数据结构与算法
DATA STRUCTURE

第二十八讲 复习

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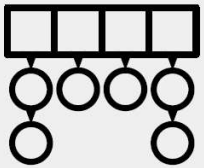
图的表示

For G = (V, E)	$\begin{bmatrix} 0 & 1 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 1 & 0 & 1 & 0 \end{bmatrix}$	
Edge Membership Is e = {u, v} in E?	$O(1)$	$O(\deg(u))$ **or** $O(\deg(v))$
Neighbor Query What are the neighbors of u?	$O(\lvert V \rvert)$	$O(\deg(v))$
Space requirements	$O(\lvert V \rvert^2)$	$O(\lvert V \rvert + \lvert E \rvert)$

图遍历算法

- 深度优先
  - 深度优先+始末时间
  - 拓扑排序
  - 检测图上的cycle
  - 强连通分支 （作业）
- 广度优先
  - 按边算的最短路径
  - 检测二分图
  - 无向图上的连通分支

最短路径算法和动态规划算法

	Dijkstra	Bellman-Ford	Floyd-Warshall	DAG上的 最短路径
Problem	Single source shortest path	Single source shortest path	All pairs shortest path	Single source
Runtime	$O(\|E\|+\|V\|\log(\|V\|))$ **worst-case** with a fibonacci heap	$O(\|V\|\|E\|)$ **worst-case**	$O(\|V\|^3)$ **worst case**	$O(\|V\|+\|E\|)$

最小生成树MST和贪婪算法

	Description	Runtime
Prim's	Grows a tree	$O(\|E\|\log(\|V\|))$ **with red-black tree** $O(\|E\|+\|V\|\log(\|V\|))$ **with Fibonacci heap**
Kruskal's	Grows a forest	$O(\|E\|\log(\|V\|))$ **with union-find** $O(\|E\|)$ **with union-find and radix sort**

FindScc

- 每个节点保存了出边集合和入边集合
- 每添加一条边，需要在两个顶点结构体里加边

```
struct Node
{
    int id;
    int key;
    // outgoing edges
    std::vector<Edge*> arcs;
    // incoming edges
    std::vector<Edge*> inArcs;

    bool active;

    Node(int value = 0)
    {
        id = getId();
        key = value;
        active = true;
    }
    int getId()
    {
        static int id = 0;
        return id++;
    }
};
```

```
void MyGraph::AddEdge(int uid, int vid)
{
    int total = _nodes.size();
    assert(uid < total && vid < total);

    Node *U = _nodes.at(uid);
    Node *V = _nodes.at(vid);

    Edge *arc = new Edge(U, V);
    U->arcs.push_back(arc);
    V->inArcs.push_back(arc);

    _edges.push_back(arc);
    assert(arc->id + 1 == (int)_edges.size());
}
```

FindScc

- 辅助函数，找下一个没访问过的节点
- 如果有按结束时间排序的traveltime，需要找结束时间最晚的第一个没访问过的节点

```
int MyGraph::GetUnvisitNode(bool *visited, pair<int, int> *sortedTravelTimes)
{
    for (size_t i = 0; i < _nodes.size(); i++)
    {
        // If travel time is not specified, check visited array.
        int id = i;
        // Otherwise, id is from travel time array.
        if (sortedTravelTimes)
        {
            id = sortedTravelTimes[i].first;
        }
        if (!visited[id])
        {
            return id;
        }
    }
    return -1;
}
```

FindScc

- 按反向边的DFS递归函数
- 因为我们在节点里加了入边，所以直接用入边进行DFS
- 但注意邻接顶点是arc里的U

```
void MyGraph::DfsReverseInternal(int uid, bool *visited, vector<int> *pScc)
{
    visited[uid] = true;
    cout << "#" << uid << ", ";
    pScc->push_back(uid);

    for (auto arc : _nodes.at(uid)->inArcs)
    {
        // Note U in arc is actually neighbor of node with uid.
        int vid = arc->U->id;
        if (!visited[vid])
        {
            DfsReverseInternal(vid, visited, pScc);
        }
    }
}
```

FindScc

```
bool PairCompare(const pair<int, int> &a, const pair<int, int> &b)
{
    return (a.second > b.second);
}

vector<vector<int>> MyGraph::FindScc()
{
    vector<vector<int>> sccs;

    cout << "\n Start DFS Time search:" << endl;
    bool *visited = new bool[_nodes.size()]();
    pair<int, int> *travelTimes = new pair<int, int>[_nodes.size()];

    // Keep dfs until all nodes are visited
    int id = -1;
    int currentTime = 0;
    while ((id = GetUnvisitNode(visited, nullptr)) >= 0)
    {
        currentTime = DfsTimeInteral(id, currentTime, visited, travelTimes);
    }

    // Reset start time in travel times to node id
    for (size_t id = 0; id < _nodes.size(); id++)
    {
        travelTimes[id].first = id;
    }

    // Sorted travel time array in order of decreasing end time
    // Note start time of node has been reset to its id
    cout << endl;
    sort(travelTimes, travelTimes + _nodes.size(), PairCompare);
    for (size_t i = 0; i < _nodes.size(); i++)
    {
        cout << "(" << travelTimes[i].first << ", " << travelTimes[i].second << ") ";
    }

    for (size_t i = 0; i < _nodes.size(); i++)
    {
        visited[i] = false;
    }

    // Starting from the first unvisited node with largest end time,
    // search for scc one by one
    cout << "Find scc" << endl;
    id = -1;
    while ((id = GetUnvisitNode(visited, travelTimes)) >= 0)
    {
        vector<int> scc;
        DfsReverseInternal(id, visited, &scc);
        sccs.push_back(scc);
        cout << endl;
    }

    delete [] visited;
    delete [] travelTimes;

    return sccs;
}
```

Dijkstra

```
bool MyGraph::Dijkstra(int s, int t)
{
    MinHeap pq; // min heap as priority queue
    int sz = (int)_nodes.size();

    // flag indicating whether a node has been visited.
    bool *visited = new bool[sz]();
    // If a shorter distance is found, parent link should be updated
    int *parent = new int[sz]();

    // Estimate distance from source node 's'.
    int *dist = new int[sz]();
    for (int id = 0; id < sz; id++)
    {
        // initialize to max, except for source node itself
        dist[id] = id == s ? 0 : INT_MAX;
        parent[id] = -1;
        // Maintain a priority queue for unvisited nodes
        pq.insertKey(id, dist[id]);
    }

    bool success = false;
    for (int i = 0; i < sz && !success; i++)
    {
        // It is difficult to implement decreaseKey, as we should
        // find out which node in heap corresponding to adjacent node.
        // Trick here: Always insert node with updated dist[].
        // Consequently, we should mark a node visited and remove min node if visited.
        while (visited[pq.getMin().id])
        {
            pq.extractMin();
        }

        // Now the min node is unvisited node.
        HeapNode hn = pq.getMin();
        pq.extractMin();

        // This indicated other unvisited nodes are not reachable.
        if (hn.priority == INT_MAX)
        {
            break;
        }
```

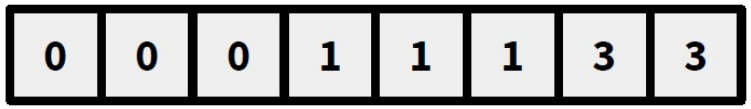
Dijkstra

```
            visited[hn.id] = true;
            for (auto arc : _nodes.at(hn.id)->arcs)
            {
                int vid = arc->V->id;
                if (visited[vid])
                {
                    continue;
                }

                // For each unvisited neighbor, update estimate dist[].
                if (dist[vid] > dist[hn.id] + arc->weight)
                {
                    // update estimate distance, parent link
                    // Trick: Insert updated value into heap, instead of decreasing key.
                    dist[vid] = dist[hn.id] + arc->weight;
                    parent[vid] = hn.id;
                    pq.insertKey(vid, dist[vid]);
                }
                if (vid == t)
                {
                    success = true;
                    break;
                }
            }
        }


        if (success)
        {
            cout << "Find path from #" << s << " to #" << t << endl;
            int curr = t;
            while (parent[curr] != -1)
            {
                cout << "#" << curr << "["<< dist[curr] << "]" << " <- ";
                curr = parent[curr];
            }
            cout << "#" << curr << "["<< dist[curr] << "]" << endl;
        }
        else
        {
            cout << "No path from #" << s << " to #" << t << endl;
        }

        return success;
    }
```

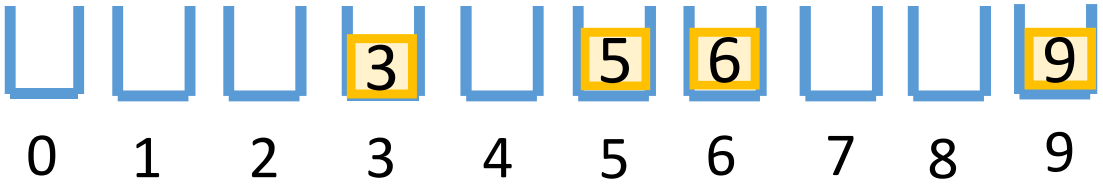
排序算法

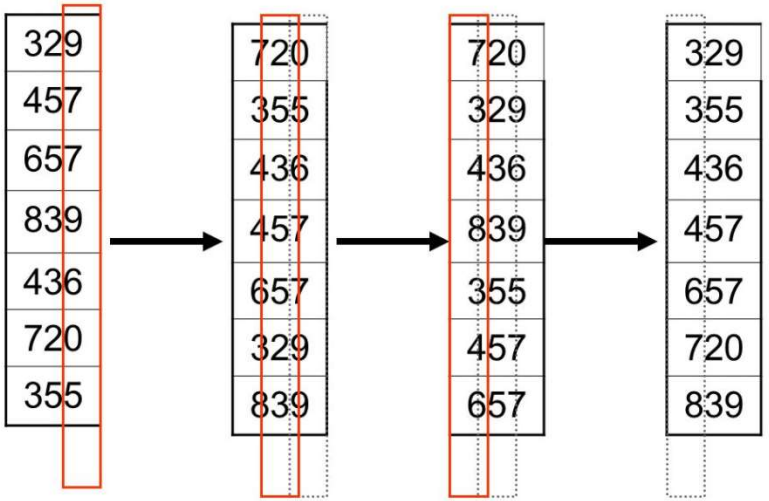
- MergeSort
- QuickSort
- CoutingSort
- RadixSort
- BucketSort

Couting sort



Bucket sort和radix sort





基本数据结构

- 数组和链表
- Stack
- Queue
- Recursion
- Heap
- HeapSort

串匹配算法

- 暴力算法
- KR算法（hashing值）
- KMP（最长前后缀匹配长度，next数组）
- BM （好后缀和坏字符规则）
- 改进的Boyer–Moore–Horspool算法（只有坏字符规则）

树

- 数的遍历
  - 中序，前序，后序
  - 遍历迭代算法
- 二叉树（线索化）
- 二叉查找树
  - AVL
  - 红黑树
- B树和B+树
- Trie
- Hash
- （习题参考Lecture20.examples.pdf）

比较

	Sorted linked lists	Sorted arrays	Balanced BSTs	Hash tables
Search	**O(n)** expected & worst-case	**O(log n)** expected & worst-case	**O(log n)** expected & worst-case O(n) worst-case for generic BSTs	**O(1)** expected **O(n)** worst-case
Insert/ Delete	**O(n)** expected & worst-case without a pointer to the element	**O(n)** expected & worst-case	**O(log n)** expected & worst case	**O(1)** expected **O(n)** worst-case without a pointer to the element

C++11新语句和STL

[C++11] for-each语句

- for语句，用循环变量student作为数组下标操作

```
int main()
{
    const int numStudents = 5;
    int scores[numStudents] = { 84, 92, 76, 81, 56 };
    int maxScore = 0;
    for (int student = 0; student < numStudents; ++student)
    {
        if (scores[student] > maxScore)
        {
            maxScore = scores[student];
        }
    }
    std::cout << "The best score was " << maxScore << '\n';

    return 0;
}
```

[C++11] for-each语句

- for-each语句

```
for (element_declaration : array)
   statement;
```

- 这里fn不是下标，fn是对fibNums[]数组元素的复制

```
int fibNums[] = { 0, 1, 1, 2, 3, 5, 8, 13, 21, 34, 55, 89 };
for (int fn : fibNums)
{
    cout << fn << " ";
}
```

- 可以直接和auto搭配使用，更方便

```
for (auto fn : fibNums)
{
    cout << fn << " ";
}
```

for-each和引用

- 如果加上引用（&），不再传值，可以修改原数组值

```
int array[5] = { 9, 7, 5, 3, 1 };
for (auto &element: array)
{
    std::cout << element << ' ';
    element = 0;
}
```

- 如果即要避免元素值复制，又不要修改原数据，用const reference

```
for (const auto &element: array)
{
    std::cout << element << ' ';
    element = 0; // ERROR: element is readonly
}
```

for-each语句

- 不能用在数组形参上，因为这个形参就是指针，不知道数组长度
- 所以参数形式最好改成指针，避免误解

```
int sumArray(int array[])
{
    int sum = 0;
    for (const auto &number : array) // ERROR!!!
    {
        sum += number;
    }
    return sum;
}
```

Template函数模板

- 如果一组重载函数仅仅是参数的类型不一样，程序的逻辑完全一样，那么这一组重载函数可以写成一个函数模板。

```
int Max(int, int);
char Max(char, char);
double Max(double, double);
```

- 所谓的函数模板就是实现类型的参数化（泛型化），即把函数中某些形式参数的类型定义成参数，称为模板参数

```
T Max(T, T);
```

- 在函数调用时，编译器根据实际参数的类型确定模板参数的值，生成不同的模板函数

函数模板的定义

- 一般的定义形式

  template<模板形式参数表>
  返回类型 FunctionName(形式参数表)
  {
  　　//函数定义体
  }

- 模板形式参数表可以包含基本数据类型，也可以包含类类型（需加前缀class或者typename）

```
template <class T>
T Max(T x, T y)
{
    T ret = (x > y) ? x : y;
    return ret;
}

int i = Max<int>(1, 5);
char c = Max<char>('c', 'p');
```

```
int i = Max(1, 5);
cout << i << endl;
double d = Max(1.5, 12.6);
cout << d << endl;
char c = Max('c', 'p');
cout << c << endl;
```

```
5
12.6
p
```

Array in C++11

- 固定长度（静态）数组。#include <array>

```
#include <array>

std::array<int, 3> myarray;
```

- 初始化

```
array<int, 5> myarray = { 9, 7, 5, 3, 1 };
array<int, 5> myarray2 { 9, 7, 5, 3, 1 };

array<int, > myarray = { 9, 7, 5, 3, 1 }; // illegal, array length must be provided
```

- 赋值

```
std::array<int, 5> myarray;
myarray = { 0, 1, 2, 3, 4 }; // okay
myarray = { 9, 8, 7 }; // okay, elements 3 and 4 are set to zero!
myarray = { 0, 1, 2, 3, 4, 5 }; // not allowed, too many elements in initializer list!
```

Array in C++11

- 用下标运算符[]取值，不检查越界

```
myarray[2] = 6;
```

- 用at()取值，检查越界

```
std::array<int, 5> myarray { 9, 7, 5, 3, 1 };
myarray.at(1) = 6; // array element 1 valid, sets array element 1 to value 6
myarray.at(9) = 10; // array element 9 is invalid, will throw error
```

- Size()取长度
- Sort()

```
std::sort(myarray.begin(), myarray.end());
```

- 遍历

```
for (const auto &element : myarray)
    std::cout << element << ' ';
```

vector in C++03

- 动态数组。#include <vector>
- 初始化

```
std::vector<int> array;
std::vector<int> array2 = { 9, 7, 5, 3, 1 };
std::vector<int> array3 { 9, 7, 5, 3, 1 }; /
```

- C++11赋值。数组是动态的

```
array = { 0, 1, 2, 3, 4 }; // okay, array length is now 5
array = { 9, 8, 7 }; // okay, array length is now 3
```

vector in C++03

- 用下标运算符[]取值，不检查越界

```
array[6] = 2; // no bounds checking
```

- 用at()取值，检查越界

```
array.at(7) = 3; // does bounds checking
```

- Size()取长度，Resize()改变长度
- 遍历

```
std::vector<int> array { 0, 1, 2 };
array.resize(5); // set size to 5

std::cout << "The length is: " << array.size() << '\n';

for (auto const &element: array)
    std::cout << element << ' ';
```

- 自动清理内存，不用担心泄露

小结

- 新的for-each语句可以使得数组遍历更容易，安全
- STL的静态数组（固定长度）：array<int, 10> stlArr
  - 可以替代C++静态数组：int arr[10]
  - 提供更多的功能，更安全
  - 注意作为函数的参数的话，不再是数组地址，而是传值，所以最好使用const reference，即传array<>的常量引用
- STL的动态数组：vector<int> stlVec
  - 可以替代C++动态数组：int *ptr = new int[length]
  - 提供更多的功能，更安全，不需要考虑内存释放的问题
  - 同样作为函数的参数的话，要注意传值的性能问题。

Q&A

Thanks!