left} + \text{right}) / 2$   
 $= (2 + 3) / 2 = 2$

6 if ( $\text{data} == \text{array}[\text{middle}]$ ) return middle index; // return 2 – index of '6'

Assume we are searching for "6"

3	5	6	9	11	18	20	25	30	45
0	1	2	3	4	5	6	7	8	9

3	5	6	9	11	18	20	25	30	45
0	1	2	3	4	5	6	7	8	9

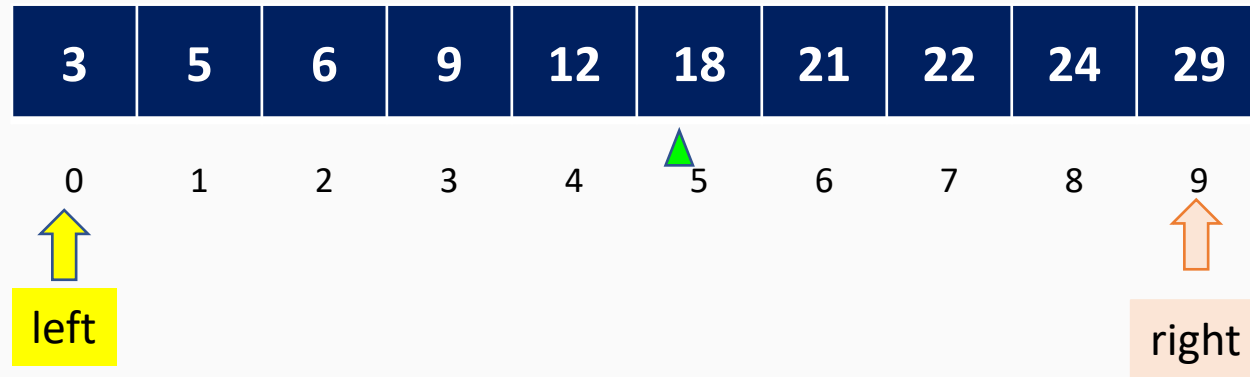
3	5	6	9
0	1	2	3

3	5	6	9
0	1	2	3

6	9
2	3



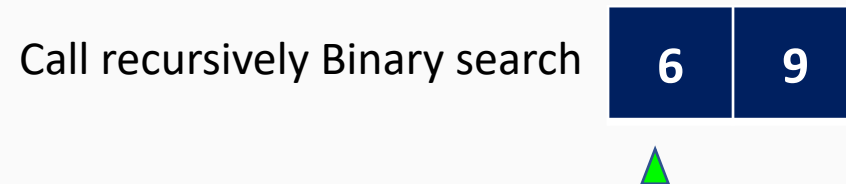
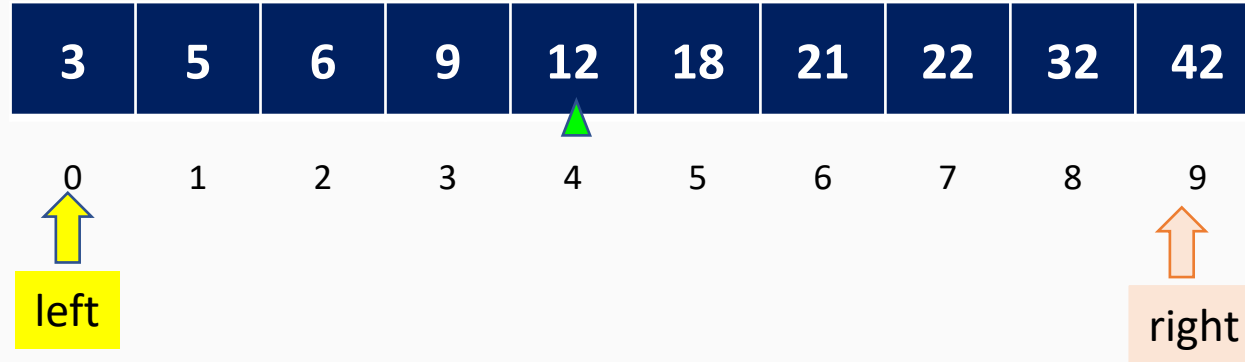
# Binary Search (Iterative)



**while (left <= right)**

- **if (array[middle] == data) return middle;**
- **if (data < array[middle]) right = middle - 1;**  
  **else left = middle + 1;**

# Binary Search (Recursive)



# Binary Search - Performance

**Time Complexity** :  $O(\log n)$

**Recursive Space Complexity** :  $O(\log n)$

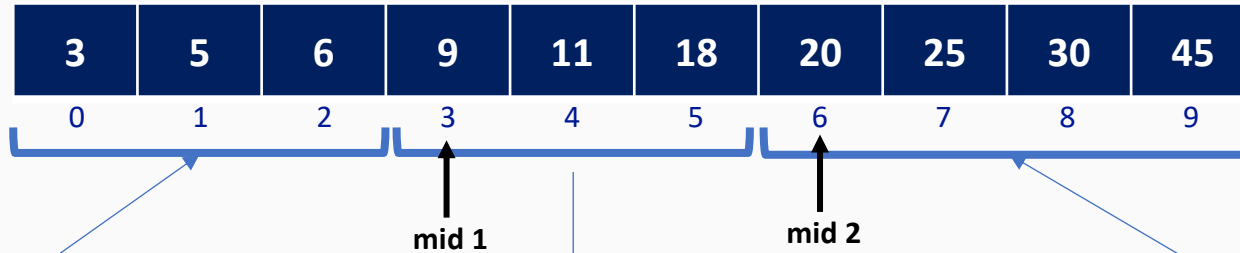
**Iterative Space Complexity** :  $O(1)$

**Lets switch to IntelliJ for implementation of  
Binary Search**



# Ternary Search

- Similar to Binary search, instead of two parts Ternary search divides into 3 parts.



$\text{partitionSize} = (\text{right} - \text{left}) / 3$

$\text{mid1} = \text{left} + \text{partitionSize}$

$\text{mid2} = \text{right} - \text{partitionSize}$

if (data < array[mid1])

ternarySearch(left, mid1 - 1)

ternarySearch(mid1+1, mid2-1)

if (data > array[mid2])

ternarySearch(mid2+1, right)

# Ternary Search-Performance

- Similar to Binary search, instead of two parts Ternary search divides into 3 parts.

## Performance Comparison

### BINARY SEARCH

$$\log_2 n$$

### TERNARY SEARCH

$$\log_3 n$$

Does not mean Ternary is faster.  
Makes too much comparison.

# Ternary Search-Implementation

- Similar to Binary Search Recursive implementation.
  1. Compare the key with the element at mid1. If equal, return mid1.
  2. If not, compare the key with mid2. If equal, return mid2.
  3. If not, then if the key is less than the element at mid1. Branch into the first part.
  4. If not, then if the key is greater than the element at mid2. Branch into the third part.
  5. If not, branch into the second (middle) part.

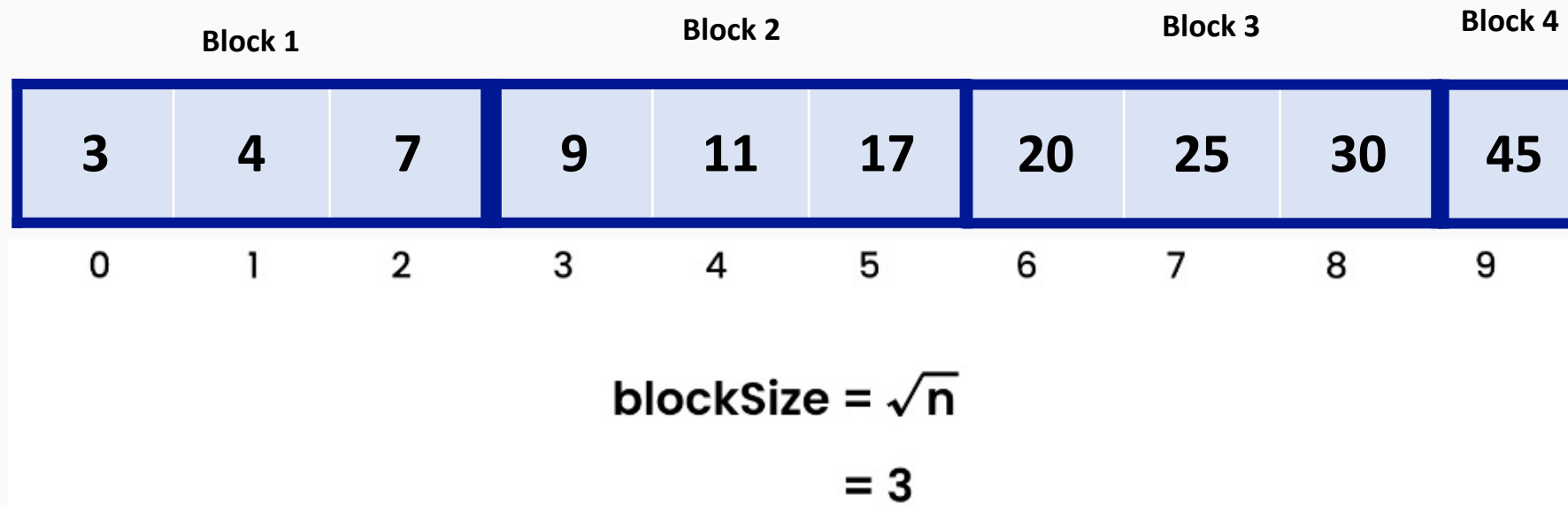
**Lets switch to IntelliJ for implementation of Ternary Search**





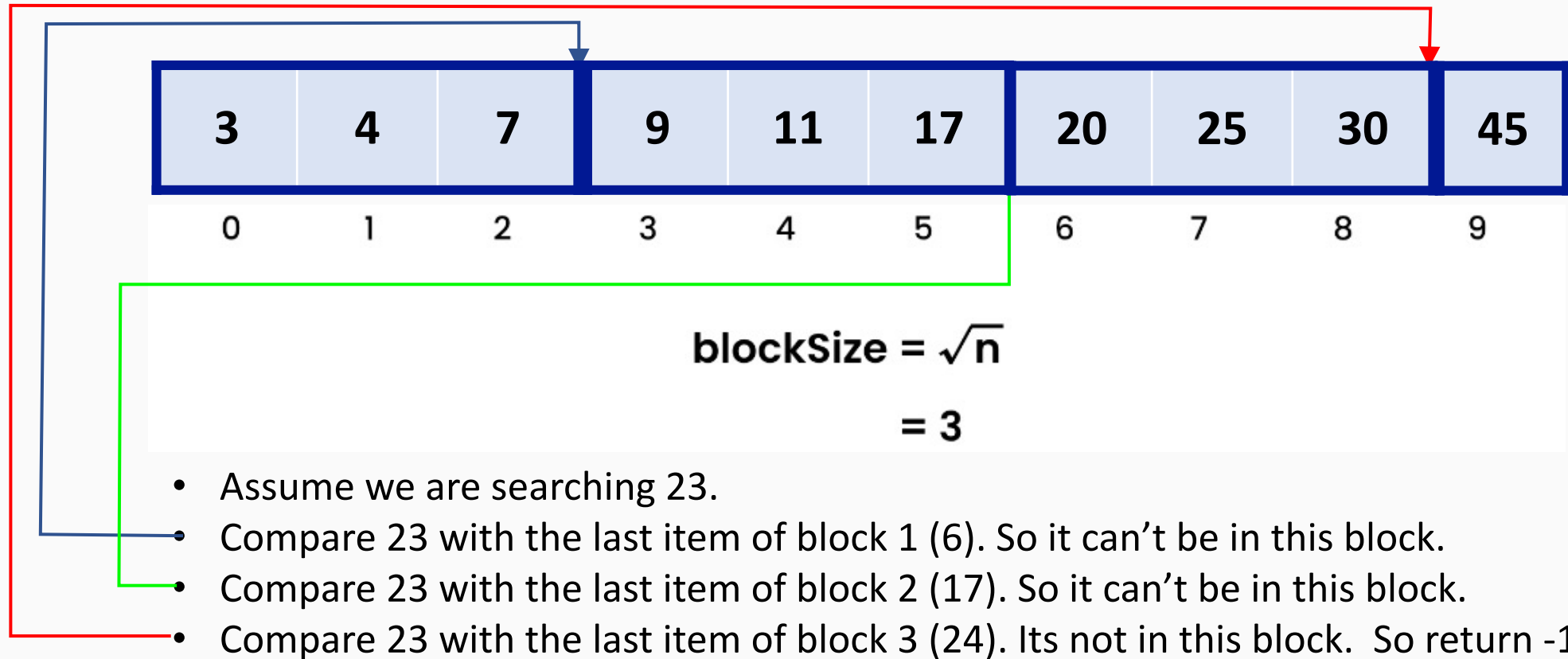
# Searching Algorithms – Jump Search

- Improvement to linear search but not as fast as Binary Search.
- For sorted arrays.
- The basic idea is to check fewer elements (than linear search) by jumping ahead by fixed steps or skipping some elements in place of searching all elements.



# Searching Algorithms – Jump Search

- Improvement to linear search but not as fast as Binary Search.



# Jump Search-Performance

**Time Complexity :  $O(\sqrt{n})$**

**Space Complexity :  $O(1)$**

# Jump Search-Implementation

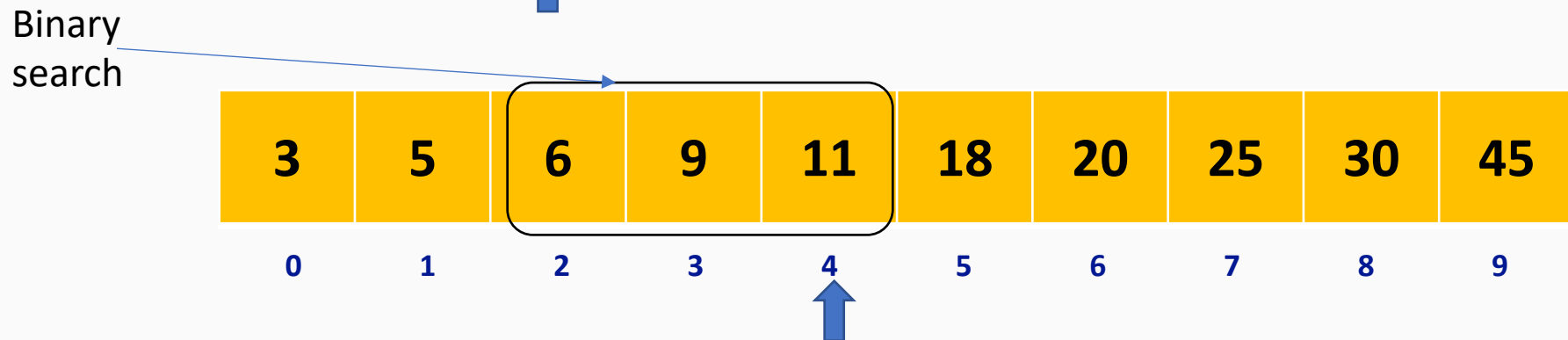
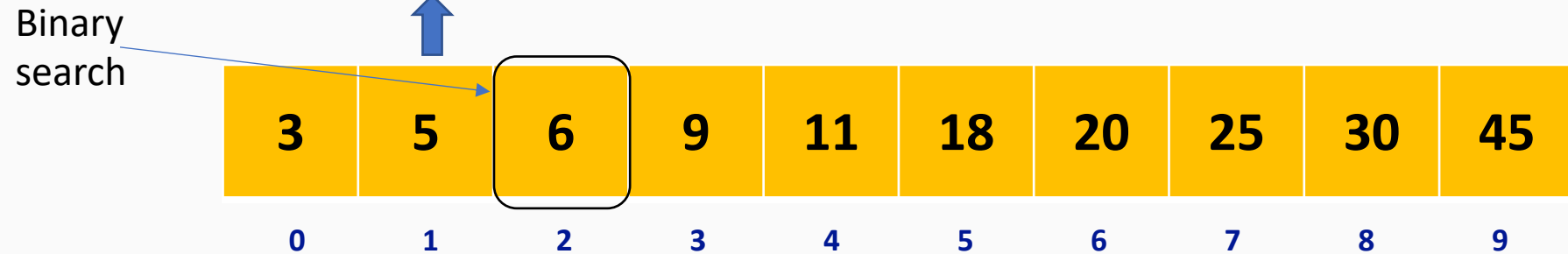
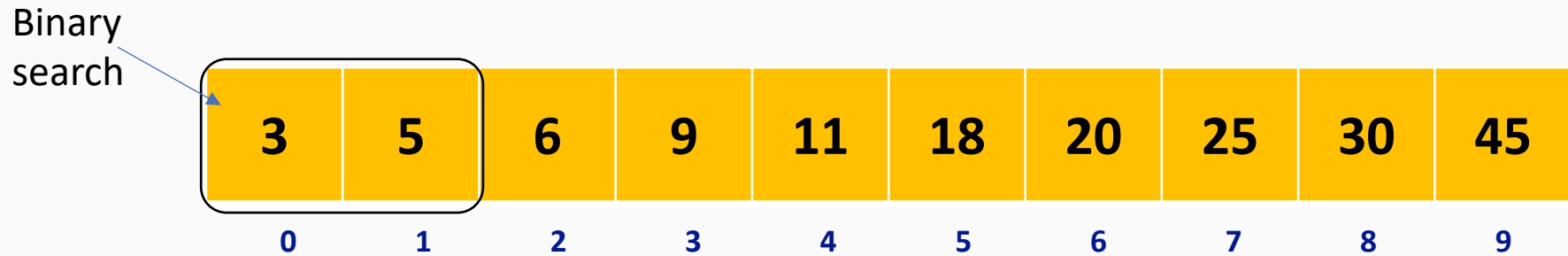
- **Calculate Block size** : `blockSize = sqrt(array.length);`
- **Set start point**: 0 for the first jump then `+=blockSize`
- **Jump Blocks**: `while (start < array.length && array[nextBlockLast - 1] < data)`  
`start=next;`  
`next=next+blockSize;`
- **If Block found** : Perform a **linear search** in that block.

**Lets switch to IntelliJ for implementation of  
Jump Search**



# Exponential Search

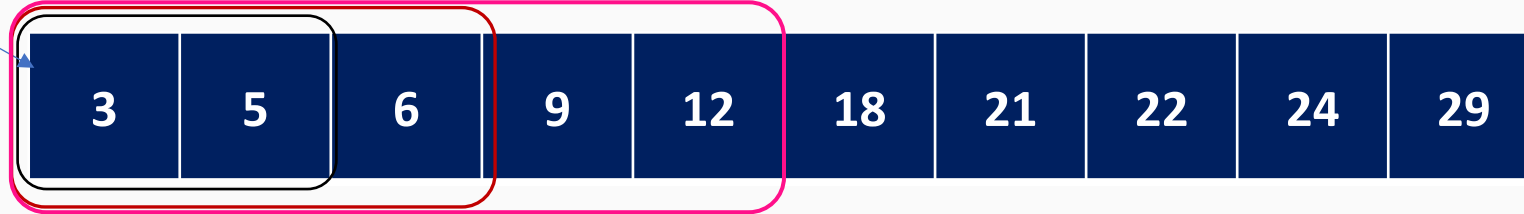
- Start with a small range, if item does not exist double the search size.



# Exponential Search-Performance

## Exponential Search

Binary  
search




Time Complexity :  $O(\log n)$

# Exponential Search-Implementation



Start here       Bound=1

- **while (bound < array.length && array[bound] < data) bound \*= 2;**
- **int left = bound / 2;**
- **int right = Math.min(bound, array.length - 1);**
- **return binarySearchRec(array, data, left, right)**  Call binary search here



**Lets switch to IntelliJ for implementation of  
Exponential Search**



# So what?

Why did we learn search algorithms, do we really need them?

# Algorithm Question-1

Given an array of integers `nums` which is sorted in ascending order, and an integer `target`, write a function to search `target` in `nums`. If `target` exists, then return its index. Otherwise, return -1.

You must write an algorithm with  $O(\log n)$  runtime complexity.

## Example 1:

**Input:** `nums = [-1,0,3,5,9,12]`, `target = 9`

**Output:** 4

**Explanation:** 9 exists in `nums` and its index is 4

## Example 2:

**Input:** `nums = [-1,0,3,5,9,12]`, `target = 2`

**Output:** -1

**Explanation:** 2 does not exist in `nums` so return -1

# Algorithm Question- Template

```
class Solution {  
    public int search(int[] nums, int target) {  
    }  
}
```

# Algorithm Question-2

You are given an  $m \times n$  integer matrix with the following two properties:

- Each row is sorted in non-decreasing order.
- The first integer of each row is greater than the last integer of the previous row.

Given an integer target, return true *if target is in matrix* or false *otherwise*.

1	3	5	7
10	11	16	20
23	30	34	60

**Input:** matrix = [[1,3,5,7],[10,11,16,20],[23,30,34,60]], target = 3

**Output:** true

**Input:** matrix = [[1,3,5,7],[10,11,16,20],[23,30,34,60]], target = 13

**Output:** false

# Algorithm Question-2 Search a 2D Matrix

**Lets switch to IntelliJ for Solution**



# Algorithm Question-2

You are given an  $m \times n$  integer matrix with the following two properties:

- Each row is sorted in non-decreasing order.
- The first integer of each row is greater than the last integer of the previous row.
- Given an integer target, return true *if target is in matrix* or false *otherwise*.
- You must write a solution in  $O(\log(m * n))$  time complexity.

1	3	5	7
10	11	16	20
23	30	34	60

**Input:** matrix = [[1,3,5,7],[10,11,16,20],[23,30,34,60]], target = 3

**Output:** true

**Input:** matrix = [[1,3,5,7],[10,11,16,20],[23,30,34,60]], target = 13

**Output:** false

# How find row/col index?

- `[[1, 3, 5, 7] , [10, 11, 16, 20], [23, 30, 34, 60]]`

	Col 0	Col 1	Col 2	Col 3
Row 0				
Row 1				
Row 2				